Assembly Language: Function Calls
Help you learn:
  • Function call problems
  • x86-64 solutions
    • Pertinent instructions and conventions
Function Call Problems

(1) Calling and returning
• How does caller function **jump** to callee function?
• How does callee function **jump back** to the right place in caller function?

(2) Passing arguments
• How does caller function pass **arguments** to callee function?

(3) Storing local variables
• Where does callee function store its **local variables**?

(4) Returning a value
• How does callee function send **return value** back to caller function?
• How does caller function access the **return value**?

(5) Optimization
• How do caller and callee function minimize memory access?
Running Example

long absadd(long a, long b)
{
    long absA, absB, sum;
    absA = labs(a);
    absB = labs(b);
    sum = absA + absB;
    return sum;
}

Calls standard C labs() function
• Returns absolute value of given long
Agenda

Calling and returning

Passing arguments
Storing local variables
Returning a value
Optimization
Problem 1: Calling and Returning

How does caller *jump* to callee?
• i.e., Jump to the address of the callee’s first instruction

How does the callee *jump back* to the right place in caller?
• i.e., Jump to the instruction immediately following the most-recently-executed call instruction

```c
... absadd(3L, -4L);
...
long absadd(long a, long b)
{
    long absA, absB, sum;
    absA = labs(a);
    absB = labs(b);
    sum = absA + absB;
    return sum;
}
```
Attempted Solution: jmp Instruction

Attempted solution: caller and callee use jmp instruction

f:
    ...
    jmp g       # Call g
fReturnPoint:
    ...

g:
    ...
    jmp fReturnPoint  # Return
Problem: callee may be called by multiple callers

f1:
...  
jmp g # Call g

f2:
...  
jmp g # Call g

f1ReturnPoint:
...  

f2ReturnPoint:
...  

g:
...  
jmp ??? # Return
 Attempted solution: Store return address in register

\begin{align*}
\text{f1:} & \quad \text{movq } $f1\text{ReturnPoint}, \%rax} \\
& \quad \text{jmp g} \quad \# \text{Call g} \\
\text{f1ReturnPoint:} & \quad \ldots
\end{align*}

\begin{align*}
\text{f2:} & \quad \text{movq } $f2\text{ReturnPoint}, \%rax} \\
& \quad \text{jmp g} \quad \# \text{Call g} \\
\text{f2ReturnPoint:} & \quad \ldots
\end{align*}

\begin{align*}
\text{g:} & \quad \ldots \\
& \quad \text{jmp } *\%rax \quad \# \text{Return}
\end{align*}

Special form of \texttt{jmp} instruction
Problem: Cannot handle nested function calls

f:

```
    movq $fReturnPoint, %rax
    jmp g    # Call g
fReturnPoint:
    ...
```

Problem if \( f() \) calls \( g() \), and \( g() \) calls \( h() \)

Return address \( g() \) -> \( f() \) is lost

\( g() \):

```
    movq $gReturnPoint, %rax
    jmp h    # Call h
gReturnPoint:
    ...
    jmp *%rax  # Return
```

\( h() \):

```
    ...
    jmp *%rax  # Return
```
Observations:

• May need to store many return addresses
  • The number of nested function calls is not known in advance
  • A return address must be saved for as long as the invocation of this function is live, and discarded thereafter
• Stored return addresses are destroyed in reverse order of creation
  • \( f() \) calls \( g() \) \( \Rightarrow \) return addr for \( g \) is stored
  • \( g() \) calls \( h() \) \( \Rightarrow \) return addr for \( h \) is stored
  • \( h() \) returns to \( g() \) \( \Rightarrow \) return addr for \( h \) is destroyed
  • \( g() \) returns to \( f() \) \( \Rightarrow \) return addr for \( g \) is destroyed
• LIFO data structure (stack) is appropriate

x86-64 solution:

• Use the STACK section of memory, usually accessed via RSP
• Via \texttt{call} and \texttt{ret} instructions
**call and ret Instructions**

**ret** instruction “knows” the return address

```
f:
  ...
  call h
  ...
  call g
  ...

h:
  ...
  ret
```

```
g:
  ...
  call h
  ...
  ret
```

Diagram:

1. Call from f to h
2. Return from h to g
3. Call from g to h
4. Return from h to f
5. Call from f to h
6. Return from h to f
### Stack operations

**RSP** (stack pointer) register points to top of stack

<table>
<thead>
<tr>
<th>Instruction</th>
<th>Equivalent to</th>
</tr>
</thead>
<tbody>
<tr>
<td>pushq src</td>
<td>subq $8, %rsp</td>
</tr>
<tr>
<td></td>
<td>movq src, (%rsp)</td>
</tr>
<tr>
<td>popq dest</td>
<td>movq (%rsp), dest</td>
</tr>
<tr>
<td></td>
<td>addq $8, %rsp</td>
</tr>
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![Diagram showing RSP register and stack](image.png)
Implementation of `call`

**RIP** (instruction pointer) register points to next instruction to be executed.

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</tr>
<tr>
<td></td>
<td><code>addq $8, %rsp</code></td>
</tr>
<tr>
<td><code>call addr</code></td>
<td><code>pushq %rip</code></td>
</tr>
<tr>
<td></td>
<td><code>jmp addr</code></td>
</tr>
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**Note:** Can’t really access RIP directly, but this is implicitly what `call` does.

`call` instruction pushes return addr (old RIP) onto stack, then jumps.
Implementation of `call`

**RIP (instruction pointer) register** points to next instruction to be executed

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| `pushq src`     | `subq $8, %rsp
movq src, (%rsp)`             |
| `popq dest`     | `movq (%rsp), dest
addq $8, %rsp`              |
| `call addr`     | `pushq %rip
jmp addr`                |

RSP (register stack pointer) is updated as follows after a call:
- Push 8 to %rsp
- Push %rip
- Jump to new address

Old RIP (instruction pointer) register points to next instruction to be executed.
Implementation of `ret`

**RIP** (instruction pointer) register points to next instruction to be executed.

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| `pushq src`   | `subq $8, %rsp`  
                | `movq src, (%rsp)` |
| `popq dest`   | `movq (%rsp), dest`  
                | `addq $8, %rsp`   |
| `call addr`   | `pushq %rip`  
                | `jmp addr`        |
| `ret`         | `popq %rip`                        |

The `ret` instruction pops the stack, thus placing the return address (old RIP) into the RIP.
Implementation of `ret`

**RIP** (instruction pointer) register points to next instruction to be executed

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| `pushq src`   | `subq $8, %rsp
movq src, (%rsp)`                           |
| `popq dest`   | `movq (%rsp), dest
addq $8, %rsp`                            |
| `call addr`   | `pushq %rip
jmp addr`                              |
| `ret`         | `popq %rip`                                      |
# long absadd(long a, long b)

absadd:

    # long absA, absB, sum
    ...
    # absA = labs(a)
    ...
    call labs
    ...
    # absB = labs(b)
    ...
    call labs
    ...
    # sum = absA + absB
    ...
    # return sum
    ...
    ret
Agenda

Calling and returning

Passing arguments

Storing local variables

Returning a value

Optimization
Problem 2: Passing Arguments

Problem:
• How does caller pass arguments to callee?
• How does callee accept parameters from caller?

```c
long absadd(long a, long b)
{
    long absA, absB, sum;
    absA = labs(a);
    absB = labs(b);
    sum = absA + absB;
    return sum;
}
```
x86-64 Solution 1: Use the Stack

Observations (déjà vu):

- May need to store many arg sets
  - The number of arg sets is not known in advance
  - Arg set must be saved for as long as the invocation of this function is live, and discarded thereafter
- Stored arg sets are destroyed in reverse order of creation
- LIFO data structure (stack) is appropriate
x86-64 Solution 2: Use Registers

x86-64 solution:
• Pass first 6 (integer or address) arguments in registers for efficiency
  • RDI, RSI, RDX, RCX, R8, R9
• More than 6 arguments ⇒
  • Pass arguments 7, 8, … on the stack
  • (Beyond scope of COS 217)
• Arguments are structures ⇒
  • Pass arguments on the stack
  • (Beyond scope of COS 217)

Callee function then saves arguments to stack
• Or maybe not!
  • See “optimization” later this lecture
• Callee accesses arguments as positive offsets vs. RSP
Running Example

```c
# long absadd(long a, long b)
absadd:
    pushq %rdi # Push a
    pushq %rsi # Push b

    # long absA, absB, sum
    ...
    # absA = labs(a)
    movq 8(%rsp), %rdi
    call labs
    ...
    # absB = labs(b)
    movq 0(%rsp), %rdi
    call labs
    ...
    # sum = absA + absB
    ...
    # return sum
    ...
    addq $16, %rsp
    ret
```
Agenda

Calling and returning
Passing arguments
**Storing local variables**
Returning a value
Optimization
Problem 3: Storing Local Variables

Where does callee function store its local variables?

```c
long absadd(long a, long b)
{
    long absA, absB, sum;
    absA = labs(a);
    absB = labs(b);
    sum = absA + absB;
    return sum;
}
```
Observations (déjà vu again!):
  • May need to store many local var sets
    • The number of local var sets is not known in advance
    • Local var set must be saved for as long as the invocation of this
      function is live, and discarded thereafter
  • Stored local var sets are destroyed in reverse order of creation
  • LIFO data structure (stack) is appropriate

x86-64 solution:
  • Use the STACK section of memory
  • Or maybe not!
    • See later this lecture
# long absadd(long a, long b)
absadd:
  pushq %rdi # Push a
  pushq %rsi # Push b

  # long absA, absB, sum
  subq $24, %rsp

  # absA = labs(a)
  movq 32(%rsp), %rdi
  call labs
  ...
  # absB = labs(b)
  movq 24(%rsp), %rdi
  call labs
  ...
  # sum = absA + absB
  movq 16(%rsp), %rax
  addq 8(%rsp), %rax
  movq %rax, 0(%rsp)
  ...
  # return sum
  ...
  addq $40, %rsp
  ret
Agenda

Calling and returning
Passing arguments
Storing local variables
Returning a value
Optimization
Problem 4: Return Values

Problem:

• How does callee function send return value back to caller function?
• How does caller function access return value?

```c
long absadd(long a, long b)
{
    long absA, absB, sum;
    absA = labs(a);
    absB = labs(b);
    sum = absA + absB;
    return sum;
}
```
x86-64 Solution: Use RAX

In principle
  • Store return value in stack frame of caller

Or, for efficiency
  • Known small size ⇒ store return value in register
  • Other ⇒ store return value in stack

x86-64 convention
  • Integer or address:
    • Store return value in RAX
  • Floating-point number:
    • Store return value in floating-point register
    • (Beyond scope of COS 217)
  • Structure:
    • Store return value on stack
    • (Beyond scope of COS 217)
Running Example

```
# long absadd(long a, long b)
absadd:
    pushq %rdi  # Push a
    pushq %rsi  # Push b

    # long absA, absB, sum
    subq $24, %rsp

    # absA = labs(a)
    movq 32(%rsp), %rdi
    call labs
    movq %rax, 16(%rsp)

    # absB = labs(b)
    movq 24(%rsp), %rdi
    call labs
    movq %rax, 8(%rsp)

    # sum = absA + absB
    movq 16(%rsp), %rax
    addq 8(%rsp), %rax
    movq %rax, 0(%rsp)

    # return sum
    movq 0(%rsp), %rax
    addq $40, %rsp
    ret
```

Complete!
Agenda

Calling and returning
Passing arguments
Storing local variables
Returning a value

Optimization
Observation: Accessing memory is expensive
- More expensive than accessing registers
- For efficiency, want to store parameters and local variables in registers (and not in memory) when possible

Observation: Registers are a finite resource
- In principle: Each function should have its own registers
- In reality: All functions share same small set of registers

Problem: How do caller and callee use same set of registers without interference?
- Callee may use register that the caller also is using
- When callee returns control to caller, old register contents may have been lost
- Caller function cannot continue where it left off
Callee-save registers

- RBX, RBP, R12, R13, R14, R15
- Callee function **must preserve** contents
- If necessary…
  - Callee saves to stack near beginning
  - Callee restores from stack near end

Caller-save registers

- RDI, RSI, RDX, RCX, R8, R9, RAX, R10, R11
- Callee function **can change** contents
- If necessary…
  - Caller saves to stack before call
  - Caller restores from stack after call
Running Example

Local variable handling in *unoptimized* version:
- At beginning, `absadd()` allocates space for local variables (`absA`, `absB`, `sum`) on stack
- Body of `absadd()` uses stack
- At end, `absadd()` pops local variables from stack

Local variable handling in *optimized* version:
- `absadd()` keeps local variables in R13, R14, R15
- Body of `absadd()` uses R13, R14, R15
- Must be careful:
  - `absadd()` cannot change contents of R13, R14, or R15
  - So `absadd()` must save R13, R14, and R15 near beginning, and restore near end
Running Example

```c
# long absadd(long a, long b)
absadd:
  pushq %r13 # Save R13, use for absA
  pushq %r14 # Save R14, use for absB
  pushq %r15 # Save R15, use for sum

  # absA = labs(a)
  pushq %rsi # Save RSI
  call labs
  movq %rax, %r13
  popq %rsi # Restore RSI

  # absB += labs(b)
  movq %rsi, %rdi
  call labs
  movq %rax, %r14

  # sum = absA + absB
  movq %r13, %r15
  addq %r14, %r15

  # return sum
  movq %r15, %rax
  popq %r15 # Restore R15
  popq %r14 # Restore R14
  popq %r13 # Restore R13
  ret
```

absadd() stores local vars in R13, R14, R15, not in memory

absadd() cannot destroy contents of R13, R14, R15

So absadd() must save R13, R14, R15 near beginning and restore near end
Running Example

Parameter handling in unoptimized version:

- `absadd()` accepts parameters (a and b) in RDI and RSI
- At beginning, `absadd()` copies contents of RDI and RSI to stack
- Body of `absadd()` uses stack
- At end, `absadd()` pops parameters from stack

Parameter handling in optimized version:

- `absadd()` accepts parameters (a and b) in RDI and RSI
- Body of `absadd()` uses RDI and RSI
- Must be careful:
  - Call of `labs()` could change contents of RDI and/or RSI
  - `absadd()` must save contents of RDI and/or RSI before call of `labs()`, and restore contents after call
Running Example

```c
# long absadd(long a, long b)
absadd:
    pushq %r13  # Save R13, use for absA
    pushq %r14  # Save R14, use for absB
    pushq %r15  # Save R15, use for sum

    # absA = labs(a)
    pushq %rsi  # Save RSI
    call labs
    movq %rax, %r13
    popq %rsi   # Restore RSI

    # absB += labs(b)
    movq %rsi, %rdi
    call labs
    movq %rax, %r14

    # sum = absA + absB
    movq %r13, %r15
    addq %r14, %r15

    # return sum
    movq %r15, %rax
    popq %r15  # Restore R15
    popq %r14  # Restore R14
    popq %r13  # Restore R13
    ret
```

**absadd()** keeps a and b in RDI and RSI, not in memory

**labs()** can change RDI and/or RSI

**absadd()** must retain contents of RSI (value of b) across 1\textsuperscript{st} call of labs()

So **absadd()** must save RSI before call and restore RSI after call
Non-Optimized vs. Optimized Patterns

Unoptimized pattern
• Parameters and local variables strictly in memory (stack) during function execution
• **Pro**: Always possible
• **Con**: Inefficient
• gcc compiler uses when invoked without –O option

Optimized pattern
• Parameters and local variables mostly in registers during function execution
• **Pro**: Efficient
• **Con**: Sometimes impossible
  • More than 6 local variables
  • Local variable is a structure or array
  • Function computes address of parameter or local variable
• gcc compiler uses when invoked with –O option, when it can!
Hybrid Patterns

Hybrids are possible

• Example
  • Parameters in registers
  • Local variables in memory (stack)

Hybrids are error prone for humans

• Example (continued from previous)
  • Step 1: Access local variable \( \leftarrow \) local var is at stack offset \( X \)
  • Step 2: Push caller-save register
  • Step 3: Access local variable \( \leftarrow \) local var is at stack offset \( X+8!!! \)
  • Step 4: Call \texttt{labs()}\n  • Step 6: Access local variable \( \leftarrow \) local var is at stack offset \( X+8!!! \)
  • Step 7: Pop caller-save register
  • Step 8: Access local variable \( \leftarrow \) local var is at stack offset \( X \)

Avoid hybrids for Assignment 4
Summary

Function calls in x86-64 assembly language

Calling and returning

• `call` instruction pushes RIP onto stack and jumps
• `ret` instruction pops from stack to RIP

Passing arguments

• Caller copies args to caller-saved registers (in prescribed order)
• Unoptimized pattern:
  • Callee pushes args to stack
  • Callee uses args as positive offsets from RSP
  • Callee pops args from stack
• Optimized pattern:
  • Callee keeps args in caller-saved registers
  • Be careful!
Summary (cont.)

Storing local variables
• Unoptimized pattern:
  • Callee pushes local vars onto stack
  • Callee uses local vars as positive offsets from RSP
  • Callee pops local vars from stack
• Optimized pattern:
  • Callee keeps local vars in callee-saved registers
  • Be careful!

Returning values
• Callee places return value in RAX
• Caller accesses return value in RAX
Putting it all together

Add up the keys of a tree

```c
struct tree {
    int key;
    struct tree *left;
    struct tree *right;
};

int sum (struct tree *t) {
    if (t==NULL)
        return 0;
    else return t->key +
            sum(t->left) +
            sum(t->right);
}
```

```
.data
sum:
.globl sum
sum: