

# COS 511: Theoretical Machine Learning

Homework #4  
Rademacher & boosting

Due:  
March 25, 2014

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## Problem 1

[15] Let  $\mathcal{F}_1, \dots, \mathcal{F}_n$  be families of real-valued functions on some space  $\mathcal{Z}$ , and let  $a_1, \dots, a_n$  be arbitrary (fixed) real numbers. Let  $\mathcal{G}$  be the class of all functions  $g$  of the form

$$g(z) = \sum_{i=1}^n a_i f_i(z)$$

where  $f_i \in \mathcal{F}_i$  for  $i = 1, \dots, n$ . For any sample  $S$ , find  $\hat{\mathcal{R}}_S(\mathcal{G})$  exactly in terms of  $a_1, \dots, a_n$ , and  $\hat{\mathcal{R}}_S(\mathcal{F}_1), \dots, \hat{\mathcal{R}}_S(\mathcal{F}_n)$ . Be sure to justify your answer.

## Problem 2

[15] Suppose, in the usual boosting set-up, that the weak learning condition is guaranteed to hold so that  $\epsilon_t \leq \frac{1}{2} - \gamma$  for some  $\gamma > 0$  which is *known* before boosting begins. Describe a modified version of AdaBoost whose final classifier is a simple (unweighted) majority vote, and show that its training error is at most  $(1 - 4\gamma^2)^{T/2}$ .

## Problem 3

Let  $\mathcal{X}_n = \{0, 1\}^n$ , and let  $\mathcal{G}_n$  be any class of boolean functions  $g : \mathcal{X}_n \rightarrow \{-1, +1\}$ . In this problem, we will see, roughly speaking, that if a function  $f$  can be written as a majority vote of polynomially many functions in  $\mathcal{G}_n$ , then under any distribution,  $f$  can be weakly approximated by some function in  $\mathcal{G}_n$ . But if  $f$  cannot be so written as a majority vote, then there exists some “hard” distribution under which  $f$  cannot be approximated by *any* function in  $\mathcal{G}_n$ .

Let  $\mathcal{M}_{n,k}$  be the class of all boolean functions that can be written as a simple majority vote of  $k$  (not necessarily distinct) functions in  $\mathcal{G}_n$ ; that is,  $\mathcal{M}_{n,k}$  consists of all functions  $f$  of the form

$$f(x) = \text{sign} \left( \sum_{j=1}^k g_j(x) \right)$$

for some  $g_1, \dots, g_k \in \mathcal{G}_n$ . Assume  $k$  is odd.

- [15] Show that if  $f \in \mathcal{M}_{n,k}$  then for all distributions  $D$  on  $\mathcal{X}_n$ , there exists a function  $g \in \mathcal{G}_n$  for which

$$\Pr_{x \sim D} [f(x) \neq g(x)] \leq \frac{1}{2} - \frac{1}{2k}.$$

- [15] Show that if  $f \notin \mathcal{M}_{n,k}$  then there exists a distribution  $D$  on  $\mathcal{X}_n$  such that for every  $g \in \mathcal{G}_n$ ,

$$\Pr_{x \sim D} [f(x) \neq g(x)] > \frac{1}{2} - \sqrt{\frac{n \ln 2}{2k}}.$$