

Have seen

- Given Inverted index, how compute the results for a query

 Merge-based algorithms
- Data structure for accessing inverted index

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- Hash table
- B+ tree







Construction of posting lists

Overview

- "document" now means preprocessed document
- One pass through collection of documents
- Gather postings for each document
- Reorganize for final set of lists: one for each term
- Look at algorithms when can't fit everything in memory
 - Main cost file page reads and writes
 - "file page" minimum unit can read from drive
 - May be multiple of "sector" device constraint

Memory- disk management

- · Have buffer in main memory
 - Size = B file pages
 - Read from disk to buffer, page at a time
 Disk cost = 1 per page
 - Write from buffer to disk, page at at time
 Disk cost = 1 per page

Sorting List on Disk - External Sorting General techique

- Divide list into size-B blocks of contiguous entries
- Read each block into buffer, sort, write out to disk
- Now have [L/B] sorted sub-lists where L is size of list in file pages
- Merge sorted sub-lists into one list – How?

Merging Lists on Disk: General technique

- · K sorted lists on disk to merge into one
- If K+1 <= B:
 - Dedicate one buffer page for output
 - Dedicate one buffer page for each list to merge input from different lists
 - Algorithm:
 - Fill 1 buffer page from each list on disk Repeat until merge complete: Merge buffer input pages to output buffer pg

When output buffer pg full, write to disk When input buffer pg empty, refill from its list

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• If K+1 > B:

- Dedicate one buffer page for output
- B-1 buffer page for input from different lists
- Define "level-0 lists": lists need to merge

If K+1 > B: Algorithm

Number of file page read/writes?

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- · Merge K sorted lists?
 - Merge-tree height = [log_{B-1}K]
 - Read/write all lists once each level.
 - Ignore breakage
 - Read/write all lists $\lceil \text{log}_{\text{B-1}} K \rceil$ times total
- External sort length L list?
 - Create [L/B] sorted sub-lists: L reads/writes
 - Merge [L/B] sorted sub-lists:
 - L *[log_{B-1} [L/B]] reads/writes
 - Total # page read/writes = O(L log_{B-1}L) 12

So far

- · Preprocessing the collection
- Sorting a list on disk (external sorting) – Cost as disk I/O

Now look at actually building

Index building Algorithm: "Block Sort-based" 1. Repeat until entire collection read: - Read documents, building (term, <attributes>, doc) tuples until buffer full • one tuple for each occurrence of a term - Sort tuples in buffer by term value as primary, doc as secondary • Tuples for one doc already together • Use sort algorithm that keeps appearance order for = keys: stable sorting - Build posting lists for each unique term in buffer • Re-writing of sorted info

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- Write partial index to disk

continuing "Blocked Sort-based"

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- 2. Merge partial indexes on disk into full index
- Partial index lists of (term:postings list) entries must be merged
- Partial postings lists for one term must be merged
 - Concatenate
 - Keep documents sorted within posting list
- If postings for one document broken across partial lists, must merge

Remarks: Index Building

- As build index:
 - Build dictionary
 - Aggregate Information on terms, e.g. document frequency
 store w/ dictionary
 - What happens if dictionary not fit in main memory as build inverted index?
- May not actually keep every term occurrence, maybe just first k.
 - Early Google did this for k=4095. Why?

What about anchor text?

- Complication
- Build separate anchor text index
 - strong relevance indicator
 - keeps index building less complicated

Other separate indexes?

Examples

- · Other strong relevance indicators
 - abstracts of documents
 compare listing abstract positions 1st in main index
 - tiered indexes based on term weights
- · types of documents
 - volatility
 - news articles
 - blogs
 - etc.