

Oat Language Specification

CIS341 – Steve Zdancewic

March 30, 2020

1 Grammar

The following grammar defines the Oat syntax. All binary operations are *left associative* with precedence levels indicated numerically. Higher precedence operators bind tighter than lower precedence ones. Here *id* ranges over lower-case identifiers and *S* indicates an upper-case “struct name.” The parts of the grammar marked with M are not part of the abstract syntax, but are included for parsing purposes (e.g. parentheses) or for use in explaining the typing judgments. Note that left-hand-sides (*lhs*) and global expressions (*gexp*) are just a subsets of expressions.

t	$::=$	types
		<code>int</code>
		<code>bool</code>
		<code>ref?</code> nullable references
		<code>ref</code> non-null references
ref	$::=$	reference types
		<code>string</code>
		<i>S</i> named (mutable) record type
		$t []$ arrays
		$(t_0, \dots, t_n) \rightarrow rt$ function pointer
		(ref) M
rt	$::=$	return types
		<code>void</code>
		t
$prog$	$::=$	prog
		ϵ
		<i>decl prog</i>
$decl$	$::=$	global declarations
		<i>gdecl</i>
		<i>fdecl</i>
		<i>tdecl</i>

<i>gdecl</i>	::=	global value declarations
		<code>global id = gexp;</code>
<i>arg</i>	::=	arg
		<i>t id</i>
<i>args</i>	::=	args
		<i>arg</i> ₁ , ..., <i>arg</i> _{<i>n</i>}
<i>fdecl</i>	::=	function declaration
		<i>rt id(args) block</i>
<i>field</i>	::=	field declarations
		<i>t x</i>
<i>fields</i>	::=	fields
		<i>field</i> ₁ ; .. ; <i>field</i> _{<i>n</i>}
		<i>fields</i> ; <i>t</i> _{<i>n</i>+1} <i>x</i> _{<i>n</i>+1} ; .. ; <i>t</i> _{<i>m</i>} <i>x</i> _{<i>m</i>}
<i>tdecl</i>	::=	struct declaration
		<code>struct S{ fields }</code>
<i>exp</i>	::=	expressions
		<code>ref null</code>
		<code>true</code>
		<code>false</code>
		<code>integer</code>
		<code>string</code>
		<i>id</i>
		<code>new t[] {exp₁, .. , exp_{<i>n</i>}}</code>
		<code>new t [exp₁] {id -> exp₂}</code>
		<i>exp</i> ₁ [<i>exp</i> ₂]
		<code>length (exp)</code>
		<code>new S{x₁=exp₁; .. ; x_{<i>n</i>}=exp_{<i>n</i>}}</code>
		<i>exp</i> . <i>id</i>
		<i>exp</i> (<i>exp</i> ₁ , .. , <i>exp</i> _{<i>n</i>})
		<i>uop exp</i>
		<i>exp</i> ₁ <i>bop</i> <i>exp</i> ₂
		(<i>exp</i>)
	M	
<i>lhs</i>	::=	left-hand-sides for assignment
		<i>id</i>
		<i>exp</i> ₁ [<i>exp</i> ₂]
		<i>exp</i> . <i>id</i>
<i>gexp</i>	::=	global initializers

	<i>integer</i> <i>string</i> <i>ref null</i> <i>true</i> <i>false</i> new <i>t</i> []{ <i>gexp</i> ₁ , .., <i>gexp</i> _{<i>n</i>} } new <i>S</i> { <i>x</i> ₁ = <i>gexp</i> ₁ ; .. ; <i>x</i> _{<i>n</i>} = <i>gexp</i> _{<i>n</i>} } <i>id</i>	64-bit integer literals C-style strings
<i>block</i>	::= { <i>stmt</i> ₁ .. <i>stmt</i> _{<i>n</i>} }	blocks
<i>vdecl</i>	::= var <i>id</i> = <i>exp</i>	local variable declarations
<i>stmt</i>	::= <i>lhs</i> = <i>exp</i> ; <i>vdecl</i> ; return <i>exp</i> ; return ; <i>exp</i> (<i>exp</i> ₁ , .., <i>exp</i> _{<i>n</i>}); if _ <i>stmt</i> for (<i>vdecls</i> ; <i>exp</i> _{opt} ; <i>stmt</i> _{opt}) <i>block</i> while (<i>exp</i>) <i>block</i>	statements assignment statement variable declaration return with value return (void) call a void-returning function conditionals
<i>if_stmt</i>	::= if (<i>exp</i>) <i>block</i> <i>else_stmt</i> if? (<i>ref id</i> = <i>exp</i>) <i>block</i> <i>else_stmt</i>	if statements standard boolean if possibly-null checked downcast
<i>else_stmt</i>	::= ϵ else <i>block</i> else if _ <i>stmt</i>	else

<i>bop</i>	::=	(left associative) binary operations
		*
		multiplication (precedence 100)
		+
		addition (precedence 90)
		-
		subtraction (precedence 90)
		<<
		shift left (precedence 80)
		>>
		shift right logical (precedence 80)
		>>>
		shift right arithmetic (precedence 80)
		<
		less-than (precedence 70)
		<=
		less-than or equal (precedence 70)
		>
		greater-than (precedence 70)
		>=
		greater-than or equal (precedence 70)
		==
		equal (precedence 60)
		!=
		not equal (precedence 60)
		&
		logical and (precedence 50)
		logical or (precedence 40)
		[&]
		bit-wise and (precedence 30)
		[]
		bit-wise or (precedence 20)

<i>uop</i>	::=	unary operations
		-
		!
		~

2 Subtyping Rules

$$H \vdash t_1 \leq t_2$$

$$\frac{}{H \vdash \text{int} \leq \text{int}} \text{ SUB_SUB_INT}$$

$$\frac{}{H \vdash \text{bool} \leq \text{bool}} \text{ SUB_SUB_BOOL}$$

$$\frac{H \vdash_r ref_1 \leq ref_2}{H \vdash ref_1? \leq ref_2?} \text{ SUB_SUB_NREF}$$

$$\frac{H \vdash_r ref_1 \leq ref_2}{H \vdash ref_1 \leq ref_2} \text{ SUB_SUB_REF}$$

$$\frac{H \vdash_r ref_1 \leq ref_2}{H \vdash ref_1 \leq ref_2?} \text{ SUB_SUB_NRREF}$$

$$H \vdash_r ref_1 \leq ref_2$$

$$\frac{}{H \vdash_r \text{string} \leq \text{string}} \text{ SUB_SUBRSTRING}$$

$$\frac{}{H \vdash_r t[] \leq t[]} \text{ SUB_SUBRARRAY}$$

$$\frac{\text{struct } S_1\{ t_1 x_1; .. ; t_n x_n; t_{n+1} x_{n+1}; .. ; t_m x_m \} \in H \quad \text{struct } S_2\{ t_1 x_1; .. ; t_n x_n \} \in H}{H \vdash_r S_1 \leq S_2} \text{ SUB_SUBRSTRUCT}$$

$$\frac{H \vdash t'_1 \leq t_1 \dots H \vdash t'_n \leq t_n \quad H \vdash_{rt} rt_1 \leq rt_2}{H \vdash_r (t_1, \dots, t_n) \rightarrow rt_1 \leq (t'_1, \dots, t'_n) \rightarrow rt_2} \text{ SUB_SUBRFUNT}$$

$$H \vdash_{rt} rt_1 \leq rt_2$$

$$\frac{}{H \vdash_{rt} \text{void} \leq \text{void}} \text{ SUB_SUBRETSVOID}$$

$$\frac{H \vdash t_1 \leq t_2}{H \vdash_{rt} t_1 \leq t_2} \text{ SUB_SUBRETRTTYP}$$

3 Well-formed types

$H \vdash t$

$$\begin{array}{c}
 \frac{}{H \vdash \text{int}} \quad \text{WF_TYPOKOKINT} \\
 \frac{}{H \vdash \text{bool}} \quad \text{WF_TYPOKOKBOOL} \\
 \frac{H \vdash_r \text{ref}}{H \vdash \text{ref}} \quad \text{WF_TYPOKOKREFT} \\
 \frac{H \vdash_r \text{ref}}{H \vdash \text{ref?}} \quad \text{WF_TYPOKOKREFTQ}
 \end{array}$$

$H \vdash_r \text{ref}$

$$\begin{array}{c}
 \frac{}{H \vdash_r \text{string}} \quad \text{WF_REFTOKOKSTRING} \\
 \frac{H \vdash t}{H \vdash_r t[]} \quad \text{WF_REFTOKOKARRAY} \\
 \frac{\text{struct } S\{ \text{fields } \} \in H}{H \vdash_r S} \quad \text{WF_REFTOKOKSTRUCT} \\
 \frac{H \vdash t_1 \dots H \vdash t_n \quad H \vdash_{rt} rt}{H \vdash_r (t_1, \dots, t_n) \rightarrow rt} \quad \text{WF_REFTOKOKFUNT}
 \end{array}$$

$H \vdash_{rt} rt$

$$\begin{array}{c}
 \frac{}{H \vdash_{rt} \text{void}} \quad \text{WF_RTYPOKVOIDOK} \\
 \frac{H \vdash t}{H \vdash_{rt} t} \quad \text{WF_RTYPOKRTPOK}
 \end{array}$$

4 Typing Rules

$\boxed{\vdash bop_1, \dots, bop_i : t}$

$$\frac{}{\vdash +, *, -, <, >, >>, [&], [|] : (\text{int}, \text{int}) \rightarrow \text{int}} \text{ TYP_INTOPS}$$

$$\frac{}{\vdash <, \leq, >, \geq : (\text{int}, \text{int}) \rightarrow \text{bool}} \text{ TYP_CMPOPS}$$

$$\frac{}{\vdash \&, | : (\text{bool}, \text{bool}) \rightarrow \text{bool}} \text{ TYP_BOOLOPS}$$

$\boxed{\vdash uop : t}$

$$\frac{}{\vdash ! : (\text{bool}) \rightarrow \text{bool}} \text{ TYP_LOGNOT}$$

$$\frac{}{\vdash \sim : (\text{int}) \rightarrow \text{int}} \text{ TYP_BITNEG}$$

$$\frac{}{\vdash - : (\text{int}) \rightarrow \text{int}} \text{ TYP_NEG}$$

$\boxed{H; G; L \vdash exp : t}$

$$\frac{H \vdash ref}{H; G; L \vdash \text{ref_null} : \text{ref?}} \text{ TYP_NULL}$$

$$\frac{}{H; G; L \vdash \text{true} : \text{bool}} \text{ TYP_BOOL_TRUE}$$

$$\frac{}{H; G; L \vdash \text{false} : \text{bool}} \text{ TYP_BOOL_FALSE}$$

$$\frac{}{H; G; L \vdash \text{integer} : \text{int}} \text{ TYP_INT}$$

$$\frac{}{H; G; L \vdash \text{string} : \text{string}} \text{ TYP_STRING}$$

$$\frac{id : t \in L}{H; G; L \vdash id : t} \text{ TYP_LOCAL}$$

$$\frac{id \notin L \quad id : t \in G}{H; G; L \vdash id : t} \text{ TYP_GLOBAL}$$

$$H \vdash t$$

$$H; G; L \vdash exp_1 : t_1 \dots H; G; L \vdash exp_n : t_n$$

$$H \vdash t_1 \leq t \dots H \vdash t_n \leq t$$

$$\frac{H \vdash t_1 \leq t \dots H \vdash t_n \leq t}{H; G; L \vdash \text{new } t[] \{exp_1, \dots, exp_n\} : t[]} \text{ TYP_CARR}$$

$$H \vdash t$$

$$H; G; L \vdash exp_1 : \text{int}$$

$$\frac{x \notin L \quad H; G; L, x : \text{int} \vdash exp_2 : t' \quad H \vdash t' \leq t}{H; G; L \vdash \text{new } t[exp_1] \{x \rightarrow exp_2\} : t[]} \text{ TYP_NEWARRAY}$$

$$H; G; L \vdash exp_1 : t[] \quad H; G; L \vdash exp_2 : \text{int}$$

$$\frac{H; G; L \vdash exp_1 : t[] \quad H; G; L \vdash exp_2 : \text{int}}{H; G; L \vdash exp_1[exp_2] : t} \text{ TYP_INDEX}$$

$$H; G; L \vdash exp : t[]$$

$$\frac{H; G; L \vdash exp : t[]}{H; G; L \vdash \text{length}(exp) : \text{int}} \text{ TYP_LENGTH}$$

$$\begin{array}{c}
\text{struct } S\{ t_1 x_1; .. ; t_n x_n \} \in H \\
H; G; L \vdash exp_1 : t'_1 .. H; G; L \vdash exp_n : t'_n \\
H \vdash t'_1 \leq t_1 .. H \vdash t'_n \leq t_n \\
\text{fields may be permuted under new} \\
\hline
H; G; L \vdash \text{new } S\{x_1=exp_1; .. ; x_n=exp_n\} : S \quad \text{TYP_STRUCTEX}
\end{array}$$

$$\frac{H; G; L \vdash exp : S \quad \text{struct } S\{ \text{fields} \} \in H \quad t \ x \in \text{fields}}{H; G; L \vdash exp.x : t} \quad \text{TYP_FIELD}$$

$$\frac{H; G; L \vdash exp : (t_1, .. , t_n) \rightarrow t \quad H; G; L \vdash exp_1 : t'_1 .. H; G; L \vdash exp_n : t'_n \quad H \vdash t'_1 \leq t_1 .. H \vdash t'_n \leq t_n}{H; G; L \vdash exp(exp_1, .. , exp_n) : t} \quad \text{TYP_CALL}$$

$$\frac{\vdash bop : (t_1, t_2) \rightarrow t \quad H; G; L \vdash exp_1 : t_1 \quad H; G; L \vdash exp_2 : t_2}{H; G; L \vdash exp_1 \ bop \ exp_2 : t} \quad \text{TYP_BOP}$$

$$\frac{H; G; L \vdash exp_1 : t_1 \quad H; G; L \vdash exp_2 : t_2 \quad H \vdash t_1 \leq t_2 \quad H \vdash t_2 \leq t_1}{H; G; L \vdash exp_1 == exp_2 : \text{bool}} \quad \text{TYP_EQ}$$

$$\frac{H; G; L \vdash exp_1 : t_1 \quad H; G; L \vdash exp_2 : t_2 \quad H \vdash t_1 \leq t_2 \quad H \vdash t_2 \leq t_1}{H; G; L \vdash exp_1 != exp_2 : \text{bool}} \quad \text{TYP_NEQ}$$

$$\frac{\vdash uop : (t) \rightarrow t \quad H; G; L \vdash exp : t}{H; G; L \vdash uop \ exp : t} \quad \text{TYP_UOP}$$

$$H;G;L_1 \vdash vdecl \Rightarrow L_2$$

$$\frac{H;G;L \vdash exp : t \quad x \notin L}{H;G;L \vdash \text{var } x = exp \Rightarrow L, x:t} \text{ TYP_DECL}$$

$$H;G;L_0 \vdash vdecls \Rightarrow L_i$$

$$\frac{H;G;L_0 \vdash vdecl_1 \Rightarrow L_1 \quad \dots \quad H;G;L_{n-1} \vdash vdecl_i \Rightarrow L_n}{H;G;L_0 \vdash vdecl_1, \dots, vdecl_n \Rightarrow L_n} \text{ TYP_VDECLS}$$

$$H;G;L_1;rt \vdash stmt \Rightarrow L_2; returns$$

$$\frac{\begin{array}{l} G;L \vdash lhs \text{ in } L \text{ or } lhs \text{ not a global function id} \\ H;G;L \vdash lhs : t \\ H;G;L \vdash exp : t' \\ H \vdash t' \leq t \end{array}}{H;G;L;rt \vdash lhs = exp; \Rightarrow L; \perp} \text{ TYP_ASSN}$$

$$\frac{H;G;L_1 \vdash vdecl \Rightarrow L_2}{H;G;L_1;rt \vdash vdecl; \Rightarrow L_2; \perp} \text{ TYP_STMTDECL}$$

$$\frac{\begin{array}{l} H;G;L \vdash exp : (t_1, \dots, t_n) \rightarrow \text{void} \\ H;G;L \vdash exp_1 : t'_1 \dots H;G;L \vdash exp_n : t'_n \\ H \vdash t'_1 \leq t_1 \dots H \vdash t'_n \leq t_n \end{array}}{H;G;L;rt \vdash exp(exp_1, \dots, exp_n); \Rightarrow L; \perp} \text{ TYP_SCALL}$$

$$\frac{\begin{array}{l} H;G;L \vdash exp : \text{bool} \\ H;G;L;rt \vdash block_1; r_1 \\ H;G;L;rt \vdash block_2; r_2 \end{array}}{H;G;L;rt \vdash \text{if}(exp) block_1 \text{ else } block_2 \Rightarrow L; r_1 \wedge r_2} \text{ TYP_IF}$$

$$\frac{\begin{array}{l} H;G;L \vdash exp : ref' ? \\ H \vdash ref' \leq ref \\ H;G;L,x:ref;rt \vdash block_1; r_1 \quad H;G;L;rt \vdash block_2; r_2 \end{array}}{H;G;L;rt \vdash \text{if?}(ref x = exp) block_1 \text{ else } block_2 \Rightarrow L; r_1 \wedge r_2} \text{ TYP_IFQ}$$

$$\frac{\begin{array}{l} H;G;L \vdash exp : \text{bool} \\ H;G;L;rt \vdash block; r \end{array}}{H;G;L;rt \vdash \text{while}(exp) block \Rightarrow L; \perp} \text{ TYP_WHILE}$$

$$\frac{\begin{array}{l} H;G;L_1 \vdash vdecls \Rightarrow L_2 \\ H;G;L_2 \vdash exp : \text{bool} \\ H;G;L_2;rt \vdash stmt \Rightarrow L_3; \perp \\ H;G;L_2;rt \vdash block; r \end{array}}{H;G;L_1;rt \vdash \text{for}(vdecls; exp_{opt}; stmt_{opt}) block \Rightarrow L_1; \perp} \text{ TYP_FOR}$$

$$\frac{H;G;L \vdash exp : t' \quad H \vdash t' \leq t}{H;G;L;t \vdash \text{return } exp; \Rightarrow L; \top} \text{ TYP_RET}$$

$$\frac{}{H;G;L;\text{void} \vdash \text{return}; \Rightarrow L; \top} \text{ TYP_RETVOID}$$

$$[H;G;L;rt \vdash block; returns]$$

$$\frac{H;G;L_0;rt \vdash_{ss} stmt_1 .. stmt_n \Rightarrow L_n;r}{H;G;L_0;rt \vdash \{stmt_1 .. stmt_n\} ;r} \quad \text{TYP_BLOCK}$$

$$[H;G;L_0;rt \vdash_{ss} stmt_1 .. stmt_n \Rightarrow L_n; returns]$$

$$H;G;L_0;rt \vdash stmt_1 \Rightarrow L_1; \perp$$

$$\dots$$

$$H;G;L_{n-2};rt \vdash stmt_{n-1} \Rightarrow L_{n-1}; \perp$$

$$H;G;L_{n-1};rt \vdash stmt_n \Rightarrow L_n;r$$

$$\frac{H;G;L_0;rt \vdash_{ss} stmt_1 .. stmt_{n-1} stmt_n \Rightarrow L_n;r}{H;G;L_0;rt \vdash_{ss} stmt_1 .. stmt_{n-1} stmt_n \Rightarrow L_n;r} \quad \text{TYP_STMTS}$$

$$[H;G \vdash_s tdecl]$$

$$\frac{H \vdash t_1 .. H \vdash t_i \ x_1 .. x_i \text{ distinct}}{H;G \vdash_s \text{struct } S\{ t_1 x_1; .. ; t_i x_i \}} \quad \text{TYP_TDECLOK}$$

$$[H;G \vdash_f fdecl]$$

$$\frac{H;G;x_1:t_1,..,x_i:t_i;rt \vdash block; \top \ x_1 .. x_i \text{ distinct}}{H;G \vdash_f rt f(t_1 x_1, .. , t_i x_i) \ block} \quad \text{TYP_FDECLOK}$$

$$[H;G \vdash prog]$$

$$\overline{H;G \vdash \epsilon} \quad \text{TYP_DEMPTY}$$

$$\frac{H;G \vdash prog}{H;G \vdash gdecl prog} \quad \text{TYP_DGDECL}$$

$$\frac{H;G \vdash_f fdecl \quad H;G \vdash prog}{H;G \vdash fdecl prog} \quad \text{TYP_DFDECL}$$

$$\frac{H;G \vdash_s tdecl \quad H;G \vdash prog}{H;G \vdash tdecl prog} \quad \text{TYP_DTDECL}$$

$$[H;G_1 \vdash_g prog \Rightarrow G_2]$$

$$\overline{H;G \vdash_g \epsilon \Rightarrow G} \quad \text{TYP_GEMPTY}$$

$$\frac{H;G_1 \vdash_g prog \Rightarrow G_2}{H;G_1 \vdash_g tdecl prog \Rightarrow G_2} \quad \text{TYP_GTDECL}$$

$$\frac{H;G_1; \cdot \vdash gexp : t \quad x \notin G_1 \quad H;G_1, x:t \vdash_g prog \Rightarrow G_2}{H;G_1 \vdash_g \text{global } x = gexp; prog \Rightarrow G_2} \quad \text{TYP_GGDECL}$$

$$\frac{H;G_1 \vdash_g prog \Rightarrow G_2}{H;G_1 \vdash_g fdecl prog \Rightarrow G_2} \quad \text{TYP_GFDECL}$$

$$[H \vdash fdecl \Rightarrow id:t]$$

$$\frac{H \vdash_{rt} rt \quad H \vdash t_1 .. H \vdash t_n}{H \vdash rt f(t_1 x_1, .. , t_n x_n) \ block \Rightarrow f:(t_1, .. , t_n) \rightarrow rt} \quad \text{TYP_FTYP}$$

$H; G_1 \vdash_f prog \Rightarrow G_2$

$$\begin{array}{c}
 \frac{}{H; G \vdash_f \epsilon \Rightarrow G} \text{ TYP_FEMPTY} \\
 \frac{H; G_1 \vdash_f prog \Rightarrow G_2}{H; G_1 \vdash_f tdecl prog \Rightarrow G_2} \text{ TYP_FTDECL} \\
 \frac{H; G_1 \vdash_f prog \Rightarrow G_2}{H; G_1 \vdash_f gdecl prog \Rightarrow G_2} \text{ TYP_FGDECL} \\
 \frac{\begin{array}{c} H \vdash fdecl \Rightarrow f:t \\ f \notin G_1 \quad H; G_1, f:t \vdash_f prog \Rightarrow G_2 \end{array}}{H; G_1 \vdash_f fdecl prog \Rightarrow G_2} \text{ TYP_FFDECL}
 \end{array}$$

$H_1 \vdash_s prog \Rightarrow H_2$

$$\begin{array}{c}
 \frac{}{H \vdash_s \epsilon \Rightarrow H} \text{ TYP_SEMPY} \\
 \frac{S \notin H_1 \quad H_1, \text{struct } S\{ \text{fields } \} \vdash_s prog \Rightarrow H_2}{H_1 \vdash_s \text{struct } S\{ \text{fields } \} \vdash_s prog \Rightarrow H_2} \text{ TYP_STDECL} \\
 \frac{H_1 \vdash_s prog \Rightarrow H_2}{H_1 \vdash_s gdecl prog \Rightarrow H_2} \text{ TYP_SGDECL} \\
 \frac{H_1 \vdash_s prog \Rightarrow H_2}{H_1 \vdash_s fdecl prog \Rightarrow H_2} \text{ TYP_SFDECL}
 \end{array}$$

$\vdash prog$

$$\frac{\cdot \vdash_s prog \Rightarrow H \quad H; G_0 \vdash_f prog \Rightarrow G_1 \quad H; G_1 \vdash_g prog \Rightarrow G_2 \quad H; G_2 \vdash prog}{\vdash prog} \text{ TYP_PROG}$$

Notes:

- The context G_0 mentioned in the rule for typechecking a complete, top-level program is the “initial context” which should contain bindings for all of the OAT built-in functions.
- The type system processes the program in several passes: (1) collect up all the structure type definitions and make sure their names don’t clash using the \vdash_s rules, (2) add all the function identifiers and their types to the global context using the \vdash_f rules, again ensuring no name clashes, (3) typecheck the global value declarations and add them to the context using the \vdash_g rules, and (4) process all the declarations one more time to examine all the struct fields to make sure their types are well formed and to typecheck the function bodies.

The rules therefore allow the types of structs to be mutually recursive, and for global values to mention function pointers as constants.

- We use \perp to indicate that a statement might not return, and \top to indicate that a statement definitely returns. When typechecking the list of statements that make up a block, only the last statement is allowed to definitely return (all the others must possibly not return). Note that TYP_FDECLOK requires that the block making up a function body definitely return.