

# COS 511: Theoretical Machine Learning

Homework #2  
Sample size bounds, growth function, VC dimension

Due:  
February 28, 2018

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## Problem 1

[10] As on Problem 1 on Homework #1, let  $X = \mathbb{R}$ , and let  $\mathcal{C}_s$  be the class of concepts defined by unions of  $s$  intervals. Compute the VC-dimension of  $\mathcal{C}_s$  exactly.

## Problem 2

[15] For  $i = 1, \dots, n$ , let  $\mathcal{G}_i$  be a space of concepts ( $\{0, 1\}$ -valued functions) defined on some domain  $X$ , and let  $\mathcal{F}$  be a space of concepts defined on  $\{0, 1\}^n$ . (That is, each  $g_i \in \mathcal{G}_i$  maps  $X$  to  $\{0, 1\}$ , and each  $f \in \mathcal{F}$  maps  $\{0, 1\}^n$  to  $\{0, 1\}$ .) Let  $\mathcal{H}$  be the space of all concepts  $h : X \rightarrow \{0, 1\}$  of the form

$$h(x) = f(g_1(x), \dots, g_n(x))$$

for some  $f \in \mathcal{F}$ ,  $g_1 \in \mathcal{G}_1, \dots, g_n \in \mathcal{G}_n$ .

Give a careful argument proving that

$$\Pi_{\mathcal{H}}(m) \leq \Pi_{\mathcal{F}}(m) \cdot \prod_{i=1}^n \Pi_{\mathcal{G}_i}(m).$$

[An **optional** continuation of this problem, applicable to feedforward networks, is given in Problem 5.]

## Problem 3

[15] Show that Sauer's Lemma is tight. That is, for each  $d = 0, 1, 2, \dots$ , give an example of a class  $\mathcal{C}$  with VC-dimension equal to  $d$  such that for each  $m$ ,

$$\Pi_{\mathcal{C}}(m) = \sum_{i=0}^d \binom{m}{i}.$$

## Problem 4

This problem explores another general method for bounding the error when the hypothesis space is infinite.

Some algorithms output hypotheses that can be represented by a small number of examples from the training set. For instance, suppose the domain is  $\mathbb{R}$  and we are learning a half-line of the form  $x \geq a$  where  $a$  defines the half-line. A simple algorithm chooses the left most positive training example  $a$  and outputs the corresponding half-line, which is clearly consistent with the data. Thus, in this case, the hypothesis can be represented by a single training example.

More formally, let  $F$  be a function mapping labeled examples to concepts, and assume that algorithm  $A$ , when given training examples  $(x_1, c(x_1)), \dots, (x_m, c(x_m))$  labeled by some unknown  $c \in \mathcal{C}$ , chooses some  $i_1, \dots, i_k \in \{1, \dots, m\}$  and outputs the consistent hypothesis  $h = F((x_{i_1}, c(x_{i_1})), \dots, (x_{i_k}, c(x_{i_k})))$ . In a sense, the algorithm has "compressed" the sample down to a sequence of just  $k$  of the  $m$  training examples. (We assume throughout that  $m > k$ .)

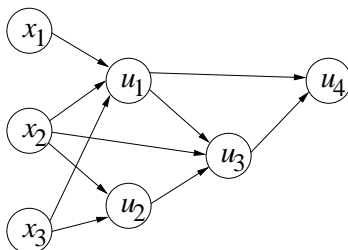
- a. [5] Give such an algorithm for axis-aligned hyper-rectangles in  $\mathbb{R}^n$  with  $k = O(n)$ . (An axis-aligned hyper-rectangle is a set of the form  $[a_1, b_1] \times \cdots \times [a_n, b_n]$ , and the corresponding concept, as usual, is the binary function that is 1 for points inside the rectangle and 0 otherwise. For  $n = 2$ , this is the class of rectangles used repeatedly as an example in class.) Your algorithm should run in time polynomial in  $m$  and  $n$ .
- b. [15] Returning to the general case, assume as usual that the examples are chosen at random from some distribution  $D$ . Also assume that the size  $k$  is fixed. Argue carefully that the error of the output hypothesis  $h$ , with probability at least  $1 - \delta$ , satisfies the bound:

$$\text{err}_D(h) \leq O\left(\frac{\ln(1/\delta) + k \ln m}{m - k}\right).$$

[Side note: A difficult, long-standing open problem asks if it is always possible to find such a “compression scheme” whose size  $k$  is equal to (or proportional to) the VC-dimension  $d$  of the target class  $\mathcal{C}$ .]

### Problem 5 – Optional (Extra Credit)

[15] This problem shows one way in which the methods we have been developing can be applied to *feedforward networks*, including (some) neural networks.



A feedforward network, as in the example above, is defined by a directed acyclic graph on a set of *input nodes*  $x_1, \dots, x_n$ , and *computation nodes*  $u_1, \dots, u_N$ . The input nodes have no incoming edges. One of the computation nodes is called the *output node*, and has no outgoing edges. Each computation node  $u_k$  is associated with a function  $f_k : \mathbb{R}^{n_k} \rightarrow \{0, 1\}$ , where  $n_k$  is  $u_k$ 's indegree (number of ingoing edges). On input  $\mathbf{x} \in \mathbb{R}^n$ , the network computes its output  $g(\mathbf{x})$  in a natural, feedforward fashion. For instance, given input  $\mathbf{x} = \langle x_1, x_2, x_3 \rangle$ , the network above computes  $g(\mathbf{x})$  as follows:

$$\begin{aligned} u_1 &= f_1(x_1, x_2, x_3) \\ u_2 &= f_2(x_2, x_3) \\ u_3 &= f_3(u_1, x_2, u_2) \\ u_4 &= f_4(u_1, u_3) \\ g(\mathbf{x}) &= u_4. \end{aligned}$$

(Here, we slightly abuse notation, writing  $x_j$  and  $u_k$  both for nodes of the network, and for the input/computed values associated with these nodes.) The number of edges in the graph is denoted  $W$ .

In what follows, we regard the underlying graph as fixed, but allow the functions  $f_k$  to vary, or to be learned from data. In particular, let  $\mathcal{F}_1, \dots, \mathcal{F}_N$  be spaces of functions. As just explained, every choice of functions  $f_1, \dots, f_N$  induces an overall function  $g : \mathbb{R}^n \rightarrow \{0, 1\}$  for the network. We let  $\mathcal{G}$  denote the space of all such functions when  $f_k$  is chosen from  $\mathcal{F}_k$  for  $k = 1, \dots, N$ .

a. Prove that

$$\Pi_{\mathcal{G}}(m) \leq \prod_{k=1}^N \Pi_{\mathcal{F}_k}(m).$$

(Note that this is a generalization of Problem 2.)

b. Let  $d_k$  be the VC-dimension of  $\mathcal{F}_k$ , and let  $d = \sum_{k=1}^N d_k$ . Assume  $m \geq d_k \geq 1$  for all  $k$ . Prove that

$$\Pi_{\mathcal{G}}(m) \leq \left( \frac{emN}{d} \right)^d.$$

c. Consider the typical case in which the functions  $f_k$  are linear threshold functions; as we know, this class of functions has VC-dimension  $d_k = n_k + 1$ . Give an exact expression for  $d$  in terms of  $N$ ,  $n$ , and  $W$ . Conclude by deriving a “big-Oh” upper bound on the generalization error of any  $g \in \mathcal{G}$  that is consistent with  $m$  random examples, assuming  $m \geq d$ . Your bound should hold with probability at least  $1 - \delta$ , and should be expressed in terms of  $N$ ,  $n$ ,  $W$ ,  $m$ , and  $\delta$ .