Concurrency in Go February 2024



https://tour.golang.org/list

https://play.golang.org

https://gobyexample.com

Today's Precept...

- 1. Two synchronization mechanisms
 - a. Locks
 - b. Channels
- 2. Mapreduce

Two synchronization mechanisms

Locks - limit access to a critical section

Channels - pass information across processes using a queue

Example: Bank account



Example: Bank account



Changes to balance are not *atomic*

```
func Deposit(amount) {
```

```
read balance
balance = balance + amount
write balance
```

Suppose the function is called in two threads, with the Thread 1 chosen to run first.

```
Thread 1
func Deposit(amount) {
```

```
read balance
balance = balance + amount
write balance
```

Thread 2

func Deposit(amount) {

read balance
balance = balance + amount
write balance

Suppose the function is called in two threads, with the Thread 1 chosen to run first.

```
Thread 1
func Deposit(amount) {
```

```
read balance
balance = balance + amount
write balance
```

Thread 2

func Deposit(amount) {

read balance
balance = balance + amount
write balance

Then, an interrupt happens, and the OS scheduler selects Thread 2 to run.

```
Thread 1

func Deposit(amount) {

func Deposit(amount) {

read balance

balance = balance + amount

write balance

tread balance

balance = balance + amount

write balance

tread balance

balance = balance + amount

tread balance

balance = balance + amount

tread balance

balance = balance + amount

tread balance

tread balance

tread balance = balance + amount

tread balance

tread balance = balance + amount

tread balance

tread balance = balance + amount

tread balance
```

Thread 1 did not write new balance to shared storage, so Thread 2 reads the old value.

```
Thread 1

func Deposit(amount) {

func Deposit(amount) {

read balance

balance = balance + amount

write balance

tread balance

balance = balance + amount

write balance

tread balance

balance = balance + amount

write balance
```

This is called a **race condition**.

```
Thread 1
```

```
func Deposit(amount) {
```

```
read balance
balance = balance + amount
write balance
```

Thread 2

func Deposit(amount) {

read balance balance = balance + amount write balance

Solution - Locks

Changes to balance are now *atomic*.

func Deposit(amount) {

```
lock balanceLock
```

read balance
balance = balance + amount
write balance
unlock balanceLock

Critical section

Good Video Explanations

Race Conditions: https://www.youtube.com/watch?v=FY9livorrJI

Deadlocks: https://www.youtube.com/watch?v=LjWug2tvSBU

Locks in Go

package account

import "sync"

}

```
type Account struct {
    balance int
```

```
func NewAccount(init int) Account {
    return Account{balance: init}
}
```

func (a *Account) CheckBalance() int {

```
return a.balance
}
func (a *Account) Withdraw(v int) {
    a.balance -= v
}
```

func (a *Account) Deposit(v int) {

```
a.balance += v
```

}

Locks in Go

package account

```
import "sync"
```

```
type Account struct {
    balance int
    mu sync.Mutex
}
```

```
func NewAccount(init int) Account {
    return Account{balance: init}
}
```

```
func (a *Account) CheckBalance() int {
    a.mu.Lock()
    defer a.mu.Unlock()
    return a.balance
}
```

```
func (a *Account) Withdraw(v int) {
    a.mu.Lock()
    defer a.mu.Unlock()
    a.balance -= v
}
```

```
func (a *Account) Deposit(v int) {
    a.mu.Lock()
    defer a.mu.Unlock()
    a.balance += v
```

}

Read Write Locks in Go

package account

import "sync"

}

```
type Account struct {
    balance int
```

```
func NewAccount(init int) Account {
    return Account{balance: init}
}
```

func (a *Account) CheckBalance() int {
 a.rwLock.RLock()

```
func (a *Account) Withdraw(v int) {
```

```
a.balance -= v
```

}

}

}

func (a *Account) Deposit(v int) {

a.balance += v

Read Write Locks in Go

package account

```
import "sync"
```

```
type Account struct {
    balance int
    rwLock sync.RWMutex
}
```

```
func NewAccount(init int) Account {
    return Account{balance: init}
}
```

```
func (a *Account) CheckBalance() int {
    a.rwLock.RLock()
    defer a.rwLock.RUnlock()
    return a.balance
}
```

```
func (a *Account) Withdraw(v int) {
    a.rwLock.Lock()
    defer a.rwLock.Unlock()
    a.balance -= v
}
```

```
func (a *Account) Deposit(v int) {
    a.rwLock.Lock()
    defer a.rwLock.Unlock()
    a.balance += v
```

}

Two Solutions to the Same Problem

Locks:

Multiple threads can reference same memory location

Use lock to ensure only one thread is updating it at any given time



Channels:

Data item initially stored in channel

Threads must request item from channel, make updates, and return item to channel



```
package account
```

```
type Account struct {
    // Fill in Here
}
```

```
func NewAccount(init int) Account {
    // Fill in Here
}
```

```
func (a *Account) CheckBalance() int {
    // What goes Here?
}
func (a *Account) Withdraw(v int) {
    // ???
}
func (a *Account) Deposit(v int) {
    // ???
}
```

```
package account
type Account struct {
    balance chan int
}
```

```
func NewAccount(init int) Account {
    a := Account{
        balance: make(chan int, 1)
    }
    a.balance <- init
    return a
}</pre>
```

```
func (a *Account) CheckBalance() int {
    // What goes Here?
}
func (a *Account) Withdraw(v int) {
    // ???
}
func (a *Account) Deposit(v int) {
    // ???
}
```

```
package account
```

```
type Account struct {
    balance chan int
}
```

```
func NewAccount(init int) Account {
    a := Account{
        balance: make(chan int, 1)
    }
    a.balance <- init
    return a
}</pre>
```

```
func (a *Account) CheckBalance() int {
    bal := <-a.balance
    a.balance <- bal
    return bal
}</pre>
```

```
func (a *Account) Withdraw(v int) {
    // ???
}
```

```
func (a *Account) Deposit(v int) {
    //???
}
```

```
package account
```

```
type Account struct {
    balance chan int
}
```

```
func NewAccount(init int) Account {
    a := Account{
        balance: make(chan int, 1)
    }
    a.balance <- init
    return a
}</pre>
```

```
func (a *Account) CheckBalance() int {
    bal := <-a.balance
    a.balance <- bal
    return bal
}</pre>
```

```
func (a *Account) Withdraw(v int) {
    bal := <-a.balance
    a.balance <- (bal - v)
}</pre>
```

```
func (a *Account) Deposit(v int) {
    //???
}
```

```
package account
```

```
type Account struct {
    balance chan int
}
```

```
func NewAccount(init int) Account {
    a := Account{
        balance: make(chan int, 1)
    }
    a.balance <- init
    return a
}</pre>
```

```
func (a *Account) CheckBalance() int {
    bal := <-a.balance
    a.balance <- bal
    return bal
}</pre>
```

```
func (a *Account) Withdraw(v int) {
    bal := <-a.balance
    a.balance <- (bal - v)
}</pre>
```

```
func (a *Account) Deposit(v int) {
    bal := <-a.balance
    a.balance <- (bal + v)
}</pre>
```

Go channels

Channels also allow us to safely communicate between *goroutines*

```
result := make(chan int, numWorkers)
```

```
// Launch workers
for i := 0; i < numWorkers; i++ {</pre>
    go func() {
         doWork()
         result <- i
    }()
}
   Wait until all worker threads have finished
for i := 0; i < numWorkers; i++ {</pre>
    handleResult(<-result)</pre>
}
fmt.Println("Done!")
```

Go channels

Easy to express asynchronous RPC

Awkward to express this using locks

```
result := make(chan int, numServers)
// Send query to all servers
for i := 0; i < numServers; i++ {
   go func() {
      resp := // ... send RPC to server
      result <- resp
   }()
}</pre>
```

// Return as soon as the first server responds
handleResponse(<-result)</pre>

Select statement

select allows a goroutine to wait on multiple channels at once

```
for {
    select {
        case money := <-dad:
            buySnacks(money)
        case money := <-mom:
            buySnacks(money)
    }
}</pre>
```

Select statement

select allows a goroutine to wait on multiple channels at once

```
for {
    select {
        case money := <-dad:
            buySnacks(money)
        case money := <-mom:
            buySnacks(money)
        case default:
            starve()
            time.Sleep(5 * time.Second)
    }
</pre>
```

Handle timeouts using select

```
// Asynchronously request an answer
// from server, timing out after X
// seconds
result := make(chan int)
```

timeout := make(chan bool)

```
// Ask server
go func() {
    response := // ... send RPC
    result <- response
}()</pre>
```

```
// Start timer
go func() {
    time.Sleep(5 * time.Second)
    timeout <- true
}()</pre>
```

```
// Wait on both channels
select {
    case res := <-result:
        handleResult(res)
    case <-timeout:
        fmt.Println("Timeout!")
}</pre>
```

```
type Lock struct {
    // ???
}
func NewLock() Lock {
    // ???
}
func (1 *Lock) Lock() {
    // ???
}
func (1 *Lock) Unlock() {
    // ???
}
```

```
type Lock struct {
     ch chan bool
}
func NewLock() Lock {
     // ???
}
func (1 *Lock) Lock() {
     // ???
}
func (1 *Lock) Unlock() {
     // ???
}
```

```
type Lock struct {
     ch chan bool
}
func NewLock() Lock {
     lock := Lock{make(chan bool, 1)}
     lock.ch <- true</pre>
     return lock
}
func (1 *Lock) Lock() {
    // ???
}
func (1 *Lock) Unlock() {
     // ???
}
```

```
type Lock struct {
     ch chan bool
}
func NewLock() Lock {
     lock := Lock{make(chan bool, 1)}
     lock.ch <- true</pre>
     return lock
}
func (1 *Lock) Lock() {
     <-lock.ch
}
func (1 *Lock) Unlock() {
     // ???
}
```

```
type Lock struct {
     ch chan bool
}
func NewLock() Lock {
     lock := Lock{make(chan bool, 1)}
     lock.ch <- true</pre>
     return lock
}
func (1 *Lock) Lock() {
     <-lock.ch
}
func (1 *Lock) Unlock() {
     lock.ch <- true</pre>
}
```

Mutexes vs. Semaphores

Mutexes allow 1 process to enter critical section at a time. Allows at most *n* concurrent accesses

Semaphores allow up to **N** processes to enter critical section simultaneously



Outline

Two synchronization mechanisms

Locks

Channels

Mapreduce

Application: Word count

How much wood would a woodchuck chuck if a woodchuck could chuck wood? how: 1, much: 1, wood: 2, would: 1, a: 2, woodchuck: 2, chuck: 2, if: 1, could: 1

Application: Word count

Locally: tokenize and put words in a hash map

How do you parallelize this?

Partition the document into *n* partitions.

Build *n* hash maps, one for each partition

Merge the *n* hash maps (by key)

How do you do this in a distributed environment?



When in the Course of human events, it becomes necessary for one people to dissolve the political bands which have connected them with another, and to assume, among the Powers of the earth, the separate and equal station to which the Laws of Nature and of Nature's God entitle them, a decent respect to the opinions of mankind requires that they should declare the causes which impel them to the separation.

Input document



When in the Course of human events, it becomes necessary for one people to

dissolve the political bands which have connected them with another, and to assume,

among the Powers of the earth, the separate and equal station to which the Laws of

Nature and of Nature's God entitle them, a decent respect to the opinions of mankind

requires that they should declare the causes which impel them to the separation.



Partition



Partition

requires that they should declare the causes which impel them

to the separation.

When in the Course of human events, it becomes necessary for one people to





Nature and of Nature's God entitle them, a decent respect to the opinions of mankind

dissolve the political bands which have connected them with another, and to assume,



among the Powers of the earth, the separate and equal station to which the Laws of requires: 1, that: 1, they: 1, should: 1, declare: 1, the: 1, causes: 1, which: 1 ...

when: 1, in: 1, the: 1, course: 1, of: 1, human: 1, events: 1, it: 1



nature: 2, and: 1, of: 2, god: 1, entitle: 1, them: 1, decent: 1, respect: 1, mankind: 1, opinion: 1 ...



dissolve: 1, the: 2, political: 1, bands: 1, which: 1, have: 1, connected: 1, them: 1 ... among: 1, the: 2, powers: 1, of: 2, earth: 1, separate: 1, equal: 1, and: 1 ...

Compute word counts locally

requires: 1, that: 1, they: 1, should: 1, declare: 1, the: 1, causes: 1, which: 1 ...

when: 1, in: 1,

the: 1, course: 1,

of: 1, human: 1, events: 1, it: 1

nature: 2, and: 1, of: 2, Now what. god: 1, entitle: 1, the decent: 1, respect: 1, god: 1, entitle: 1, them: 1,

mankind: 1, opinion: 1 ...

How to merge results?

dissolve: 1, the: 2,

political: 1, bands: 1,

which: 1, have: 1,

connected: 1, them: 1 ...

among: 1, the: 2,

powers: 1, of: 2,

earth: 1, separate: 1,

equal: 1, and: 1 ...

Compute word counts locally

Merging results computed locally

Several options

Don't merge — requires additional computation for correct results

Send everything to one node — what if data is too big? Too slow...

Partition key space among nodes in cluster (e.g. [a-e], [f-j], [k-p] ...)

- 1. Assign a key space to each node
- 2. Split local results by the key spaces
- 3. Fetch and merge results that correspond to the node's key space

requires: 1, that: 1, they: 1, should: 1, declare: 1, the: 1, causes: 1, which: 1 ...

when: 1, in: 1, the: 1, course: 1, of: 1, human: 1, events: 1, it: 1



nature: 2, and: 1, of: 2, god: 1, entitle: 1, them: 1, decent: 1, respect: 1, mankind: 1, opinion: 1 ...



dissolve: 1, the: 2, political: 1, bands: 1, which: 1, have: 1, connected: 1, them: 1 ... among: 1, the: 2,
powers: 1, of: 2,
earth: 1, separate: 1,
equal: 1, and: 1 ...



Split local results by key space



All-to-all shuffle





Merge results received from other nodes

Mapreduce

Partition dataset into many chunks

Map stage: Each node processes one or more chunks locally

Reduce stage: Each node fetches and merges partial results from all other nodes

Mapreduce Interface

map(key, value) -> list(<k', v'>)

Apply function to (key, value) pair

Outputs list of intermediate pairs

reduce(key, list<value>) -> <k', v'>

Applies aggregation function to values Outputs result

Mapreduce: Word count

map(key, value):

// key = document name
// value = document contents
for each word w in value:
 emit (w, 1)

reduce(key, values):

```
// key = the word
// values = number of occurrences of that word
count = sum(values)
emit (key, count)
```

Mapreduce: Word count



Why is implementing MapReduce hard?

- Failure is common
 - Even if each machine is available p = 99.999% of the time, a datacenter with n = 100,000 machines still encounters failures $(1-p^n) = 63\%$ of the time
- Data skew causes unbalanced performance across cluster
- \rightarrow Problems occur at scale.
- → Hard to debug!

