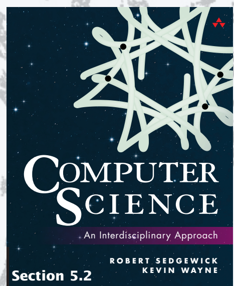


http://introcs.cs.princeton.edu

17. A Computing Machine

COMPUTER SCIENCE
SEDGEWICK / WAYNE
PART II: ALGORITHMS, MACHINES, and THEORY

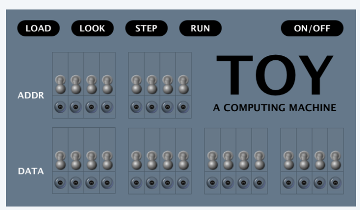
17. A Computing Machine

- Overview
- Data types
- Instructions
- Operating the machine
- Machine language programming

CS.17.A.MachineI.Overview

A TOY computing machine

TOY is an imaginary machine similar to:
- Ancient computers.
- Today's smartphone processors.
- Countless other devices designed and built over the past 50 years.



Smartphone processor, 2010s


PDP-8, 1970s

Reasons to study TOY

Prepare to learn about computer architecture
- How does your computer's processor work?
- What are its basic components?
- How do they interact?

Learn about machine-language programming.
- How do Java programs relate to your computer?
- Key to understanding Java references.
- Intellectual challenge of a new programming regime.
- Still necessary in some modern applications.

multimedia, computer games, embedded devices, scientific computing,...

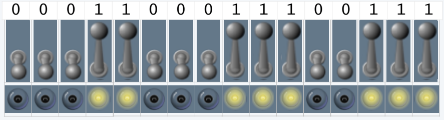
Learn fundamental abstractions that have informed processor design for decades.

Bits and words

Everything in TOY is encoded with a sequence of *bits* (value 0 or 1).

- Why? Easy to represent two states (on and off) in real world.
- Bits are organized in 16-bit sequences called *words.*

0 0 0 1 1 0 0 0 1 1 1 0 0 1 1 1



More convenient for humans: *hexadecimal notation* (base 16)

- 4 *hex digits* in each word.
- Convert to and from binary 4 bits at a time.

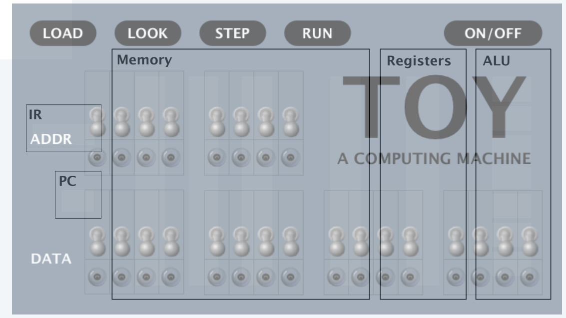
0	0	0	1	1	0	0	0	1	1	1	0	0	1	1	1
1				8				E				7			

binary	*hex*
0000	0
0001	1
0010	2
0011	3
0100	4
0101	5
0110	6
0111	7
1000	8
1001	9
1010	A
1011	B
1100	C
1101	D
1110	E
1111	F

Inside the box

Components of TOY machine

- Memory
- Registers
- Arithmetic and logic unit (ALU)
- Program counter (PC)
- Instruction register (IR)



Memory

Holds data and instructions

- 256 words
- 16 bits in each word
- Connected to registers
- Words are *addressable*

Use *hexadecimal* for addresses

- Number words from 00 to FF
- *Think in hexadecimal*
- Use array notation
- Example: `M[2A]` = **C024**

Memory

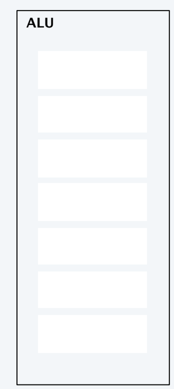
00	0000	10	8A01	20	7101		F0	F0F0
01	FFFE	11	8B02	21	8AFF		F1	0505
02	000D	12	1CAB	22	7680		F2	000D
03	0003	13	9C03	23	7B00		F3	1000
04	0001	14	0001	24	CA2B		F4	0101
05	0000	15	0010	25	8CFF		F5	0010
06	0000	16	0100	26	156B		F6	0001
07	0000	17	1000	27	BC05	...	F7	0010
08	0000	18	0100	28	2AA1		F8	0100
09	0000	19	0010	29	2BB1		F9	1000
0A	0000	1A	0001	2A	C024		FA	0100
0B	0000	1B	0010	2B	0000		FB	0010
0C	0000	1C	0100	2C	0000		FC	0001
0D	0000	1D	1000	2D	0000		FD	0010
0E	0000	1E	0100	2E	0000		FE	0100
0F	0000	1F	0010	2F	0000		FF	0100

Table of 256 words *completely specifies* contents of memory.

Arithmetic and logic unit (ALU)

ALU

- TOY's computational engine
- A *calculator*, not a computer
- Hardware that implements *all* data-type operations
- How? Stay tuned for computer architecture lectures



Registers

Registers

- 16 words, addressable in hex from 0 to F (use names R[0] through R[F])
- Scratch space for calculations and data movement.
- Connected to memory and ALU
- *By convention,* R[0] *is always* 0. ← often simplifies code (stay tuned) In our code, we often also keep 0001 in R[1].

Q. Why not just connect memory directly to ALU?

A. Too many different memory names (addresses).

Q. Why not just connect memory locations to one another?

A. Too many different connections.

Table of 16 words *completely specifies* contents of registers.

Registers	
R[0]	0 0 0 0
R[1]	0 0 0 1
R[2]	F F F E
R[3]	1 C A B
R[4]	0 0 0 1
R[5]	0 0 0 0
R[6]	F A C E
R[7]	0 0 0 0
R[8]	F 0 0 1
R[9]	0 0 0 0
R[A]	0 0 0 5
R[B]	0 0 0 8
R[C]	0 0 0 D
R[D]	0 0 0 0
R[E]	0 0 0 0
R[F]	0 0 0 0

Program counter and instruction register

TOY operates by executing a sequence of instructions.

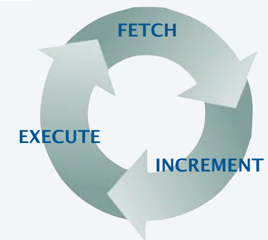
Critical abstractions in making this happen

- **Program Counter (PC).** Memory address of next instruction.
- **Instruction Register (IR).** Instruction being executed.

PC	IR
10	1 C A B

Fetch-increment-execute cycle

- Fetch: Get instruction from memory into IR.
- Increment: Update PC to point to *next* instruction.
- Execute: Move data to or from memory, change PC, or perform calculations, as specified by IR.



The state of the machine

Contents of memory, registers, and PC at a particular time

- Provide a record of what a program has done.
- Completely determines what the machine will do.

ALU and IR hold intermediate states of computation

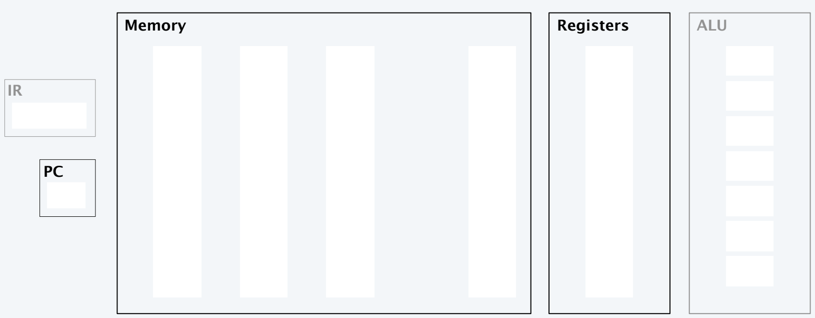


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Image sources

http://pixabay.com/en/man-flashlight-helmet-detective-308611/
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CS.17.B.MachineI.Types

TOY data type

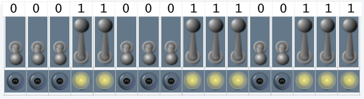
A data type is a set of values and a set of operations on those values.

TOY's data type is 16-bit *two's complement* integers.

Two kinds of operations
- Arithmetic.
- Bitwise.

All other types of data must be implemented with *software*
- 32-bit and 64-bit integers.
- 32-bit and 64-bit floating point values.
- Characters and strings.
- ...

All values are represented in 16-bit words.



TOY data type (original design): Unsigned integers

Values. 0 to $2^{16}-1$, encoded in binary (or, equivalently, hex).

Example. 6375_{10}.

	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
binary	0	0	0	1	1	0	0	0	1	1	1	0	0	1	1	1
				2^{12}	$+2^{11}$				$+2^7$	$+2^6$	$+2^5$			$+2^2$	$+2^1$	$+2^0$

hex	1	8	E	7
	1×16^3	$+ 8 \times 16^2$	$+ 14 \times 16$	$+ 7$
	4096	+ 2048	+ 224	+ 7

Operations.
- Add.
- Subtract.
- Test if 0.

Example. 18E7 + 18E7 = 31CE

	0	0	0	1	1	0	0	0	1	1	1	0	0	1	1	1
+	0	0	0	1	1	0	0	0	1	1	1	0	0	1	1	1
=	0	0	1	1	0	0	0	1	1	1	0	0	1	1	1	0

Warning. TOY ignores overflow.

TOY data type (better design): two's complement

Values. -2^{15} to $2^{15}-1$, encoded in *16-bit two's complement*.

includes negative integers!

Operations.
- Add.
- Subtract.
- Test if positive, negative, or 0.

16 bit two's complement
- 16-bit binary representation of x for positive x.
- 16-bit binary representation of $2^{16} - |x|$ for negative x.

Useful properties
- Leading bit (bit 15) signifies sign.
- 0000000000000000 represents zero.
- Add/subtract is *the same* as for unsigned.

decimal	*hex*	*binary*
+32,767	7FFF	0111111111111111
+32,766	7FFE	0111111111111110
+32,765	7FFD	0111111111111101
...		
+3	0003	0000000000000011
+2	0002	0000000000000010
+1	0001	0000000000000001
0	0000	0000000000000000
−1	FFFF	1111111111111111
−2	FFFE	1111111111111110
−3	FFFD	1111111111111101
...		
−32,766	8002	1000000000000010
−32,767	8001	1000000000000001
−32,768	8000	1000000000000000

slight annoyance: one extra negative value

Two's complement: conversion

To convert from decimal to two's complement

- If greater than +32,767 or less than −32,768 report error.
- Convert to 16-bit binary.
- If not negative, done.
- If negative, *flip all bits and add 1*.

Examples

$+13_{10}$	0000000000001101	000D
-13_{10}	1111111111110011	FFF3
$+256_{10}$	0000000100000000	0100
-256_{10}	1111111100000000	FF00

To convert from two's complement to decimal

- If sign bit is 1, *flip all bits and add 1* and output minus sign.
- Convert to decimal.

Examples

0001	0000000000000001	1_{10}
FFFF	1111111111111111	-1_{10}
FF0D	1111111100001101	-243_{10}
00F3	0000000011110011	$+243_{10}$

To add/subtract

- Use same rules as for unsigned binary.
- (Still) ignore overflow.

Example

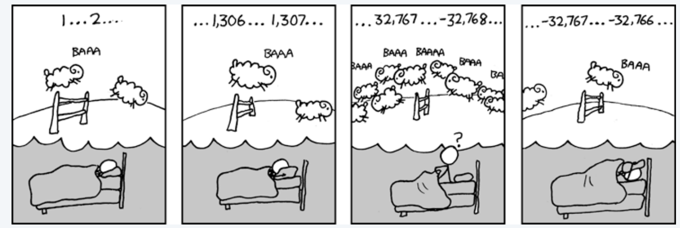
-256_{10}	1111111100000000	FF00
$+13_{10}$	+0000000000001101	+000D
$=-243_{10}$	=1111111100001101	=FF0D

Overflow in two's complement

$32{,}767_{10} = 2^{15}-1$	0111111111111111	7FFF		
+1	+ 0000000000000001	+ 0001		
	= 1000000000000000	= 8000	$= -2^{15} = -32{,}768_{10}$	

largest (positive) number

smallest (negative) number



http://xkcd.com/571/

TOY data type: Bitwise operations

Operations

- Bitwise AND.
- Bitwise XOR.
- Shift left.
- Shift right.

	0	1	0	1	1	0	0	1	0	1	0	0	1	0	0	0
AND	0	0	0	1	1	1	1	1	0	0	0	0	0	1	0	1
=	0	0	0	1	1	0	0	1	0	0	0	0	0	0	0	0

x	*y*	*x* AND *y*
0	0	0
0	1	0
1	0	0
1	1	1

	0	1	0	1	1	0	0	1	0	1	0	0	1	0	0	0
XOR	0	0	0	1	1	1	1	1	0	0	0	0	0	1	0	1
=	0	1	0	0	0	1	1	0	0	1	0	0	1	1	0	1

x	*y*	*x* XOR *y*
0	0	0
0	1	1
1	0	1
1	1	0

	0	0	0	0	1	1	1	0	0	0	0	0	1	0	0	1
Shift left 3	0	1	1	1	0	0	0	0	0	1	0	0	1	0	0	0

fill with 0s

	0	0	0	0	1	1	1	0	0	0	0	0	1	0	0	1
Shift right 3	0	0	0	0	0	0	0	1	1	1	0	0	0	0	0	1

fill with 0s

Special note: Shift left/right operations also implement multiply/divide by powers of 2 for integers.

shift right fills with 1s if leading bit is 1

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Image sources

http://pixabay.com/en/network-media-binary-computer-65923/

https://xkcd.com/571/

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CS.17.C.MachineI.Instructions

TOY instructions

ANY 16-bit (4 hex digit) value defines a TOY instruction.

First hex digit specifies which instruction.

Each instruction changes machine state in a well-defined way.

category	*opcodes*	*implements*	*changes*
operations	1 2 3 4 5 6	data-type operations	registers
data movement	7 8 9 A B	data moves between registers and memory	registers, memory
flow of control	0 C D E F	conditionals, loops, and functions	PC

opcode	*instruction*
0	halt
1	add
2	subtract
3	bitwise and
4	bitwise xor
5	shift left
6	shift right
7	load address
8	load
9	store
A	load indirect
B	store indirect
C	branch if zero
D	branch if positive
E	jump register
F	jump and link

Encoding instructions

ANY 16-bit (4 hex digit) value defines a TOY instruction.

Two different instruction formats

- Type RR: Opcode and 3 registers.

15 14 13 12 | 11 10 9 8 | 7 6 5 4 | 3 2 1 0
opcode | *destination d* | *source s* | *source t*

- Type A: Opcode, 1 register, and 1 memory address.

15 14 13 12 | 11 10 9 8 | 7 6 5 4 3 2 1 0
opcode | *destination d* | *address addr*

Examples

1 C A B	*Add* R[A] *to* R[B] *and put result in* R[C].
8 A 1 5	*Load into* R[A] *data from* M[15].

opcode		*instruction*
0	RR	halt
1	RR	add
2	RR	subtract
3	RR	bitwise and
4	RR	bitwise xor
5	RR	shift left
6	RR	shift right
7	A	load address
8	A	load
9	A	store
A	RR	load indirect
B	RR	store indirect
C	A	branch if zero
D	A	branch if positive
E	RR	jump register
F	A	jump and link

A TOY program

Add two integers
- Load operands from memory into registers.
- Add the registers.
- Put result in memory.

Load into R[A] data from M[15]
Load into R[B] data from M[16]
Add R[A] and R[B] and put result into R[C]
Store R[C] into M[17]
Halt

PC: 10 → 11, 12, 13, 14

Memory		
10	8 A 1 5	R[A] ← M[15]
11	8 B 1 6	R[B] ← M[16]
12	1 C A B	R[C] ← R[A] + R[B]
13	9 C 1 7	M[17] ← R[C]
14	0 0 0 0	halt
15	0 0 0 8	
16	0 0 0 5	
17	0 0 0 D	
. . .		

Registers	
. . .	
A	0 0 0 8
B	0 0 0 5
C	0 0 0 D
. . .	

Q. How can you tell whether a word is an instruction?

A. If the PC has its address, it *is* an instruction!

Same program with different data

Add two integers
- Load operands from memory into registers.
- Add the registers.
- Put result in memory.

Load	into R[A] data from M[15]
Load	into R[B] data from M[16]
Add	R[A] and R[B] and put result into R[C]
Store	R[C] into M[17]
Halt	

Q. How can you tell whether a word is data ?

A. If it is added to another word, it *is* data !

PC: 10 → (11, 12, 13, 14)

Memory

10	8 A 1 5	R[A] ← M[15]
11	8 B 1 6	R[B] ← M[16]
12	1 C A B	R[C] ← R[A] + R[B]
13	9 C 1 7	M[17] ← R[C]
14	0 0 0 0	halt
15	8 B 1 6	
16	1 C A B	
17	A 7 C 1	
...		

instruction (M[12] 1 C A B)
data (M[16] 1 C A B)

Registers

...		
A	8 B 1 6	$-29,930_{10}$
B	1 C A B	$7,339_{10}$
C	A 7 C 1	$-22,591_{10}$
...		

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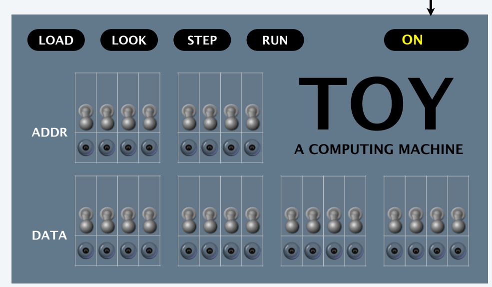
Outside the box

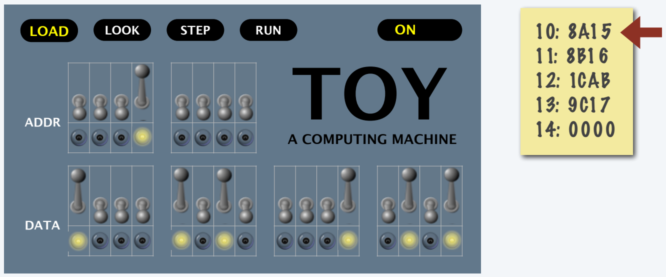
User interface
- Switches.
- Lights.
- Control Buttons.

First step: Turn on the machine!



Loading a program into memory

To load an instruction
- Set 8 memory address switches.
- Set 16 data switches to instruction encoding.
- Press LOAD to load instruction from switches into addressed memory word.



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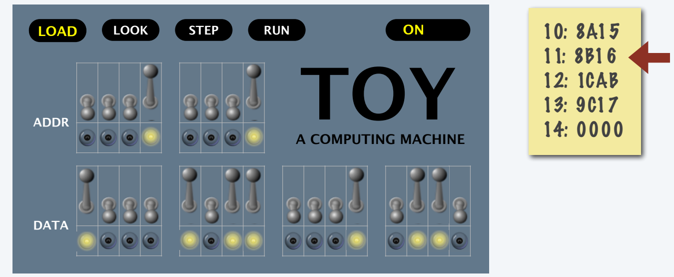

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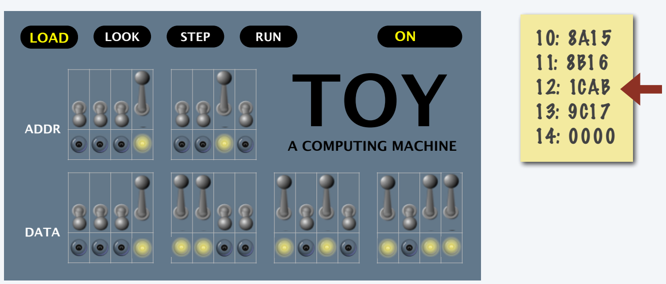

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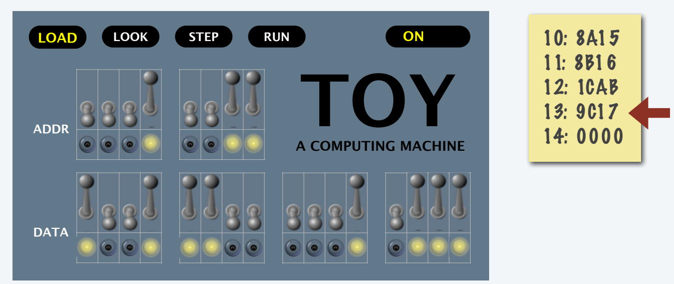

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To load an instruction

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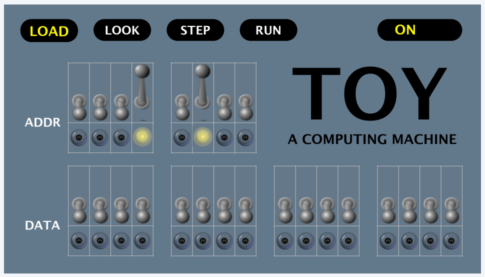
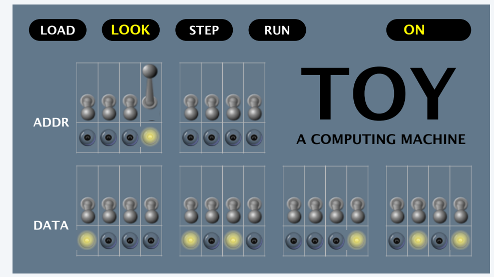


10: 8A15
11: 8B16
12: 1CAB
13: 9C17
14: 0000

Looking at what's in the memory

To double check that you loaded the data correctly

- Set 8 memory address switches.
- Press LOOK to examine the addressed memory word.

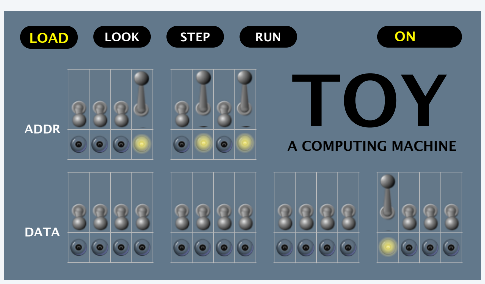


10: 8A15
11: 8B16
12: 1CAB
13: 9C17
14: 0000

Loading data into memory

To load data, use the *same* procedure as for instructions

- Set 8 memory address switches.
- Set 16 data switches to *data* encoding.
- Press LOAD to load *data* from switches into addressed memory word.



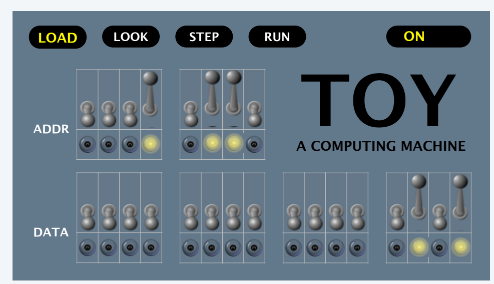
10: 8A15
11: 8B16
12: 1CAB
13: 9C17
14: 0000

15: 0008
16: 0005

Loading data into memory

To load data, use the *same* procedure as for instructions

- Set 8 memory address switches.
- Set 16 data switches to *data* encoding.
- Press LOAD to load *data* from switches into addressed memory word.



10: 8A15
11: 8B16
12: 1CAB
13: 9C17
14: 0000

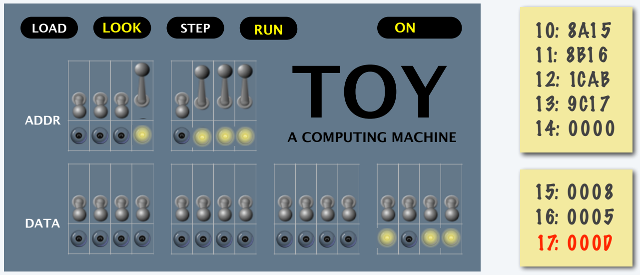
15: 0008
16: 0005

Running a program

To run a program, set the address switches to the address of first instruction and press RUN.

[data lights may flash, but all (and RUN light) go off when HALT instruction is reached]

To see the output, set the address switches to the address of expected result and press LOOK.

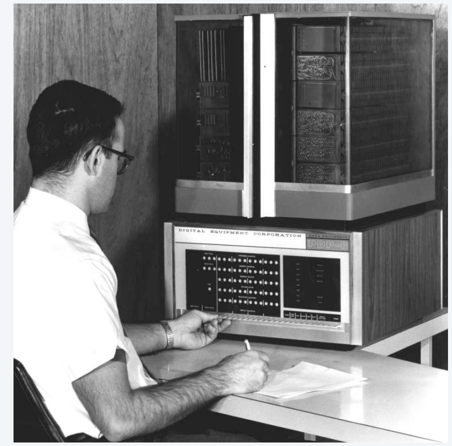


To run the program again, enter different data and press RUN again.

Switches and lights

Q. Did people really program this way?

A. Yes!



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CS.17.E.MachineI.Programming

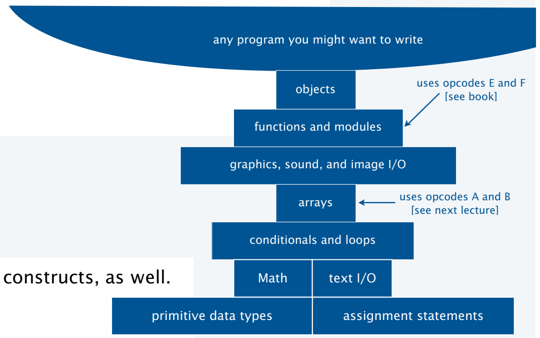
Machine language programming

TOY instructions support the same basic programming constructs as Java.

- Primitive data types.
- Assignment statements.
- Conditionals and loops.
- Arrays (next lecture).
- Standard input and output (next).

and can support advanced programming constructs, as well.

- Functions and libraries (see text).
- Linked structures (see text).



Conditionals and loops

To control the flow of instruction execution

- Test a register's value.
- Change the PC, depending on the value.

opcode	*instruction*
C	branch if zero
D	branch if positive

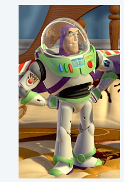
Example: Absolute value of R[A]

10	**DA12**	If R[A] > 0 set PC to 12 (skip 11)
11	**2A0A**	Subtract R[A] from 0 (R[0]) and put result into R[A]
12	**...**	

Example: Typical `while` loop (assumes R[1] is 0001)

10	**CA15**	If R[A] is 0 set PC to 15
11	**...**	
12	**...**	
13	**2AA1**	Decrement R[A] by 1
14	**C010**	Set PC to 10
15	**...**	

```
while (a != 0) {
  ...
  ...
    a--;
}
```

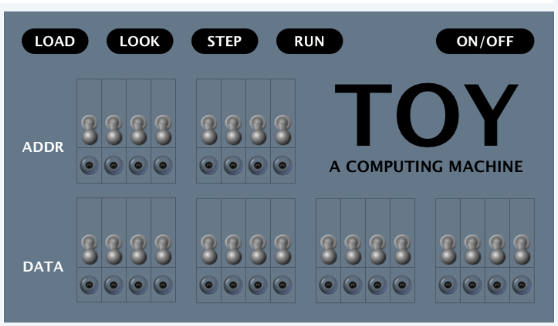


To infinity and beyond!

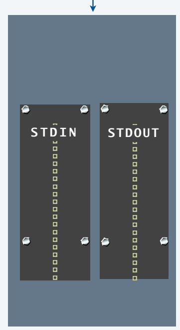
Standard input and output

An immediate problem

- We're not going to be able to address real-world problems with just switches and lights for I/O!
- One solution: Paper tape.



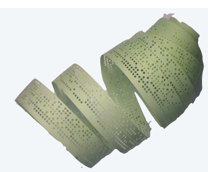
Need to bolt new I/O devices to the side of the machine.

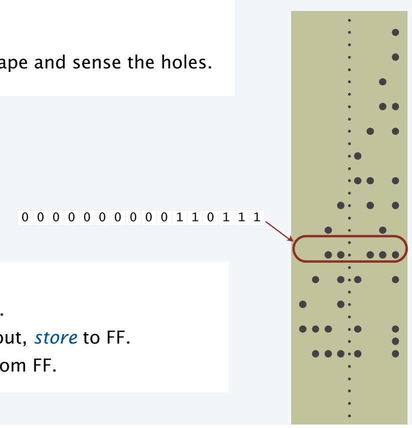


Standard input and output

Punched paper tape

- Encode 16-bit words in two 8-bit rows.
- To *write* a word, *punch a hole* for each 1.
- To *read* a word, shine a light behind the tape and sense the holes.





TOY mechanism

- Connect hardware to memory location FF.
- To *write* the contents of a register to stdout, *store* to FF.
- To *read* from stdin into a register, *load* from FF.

Flow of control and standard output example: Fibonacci numbers

Register trace

A	1	1	2	3	5	8	13	21	34	55	89
B	1	2	3	5	8	13	21	34	55	89	144
C	2	3	3	5	8	13	21	34	55	89	144
9	A	9	8	7	6	5	4	3	2	1	0

PC →	40	7 1 0 1	R[1] = 1
	41	7 A 0 1	R[A] = 1
	42	7 B 0 1	R[B] = 1
	43	8 9 4 C	R[9] = M[4C]
	44	C 9 4 B	if (R[9] == 0) PC = 4B
	45	9 A F F	write R[A] to stdout
	46	1 C A B	R[C] = R[A] + R[B]
	47	1 A B 0	R[A] = R[B]
	48	1 B C 0	R[B] = R[C]
	49	2 9 9 1	R[9] = R[9] - 1
	4A	C 0 4 4	PC = 44
	4B	0 0 0 0	halt
	4C	0 0 0 A	

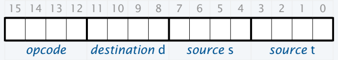
...

```
a = 1;
b = 1;
i = N;
while (i != 0) {
   StdOut.print(a);
   c = a + b;
   a = b;
   b = c;
   i = i - 1;
}
```

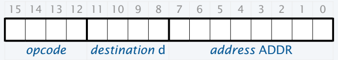
TOY reference card

opcode	operation	format	pseudo-code
0	halt	—	halt
1	add	RR	R[d] = R[s] + R[t]
2	subtract	RR	R[d] = R[s] - R[t]
3	bitwise and	RR	R[d] = R[s] & R[t]
4	bitwise xor	RR	R[d] = R[s] ^ R[t]
5	shift left	RR	R[d] = R[s] << R[t]
6	shift right	RR	R[d] = R[s] >> R[t]
7	load addr	A	R[d] = addr
8	load	A	R[d] = M[addr]
9	store	A	M[addr] = R[d]
A	load indirect	RR	R[d] = M[R[t]]
B	store indirect	RR	M[R[t]] = R[d]
C	branch zero	A	if (R[d] == 0) PC = addr
D	branch positive	A	if (R[d] > 0) PC = addr
E	jump register	RR	PC = R[d]
F	jump and link	A	R[d] = PC + 1; PC = addr

Format RR



Format A



ZERO R[0] is always 0.

STANDARD INPUT Load from FF.

STANDARD OUTPUT Store to FF.

Pop quiz 1 on TOY

Q. What is the interpretation of

1A75 as a TOY instruction?

1A75 as a two's complement integer value?

0FFF as a TOY instruction?

0FFF as a two's complement integer value?

8888 as a TOY instruction?

8888 as a two's complement integer value? (Answer in base 16).

Pop quiz 2 on TOY

Q. How does one flip all the bits in a TOY register ?

Pop quiz 3 on TOY

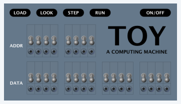
Q. What does the following TOY program leave in R[2] ?

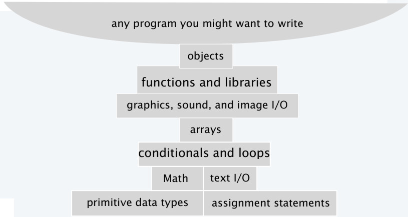
```
10  7 C 0 A    R[C] = 10₁₀
11  7 1 0 1    R[1] = 1
12  7 2 0 1    R[2] = 1
13  1 2 2 2    R[2] = R[2] + R[2]
14  2 C C 1    R[C] = R[C] - 1
15  D C 1 3    if (R[C] > 0) PC = 13
16  0 0 0 0    HALT
```


TOY vs. your laptop

Two different computing machines

- Both implement basic data types, conditionals, loops, and other low-level constructs.
- Both can have arrays, functions, and other high-level constructs.
- Both have infinite input and output streams.





Q. Is 256 words enough to do anything useful?

A. Yes! (See book, and stay tuned for next lecture.)

A. Yes! It is a Turing Machine, with a read/write I/O device (see theory lectures).

OK, we definitely want a faster version with more memory when we can afford it...

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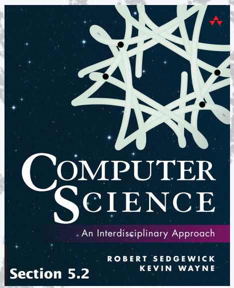

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17. A Computing Machine