

Dynamo / Bayou

Feb 23rd & 24th, 2022

[Adapted from Andrew Or's]

Some context...

Dynamo and Bayou both offer **high availability** and **weak consistency**

Most traditional databases offer **strong consistency** and **low availability**
Not suitable for modern applications with super high demands

What are some example applications of each?

- Flight ticket booking (HA)

- Amazon shopping carts (HA)

- Offline edits (HA)

- Billing services (SC)

- Bank accounts (SC)

Availability is important



Tens of millions of customers at peak times

Tens of millions of shopping cart requests, 3 million checkouts per day

Hundreds of thousands of concurrently active sessions

Strict Service-Level Agreements (SLAs) translate to business value

Dynamo

Fully decentralized, highly available key-value store

Always writeable, resolve conflicts during reads --- Eventual Consistency

API for clients to specify requirements (99.9th percentile)

Departure from RDBMS: simpler functionality, fewer guarantees, runs on commodity hardware (low-end, broadly compatible, non-specialized machines)

(Read the [original paper](#), especially Section 4)

Techniques for achieving availability

Consistent hashing for partitioning key space

Vector clocks for reconciling conflicts during reads

Sloppy quorums for handling temporary failures

Anti-entropy using Merkle trees for syncing key-value pairs

Gossip-based protocol for membership notifications

Table 1: Summary of techniques used in *Dynamo* and their advantages.

Problem	Technique	Advantage
Partitioning	Consistent Hashing	Incremental Scalability
High Availability for writes	Vector clocks with reconciliation during reads	Version size is decoupled from update rates.
Handling temporary failures	Sloppy Quorum and hinted handoff	Provides high availability and durability guarantee when some of the replicas are not available.
Recovering from permanent failures	Anti-entropy using Merkle trees	Synchronizes divergent replicas in the background.
Membership and failure detection	Gossip-based membership protocol and failure detection.	Preserves symmetry and avoids having a centralized registry for storing membership and node liveness information.

Techniques for achieving availability

Consistent hashing for partitioning key space

Vector clocks for reconciling conflicts during reads

Sloppy quorums for handling temporary failures

Anti-entropy using Merkle trees for syncing key-value pairs

Gossip-based protocol for membership notifications

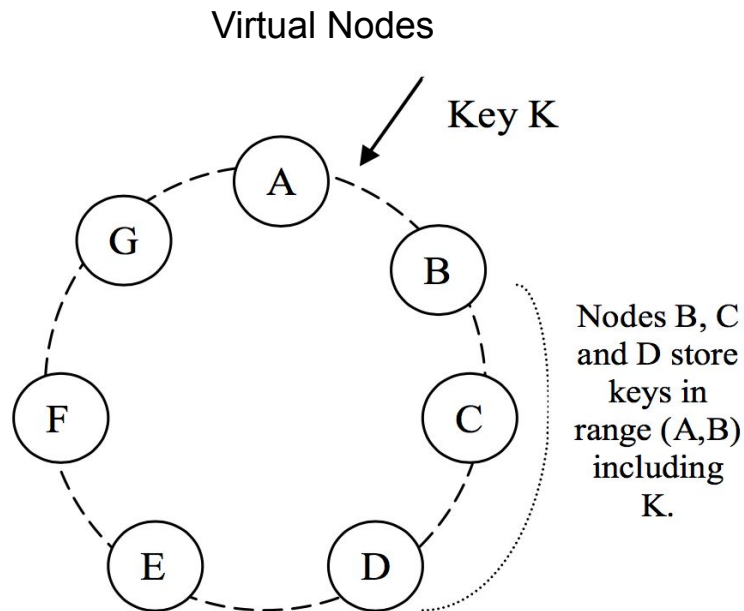
Consistent Hashing

Assign each node a random position on the ring

Node owns the preceding key range

For fault tolerance, replicate each key at N successor nodes in the ring

Virtual nodes: each physical node gets assigned multiple nodes on the ring (e.g. B, D, F)



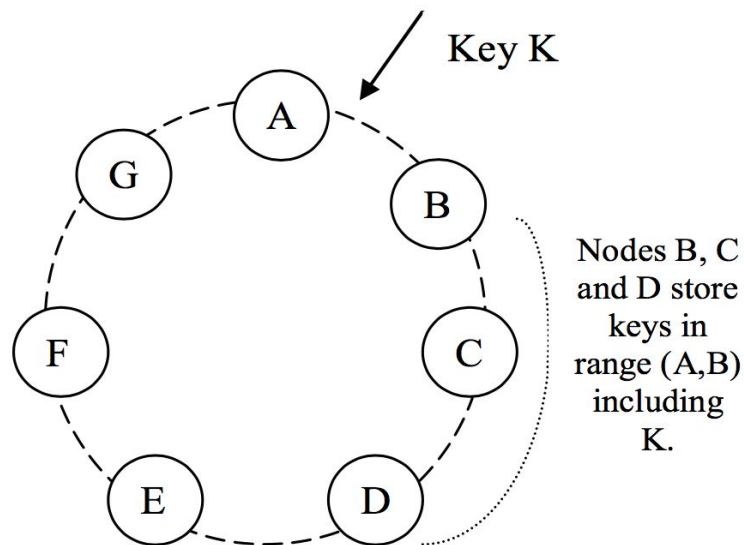
Consistent Hashing

Desirable properties?

Uniform distribution of load

Minimum object movements when nodes join or leave the ring

Number of virtual nodes can be adjusted for device heterogeneity



Techniques for achieving availability

Consistent hashing for partitioning key space

Vector clocks for reconciling conflicts during reads

Sloppy quorums for handling temporary failures

Anti-entropy using Merkle trees for syncing key-value pairs

Gossip-based protocol for membership notifications

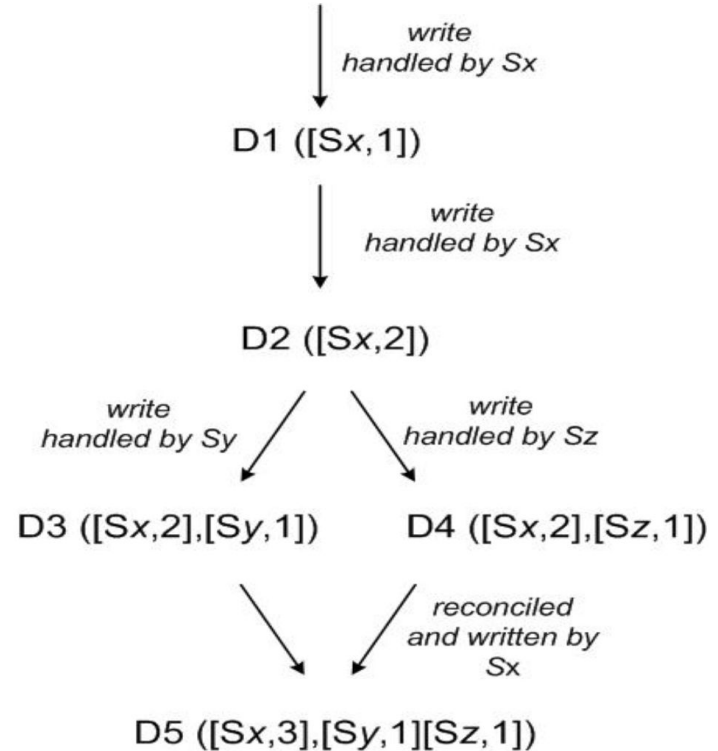
Conflict resolution

Two machines write different values to the same key

Vector clocks: list of (node, count) pairs where count is incremented on write

If one vector clock subsumes another, discard older value

Else, return all conflicting values to client



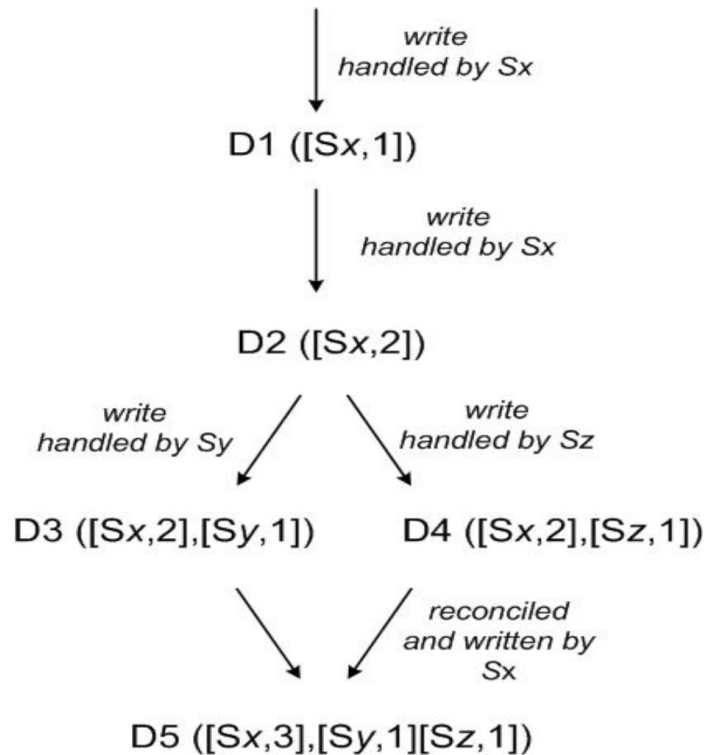
Context contains vector clocks

Dynamo client API is simple:

get(key) (value, context)

put(key, value, context)

Common pattern: put after get



Shopping Cart



Cat-Opoly by LatefortheSky

In Stock

✓Prime

This is a gift [Learn more](#)

[Delete](#) | [Save for later](#)



Fancy Feast Wet Cat Food, Grilled, Seafood Feast Variety Pack, 3-Ounce Can, Pack of 24 by Purina Fancy Feast

In stock. Usually ships within 3 to 4 days.

Shipped from: [Connect Buy](#)

Gift options not available. [Learn more](#)

[Delete](#) | [Save for later](#)

Shopping Cart



Furhaven Orthopedic Mattress Pet Bed, Large, Chocolate, for Dogs and Cats by Furhaven Pet

In Stock

✓Prime

This is a gift [Learn more](#)

[Delete](#) | [Save for later](#)

Shopping Cart



Cat-Opoly by LatefortheSky

In Stock

✓Prime

This is a gift [Learn more](#)

[Delete](#) | [Save for later](#)



Fancy Feast Wet Cat Food, Grilled, Seafood Feast Variety Pack, 3-Ounce Can, Pack of 24 by Purina Fancy Feast

In stock. Usually ships within 3 to 4 days.

Shipped from: [Connect Buy](#)

Gift options not available. [Learn more](#)

[Delete](#) | [Save for later](#)



Furhaven Orthopedic Mattress Pet Bed, Large, Chocolate, for Dogs and Cats by Furhaven Pet

In Stock

✓Prime

This is a gift [Learn more](#)

[Delete](#) | [Save for later](#)

Shopping Cart



Cat-Opoly by LatefortheSky

In Stock

Prime

This is a gift [Learn more](#)

[Delete](#) | [Save for later](#)



Fancy Feast Wet Cat Food, Grilled, Seafood Feast Variety Pack, 3-Ounce Can, Pack of 24 by Purina Fancy Feast

In stock. Usually ships within 4 days.

Shipped from: [Connect Buy](#)
Gift options not available. [Learn more](#)

[Delete](#) | [Save for later](#)

DELETE

Shopping Cart



Fancy Feast Wet Cat Food, Grilled, Seafood Feast Variety Pack, 3-Ounce Can, Pack of 24 by Purina Fancy Feast

In stock. Usually ships within 3 to 4 days.

Shipped from: [Connect Buy](#)
Gift options not available. [Learn more](#)

[Delete](#) | [Save for later](#)



Furhaven Orthopedic Mattress Pet Bed, Large, Chocolate, for Dogs and Cats by Furhaven Pet

In Stock

Prime

This is a gift [Learn more](#)

[Delete](#) | [Save for later](#)

Shopping Cart



Cat-Opoly by LatefortheSky

In Stock

Prime

This is a gift [Learn more](#)

[Delete](#) | [Save for later](#)



Fancy Feast Wet Cat Food, Grilled, Seafood Feast Variety Pack, 3-Ounce Can, Pack of 24 by Purina Fancy Feast

In stock. Usually ships within 3 to 4 days.

Shipped from: [Connect Buy](#)
Gift options not available. [Learn more](#)

[Delete](#) | [Save for later](#)



Furhaven Orthopedic Mattress Pet Bed, Large, Chocolate, for Dogs and Cats by Furhaven Pet

In Stock

Prime

This is a gift [Learn more](#)

[Delete](#) | [Save for later](#)

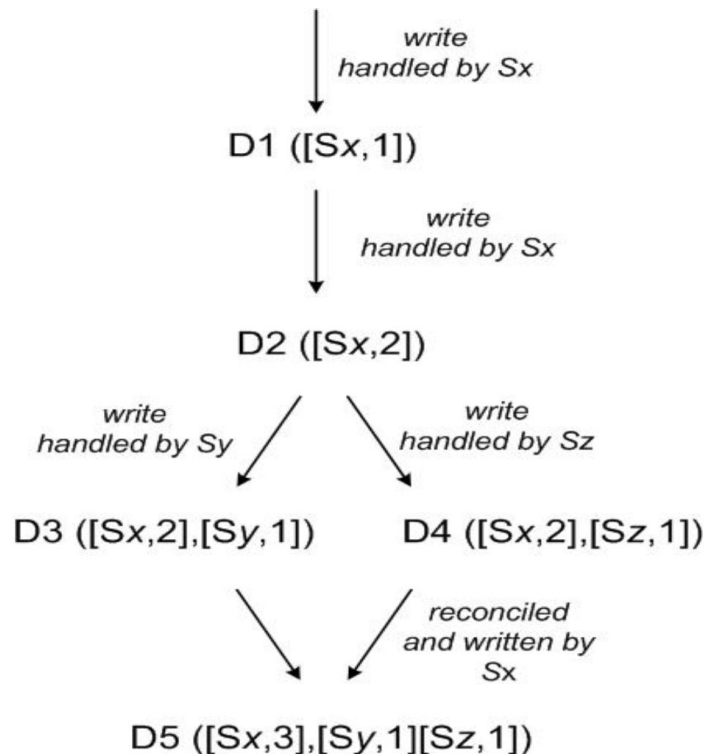
Conflict resolution

Two machines write different values to the same key

Vector clocks: list of (node, count) pairs where count is incremented on write

If one vector clock subsumes another, discard older value

Else, return all conflicting values to client



Techniques for achieving availability

Consistent hashing for partitioning key space

Vector clocks for reconciling conflicts during reads

Sloppy quorums for handling temporary failures

Anti-entropy using Merkle trees for syncing key-value pairs

Gossip-based protocol for membership notifications

Sloppy Quorums

Write to N nodes, return success when $W < N$ nodes respond

Read from N nodes, return value(s) from $R < N$ nodes

Typically, $W+R > N$ means at least one writer and one reader overlap, so values are consistent

Sloppy here means skip nodes that have failed, such that even if $W+R > N$, the readers and writers may not overlap = **not consistent!**

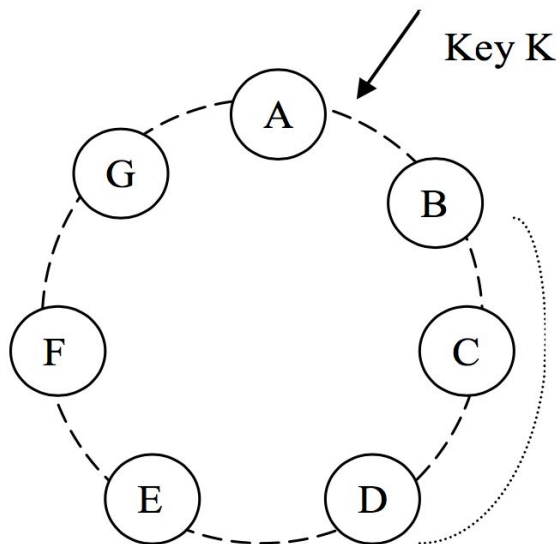
Sloppy Quorums

Example:

Typical values are $N = 3$, $W = R = 2$

Nodes C and D have failed, so key k is written to E and F instead

Nodes C and D recover, and now client tries to read from C and D = **stale value**



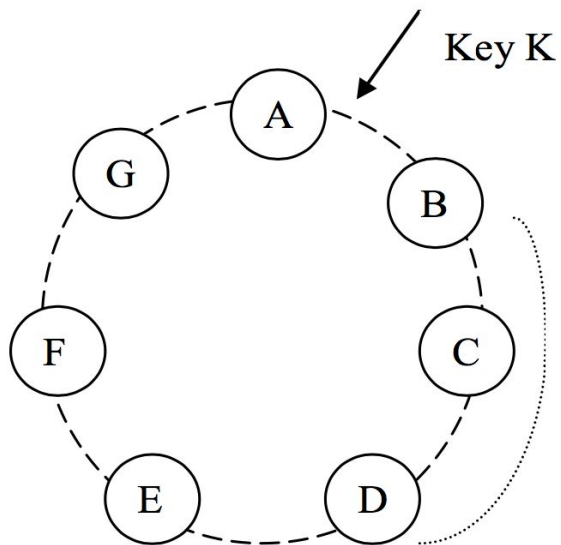
Hinted Handoff

“Hint” refers to the node the data originally belongs to

Example:

Nodes E and F remember they are writing on behalf of C and D

As soon as C and D recovers, E and F transfer their values for k to C and D



Sloppy Quorums

Write to N nodes, return success when $W < N$ nodes respond

Read from N nodes, return value(s) from $R < N$ nodes

Typically, $W+R > N$ means at least one writer and one reader overlap, so values are consistent

Sloppy here means skip nodes that have failed, such that even if $W+R > N$, the readers and writers may not overlap = **not consistent!**

Techniques for achieving availability

Consistent hashing for partitioning key space

Vector clocks for reconciling conflicts during reads

Sloppy quorums for handling temporary failures

Anti-entropy using Merkle trees for syncing key-value pairs

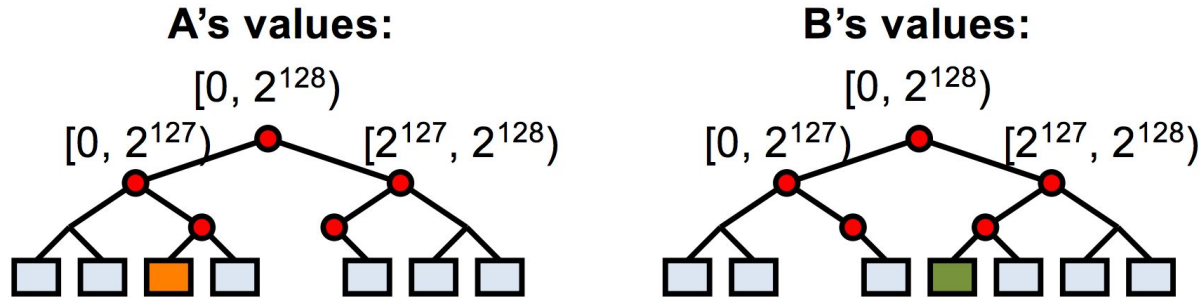
Gossip-based protocol for membership notifications

Anti-entropy using Merkle trees

Goal: minimize durability loss from above techniques

Nodes responsible for the same key spaces exchange Merkle trees

Find differences quickly while exchanging little information



Techniques for achieving availability

Consistent hashing for partitioning key space

Vector clocks for reconciling conflicts during reads

Sloppy quorums for handling temporary failures

Anti-entropy using Merkle trees for syncing key-value pairs

Gossip-based protocol for membership notifications

Membership notification

Gossip-based protocol to propagate membership changes

Each node learns the key spaces handled by all other nodes

Result: zero-hop distributed hash table (DHT)

Clearly not infinitely scalable, but storage requirement not a problem in practice

Bayou

What is it?

- Weakly consistent, replicated storage system

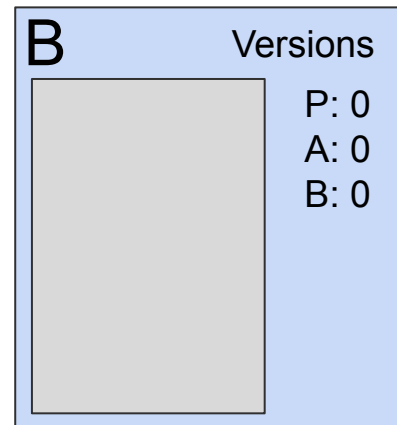
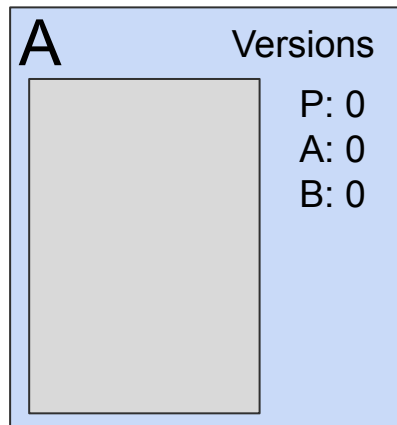
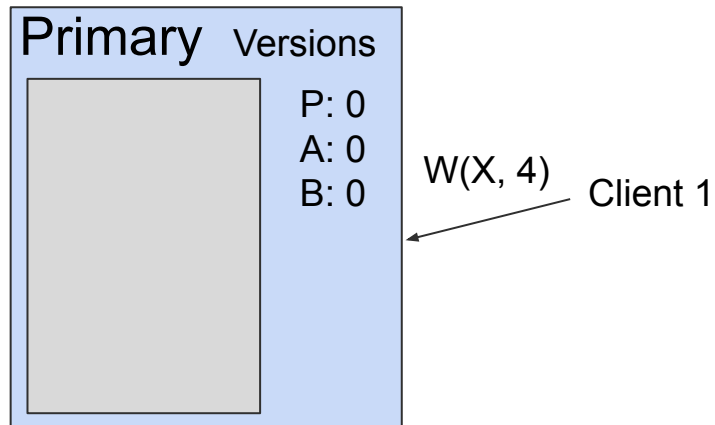
Goals:

- Maximize availability, *support offline collaboration*
- Minimize network communication
- Agree on all values (eventually)

Bayou Writes

Legend

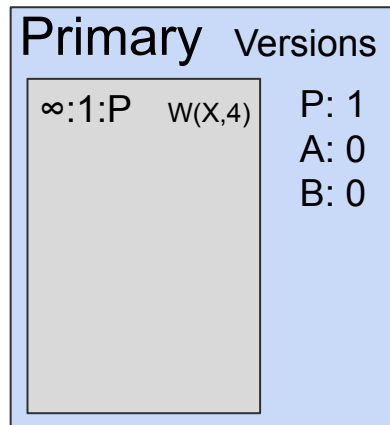
Commit Timestamp:Write Timestamp:Write Server



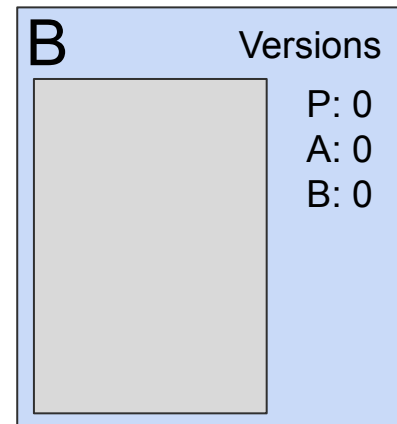
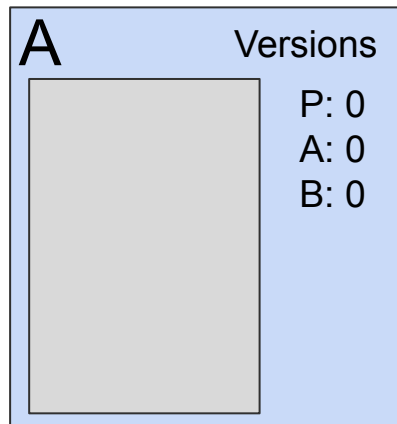
Bayou Writes

Legend

Commit Timestamp:Write Timestamp:Write Server



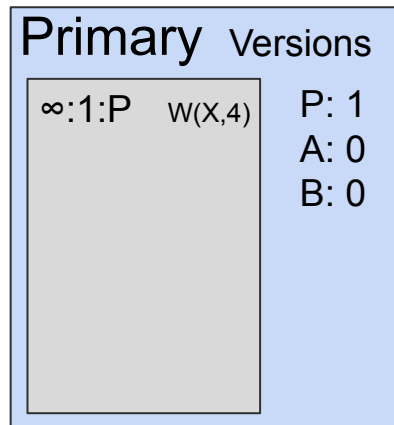
Client 1



Bayou Writes

Legend

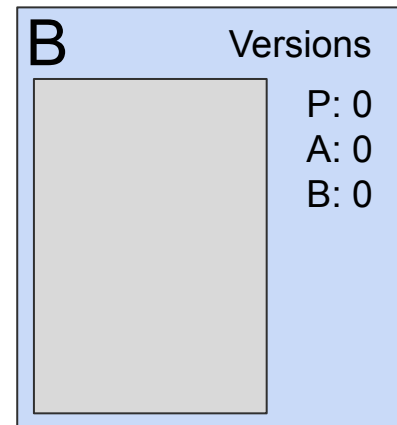
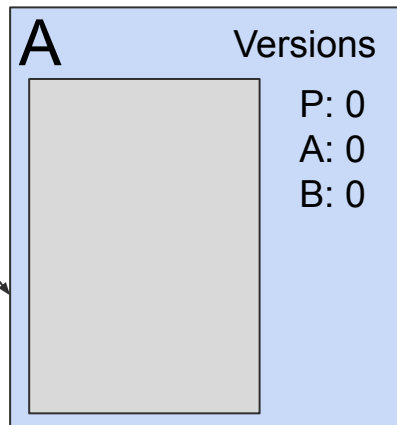
Commit Timestamp:Write Timestamp:Write Server



Client 1

W(Y, 8)

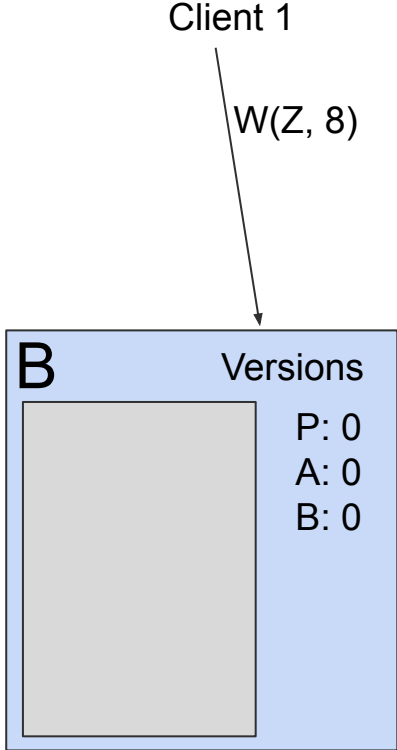
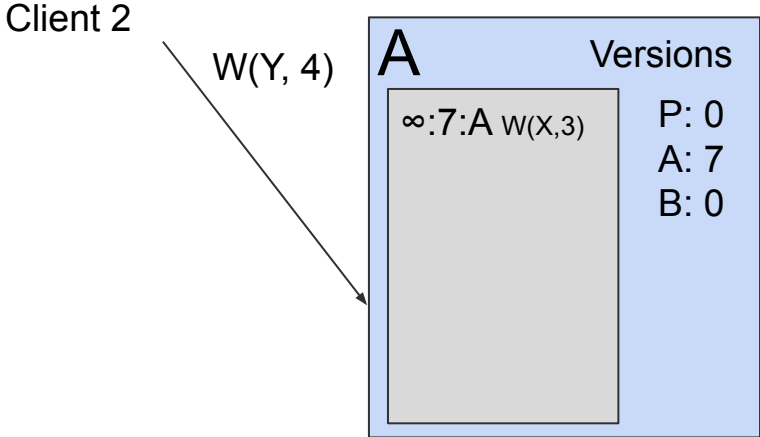
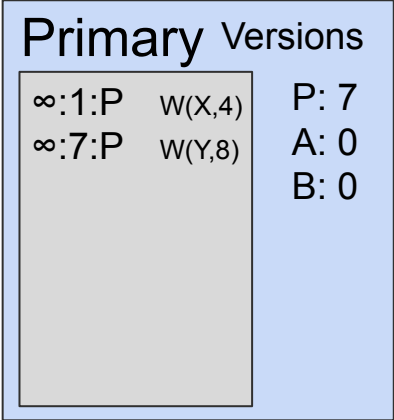
Client 2
W(X, 3)



Bayou Writes

Legend

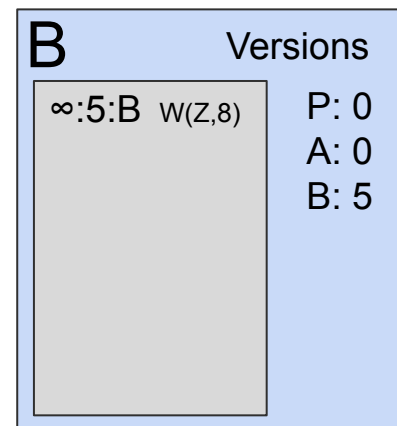
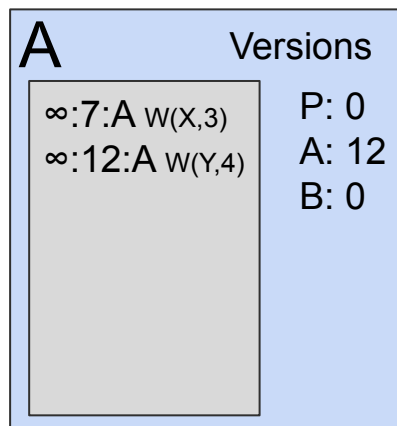
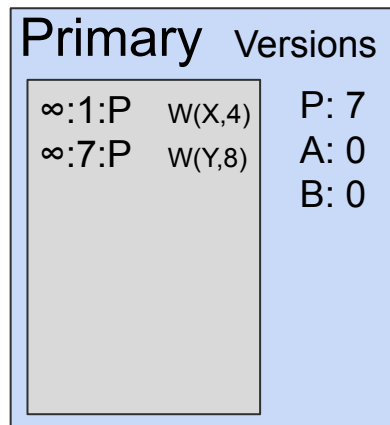
Commit Timestamp:Write Timestamp:Write Server



Bayou Writes

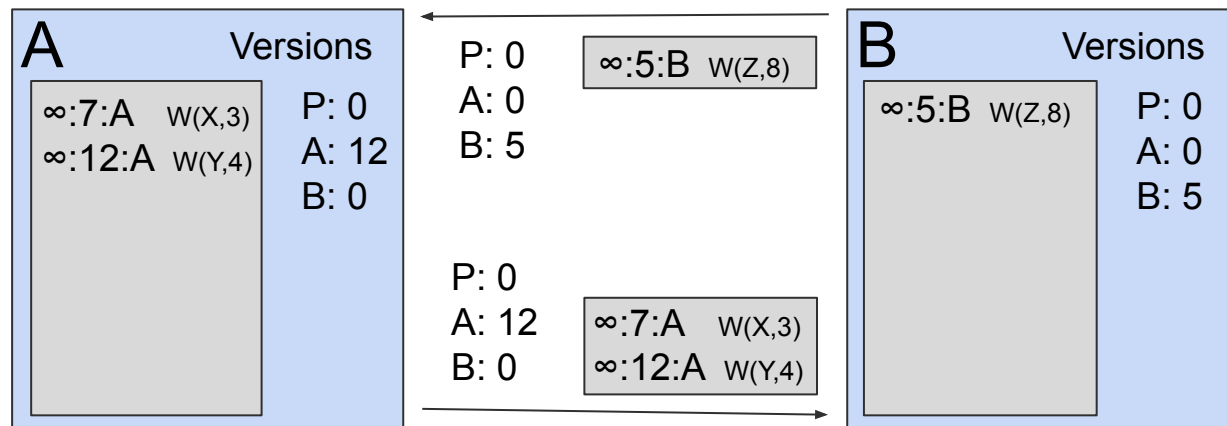
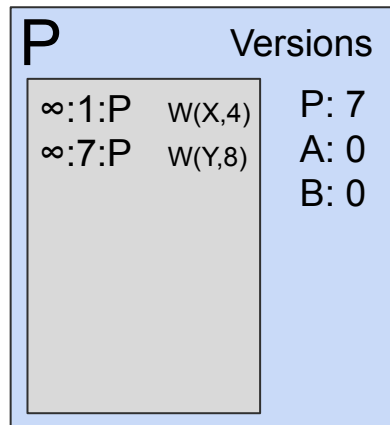
Legend

Commit Timestamp:Write Timestamp:Write Server

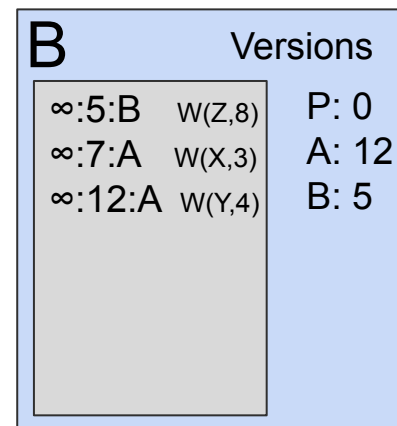
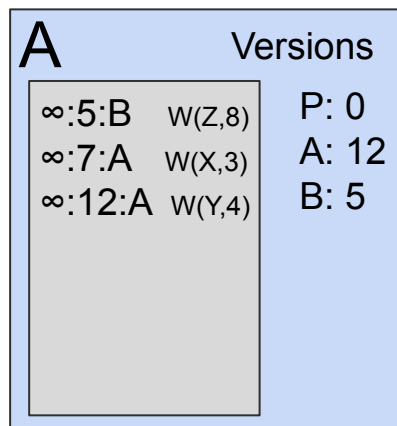
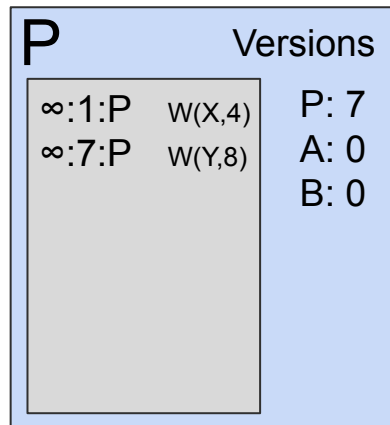


Bayou Anti-Entropy

Anti-Entropy Session
A & B

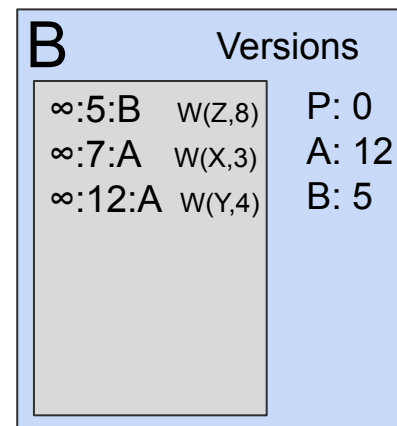
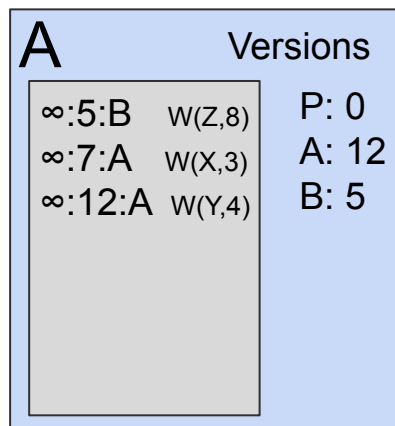
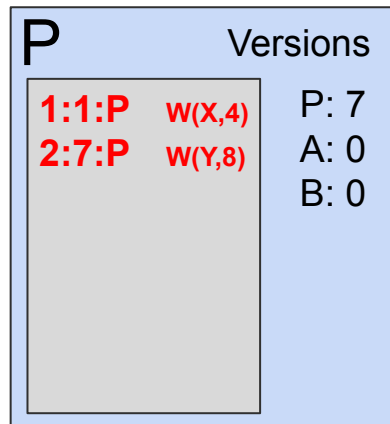


Bayou Anti-Entropy



Bayou Commit

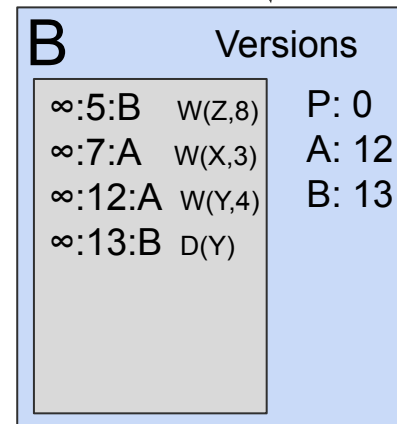
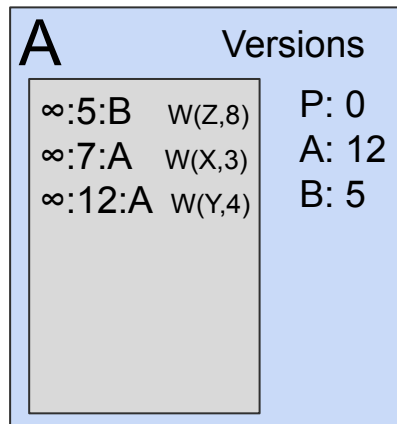
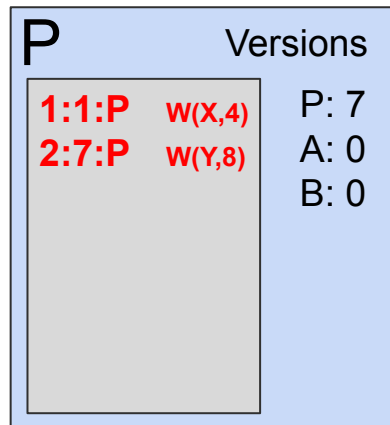
Primary **commits** its entries



Bayou Write

Write after anti-entropy session

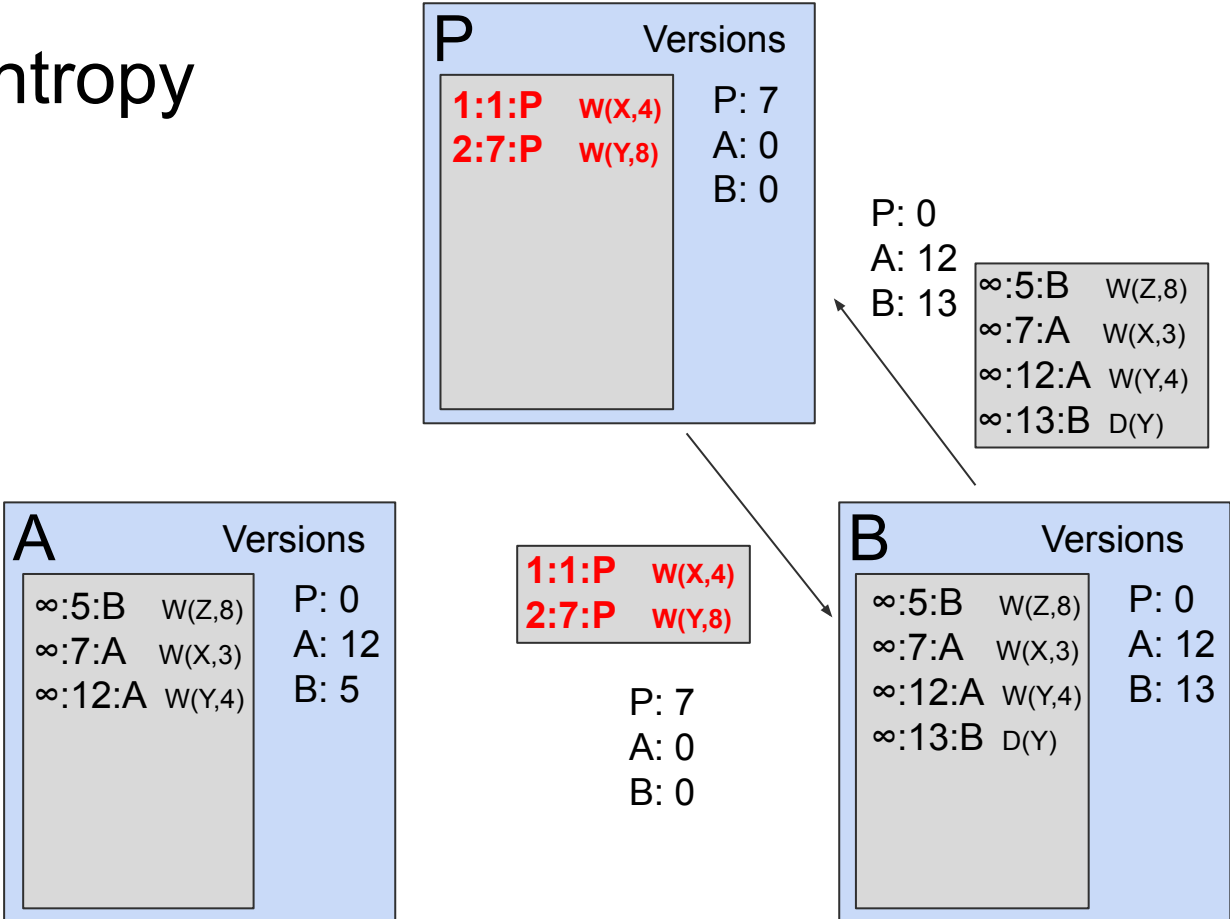
Write timestamp = $\max(\text{clock}, \max(\text{TS})+1)$



Client 1
D(Y)

Bayou Anti-Entropy

Anti-Entropy Session
P & B



Bayou Anti-Entropy

Anti-Entropy Session

P & B

Primary respects causality

P		Versions
1:1:P	w(x,4)	P: 7
2:7:P	w(y,8)	A: 12
∞ :5:B	w(z,8)	B: 13
∞ :7:A	w(x,3)	
∞ :12:A	w(y,4)	
∞ :13:B	D(Y)	

A		Versions
∞ :5:B	w(z,8)	P: 0
∞ :7:A	w(x,3)	A: 12
∞ :12:A	w(y,4)	B: 5

B		Versions
1:1:P	w(x,4)	P: 7
2:7:P	w(y,8)	A: 12
∞ :5:B	w(z,8)	B: 13
∞ :7:A	w(x,3)	
∞ :12:A	w(y,4)	
∞ :13:B	D(Y)	

Bayou Commit

Primary **commits** Its entries

P		Versions
1:1:P	w(X,4)	P: 7
2:7:P	w(Y,8)	A: 12
3:5:B	w(Z,8)	B: 13
4:7:A	w(X,3)	
5:12:A	w(Y,4)	
6:13:B	D(Y)	

A		Versions
∞ :5:B	w(Z,8)	P: 0
∞ :7:A	w(X,3)	A: 12
∞ :12:A	w(Y,4)	B: 5

B		Versions
1:1:P	w(X,4)	P: 7
2:7:P	w(Y,8)	A: 12
∞ :5:B	w(Z,8)	B: 13
∞ :7:A	w(X,3)	
∞ :12:A	w(Y,4)	
∞ :13:B	D(Y)	

Bayou

After a number of commits
and anti-entropy sessions
(without further writes)

P		Versions
1:1:P	w(X,4)	P: 7
2:7:P	w(Y,8)	A: 12
3:5:B	w(Z,8)	B: 13
4:7:A	w(X,3)	
5:12:A	w(Y,4)	
6:13:B	D(Y)	

A		Versions
1:1:P	w(X,4)	P: 7
2:7:P	w(Y,8)	A: 12
3:5:B	w(Z,8)	B: 13
4:7:A	w(X,3)	
5:12:A	w(Y,4)	
6:13:B	D(Y)	

B		Versions
1:1:P	w(X,4)	P: 7
2:7:P	w(Y,8)	A: 12
3:5:B	w(Z,8)	B: 13
4:7:A	w(X,3)	
5:12:A	w(Y,4)	
6:13:B	D(Y)	

Bayou and Dynamo similarities

Anti-entropy to achieve eventual consistency

Exchange vector clocks to determine order of operations

Expose conflict resolution to application

High availability!