

Not So Predictable Mining Pools

Attacking Solo Mining Pools by Bagging Blocks and Conning Competitors

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Outline

- Predictable Solo Mining is a new payout scheme being used in real-world cryptocurrency mining pools
- Our work examines the security of the Predictable Solo Mining payout scheme
- We introduce three attacks on the payout scheme
 - One attack exploiting cheap rewards in the pool
 - Two attacks increasing the cost others pay for rewards

BACKGROUND

Mining Pools

- The number of miners means solo mining is realistically unprofitable due to variability in profits
- Variability in profits goes down with larger miner hashrates
- Mining pools aggregate computational power, receive more consistent rewards, and distribute rewards to the members of the pool

Payout Scheme

- Determines how to allocate the pools revenue between individual miners
- Ideally we want a mining pool scheme to exhibit:
 - Incentive Compatibility
 - Proportional Fairness
- Mining pool operators want competitive advantage, leading to different payout schemes being used that aren't vetted

Payout Schemes: Details

- Users submit partial proofs of work to receive “shares”
 - Higher difficulty proofs of work worth more shares
- Example: Pay Per Last N Shares (PPLNS)
 - Only the last N shares submitted are considered when calculating rewards after a block is found
- More in use today, prior work shows that some violate incentive compatibility and fairness properties

Predictable Solo Mining (PSM)

- Each submitted share will increase the credit of the miner who submitted the share by the share difficulty
 - Miners with higher hash rates move up the leaderboard faster
- PSM is unique in that it **does not** divide the block reward to the pool
 - Share leader receives entirety of the reward
- *Post Reward Shares = Pre Reward Shares – Runner Up Shares*

Simple PSM Example

Shares Pre Block

- 1. A - 10,000
- 2. B - 8,000
- 3. C - 4,000
- 4. D - 3,000
- 5. E - 1,000



Shares Post-Block

- 1. B - 8,000
- 2. C - 4,000
- 3. D - 3,000
- 4. A - 2,000 ← Previous Leader
- 5. E - 1,000

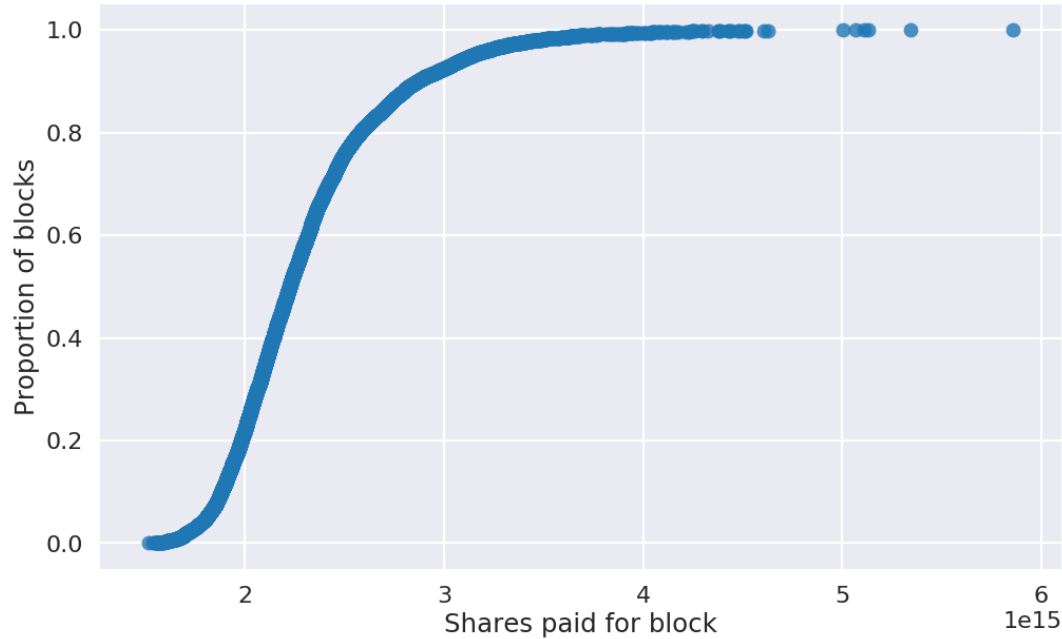
ATTACKS

Key Insights

- “Cost” of a block reward can be characterized by the number of shares held by the second place miner
- PSM Claim: The average block cost is equal to the network difficulty
- The amount of shares expended winning two different blocks, which have the *same monetary value*, varies by up to a factor of *four*

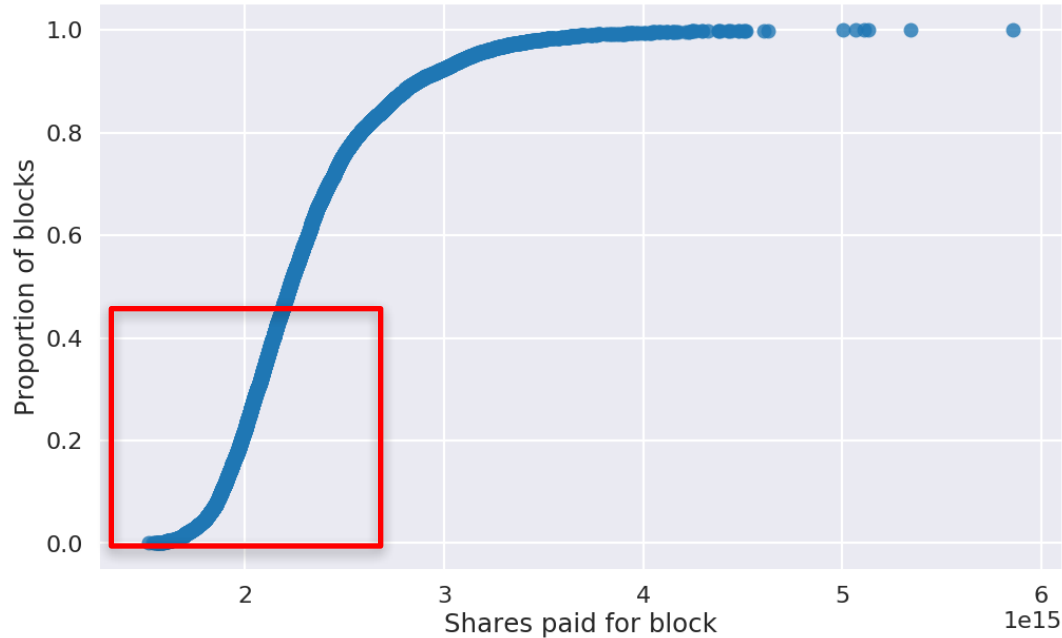
Cost of Blocks in PSM

Distribution of block costs with 100 miners over 10000 rounds



Exploiting Cheap Blocks

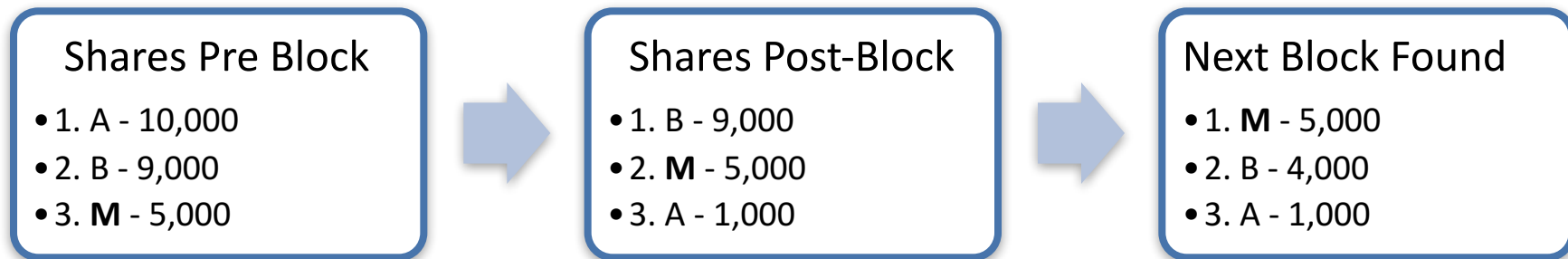
Distribution of block costs with 100 miners over 10000 rounds



Share-Cost Minimization

- Honest miners submit all of their work to the pool, driving themselves up the leaderboard
- Attacker only wants to win “cheap” blocks
- A malicious miner can refuse to place any more than n shares into their account, and only win blocks at a cost of ***at most n*** shares
- Violates proportional fairness

Share-Cost Minimization: Example

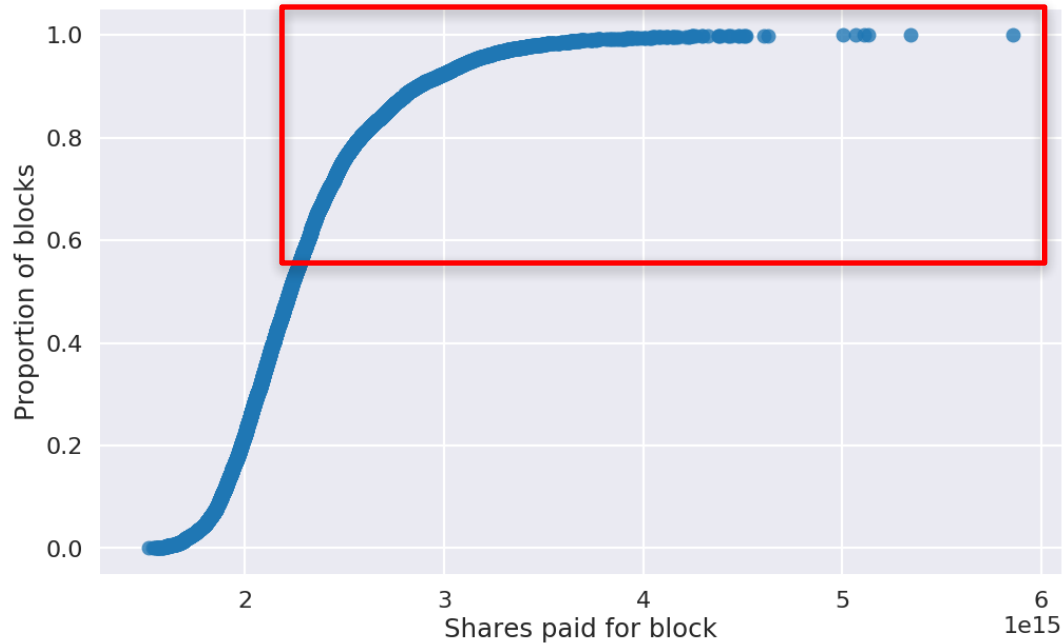


Leftover Computing Power

- Only submitting a set number of shares to the pool leads to leftover computing power
 - Spend this computational power in the same pool
 - Spend this computational power in other pools
- Violates incentive compatibility

Exploiting Expensive Blocks

Distribution of block costs with 100 miners over 10000 rounds



Malicious Share Donation

- Many pools do not authenticate share submissions
- A malicious miner can submit shares to the 2nd place miner to minimize the gap between 1st and 2nd place
- Effectively maximizes the average cost the target miner pays for each block

Malicious Share Donation: Example

Shares Pre-Donation

- 1. T - 10,000
- 2. B - 9,000
- 3. C - 4,000



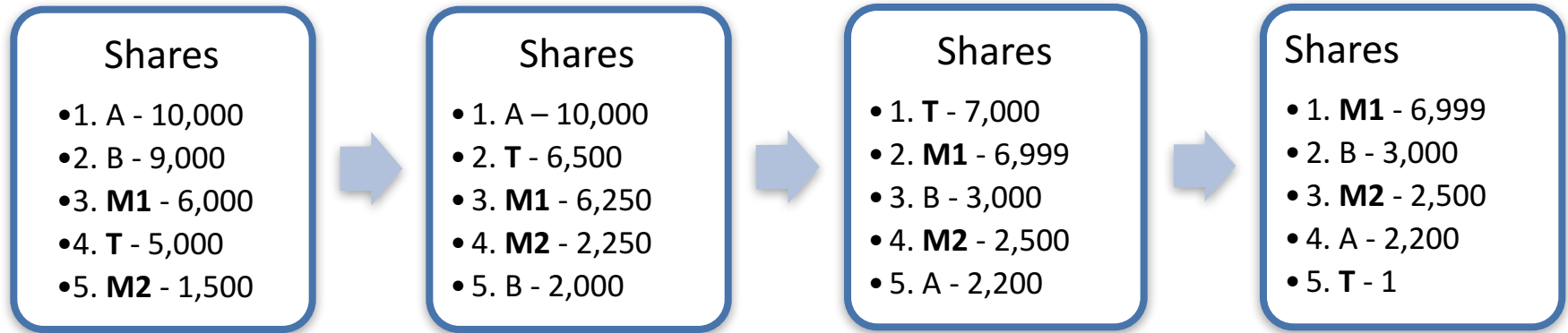
Shares Post-Donation

- 1. T - 10,000
- 2. B - 9,999
- 3. C - 4,000

Multiple Account Idling

- Share Donation Attack is intuitive and effective
 - Relies on lack of authentication in pools
- We can increase the average block cost for a target miner in pools with authentication
 - Do not need to donate to other miners
- Use multiple accounts, idle one account until target miner in range

Multiple Account Idling: Example

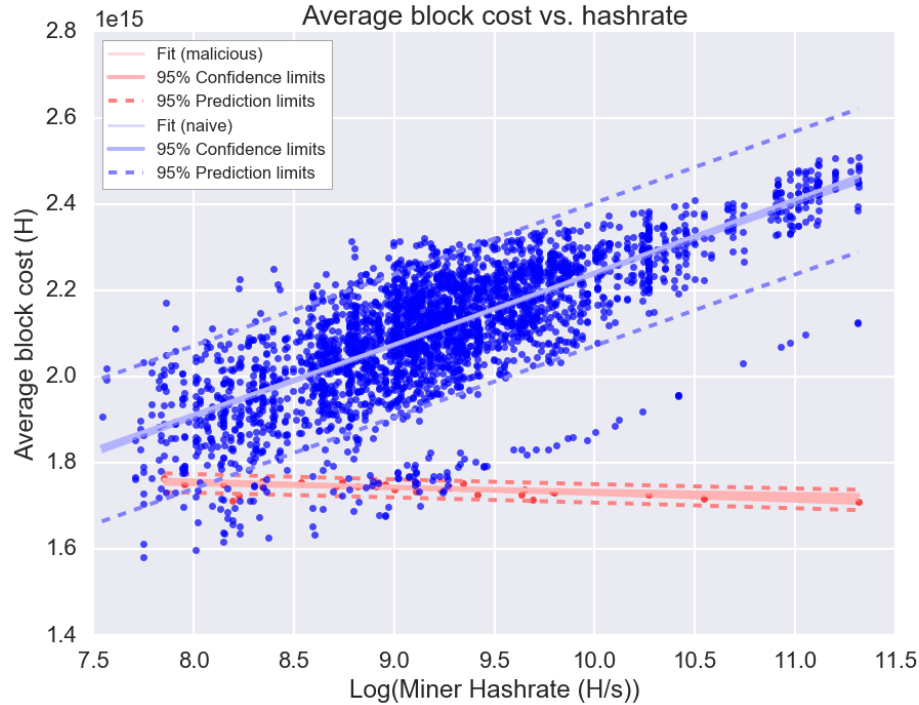


EVALUATION

Simulation

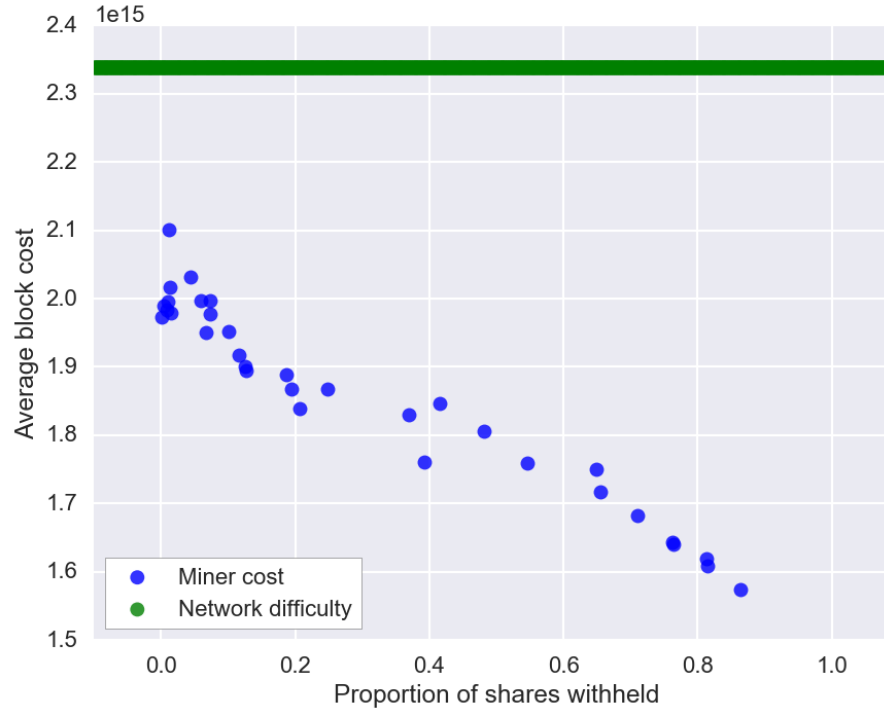
- Important to test attacks with real-world pool hashrates
- Collected active miners via Ethpool and Ethermine API
- Built discrete mining pool simulator from collected hashrates
- Mining pool simulator runs with both honest and malicious miners using current network difficulty
- Code Available at: <https://github.com/VolSec/amingpoolsimulator>

Share-Cost Minimization

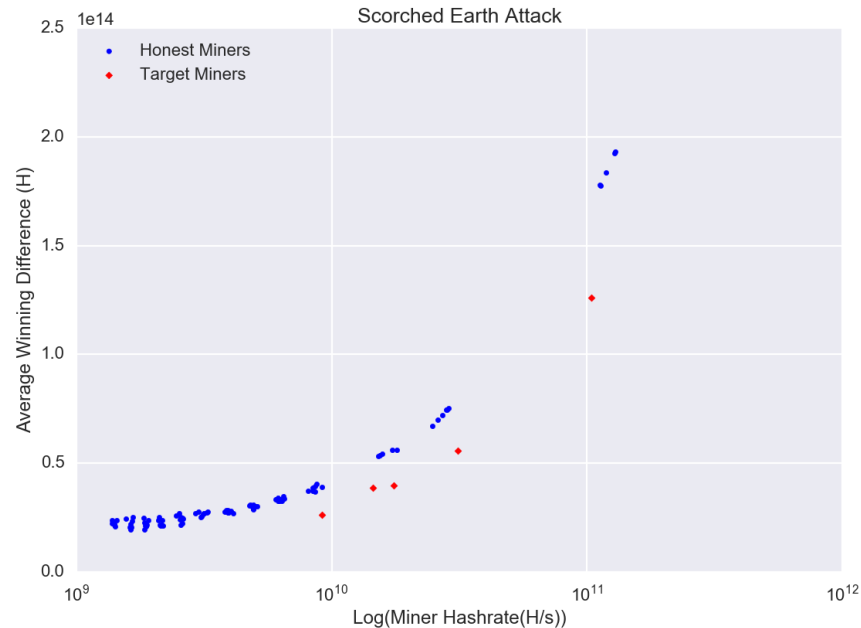
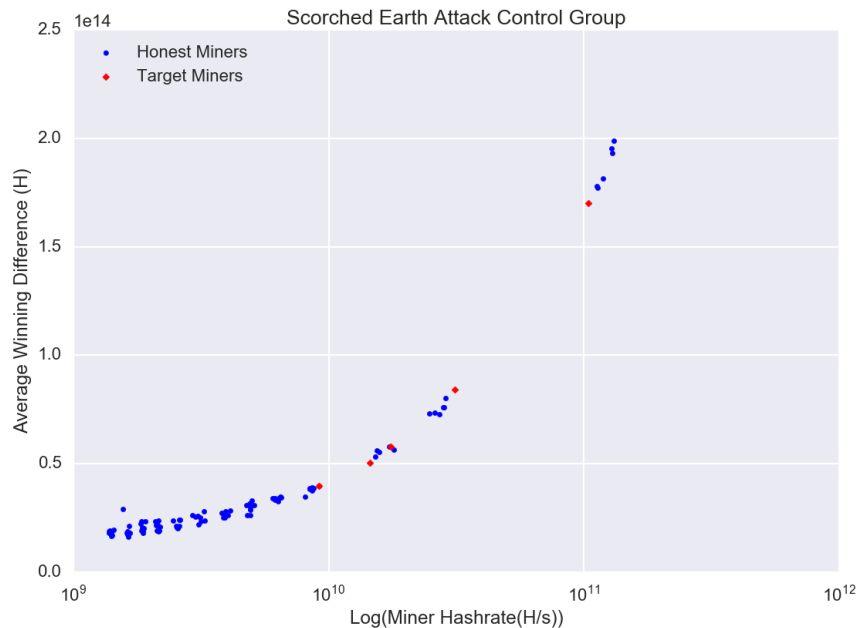


Share-Cost Minimization

Attack Results for Typical Mid-range Miner



Malicious Share Donation



Multiple Account Idling

Attacker / Target Ratio	% Decrease in Average Winning Difference
1.2	.03
4.2	5.02
7.5	6.31
9.0	5.6
14.2	8.36

Conclusions

- Payout schemes need to be vetted for incentive compatibility and fairness before being used in practice
- In any payout scheme, a single miner should not be able to influence the price of the reward of another miner
- Authentication in pools can help reduce future attacks

Questions?

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