

Light at the
end of the tunnel



Course Overview

Lecture 23

COS 461: Computer Networks

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Key Concepts in Networking

Some Key Concepts

- **Course was organized around protocols**
 - But a small set of concepts recur in many protocols
- **General CS concepts**
 - Hierarchy, indirection, caching, randomization
- **Networking-specific concepts**
 - Soft state, layering, (de)multiplexing
 - End-to-end argument

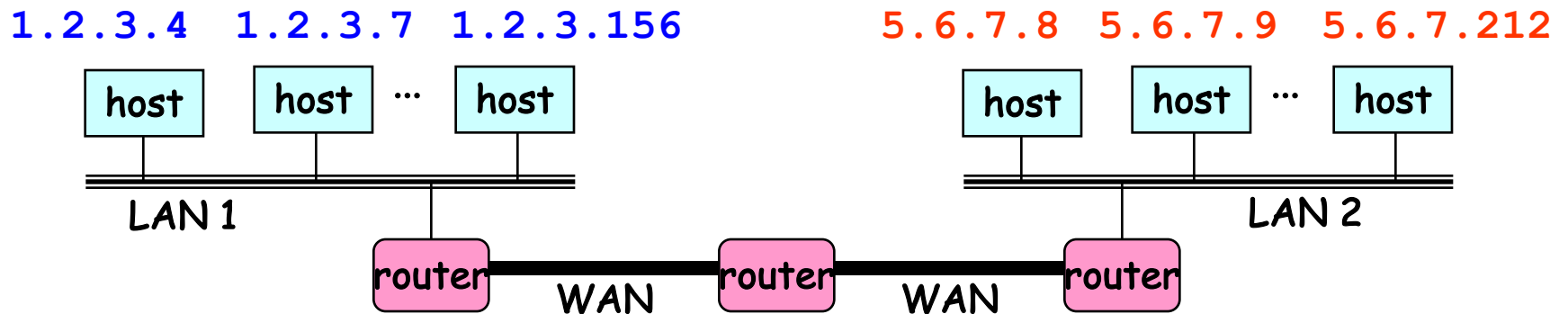
Hierarchy

- Scalability of large systems
 - Cannot store all information everywhere
 - Cannot centrally coordinate everything
- Hierarchy to manage scale
 - Divide system into smaller pieces
- Hierarchy to divide control
 - Decentralized management
- Examples in the Internet
 - IP addresses, routing protocols, DNS, P2P



Hierarchy: IP Address Blocks

- Number related hosts from a common subnet
 - 1.2.3.0/24 on the left LAN
 - 5.6.7.0/24 on the right LAN



1.2.3.0/24	←
5.6.7.0/24	→

forwarding table

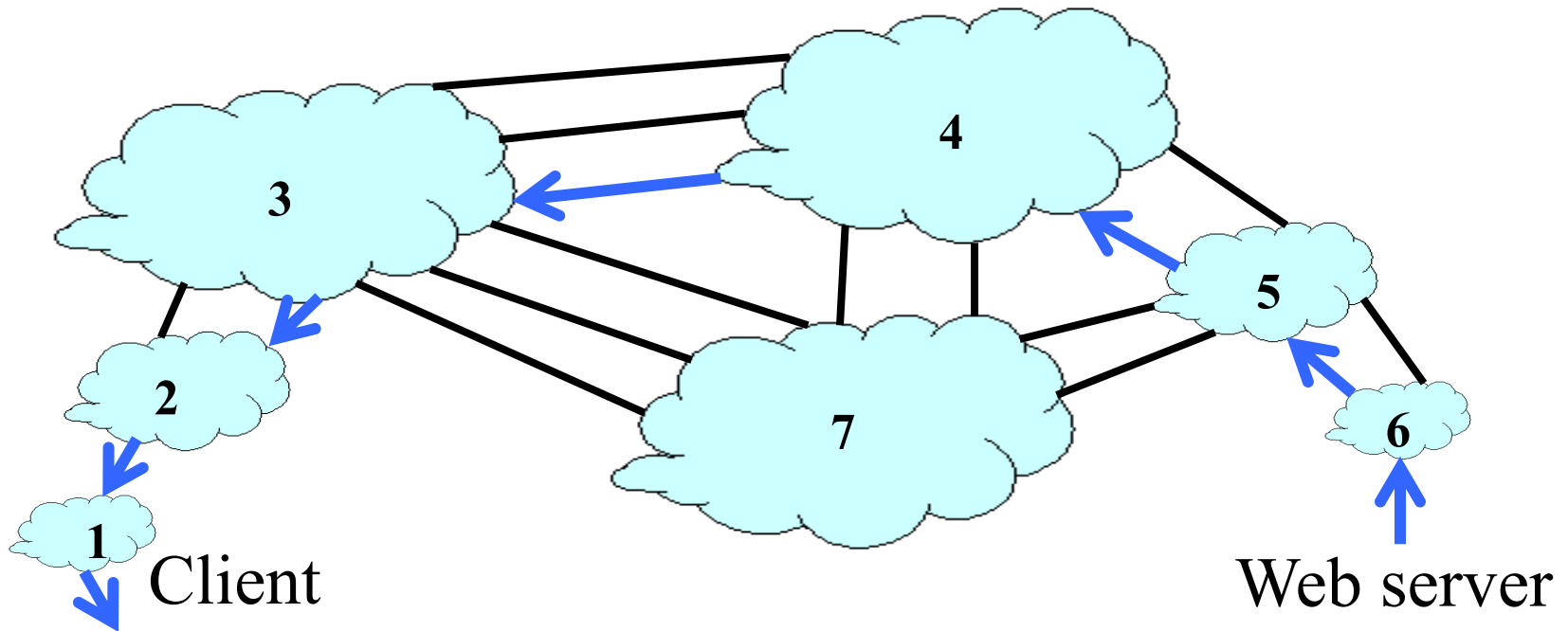
Hierarchy: IP Address Blocks

- **Separation of control**
 - Prefix: assigned to an institution
 - Addresses: assigned by institution to its nodes
- **Who assigns prefixes?**
 - Internet Corporation for Assigned Names & Numbers
 - Regional Internet Registries (RIRs)
 - Internet Service Providers (ISPs)
 - Stub networks
 - Regions within an enterprise

Hierarchy: Routing Protocols

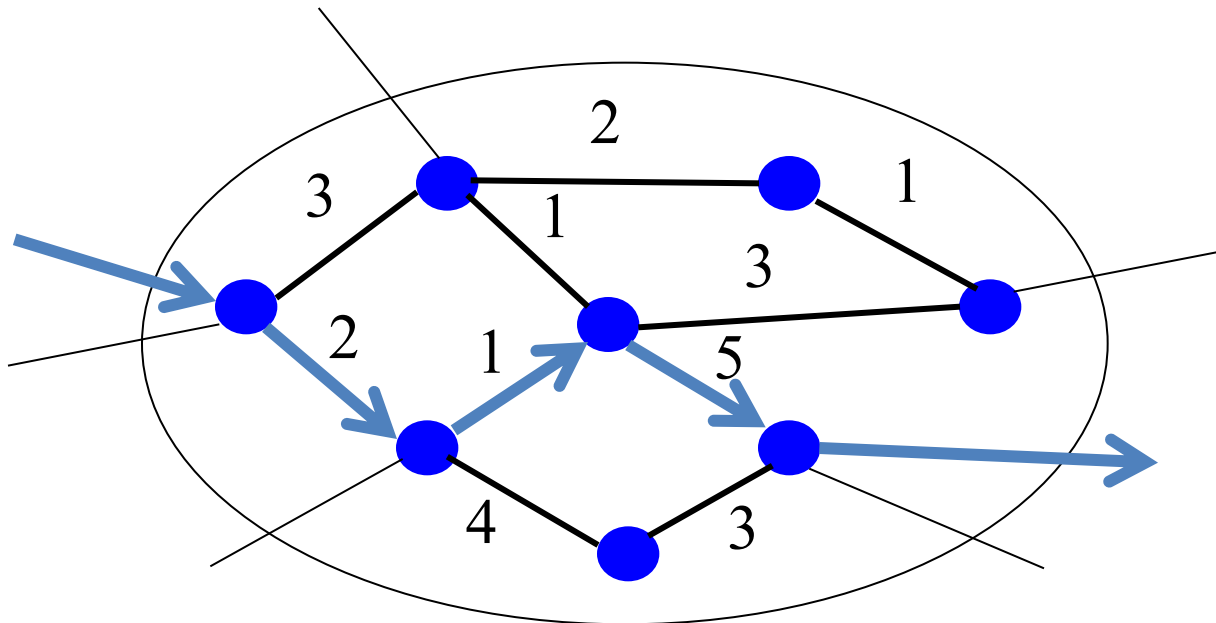
- **AS-level topology**

- Nodes are Autonomous Systems (ASes)
- Edges are links and business relationships
- Hides the detail within each AS's network



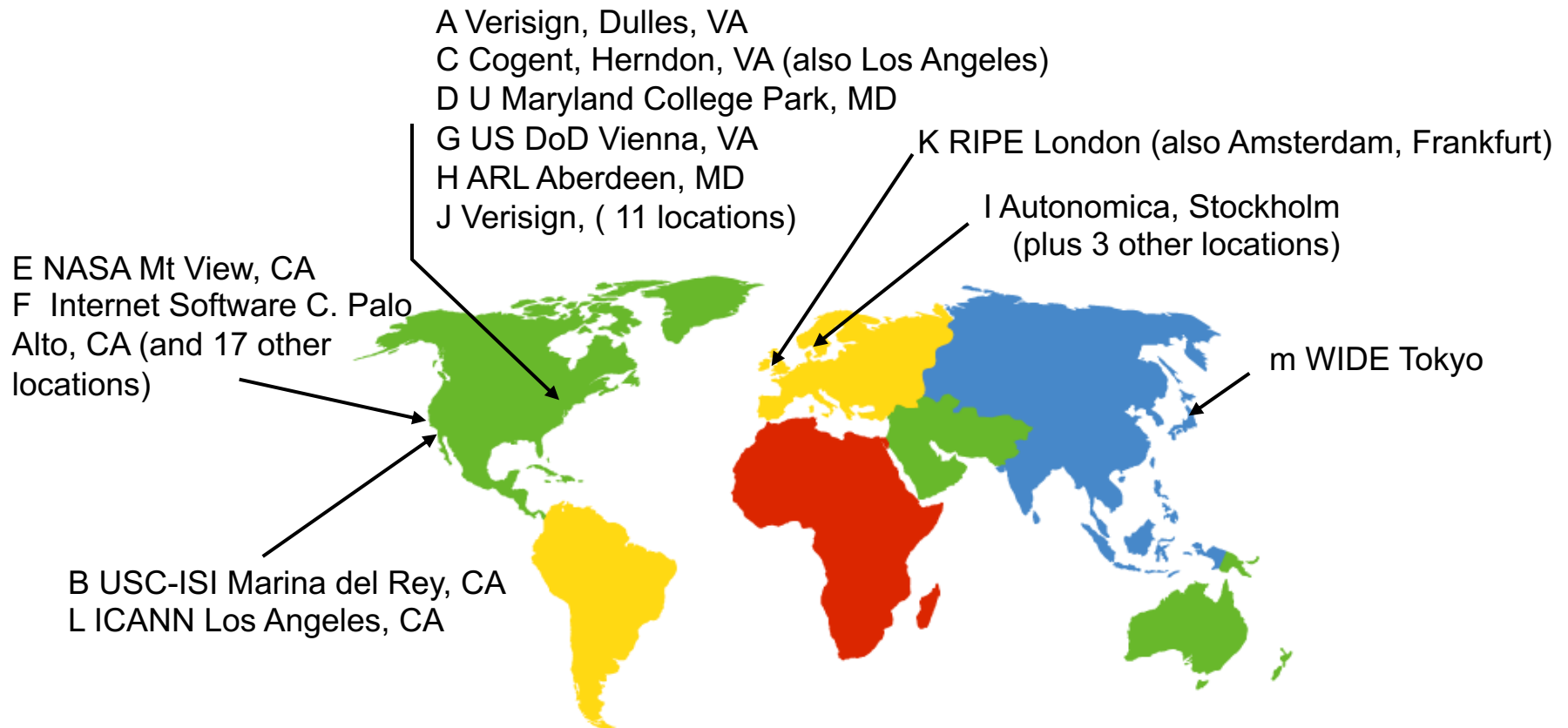
Hierarchy: Routing Protocols

- Interdomain routing ignores details in an AS
 - Routers flood information to learn the topology
 - Routers determine "next hop" to other routers...
 - By computing shortest paths based on link weights

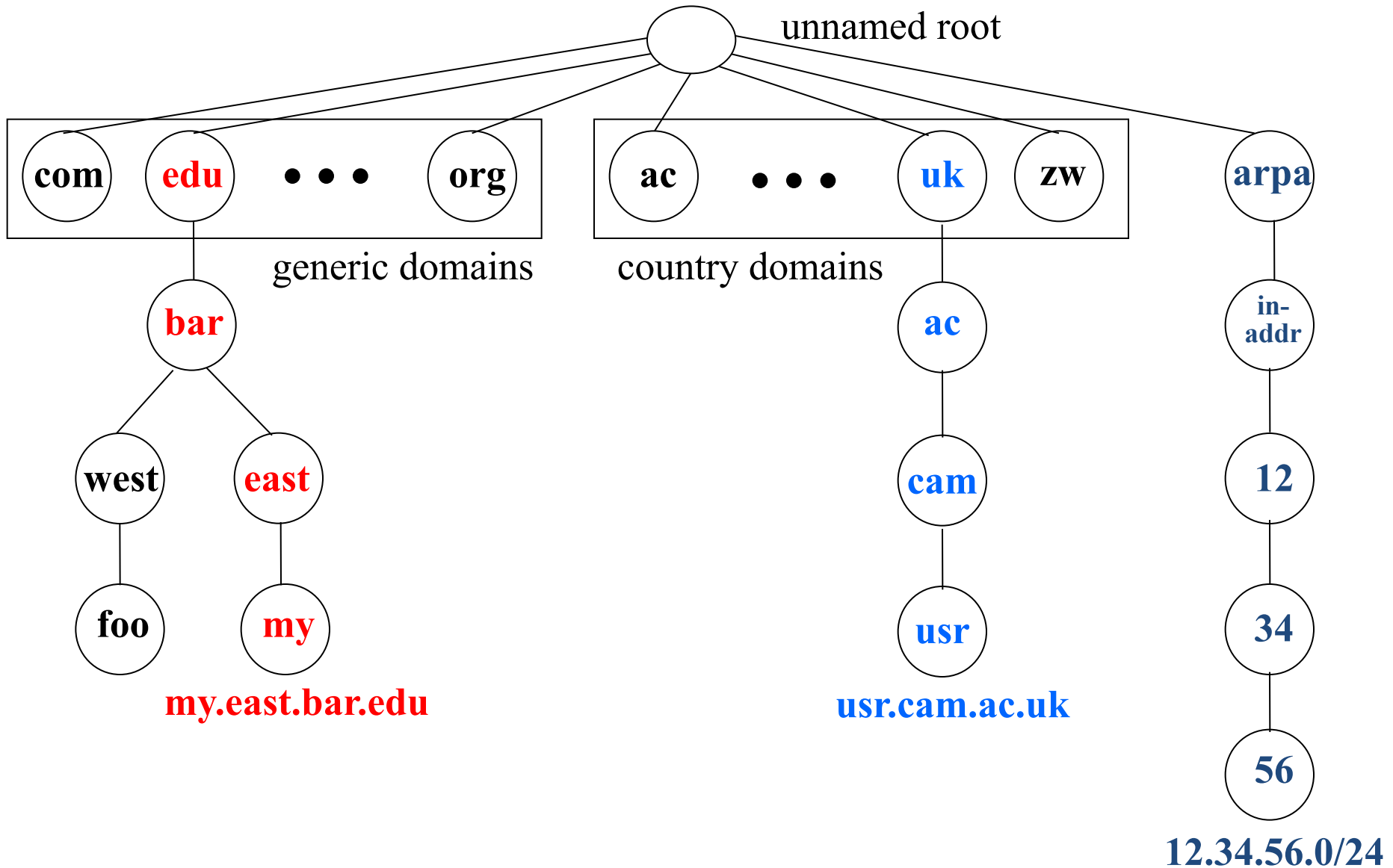


Hierarchy: Domain Name System

- 13 root servers (see <http://www.root-servers.org/>)
- Labeled A through M

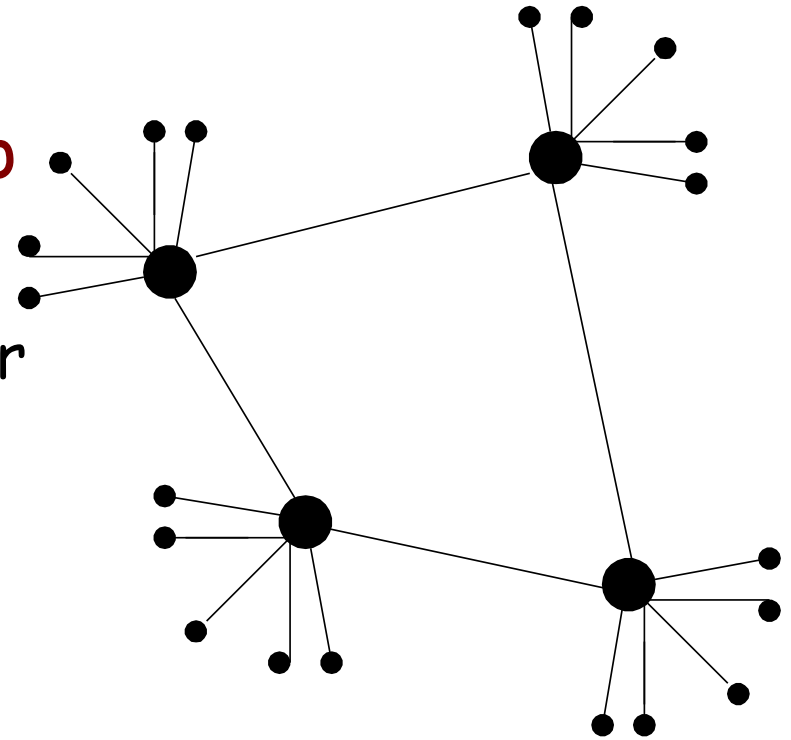


Hierarchy: Domain Name System



Hierarchy: Super Peers in P2P (KaZaA)

- Each peer is either group leader or assigned to group leader
 - TCP connection between peer and its group leader
 - TCP connections between some pairs of group leaders
- Group leader tracks the content in all its children



- ordinary peer
- group-leader peer
- neighboring relationships in overlay network

Indirection

- **Referencing by name**
 - Rather than the value itself
 - E.g., manipulating a variable through a pointer
- **Benefits of indirection**
 - Human convenience
 - Reducing overhead when things change
- **Examples of indirection in the Internet**
 - Names vs. addresses
 - Mobile IP

Indirection: Names vs. Addresses

- **Host name to IP address**
 - Mnemonic names to location-dependent addresses
 - E.g., from `www.cnn.com` to `64.236.16.20`
 - Using the Domain Name System (DNS)
- **From IP address to MAC address**
 - From hierarchical global address to interface card
 - E.g., from `64.236.16.20` to `00-15-C5-49-04-A9`
 - Using the Address Resolution Protocol (ARP)

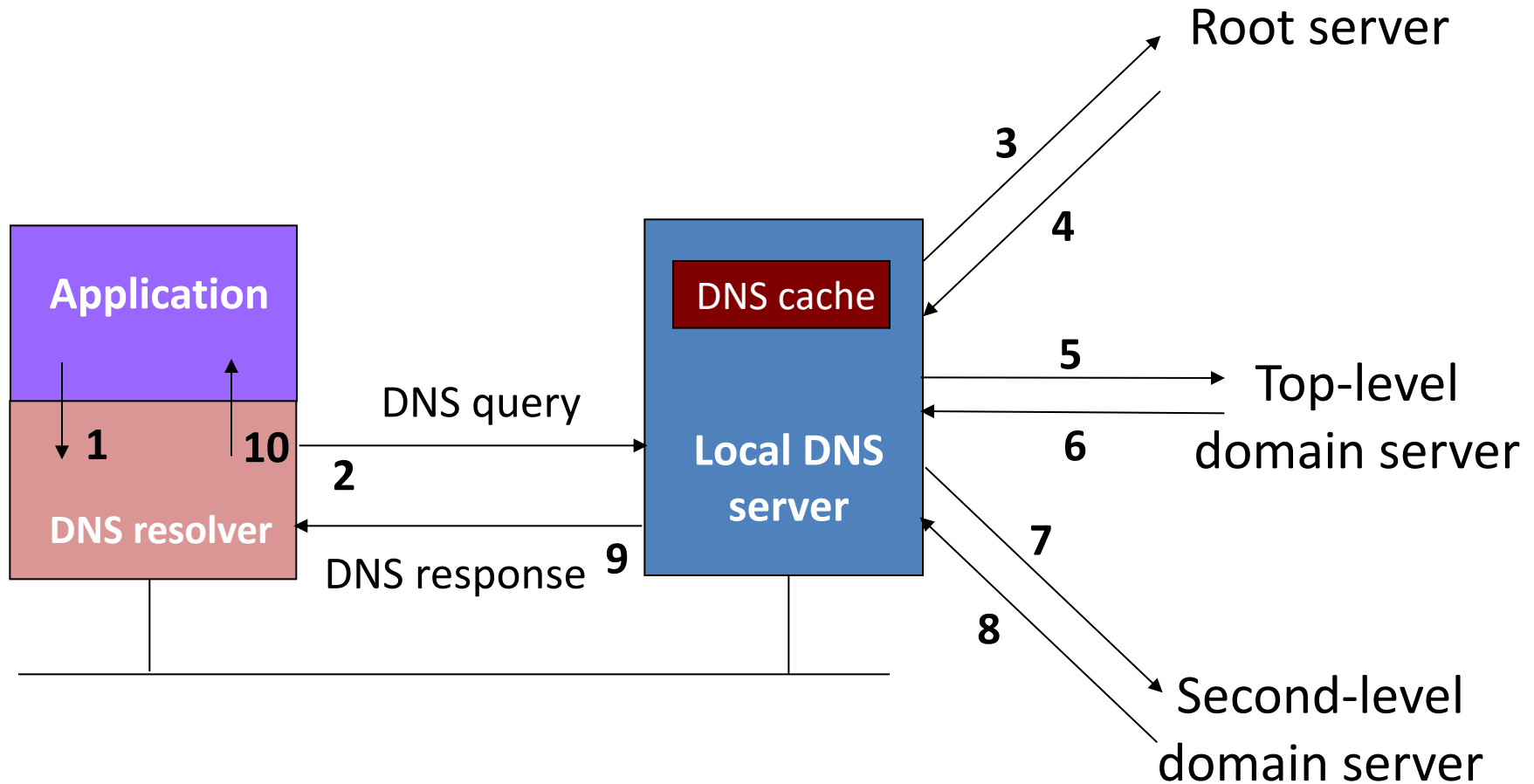
Indirection: Load Balancers & Switches

- Not fixed binding of IPs or MAC address to physical machine
 - NAT allows multiple machines to share single public IP address
 - Load balancers: Machines share IP address, LB maps to physical machine by network flow
 - VM can migrate across L2 network through gratuitous ARP

Caching

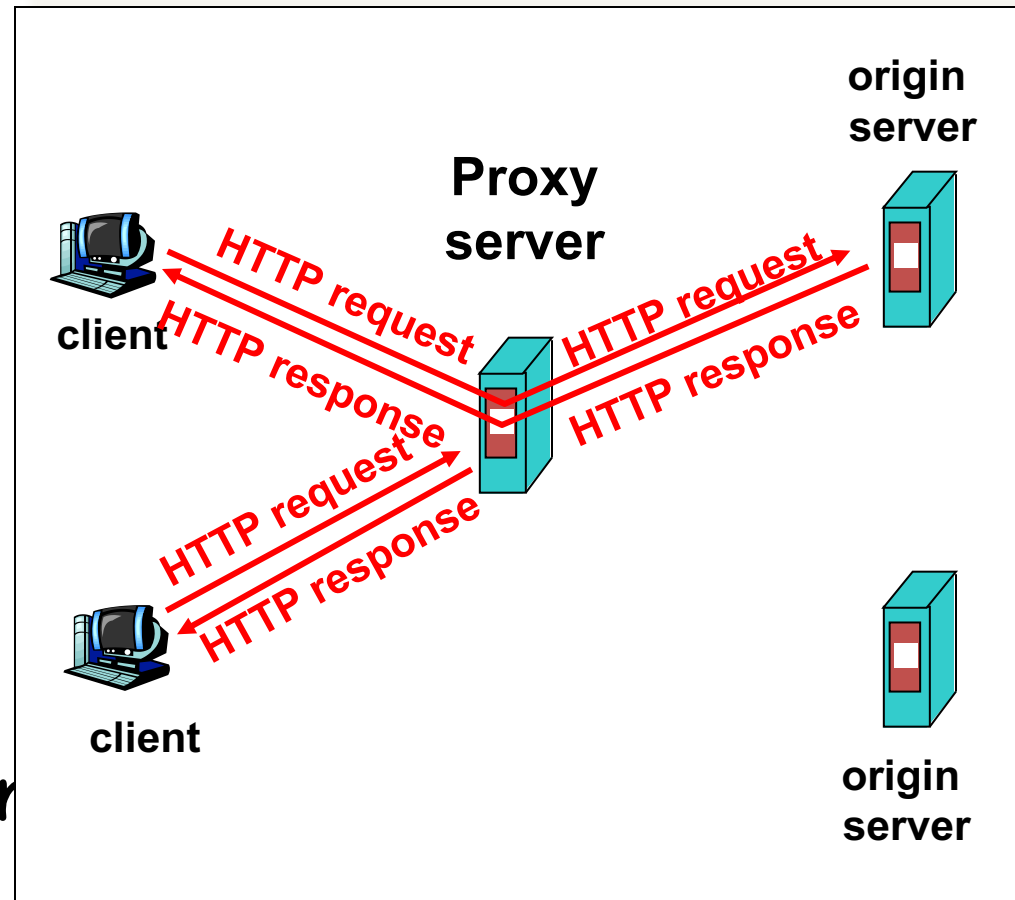
- **Duplicating data stored elsewhere**
 - To reduce latency for accessing the data
 - To reduce resources consumed
- **Caching is often quite effective**
 - Speed difference between cache and primary copy
 - Locality of reference, and small set of popular data
- **Examples from the Internet**
 - DNS caching, Web caching

Caching: DNS Caching



Caching: Web Caching

- **Caching location**
 - Proxy cache
 - Browser cache
- **Better performance**
 - Lower RTT
 - Existing connection
 - Less network load

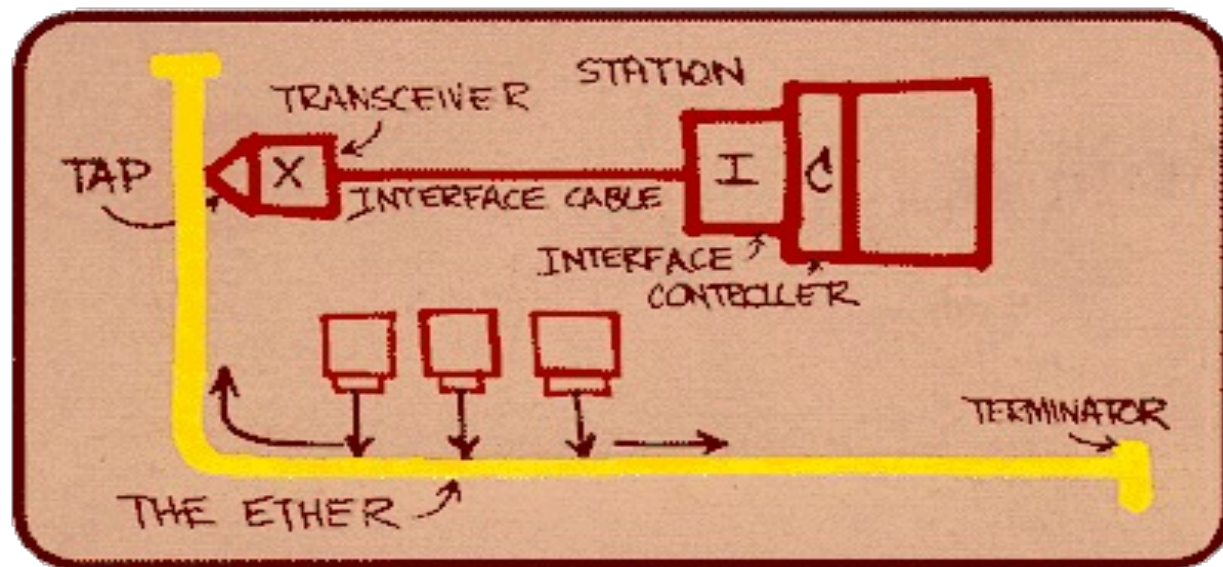


Randomization

- Distributed adaptive algorithms
 - Multiple distributed parties
 - Adapting independently
- Risk of synchronization
 - Many parties reacting at the same time
 - Leading to bad aggregate behavior
- Randomization can desynchronize
 - Ethernet back-off, Random Early Detection
- Rather than imposing centralized control

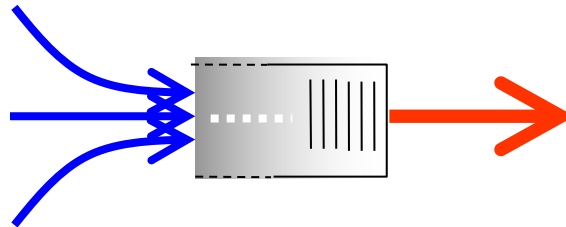
Randomization: Ethernet Back-off

- **Random access: exponential back-off**
 - After collision, wait random time before retrying
 - After m^{th} , choose K randomly from $\{0, \dots, 2^m - 1\}$
 - Wait for $K * 512$ bit times before trying again



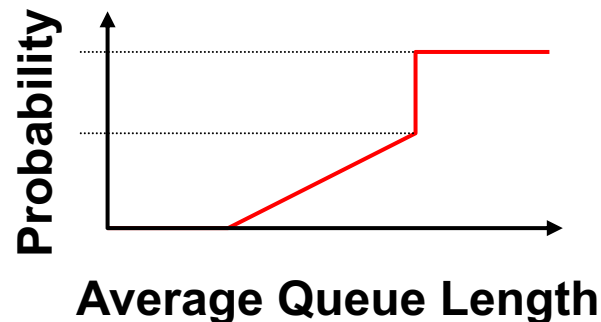
Randomization: Dropping Packets Early

- Congestion on a link
 - Eventually the queue becomes full
 - And new packets must be dropped
- Drop-tail queuing leads to bursty loss
 - Many packets encounter a full queue
 - Many TCP senders reduce their sending rates



Randomization: Dropping Packets Early

- **Better to give early feedback**
 - Get a few connections to slow down
 - ... before it is too late
- **Random Early Detection (RED)**
 - Randomly drop packets when queue (near) full
 - Drop rate increases as function of queue length



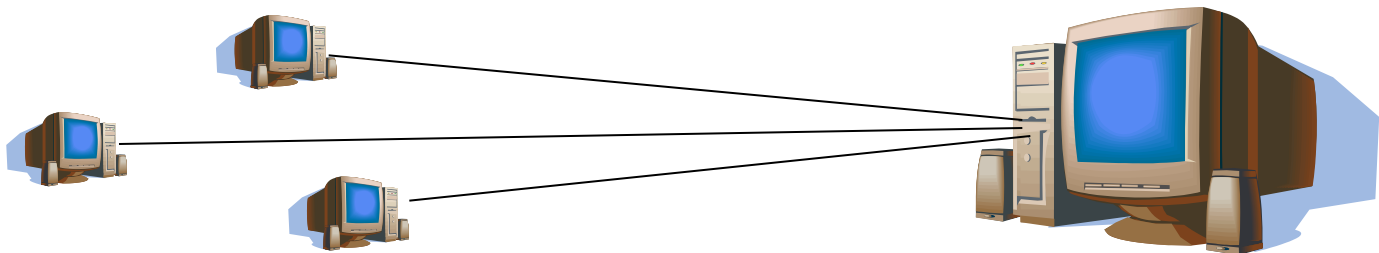
Soft State

- **State: stored in nodes by network protocols**
 - Installed by receiver of a set-up message
 - Updated when conditions change
- **Hard state: valid unless told otherwise**
 - Removed by receiver of tear-down message
 - Requires error handling to deal with sender failure
- **Soft state: invalid if not told to refresh**
 - Periodically refreshed, removed by timeout
- **Soft state reduces complexity**
 - DNS caching, DHCP leases



Soft State: DNS Caching

- **Cache consistency is a hard problem**
 - Ensuring the cached copy is not out of date
- **Strawman: explicit revocation or updates**
 - Keep track of everyone who has cached information
 - If name-to-host mapping changes, update caches
- **Soft state solution**
 - DNS responses include a "time to live" (TTL) field
 - Cached entry is deleted after TTL expires

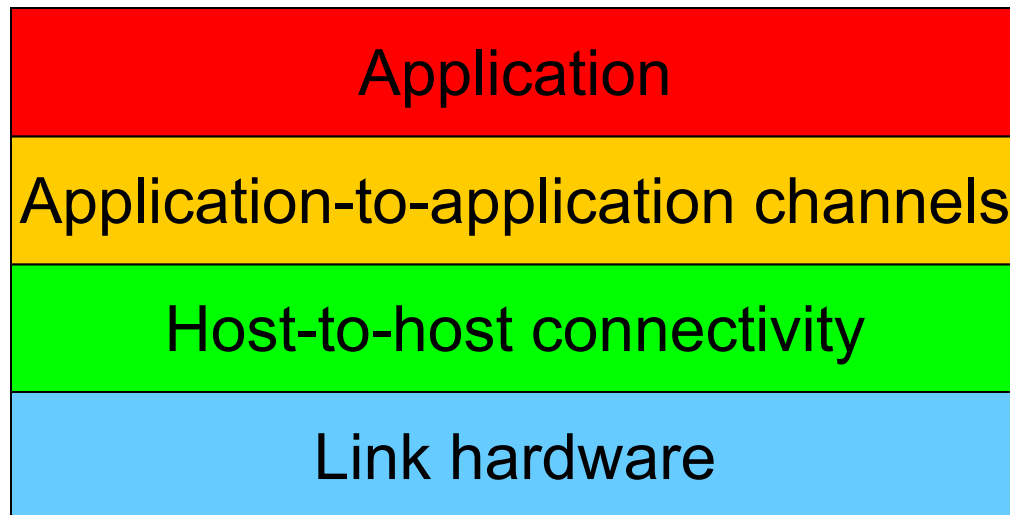


Soft State: DHCP Leases

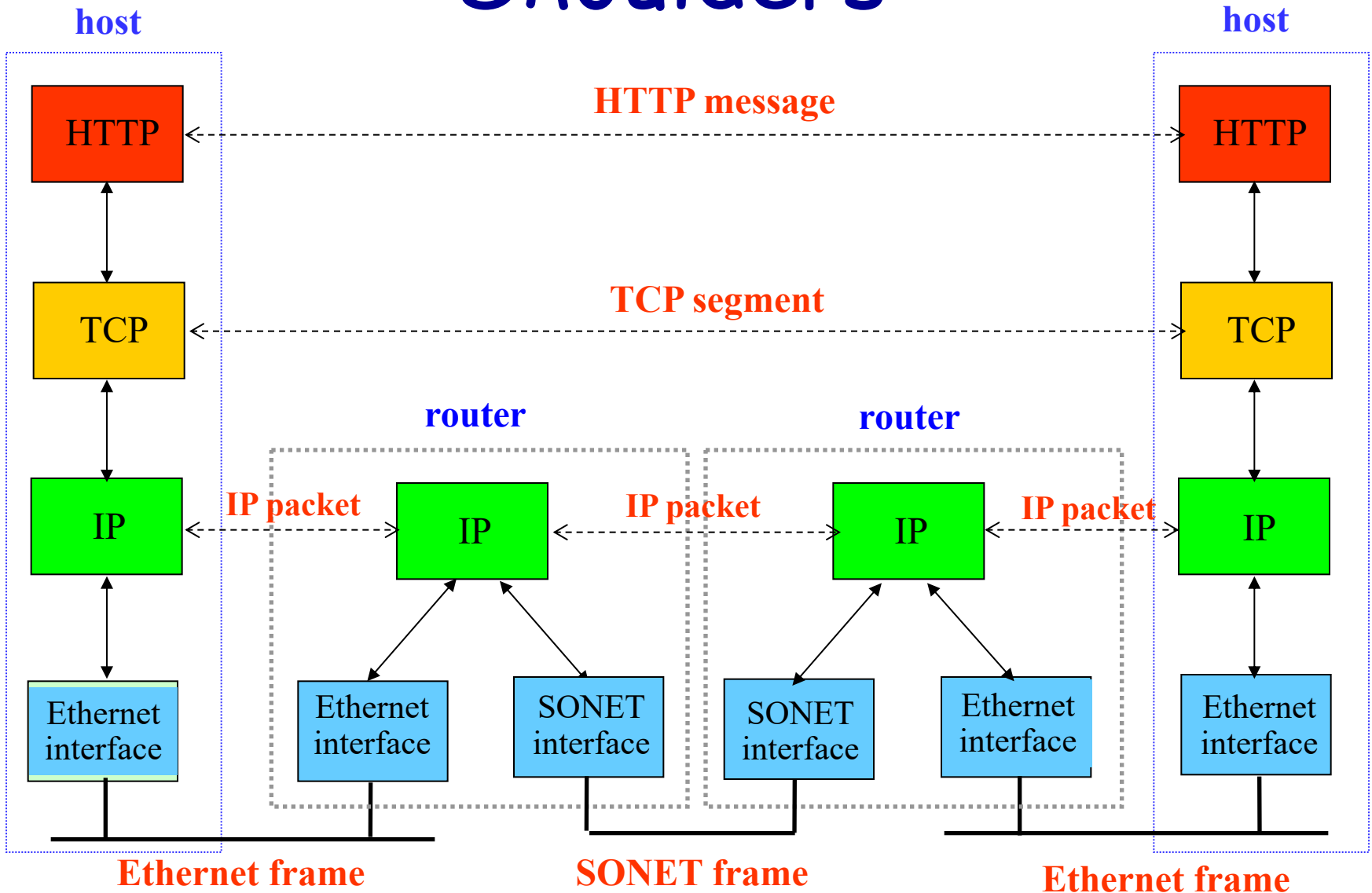
- DHCP "offer message" from the server
 - Configuration parameters (proposed IP address, mask, gateway router, DNS server, ...)
 - Lease time (the time information remains valid)
- Why is a lease time necessary?
 - Client can release address (DHCP RELEASE)
 - E.g., "ipconfig /release" or clean shutdown of computer
 - But, the host might not release the address
 - E.g., the host crashes or buggy client software
 - You don't want address to be allocated forever

Layering: A Modular Approach

- **Sub-divide the problem**
 - Each layer relies on services from layer below
 - Each layer exports services to layer above
- **Interface between layers defines interaction**
 - Hides implementation details
 - Layers can change without disturbing other layers

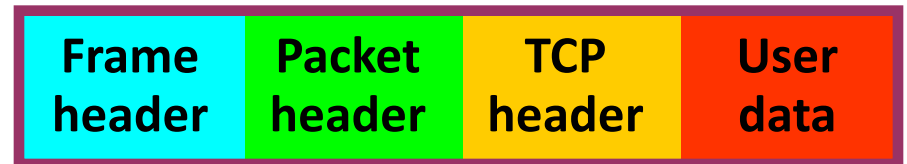
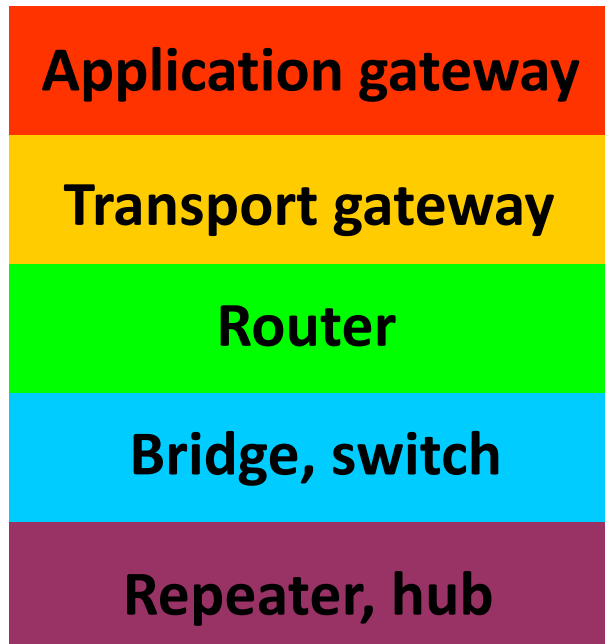


Layering: Standing on Shoulders



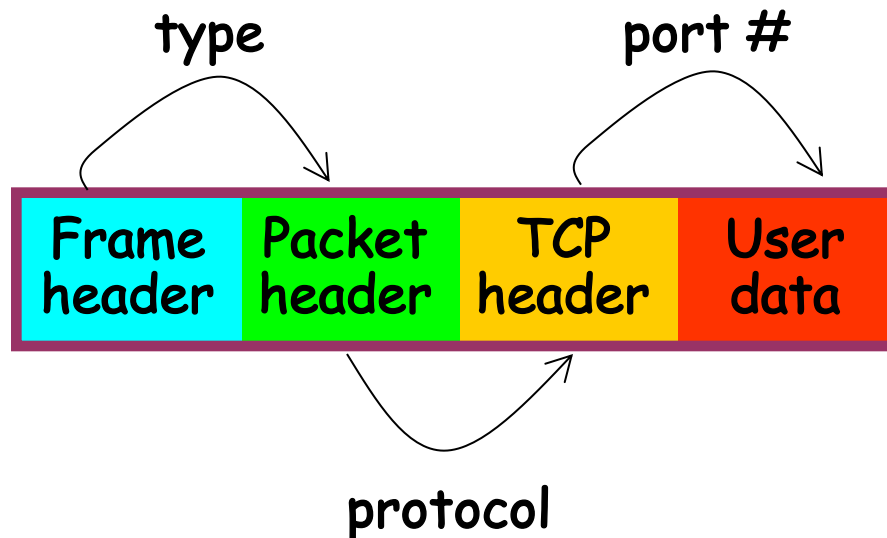
Layering: Encapsulation of Data

- Different devices switch different things
 - Physical layer: electrical signals (repeaters and hubs)
 - Link layer: frames (bridges and switches)
 - Network layer: packets (routers)

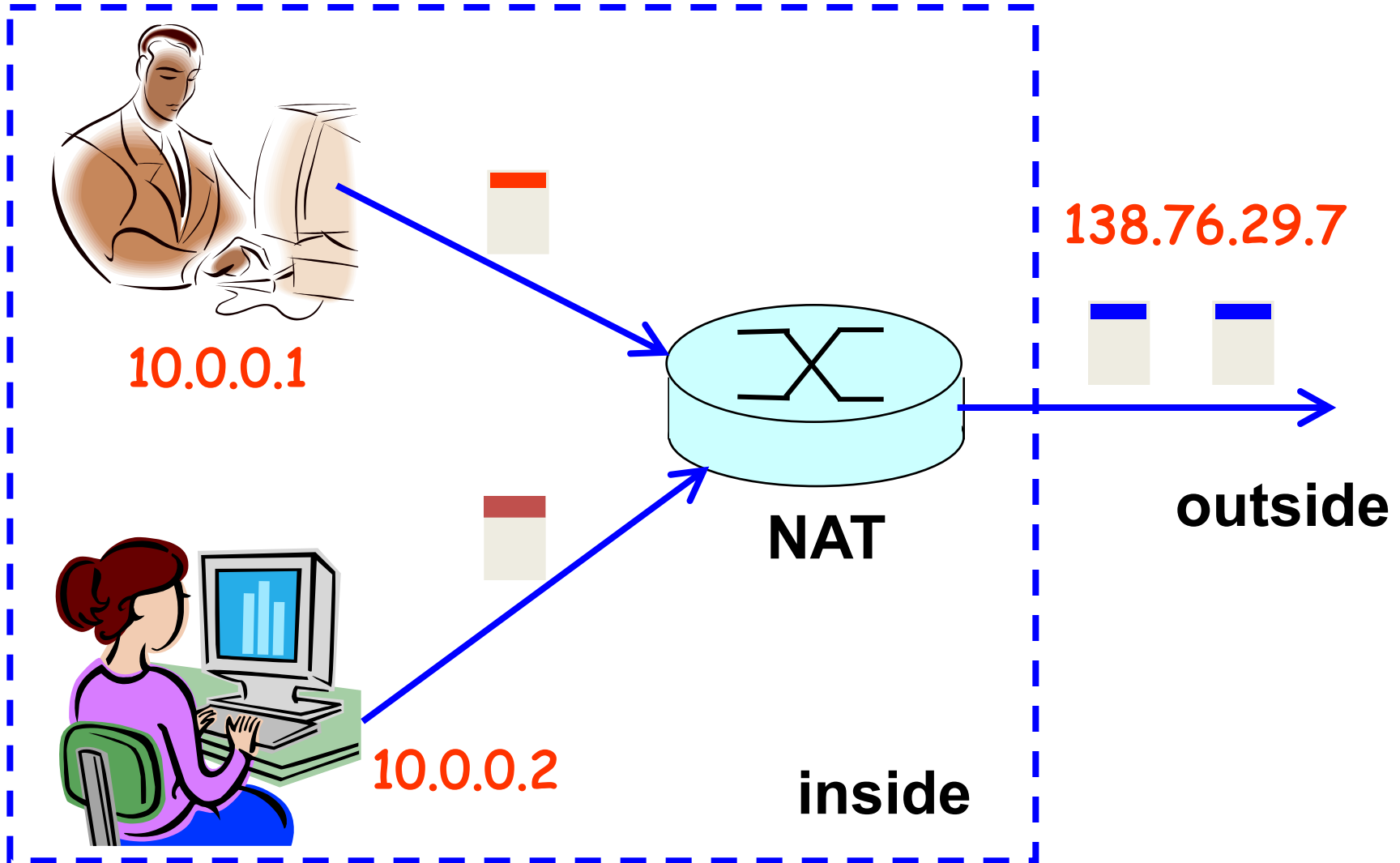


Demultiplexing

- Separating multiple streams out of one
 - Recognizing the separate streams
 - Treating the separate streams accordingly
- Examples in the Internet



(De)multiplexing: With a NAT



Power at the End Host

End-to-End Principle

Whenever possible, communications protocol operations should be defined to occur at the **end-points** of a communications system.

Programmability

With programmable end hosts, new network services can be added at **any time, by anyone**.

What Will Happen to the Internet

No Strict Notions of Identity



"On the Internet, nobody knows you're a dog."

- Leads to
 - Spam
 - Spoofing
 - Denial-of-service
 - Route hijacking

Protocols Designed Based on Trust

- That you don't spoof your addresses
 - MAC spoofing, IP address spoofing, spam, ...
- That port numbers correspond to applications
 - Rather than being arbitrary, meaningless numbers
- That you adhere to the protocol
 - Ethernet exponential back-off after a collision
 - TCP additive increase, multiplicative decrease
- That protocol specifications are public
 - So others can build interoperable implementations

Nobody in Charge

- Traffic traverses many Autonomous Systems
 - Who's fault is it when things go wrong?
 - How do you upgrade functionality?
- Implicit trust in the end host
 - What if some hosts violate congestion control?
- Anyone can add any application
 - Whether or not it is legal, moral, good, etc.
- Spans many countries
 - So no one government can be in charge

Challenging New Requirements

- Disseminating data
- Mobile, multi-homed hosts
- Sometimes-connected hosts
- Large number of hosts
- Real-time applications

The Internet of the Future

- Can we fix what ails the Internet
 - Security, performance, reliability
 - Upgradability, managability
 - <Your favorite gripe here>
- Without throwing out baby with bathwater
 - Ease of adding new hosts
 - Ease of adding new services
 - Ease of adding new link technologies
- An open technical and policy question...