

## Degree of the OR function

**Definition 6.1.** A polynomial  $g \in \mathbb{Z}_m[\mathbf{x}_1, \dots, \mathbf{x}_n]$  represents a Boolean function  $f : \{0, 1\}^n \rightarrow \{0, 1\} \pmod m$  if

$$\forall \mathbf{x} \in \{0, 1\}^n, g(\mathbf{x}) \begin{cases} = 0 \pmod m & \text{if } f(\mathbf{x}) = 0 \\ \neq 0 \pmod m & \text{if } f(\mathbf{x}) = 1 \end{cases}$$

Note that without loss of generality we can take  $g$  to be multilinear, that is,  $\forall i \in [n], \deg(\mathbf{x}_i) \leq 1$  in  $g$ . This is because if  $a \in \{0, 1\}$ , then  $\forall h \geq 1, a^h = a$ .

**Definition 6.2.** Let  $\deg_p(f) = \min\{\deg(g) \mid g \text{ represents } f \pmod p\}$

**Lemma 6.1.** Suppose  $f \in \mathbb{F}_p[\mathbf{x}_1, \dots, \mathbf{x}_n]$ , where  $p$  is prime, is a multilinear non-zero polynomial of degree at most  $d$ , where  $d > 0$ . Then:

$$|\{\mathbf{a} \in \{0, 1\}^n \mid f(\mathbf{a}) = 0\}| \leq 2^n - 2^{n-d},$$

where equality is achieved for  $f = \prod_{i=1}^d \mathbf{x}_i$ .

*Proof.* We can prove the lemma by induction on  $n$ . In the base case,  $n = 1$ , we have  $d = 1$ , so  $f$  is an expression linear in  $\mathbf{x}_1$ , and thus can have at most one zero, so the inequality holds. Suppose we have shown the inequality holds for polynomials of  $n - 1$  variables. We show that it also holds for any polynomial  $f$  with of  $n$  variables. We can express  $f$  in the following way:  $f(\mathbf{x}_1, \dots, \mathbf{x}_n) = \mathbf{x}_1 g(\mathbf{x}_2, \dots, \mathbf{x}_n) + h(\mathbf{x}_2, \dots, \mathbf{x}_n)$ . Notice that the polynomial  $g$  is of degree at most  $d - 1$ . Then, if we let  $K = |\{\mathbf{a} \in \{0, 1\}^{n-1} \mid g(\mathbf{a}) = 0\}|$ , we have:

$$\begin{aligned} |\{\mathbf{a} \in \{0, 1\}^n \mid f(\mathbf{a}) = 0\}| &\leq 2|\{\mathbf{a} \in \{0, 1\}^{n-1} \mid g(\mathbf{a}) = 0\}| + |\{\mathbf{a} \in \{0, 1\}^{n-1} \mid g(\mathbf{a}) \neq 0\}| \\ &= 2K + (2^{n-1} - K) \\ &= K + 2^{n-1} \\ &\leq 2^{n-1} - 2^{n-1-d+1} + 2^{n-1} \\ &= 2^n - 2^{n-d}, \end{aligned}$$

where the first inequality holds because if  $g(\mathbf{x}_2, \dots, \mathbf{x}_n) \neq 0$ , there is exactly one value of  $\mathbf{x}_1$ , for which  $f(\mathbf{x}) = 0$ , but if  $g(\mathbf{x}_2, \dots, \mathbf{x}_n) = 0$ , there might be up to two values of  $\mathbf{x}_1$  such that  $f(\mathbf{x}) = 0$ . The second inequality holds by the inductive step.  $\square$

**Example 6.1.** Let  $f(\mathbf{x}) = OR(\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_n) = \begin{cases} 0 & \text{if } \forall i \in [n], \mathbf{x}_i = 0 \\ 1 & \text{otherwise} \end{cases}$

**Claim 6.1.** If  $p$  is prime, then  $\deg_p(f) \geq \frac{n}{p-1}$ .

*Proof.* Suppose that  $g \in \mathbb{F}_p[\mathbf{x}_1, \dots, \mathbf{x}_n]$  represents  $f$ . Take  $h = 1 - g^{p-1}(\mathbf{x})$ . Then

$$h(\mathbf{x}) = \begin{cases} 1 & \text{if } \forall i \in [n], \mathbf{x}_i = 0 \\ 0 & \text{otherwise} \end{cases}.$$

Let  $\tilde{h}(\mathbf{x})$  be the multilinear polynomial such that  $\forall \mathbf{a} \in \{0, 1\}^n, \tilde{h}(\mathbf{a}) = h(\mathbf{a})$ . We can get  $\tilde{h}$  from  $h$  by replacing all occurrences of  $\mathbf{x}_i^k$  with  $\mathbf{x}_i$  for any  $i \in [n]$  and  $k \geq 1$ . Notice that  $\deg(\tilde{h}) \leq \deg(h) \leq (p-1)\deg(g)$ . Now by Lemma 6.1,  $|\mathbf{a} \in \{0, 1\}^n | \tilde{h}(\mathbf{a}) = 0| \leq 2^n - 2^{n-\deg(\tilde{h})}$ . On the other hand, we know that  $\tilde{h}$  has exactly  $2^n - 1$  zeros, so we get that  $2^n - 1 \leq 2^n - 2^{n-\deg(\tilde{h})}$ , therefore  $\deg(\tilde{h}) \geq n$ . From here and from  $(p-1)\deg(g) \geq \deg(\tilde{h})$  we get that  $\deg(g) \geq \frac{n}{p-1}$ .  $\square$

Note that Claim 6.1 also holds if  $p$  is a power of some prime number [TB98].

Surprisingly, replacing  $\mathbb{F}_p$  with  $\mathbb{Z}_m$ , where  $m$  is composite, allows a representation of OR to have much smaller degree.

**Theorem 6.1** ([BBR92]). There exists a polynomial  $g \in \mathbb{Z}_6[\mathbf{x}_1, \dots, \mathbf{x}_n]$  that represents the OR function mod 6 with  $\deg(g) \leq O(\sqrt{n})$ .

*Proof.* We are going to use the following theorem:

**Theorem 6.2** (Chinese Remainder Theorem [Ste08]). Let  $m = pq$ , where  $p$  and  $q$  are different primes. Then  $\mathbb{Z}_m \cong \mathbb{Z}_p \times \mathbb{Z}_q$ , where the isomorphism  $\varphi : \mathbb{Z}_m \rightarrow \mathbb{Z}_p \times \mathbb{Z}_q$  is given by  $\varphi(k) = (k \bmod p, k \bmod q)$ .

**Corollary 6.1.** Let  $m = pq$  for  $p$  and  $q$  distinct primes. If  $k = 0 \bmod p$  and  $k = 0 \bmod q$ , then  $k = 0 \bmod m$ .

**Example 6.2.** Let  $m = 6, p = 2, q = 3$ . Then  $\mathbb{Z}_6 \cong \mathbb{Z}_2 \times \mathbb{Z}_3$ . The table below gives the isomorphism from  $\mathbb{Z}_6$  to  $\mathbb{Z}_2 \times \mathbb{Z}_3$ .

$k$	$\varphi(k)$
0	(0, 0)
1	(1, 1)
2	(0, 2)
3	(1, 0)
4	(0, 1)
5	(1, 2)

Notice that addition and multiplication is preserved under  $\varphi$ , that is  $\varphi(k_1 + k_2) = \varphi(k_1) + \varphi(k_2)$ . For instance,  $2 + 4 = 0 \pmod 6$  in  $\mathbb{Z}_6$ , and  $(0, 2) + (0, 1) = (0, 3) = (0, 0)$  in  $\mathbb{Z}_2 \times \mathbb{Z}_3$ .

**Definition 6.3.**  $k \in \mathbb{Z}_6 \setminus \{0\}$  is a *zero divisor* if  $\exists \ell \in \mathbb{Z}_6 \setminus \{0\}$  such that  $k \cdot \ell = 0$ .

Notice that  $k$  is a zero divisor if and only if at least one of the two elements of the pair  $\varphi(k)$  is zero.

The Chinese Remainder Theorem extends to polynomials, that is, if  $p$  and  $q$  are prime numbers with  $m = pq$ , then  $\mathbb{Z}_m[\mathbf{x}_1, \dots, \mathbf{x}_n] \cong \mathbb{Z}_p[\mathbf{x}_1, \dots, \mathbf{x}_n] \times \mathbb{Z}_q[\mathbf{x}_1, \dots, \mathbf{x}_n]$ , where the isomorphism  $\varphi : \mathbb{Z}_m[\mathbf{x}_1, \dots, \mathbf{x}_n] \rightarrow \mathbb{Z}_p[\mathbf{x}_1, \dots, \mathbf{x}_n] \times \mathbb{Z}_q[\mathbf{x}_1, \dots, \mathbf{x}_n]$  is given by  $\varphi(f(x)) = (f(x) \pmod p, f(x) \pmod q)$ . The reason why this extension of the  $\varphi$  we defined above works is because a polynomial on  $n$  variables modulo some number  $t$  is an expression that uses sums and products and so we can simply apply  $\varphi$  on all the coefficients.

**Example 6.3.** If  $m = 6, p = 2, q = 3$ , take  $f(x) = 4\mathbf{x}_1^2 + 3\mathbf{x}_1\mathbf{x}_2 + 5\mathbf{x}_1 + 1 \in \mathbb{Z}_6[\mathbf{x}_1, \mathbf{x}_2]$ . Then we have  $\varphi(f(x)) = (\mathbf{x}_1\mathbf{x}_2 + \mathbf{x}_1 + 1, 2\mathbf{x}_1^2 + 2\mathbf{x}_1 + 1)$ .

Moreover, the extension works in the other direction as well: given two polynomials  $f_p \in \mathbb{Z}_p[\mathbf{x}_1, \dots, \mathbf{x}_n]$  and  $f_q \in \mathbb{Z}_q[\mathbf{x}_1, \dots, \mathbf{x}_n]$ , there exists a unique polynomial  $f \in \mathbb{Z}_m[\mathbf{x}_1, \dots, \mathbf{x}_n]$  such that  $f \pmod p = f_p$  and  $f \pmod q = f_q$  and  $\deg(f) = \max\{\deg(f_p), \deg(f_q)\}$ .

We give an outline of the proof of Theorem 6.1 here. For any  $\mathbf{a} \in \{0, 1\}^n$ , let  $\|\mathbf{a}\| = |\{i \in [n] \mid \mathbf{a}_i = 1\}|$ , also referred to as the Hamming weight of  $\mathbf{a}$ . We construct two polynomials  $f_2 \in \mathbb{Z}_2[\mathbf{x}_1, \dots, \mathbf{x}_n]$  and  $f_3 \in \mathbb{Z}_3[\mathbf{x}_1, \dots, \mathbf{x}_n]$  such that, if  $2^d \approx \sqrt{n}$  and  $3^e \approx \sqrt{n}$ , we have that  $\forall \mathbf{a} \in \{0, 1\}^n$ :

$$\begin{aligned} f_2(\mathbf{a}) = 0 &\Leftrightarrow \|\mathbf{a}\| = 0 \pmod{2^d} \\ f_3(\mathbf{a}) = 0 &\Leftrightarrow \|\mathbf{a}\| = 0 \pmod{3^e} \end{aligned}$$

We make sure that  $f_2$  and  $f_3$  have degree approximately  $\sqrt{n}$ . Intuitively, the reason we can do this is because by the Schwartz-Zippel lemma, the number of zeros of  $f_2$  and  $f_3$ , if they have degree approximately  $\sqrt{n}$ , is at most  $\sqrt{n}2^{n-1}$ , and the number of values of  $\mathbf{a}$  such that the above condition requires that  $f_2(\mathbf{a}) = 0$  or  $f_3(\mathbf{a}) = 0$  is far smaller than  $\sqrt{n}2^{n-1}$  – it's approximately  $\sum_{i=1}^{\sqrt{n}} \binom{n}{i\sqrt{n}} \leq \sum_{i=1}^{\sqrt{n}} 2^{n-1} = \sqrt{n}2^{n-1}$ . Then, using the Chinese Remainder Theorem, we combine  $f_2$  and  $f_3$  to get  $f = (f_2, f_3) \in \mathbb{Z}_6[\mathbf{x}_1, \dots, \mathbf{x}_n]$ . This gives

us

$$\begin{aligned} f(\mathbf{a}) = 0 &\Leftrightarrow f_2(\mathbf{a}) = 0 \text{ and } f_3(\mathbf{a}) = 0 \\ &\Leftrightarrow \|\mathbf{a}\| = 0 \pmod{2^d} \text{ and } \|\mathbf{a}\| = 0 \pmod{3^e}, \end{aligned}$$

and using the Chinese Remainder Theorem on  $\mathbb{Z}_{2^d} \times \mathbb{Z}_{3^e} \cong \mathbb{Z}_{2^d 3^e}$ , we have that

$$\|\mathbf{a}\| = 0 \pmod{2^d} \text{ and } \|\mathbf{a}\| = 0 \pmod{3^e} \Leftrightarrow \|\mathbf{a}\| = 0 \pmod{2^d 3^e}.$$

If we make sure that  $2^d 3^e > n$ , then we get

$$\begin{aligned} \|\mathbf{a}\| = 0 \pmod{2^d 3^e} &\Leftrightarrow \|\mathbf{a}\| = 0 \\ &\Leftrightarrow \mathbf{a} = \mathbf{0}, \end{aligned}$$

where we use  $\mathbf{a} = \mathbf{0}$  to denote  $\forall i \in [n], \mathbf{a}_i = 0$ . Putting it all together, we get  $f(\mathbf{a}) = 0 \Leftrightarrow \mathbf{a} = \mathbf{0}$ , which is what we wanted to show.

Now we give the details of this intuition. We use the following lemma.

**Lemma 6.2.** If  $p$  is a fixed prime number, then  $\forall d \in \mathbb{N}$ , there exists a polynomial  $f_p \in \mathbb{Z}_p[\mathbf{x}_1, \dots, \mathbf{x}_n]$  such that  $\deg(f) \leq p^d - 1$  and  $\forall \mathbf{a} \in \{0, 1\}^n, f_p(\mathbf{a}) = 0 \Leftrightarrow \|\mathbf{a}\| = 0 \pmod{p^d}$ .

*Proof.* Let  $w = \|\mathbf{a}\| \in \{0, 1, \dots, n\}$ . Write  $w$  in its base  $p$  expansion:  $w = \sum_{j=0}^{\infty} \mathbf{w}_j p^j$ , with each  $\mathbf{w}_j \in \mathbb{Z}_p = \{0, 1, \dots, p-1\}$ . Then we define  $f_p$  as a symmetric function in  $\mathbf{w}_0, \mathbf{w}_1, \dots, \mathbf{w}_{d-1}$  of degree at most  $d(p-1)$ . More specifically, let  $f_p(\mathbf{a}) = \prod_{j=0}^{d-1} (1 - \mathbf{w}_j^{p-1}) - 1 \pmod{p}$ . Now we show that  $f_p(\mathbf{a}) = 0 \Leftrightarrow \|\mathbf{a}\| = 0 \pmod{p^d}$ .

- First, suppose that  $\|\mathbf{a}\| = w = 0 \pmod{p^d}$ . This means that  $\forall j \in \{0, 1, \dots, d-1\}, \mathbf{w}_j = 0$ . Then  $f_p(\mathbf{a}) = \prod_{j=0}^{d-1} (1 - \mathbf{w}_j^{p-1}) - 1 \pmod{p} = \prod_{j=0}^{d-1} 1 - 1 \pmod{p} = 1 - 1 \pmod{p} = 0$ .
- Now suppose that  $f_p(\mathbf{a}) = 0$ . Suppose that for some  $j \in \{0, 1, \dots, d-1\}, \mathbf{w}_j \neq 0$ . Then by Fermat's Little Theorem  $\mathbf{w}_j^{p-1} = 1 \pmod{p}$ , so  $f_p(\mathbf{a}) = \prod_{j=0}^{d-1} (1 - \mathbf{w}_j^{p-1}) - 1 \pmod{p} = 0 - 1 \pmod{p} = p - 1 \neq 0$ . We have reached a contradiction assuming that for some  $j \in \{0, 1, \dots, d-1\}, \mathbf{w}_j \neq 0$ . Therefore,  $\forall j \in \{0, 1, \dots, d-1\}, \mathbf{w}_j = 0$ , which implies that  $w = 0 \pmod{p^d}$ .

Now we need to show that  $\forall i \in \{0, 1, \dots, d-1\}$ , we can write  $\mathbf{w}_j(\mathbf{a})$  as a low degree polynomial in  $\mathbb{Z}_p[\mathbf{a}_1, \mathbf{a}_2, \dots, \mathbf{a}_n]$ , so that  $f_p(\mathbf{a})$  has degree at most  $p^d - 1$ . We make use of Lucas' Theorem.

**Theorem 6.3** (Lucas' Theorem [AL12]). Let  $p$  be a prime number and let  $r, s \in \mathbb{N}$ . If the base  $p$  expansions of  $r$  and  $s$  are  $r = \sum_{j=0}^{\infty} \mathbf{r}_j p^j$  and  $s = \sum_{j=0}^{\infty} \mathbf{s}_j p^j$ , then  $\binom{r}{s} = \prod_{j=0}^{\infty} \binom{\mathbf{r}_j}{\mathbf{s}_j} \pmod{p}$ , where we define  $\binom{m}{n} = 0$  in case  $m < n$ .

We use Lucas' Theorem to show the following claim, which in turn will help us finish the proof of Lemma 6.2.

**Claim 6.2.** For any  $j \in \{0, 1, \dots, d-1\}$ ,  $\mathbf{w}_j(\mathbf{x}) = \text{Sym}_{p^j}(\mathbf{x}) \pmod p$ , where  $\text{Sym}_s(\mathbf{x})$ ,  $s \in [n]$ , is the  $s$ -th elementary symmetric polynomial, defined in the following way:

$$\text{Sym}_s(\mathbf{x}) = \sum_{S \subseteq [n], |S|=s} \prod_{j \in S} \mathbf{x}_j.$$

*Proof.* Notice that  $\text{Sym}_{p^j}(\mathbf{a}) = \binom{w}{p^j}$ , and by Theorem 6.3,  $\binom{w}{p^j} = \binom{\mathbf{w}_j}{1} \prod_{j' \neq j} \binom{\mathbf{w}_{j'}}{0} = \binom{\mathbf{w}_j}{1} = \mathbf{w}_j$ .

□

Using Claim 6.2, we get that  $f_p(\mathbf{x}) = \prod_{j=0}^{d-1} (1 - \text{Sym}_{p^j}(\mathbf{x})) - 1$ . This gives us  $\deg(f_p) \leq \sum_{j=0}^{d-1} \deg(\text{Sym}_{p^j}(\mathbf{x})) (p-1) = \sum_{j=0}^{d-1} p^j (p-1) = p^d - 1$ . Thus, we have a polynomial  $f_p$  of degree at most  $p^d - 1$  such that  $f_p(\mathbf{a}) = 0 \pmod p$  if and only if  $\|\mathbf{a}\| = 0 \pmod p^d$ , which completes the proof of Lemma 6.2.

□

Now we finish the proof of Theorem 6.1. Take  $d$  and  $e$  such that  $\sqrt{n} < 2^d \leq 2\sqrt{n}$  and  $\sqrt{n} < 3^e \leq 3\sqrt{n}$ . By Lemma 6.2, there are two polynomials  $f_2 \in \mathbb{Z}_2[\mathbf{x}_1, \dots, \mathbf{x}_n]$  and  $f_3 \in \mathbb{Z}_3[\mathbf{x}_1, \dots, \mathbf{x}_n]$  such that  $\deg(f_2) \leq 2^d - 1$ ,  $\deg(f_3) \leq 3^e - 1$ , and  $f_2(\mathbf{a}) = 0$  iff  $\|\mathbf{a}\| = 0 \pmod{2^d}$ ,  $f_3(\mathbf{a}) = 0$  iff  $\|\mathbf{a}\| = 0 \pmod{3^e}$ . Now by the Chinese Remainder Theorem for polynomials,  $f = (f_2, f_3) \in \mathbb{Z}_6[\mathbf{x}_1, \dots, \mathbf{x}_n]$  has degree at most  $\max(\deg(f_2), \deg(f_3))$ , which is at most  $3\sqrt{n}$ , and  $f(\mathbf{a}) = 0$  iff  $\|\mathbf{a}\| = 0 \pmod{2^d 3^e}$ . Since  $2^d 3^e > n$ ,  $\|\mathbf{a}\| = 0 \pmod{2^d 3^e}$  iff  $\|\mathbf{a}\| = 0$  iff  $\forall i \in [n], \mathbf{a}_i = 0$ . So  $f(\mathbf{a}) = 0$  iff  $\forall i \in [n], \mathbf{a}_i = 0$ . Thus,  $f$  represents the OR function.

□

## Obtaining a Matching Vector family

Using Theorem 6.1, we can get a Matching Vector family.

**Theorem 6.4** ([Gro00]). There exists a Matching Vector family over  $\mathbb{Z}_6^\ell$  of size  $\ell^{\frac{C \log \ell}{\log^2 \log \ell}}$ , where  $C = \frac{1}{81}$ .

*Proof.* We will use Theorem 6.1. If  $h$  is the number of variables in the OR function we are looking at, then the polynomial  $f \in \mathbb{Z}_6[\mathbf{x}_1, \dots, \mathbf{x}_h]$ , which represents OR has  $\deg(f) = O(\sqrt{h})$ . Define the following  $2^h$  polynomials:  $\forall \mathbf{b} \in \{0, 1\}^h$ , let  $g_{\mathbf{b}}(\mathbf{x}) = f((\mathbf{b} - \mathbf{x})^2)$ , where  $(\mathbf{a}_1, \dots, \mathbf{a}_h)^2 = (\mathbf{a}_1^2, \dots, \mathbf{a}_h^2)$ . We will use the following claim:

**Claim 6.3.**

$$\forall \mathbf{a} \in \{0, 1\}^h, g_{\mathbf{b}}(\mathbf{a}) \begin{cases} = 0 \pmod{6} & \text{if } \mathbf{a} = \mathbf{b} \\ \neq 0 \pmod{6} & \text{if } \mathbf{a} \neq \mathbf{b} \end{cases}$$

*Proof.* If  $\mathbf{a} = \mathbf{b}$ , then  $g_{\mathbf{b}}(\mathbf{a}) = f(\mathbf{0}) = 0$ . Otherwise,  $(\mathbf{a} - \mathbf{b})^2 \neq \mathbf{0}$ , so  $f((\mathbf{a} - \mathbf{b})^2) \neq 0 \pmod{6}$ .  $\square$

We now have two sets of vectors of size  $2^h$ ,  $\mathbf{a} \in \{0, 1\}^h$  and  $\mathbf{b} \in \{0, 1\}^h$ , with the property outlined in Claim 6.3. This has a similar structure to Matching Vector families, but is not quite what we want – we want the dot product of  $\mathbf{a}$  and  $\mathbf{b}$  to have the property that  $g_{\mathbf{b}}(\mathbf{a})$  has. To achieve this, we will "linearize" the two sets of vectors. We will need the following definitions.

**Definition 6.4.** For any polynomial  $g \in \mathbb{Z}_6[\mathbf{x}_1, \dots, \mathbf{x}_h]$  with  $\deg(g) \leq d$ , let  $\text{Coeff}(g) \in \mathbb{Z}_6^{\binom{h+d}{d}}$  be the coefficient vector of  $g$  under some fixed order of monomials.

**Definition 6.5.** Let  $V_{\leq d} : \mathbb{Z}_6^h \rightarrow \mathbb{Z}_6^{\binom{h+d}{d}}$ , the degree  $\leq d$  Veronese embedding, be such that  $\forall g \in \mathbb{Z}_6[\mathbf{x}_1, \dots, \mathbf{x}_h]$  with  $\deg(g) \leq d$ , and  $\forall \mathbf{a} \in \mathbb{Z}_6^h$ , we have  $g(\mathbf{a}) = \langle \text{Coeff}(g), V_{\leq d}(\mathbf{a}) \rangle$

**Example 6.4.** Let  $h = 2$ ,  $d = 2$ . If  $\mathbf{a} = (\mathbf{a}_1, \mathbf{a}_2)$ , then  $V_{\leq 2}(\mathbf{a}) = (1, \mathbf{a}_1, \mathbf{a}_2, \mathbf{a}_1\mathbf{a}_2, \mathbf{a}_1^2, \mathbf{a}_2^2)$  and  $\text{Coeff}(g) = (g_0, g_1, g_2, g_3, g_4, g_5)$ , where  $g(\mathbf{x}) = g_0 + g_1\mathbf{x}_1 + g_2\mathbf{x}_2 + g_3\mathbf{x}_1\mathbf{x}_2 + g_4\mathbf{x}_1^2 + g_5\mathbf{x}_2^2$ .

Now we can set  $(\mathbf{v}_{\mathbf{b}} = \text{Coeff}(g_{\mathbf{b}}))_{\mathbf{b} \in \{0,1\}^h}$  and  $(\mathbf{u}_{\mathbf{a}} = V_{\leq d}(\mathbf{a}))_{\mathbf{a} \in \{0,1\}^h}$ . Then  $((\mathbf{v}_{\mathbf{b}})_{\mathbf{b} \in \{0,1\}^h}, (\mathbf{u}_{\mathbf{a}})_{\mathbf{a} \in \{0,1\}^h})$  is a Matching Vector family of size  $2^h$ . This is because  $\forall \mathbf{a}, \mathbf{b} \in \{0, 1\}^h$ ,  $\langle \mathbf{u}_{\mathbf{a}}, \mathbf{v}_{\mathbf{b}} \rangle = g_{\mathbf{b}}(\mathbf{a})$ , which is 0 modulo 6 if and only if  $\mathbf{a} = \mathbf{b}$ . We have that  $\forall \mathbf{a}, \mathbf{b} \in \{0, 1\}^h$ ,  $\mathbf{v}_{\mathbf{b}}, \mathbf{u}_{\mathbf{a}} \in \mathbb{Z}_6^{\binom{h+d}{d}}$ . Note that  $d \leq 3\sqrt{h}$ . Set  $l = \binom{h+d}{d}$ . Then we have that  $l = \frac{(h+d)!}{h!d!} = \frac{(h+1) \cdot (h+1) \cdots (h+d)}{1 \cdot 2 \cdots d} \leq (2h)^d \leq 2h^{3\sqrt{h}}$ . Now the size of our Matching Vector family is  $2^h$ , and we'll show that  $2^h \geq l^{\frac{C \log l}{\log^2 \log l}}$ , where  $C = \frac{1}{81}$ . Since  $l^{\frac{C \log l}{\log^2 \log l}} = 2^{\frac{C \log^2 l}{\log^2 \log l}}$ , it is enough to show that  $h \geq \frac{C \log^2 l}{\log^2 \log l}$ .

$$\begin{aligned}
& \frac{C \log^2 l}{\log^2 \log l} \leq \\
& \frac{\left(3\sqrt{h} \log(2h)\right)^2}{81 \log^2 \left(3\sqrt{h} \log(2h)\right)} \leq \\
& \frac{9h \log^2(2h)}{81 \log^2(\sqrt{h})} = \\
& \frac{h \left(2 \log(\sqrt{h}) + 1\right)^2}{9 \log^2(\sqrt{h})} \leq \\
& \frac{9h \log^2(\sqrt{h})}{9 \log^2(\sqrt{h})} = h
\end{aligned}$$

Thus, we have a Matching Vector family of size  $l^{\frac{\log l}{81 \log^2 \log l}}$ .

□

**Exercise 6.1.** Extend the construction given in the proof of Theorem 6.4 to  $\mathbb{Z}_m$ , where  $m = p_1 \dots p_t$ , where  $p_1, \dots, p_t$  are distinct primes,  $t > 2$ . What is the size of the Matching Vector family we can get in the case of  $t$  primes?

**Exercise 6.2.** Show that if  $p$  is prime, then any function  $f : \mathbb{Z}_p^n \rightarrow \{0, 1\}$  can be computed exactly by a polynomial  $g \in \mathbb{Z}_p[\mathbf{x}_1, \dots, \mathbf{x}_n]$ . Show that this is not true if we replace  $\mathbb{Z}_p$  with  $\mathbb{Z}_{pq}$  for  $p$  and  $q$  prime.

**Exercise 6.3.** Let  $f : \{0, 1\}^{2n} \rightarrow \{0, 1\}$  be defined in the following way:

$$\forall \mathbf{a}, \mathbf{b} \in \{0, 1\}^n, f(\mathbf{a}, \mathbf{b}) = \begin{cases} 0 & \text{if } \mathbf{a} = \mathbf{b} \\ 1 & \text{if } \mathbf{a} \neq \mathbf{b} \end{cases}$$

Find a polynomial  $g \in \mathbb{Z}_6[\mathbf{x}_1, \dots, \mathbf{x}_n, \mathbf{y}_1, \dots, \mathbf{y}_n]$  that represents  $f$  with  $\deg(g) = O(\sqrt{n})$ . Is there a polynomial  $g' \in \mathbb{Z}_p[\mathbf{x}_1, \dots, \mathbf{x}_n, \mathbf{y}_1, \dots, \mathbf{y}_n]$  for  $p$  prime with  $\deg(g') = O(\sqrt{n})$  that represents  $f$ ?

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