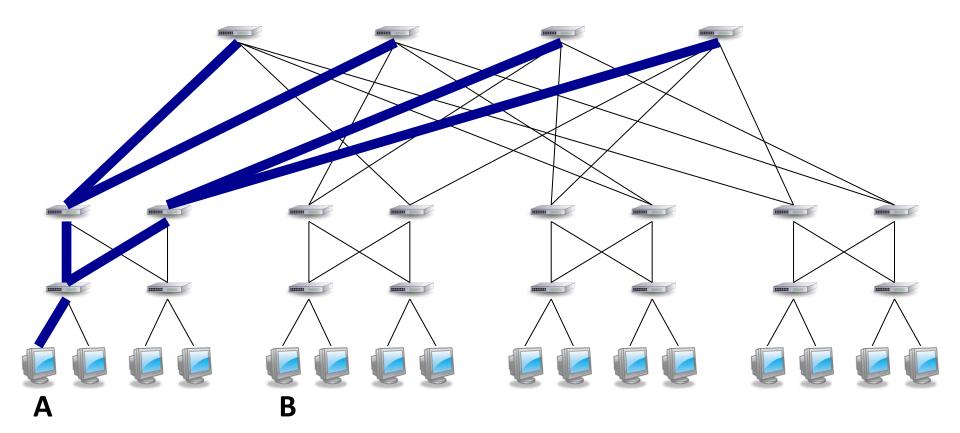
Bridging the Theory-Practice Gap in Multi-Commodity Flow Routing

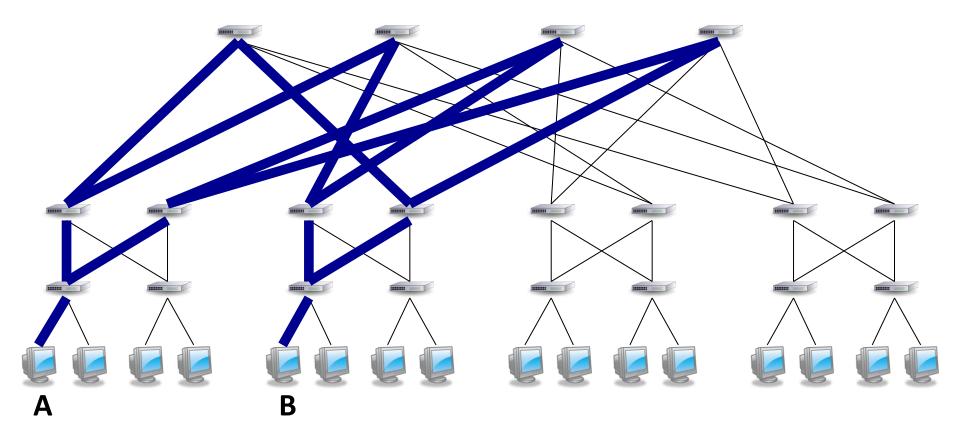
Siddhartha Sen, DISC 2011

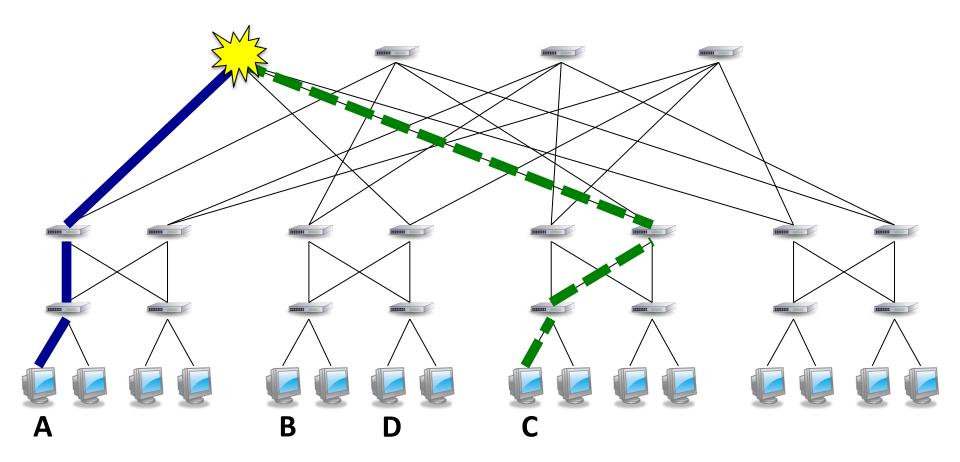
Joint work with Sunghwan Ihm, Kay Ousterhout, and Mike Freedman Princeton University

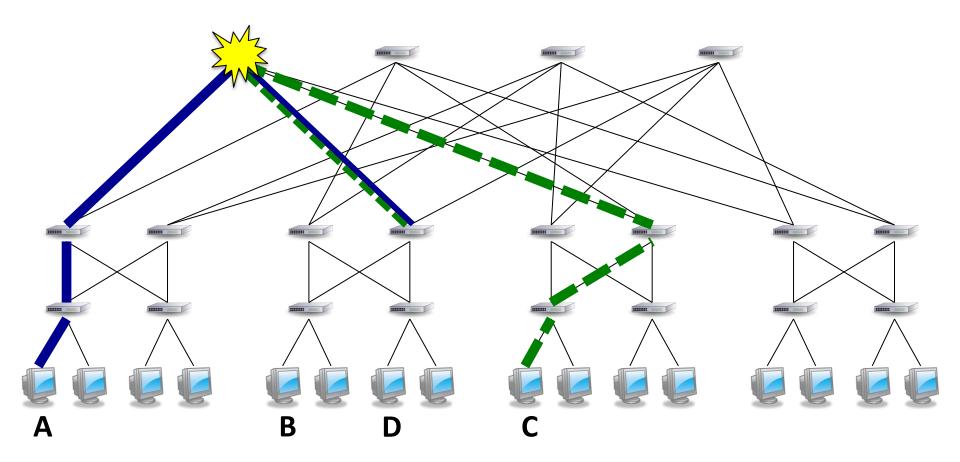
Routing flows in data center networks

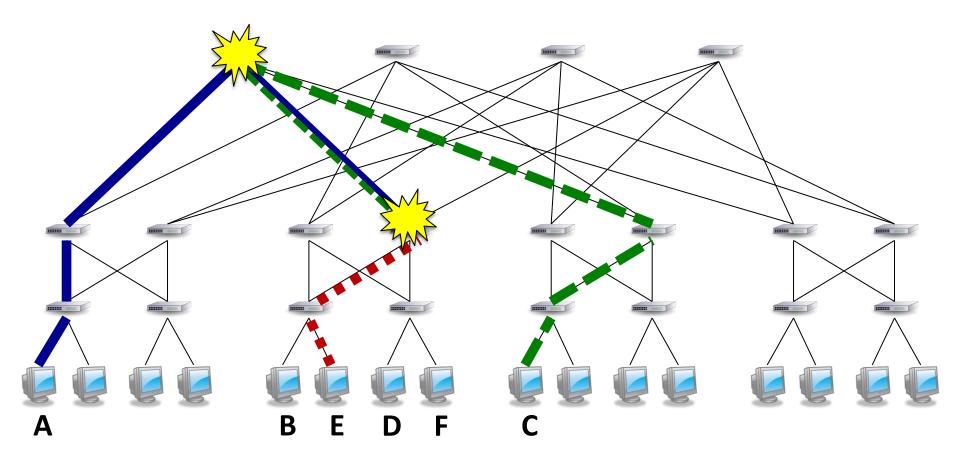


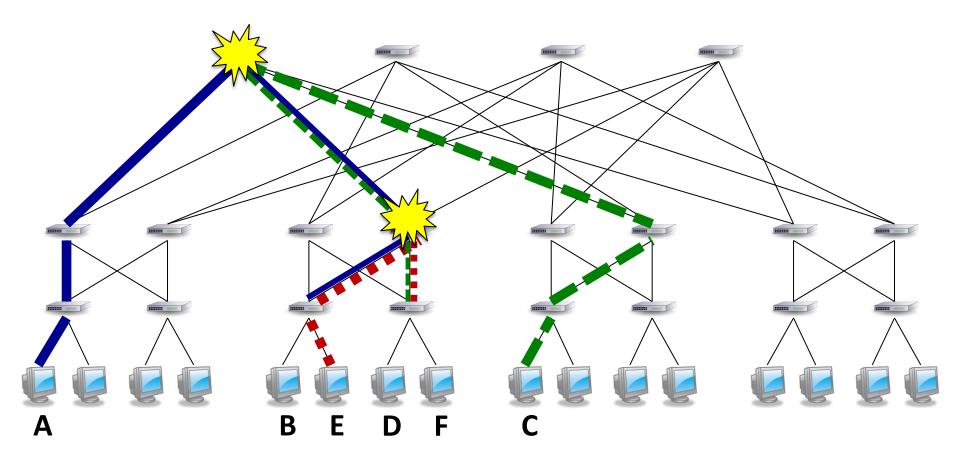
Routing flows in data center networks

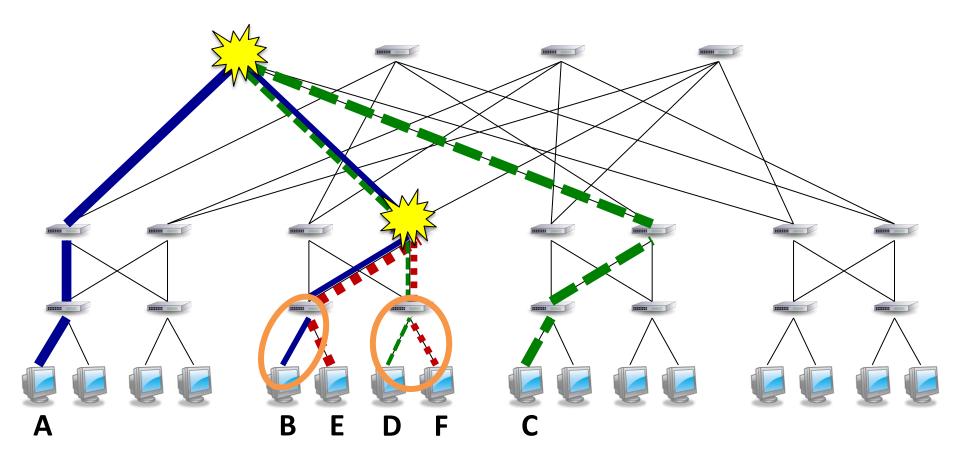




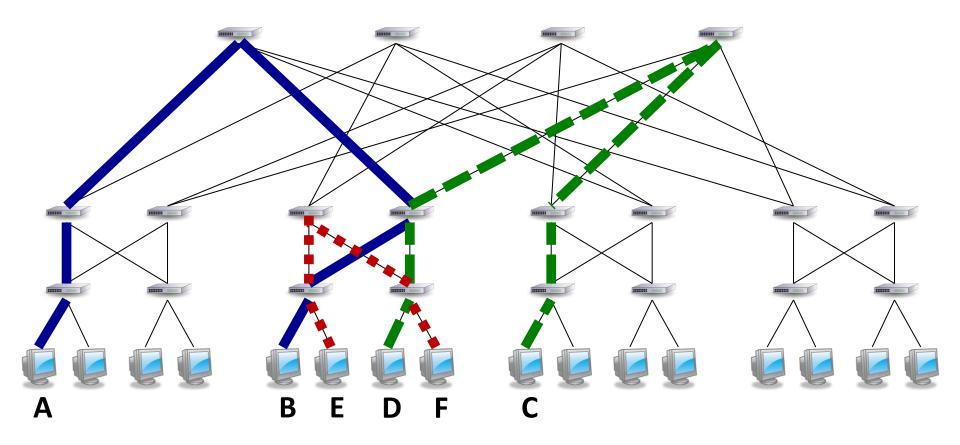








... but there is available capacity!



... but there is available capacity!

Must compute routes repeatedly: real workloads are dynamic (~ms)!

Multi-commodity flow problem

- Input: Network G = (V,E) of switches and links
 Flows K = {(s_i,t_i,d_i)} of source, target,
 demand tuples
- **Goal:** Compute flow that maximizes minimum **fraction** of any *d_i* routed
- Requires fractionally splitting flows, otherwise no O(1)-factor approximation

Prior solutions

• Sequential model

• Billboard model

• Routers model

more decentralized

Prior solutions

- Sequential model
 - Theory: [Vaidya89, PlotkinST95, GargK07, ...]
 - Practice: [BertsekasG87, BurnsOKM03, Hedera10, ...]
- Billboard model
 - Theory: [AwerbuchKR07, AwerbuchK09, ...]
 - Practice: [MATE01, TeXCP05, MPTCP11, …]
- Routers model
 - Theory: [AwerbuchL93, AwerbuchL94, AwerbuchK07, ...]
 - Practice: [REPLEX06, COPE06, FLARE07, ...]

Prior solutions

- Sequential model
 - Theory: [Vaidya89, PlotkinST95, GargK07, ...]

Theory-practice gap:

Models unsuitable for dynamic workloads
 Splitting flows difficult in practice

Routers model

- Theory: [AwerbuchL93, AwerbuchL94, AwerbuchK07, ...]
- Practice: [REPLEX06, COPE06, FLARE07, ...]

Contributions

- Demonstrate why prior solutions fail, both theoretically and practically
- Propose theoretical + practical fixes
- Devise algorithms in new framework

Contributions

- Demonstrate why prior solutions fail, both theoretically and practically
- Propose theoretical + practical fixes
- Devise algorithms in new framework

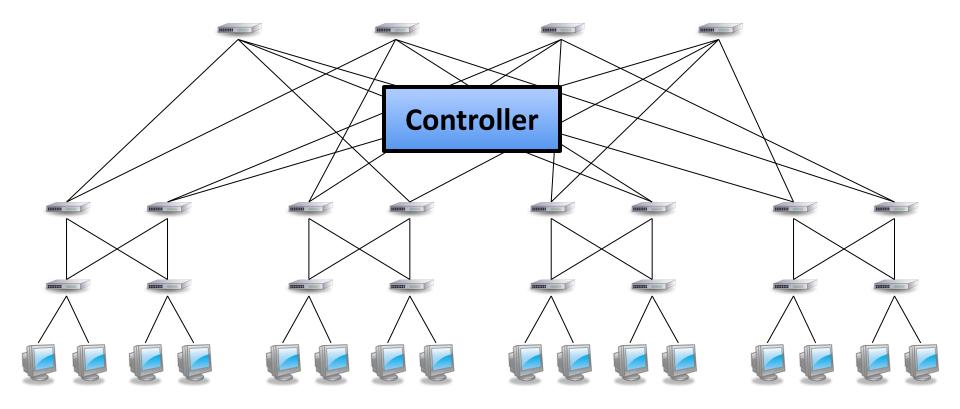
Goal: Provably optimal + practical multi-commodity flow routing

Problems

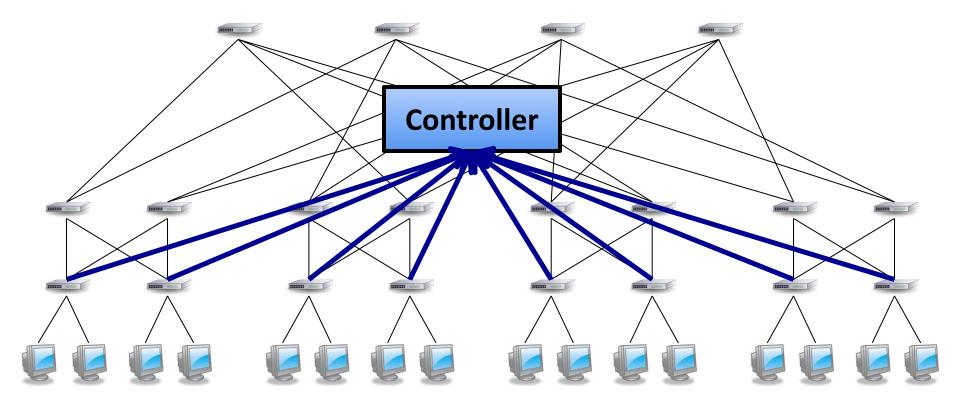
1. Dynamic workloads

Solutions

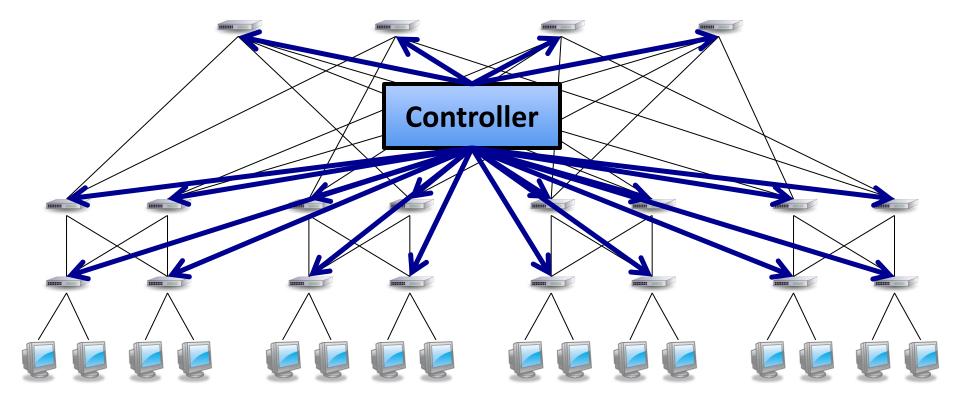
Sequential solutions don't scale



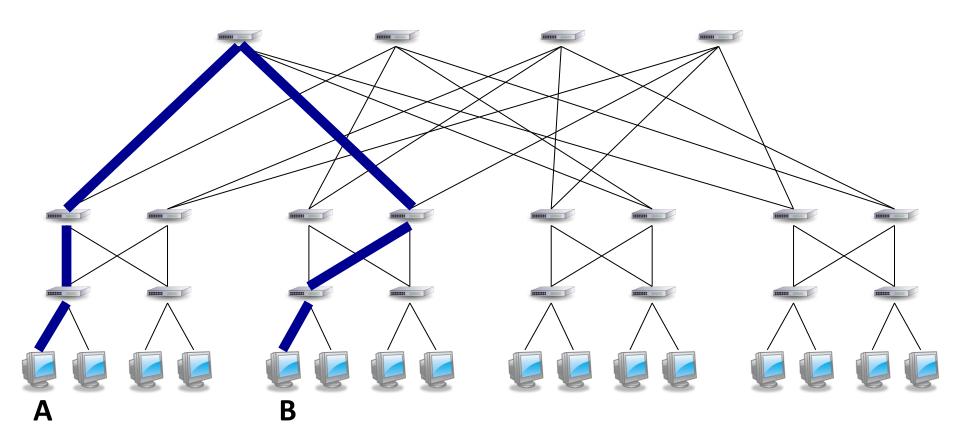
Sequential solutions don't scale



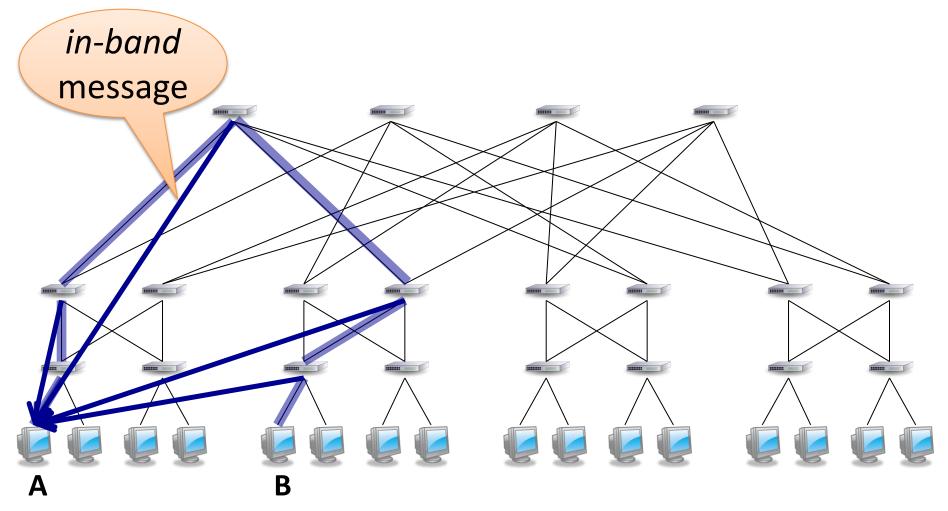
Sequential solutions don't scale



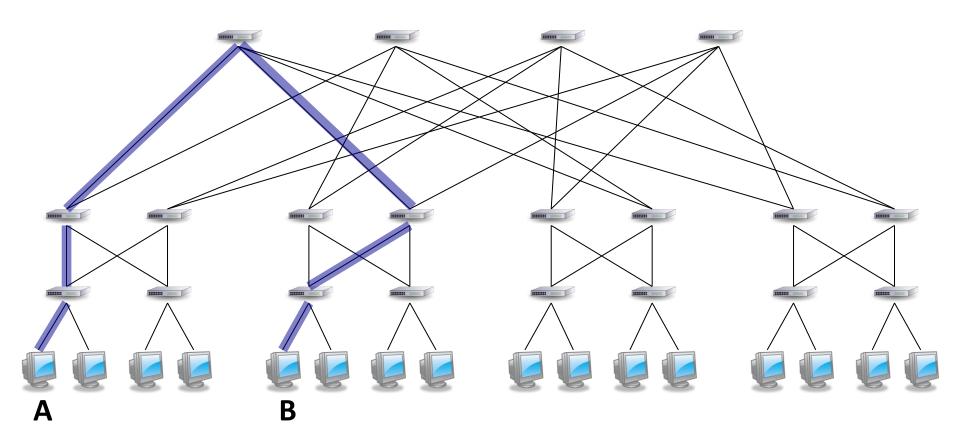
Billboard solutions require link utilization information...



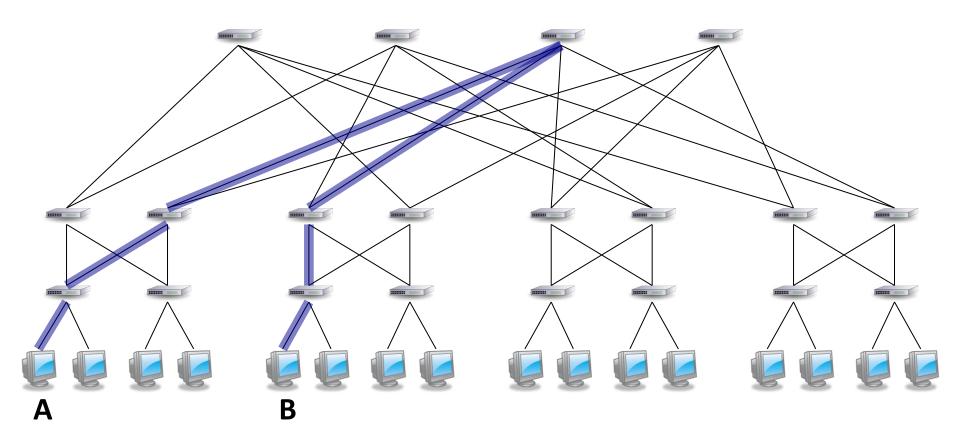
Billboard solutions require link utilization information...



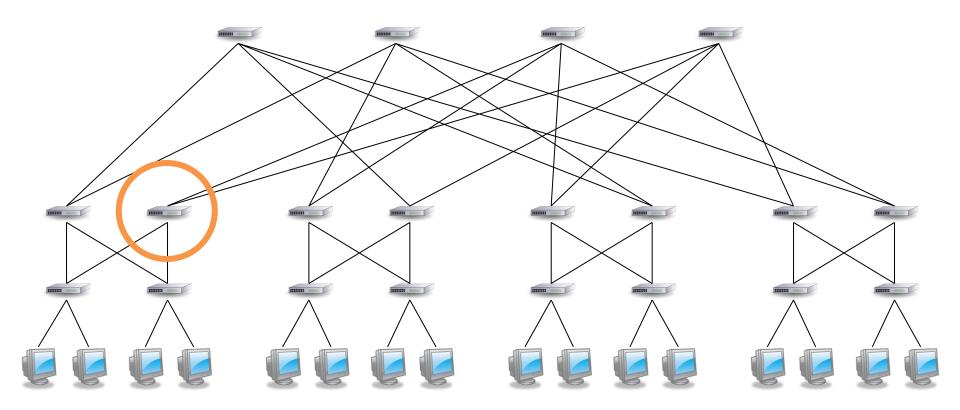
... and use path-based (exponential) representations



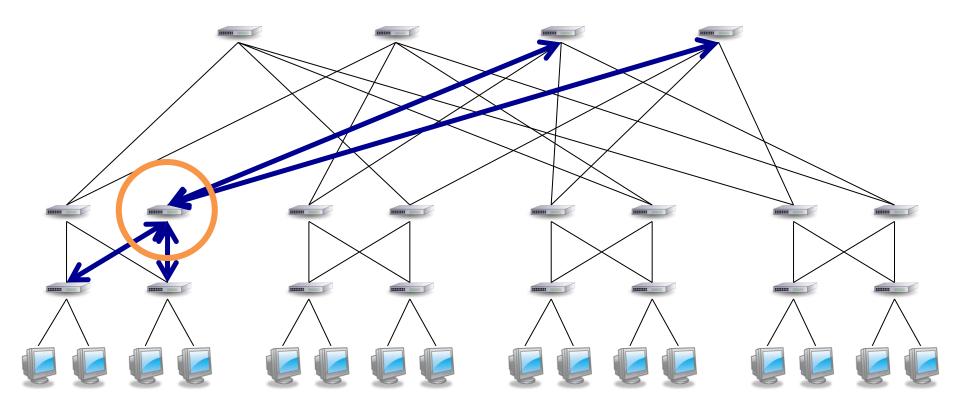
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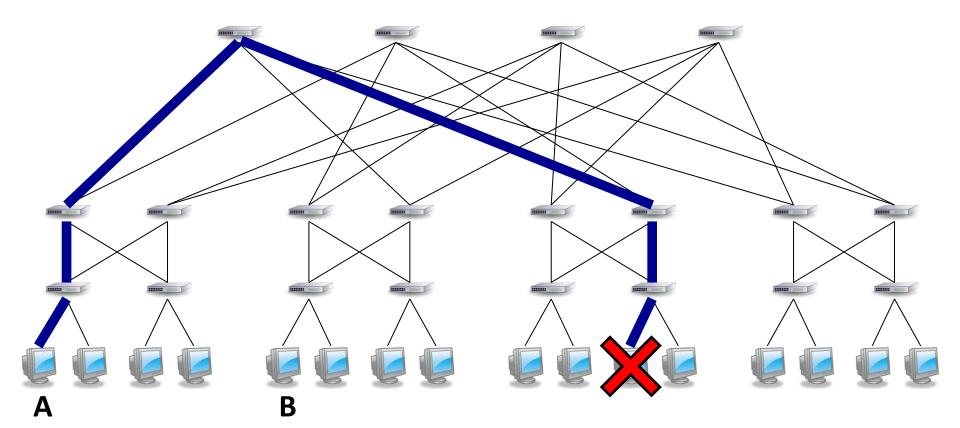
Routers solutions are local and hence scalable...



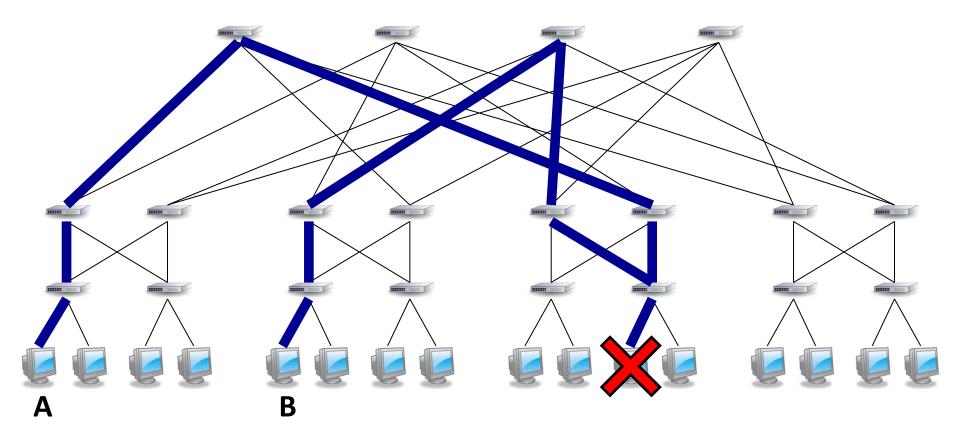
Routers solutions are local and hence scalable...



... but lack knowledge of global routes



... but lack knowledge of global routes



Problems

1. Dynamic workloads

Solutions

- Routers Plus Preprocessing (RPP) model
 - Poly-time preprocessing is free
 - In-band messages are free

Problems

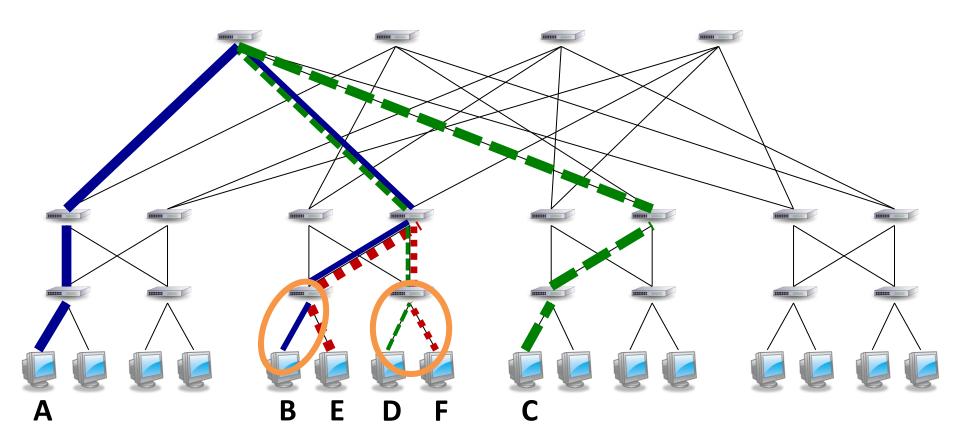
1. Dynamic workloads

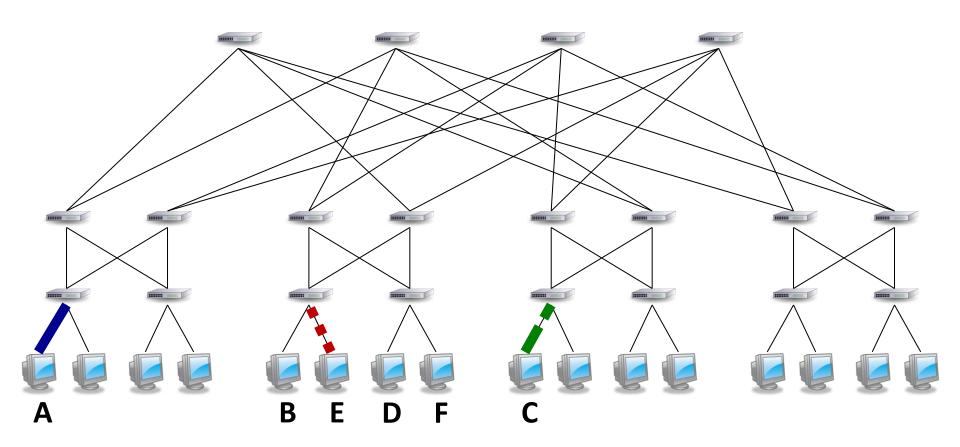
2. Splitting flows

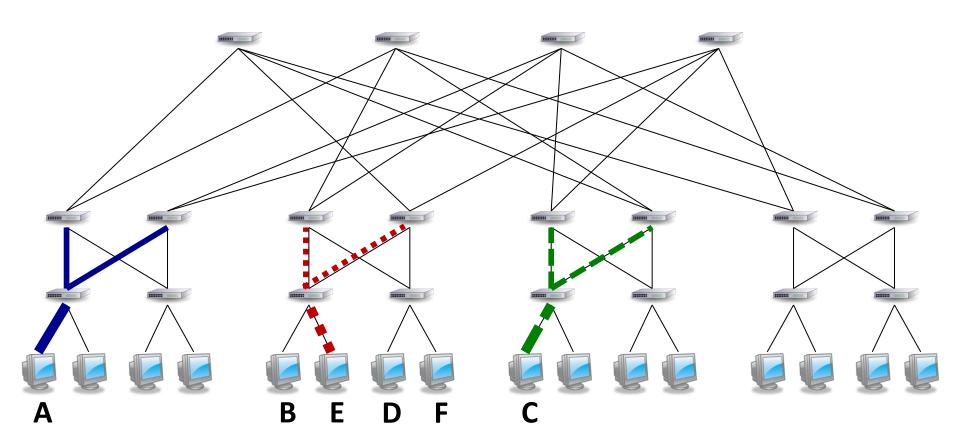
Solutions

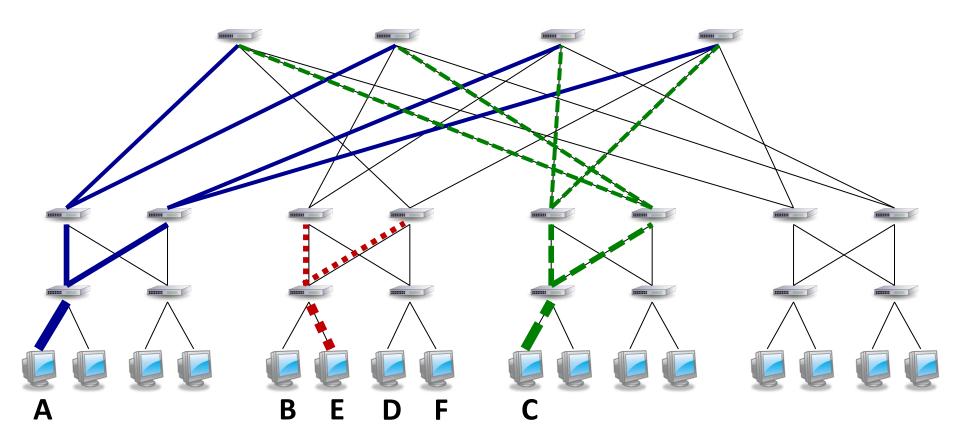
- 1. Routers Plus Preprocessing (RPP) model
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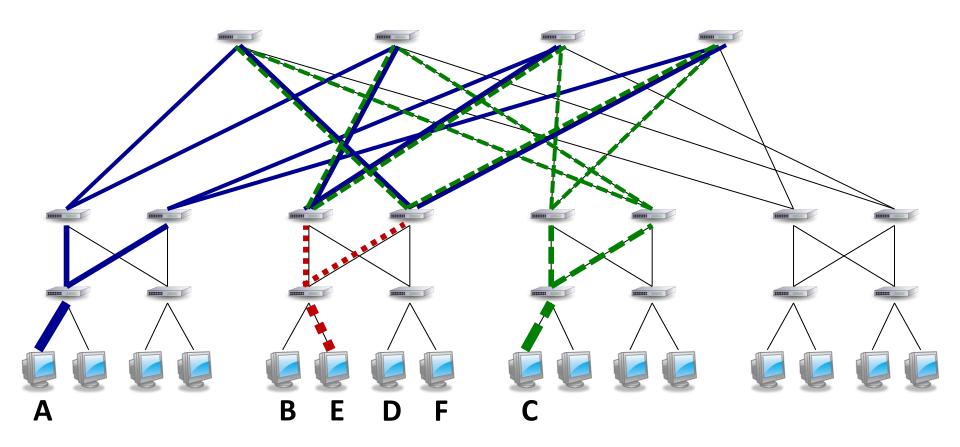
No splitting + collisions

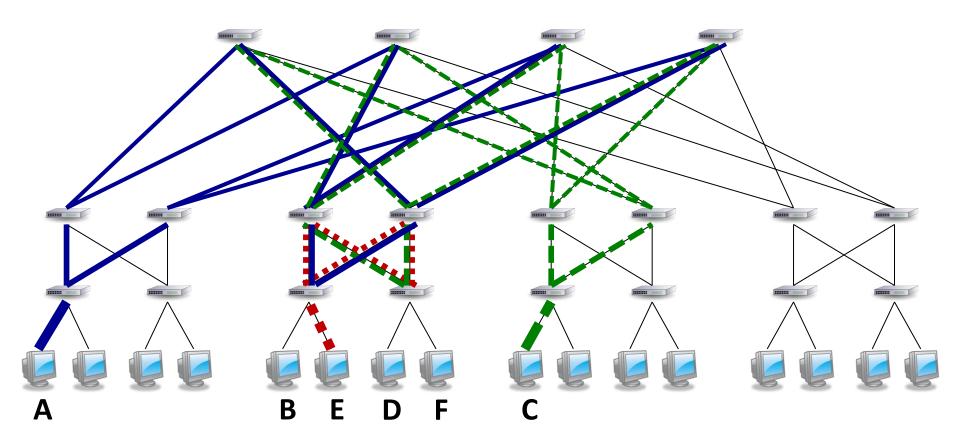




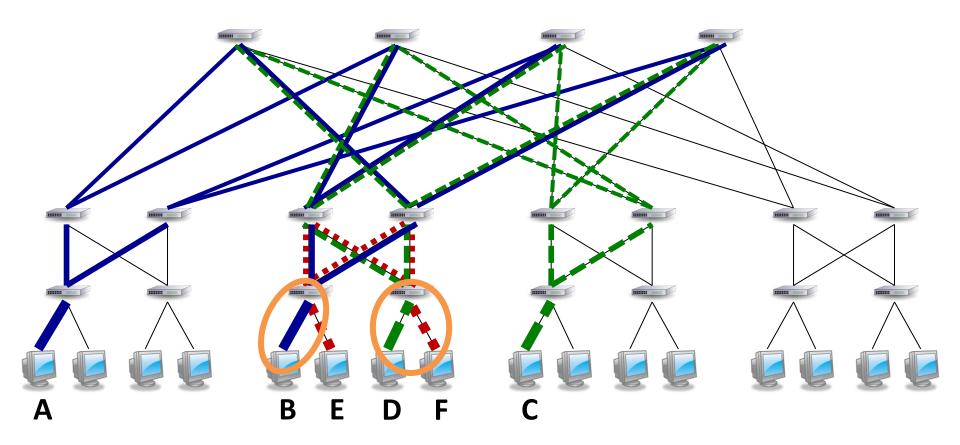








Proactive splitting



Proactive splitting



- Split every flow?
 - Inefficient
- What granularity to split at?
 - Per-packet \Rightarrow too much reordering

Problems

1. Dynamic workloads

2. Splitting flows

Solutions

- Routers Plus Preprocessing (RPP) model
 - Poly-time preprocessing is freeIn-band messages are free

2. Splitting technique

- Group flows by target, split aggregate flow
- Group contiguous packets into *flowlets* to reduce reordering

Problems

1. Dynamic workloads

2. Splitting flows

3. Switch \neq end host

 Limited processing, high-speed matching on packet headers

Solutions

- 1. Routers Plus Preprocessing (RPP) model
 - Poly-time preprocessing is freeIn-band messages are free

2. Splitting technique

- Group flows by target, split aggregate flow
- Group contiguous packets into *flowlets* to reduce reordering
- 3. Add forwarding table rules to programmable switches
 - Match TCP seq num header, use bit tricks to create flowlets

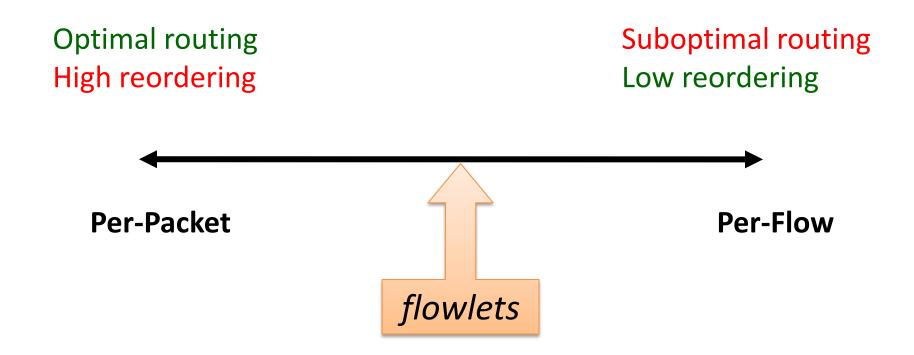
Conclusion

 Both theoretical and practical innovations needed to bridge theory-practice gap

• LOCALFLOW: optimal algorithm in new framework for data center networks

Additional slides

Granularity of splitting



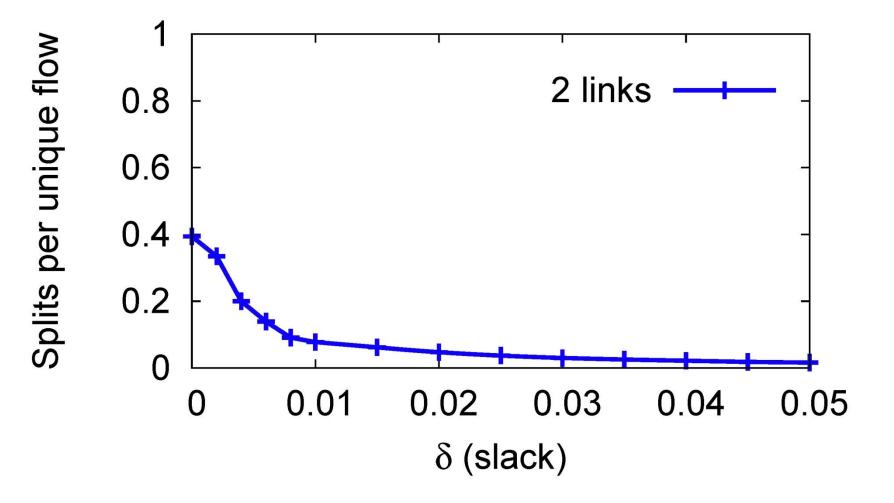
Line rate splitting

Flow	TCP seq num	Link
$A \rightarrow B$	*0*****	1
$A \rightarrow B$	*10****	2
$A \rightarrow B$	*11****	3
	flowlet = 16 packets	

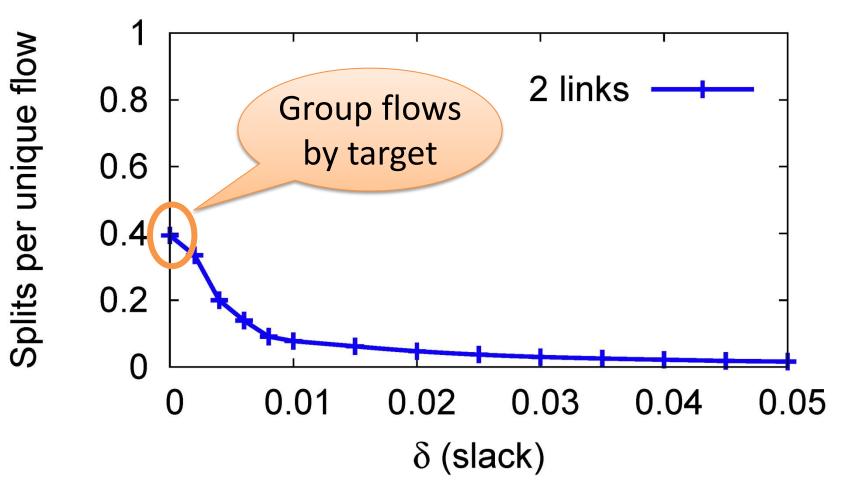
Line rate splitting

1/2	TCP seq num	Link
$A \rightarrow B$	*0*****	1
. 1/4	* 10****	2
A	*11****	3
1/4		
	flowlet = 16 packets	

LOCALFLOW: Frequency of splitting



LOCALFLOW: Frequency of splitting



LOCALFLOW: Frequency of splitting

Splits per unique flow

