

## Average number of compares for QUICKSORT with distinct keys

**Recurrence from recursive program**

$$C_N = N - 1 + \frac{1}{N} \sum_{1 \leq j \leq N} (C_{j-1} + C_{N-j})$$

**Change  $j$  to  $N + 1 - j$  in second sum**

$$C_N = N - 1 + \frac{2}{N} \sum_{1 \leq j \leq N} C_{j-1}.$$

**Multiply both sides by  $N$**

$$NC_N = N(N - 1) + \frac{2}{N} \sum_{1 \leq j \leq N} C_{j-1}.$$

**Subtract same equation for  $N - 1$**

$$NC_N - (N - 1)C_{N-1} = 2N + 2C_{N-1}$$

**Rearrange terms**

$$NC_N = (N + 1)C_{N-1} + 2N$$

**Divide by  $N(N + 1)$**

$$\frac{C_N}{N + 1} = \frac{C_{N-1}}{N} + \frac{2}{N + 1}$$

**Telescope**

$$\frac{C_N}{N + 1} = 2(H_{N+1} - 1)$$

**Approximate**

$$C_N \approx 2N \ln N$$

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## Average number of compares for QUICKSORT with equal keys

### Recurrence for average number of comparisons

$$C(x_1, \dots, x_n) = N + 1 + \frac{1}{N} \sum_{1 \leq j \leq N} x_j (C(x_1, \dots, x_{j-1}) + C(x_{j+1} \dots x_n))$$

**Multiply both sides by  $N = x_1 + \dots + x_n$**

$$NC(x_1, \dots, x_n) = N(N - 1) + \sum_{1 \leq j \leq N} x_j C(x_1, \dots, x_{j-1}) + \sum_{1 \leq j \leq N} x_j C(x_{j+1}, \dots, x_n).$$

**Subtract same equation for  $x_2, \dots, x_n$  (with  $D(x_1 \dots x_n) \equiv C(x_1, \dots, x_n) - C(x_2, \dots, x_n)$ )**

$$(x_1 + \dots + x_n)D(x_1, \dots, x_n) = x_1^2 - x_1 + 2x_1(x_2 + \dots + x_n) + \sum_{2 \leq j \leq n} x_j D(x_1, \dots, x_{j-1})$$

**Subtract same equation for  $x_1, \dots, x_{n-1}$**

$$(x_1 + \dots + x_n)D(x_1, \dots, x_n) - (x_1 + \dots + x_{n-1})D(x_1, \dots, x_{n-1}) = 2x_1x_n + x_n D(x_1, \dots, x_{n-1})$$

**Simplify, divide by  $N$**

$$D(x_1, \dots, x_n) = D(x_1, \dots, x_{n-1}) + \frac{2x_1x_n}{x_1 + \dots + x_n}$$

**Telescope (twice)**

$$C(x_1, \dots, x_n) = N - n + 2 \sum_{1 \leq k \leq j \leq n} \frac{x_k x_j}{x_k + \dots + x_j}$$

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## Upper bound on QUICKSORT entropy

### Quicksort entropy definition

$$Q = \sum_{1 \leq k < j \leq n} \frac{p_k p_j}{p_k + \dots + p_j}$$

### Separate double sum

$$Q = \sum_{1 \leq k < n} p_k \sum_{k < j \leq n} \frac{p_j}{p_k + \dots + p_j}$$

**Substitute**  $q_{ij} = (p_i + \dots + p_j / p_i)$  (**note:**  $1 = q_{ii} \leq q_{i(i+1)} \leq \dots \leq q_{in} < 1/p_i$ )

$$Q = \sum_{1 \leq k < n} p_k \sum_{k < j \leq n} \frac{q_{kj} - q_{k(j-1)}}{q_{kj}}$$

### Bound with integral

$$Q < \sum_{1 \leq k < n} p_k \int_{q_{kk}}^{q_{kn}} \frac{1}{x} dx$$

### Simplify

$$Q < \sum_{1 \leq k < n} p_k \ln q_{kn} \leq \sum_{1 \leq k < n} p_k (-\ln p_k) = H \ln 2$$