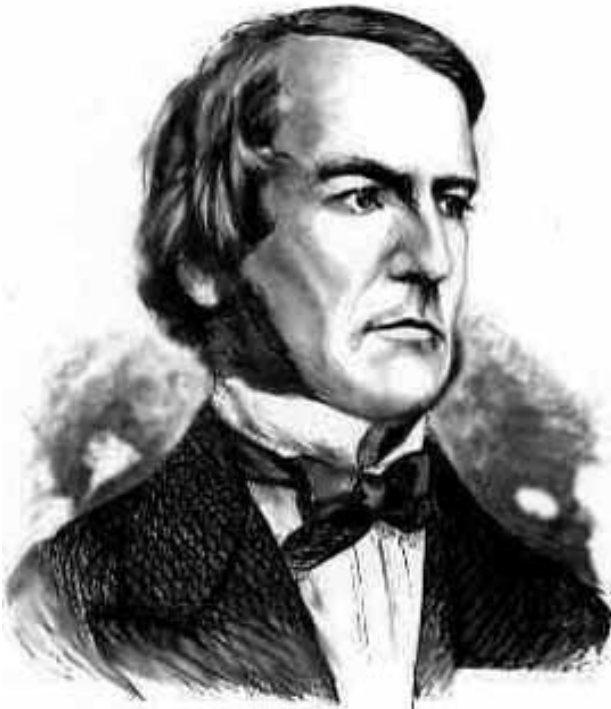
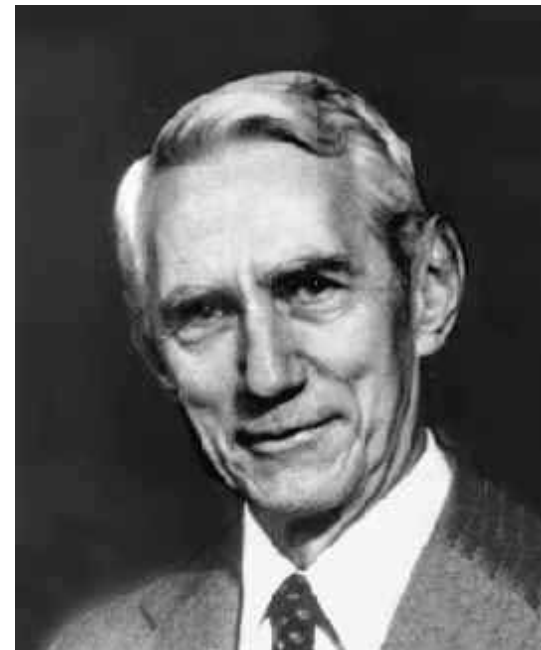


6.1 Combinational Circuits



George Boole (1815 - 1864)



Claude Shannon (1916 - 2001)

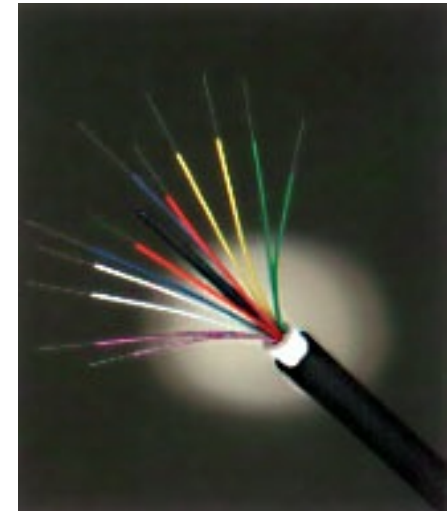
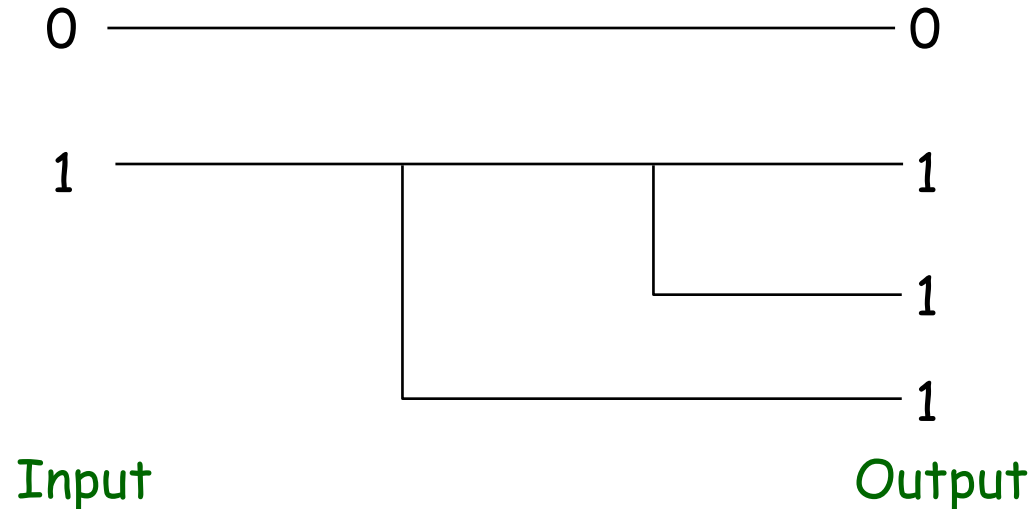
Signals and Wires

Digital signals

- Binary (or "logical") values: 1 or 0, on or off, high or low voltage

Wires.

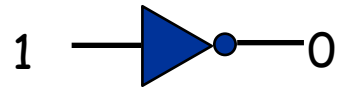
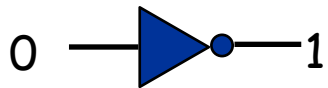
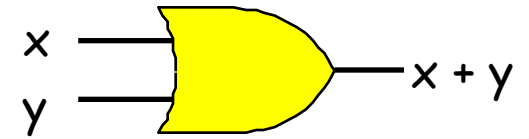
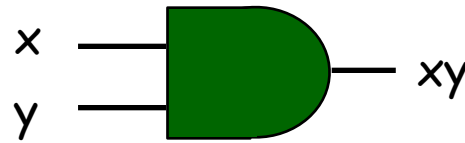
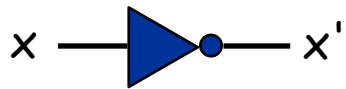
- Propagate logical values from place to place.
- Signals "flow" from left to right.
 - A drawing convention, sometimes violated
 - Actually: flow from producer to consumer(s) of signal



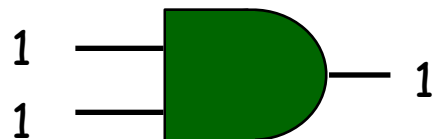
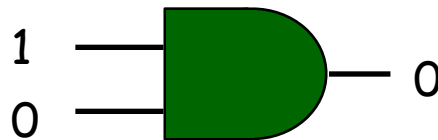
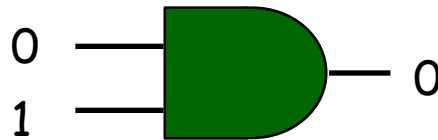
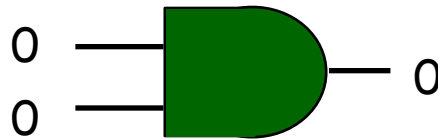
Logic Gates

Logical gates.

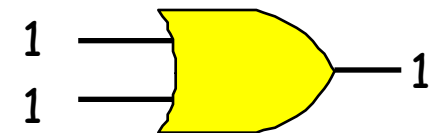
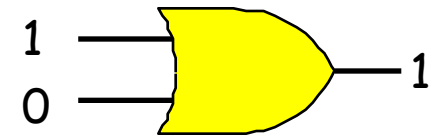
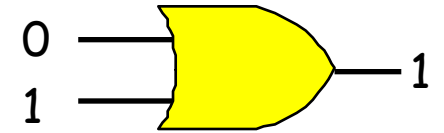
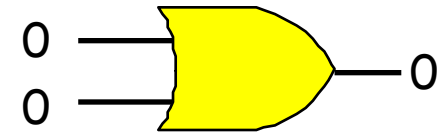
- Fundamental building blocks.



NOT



AND

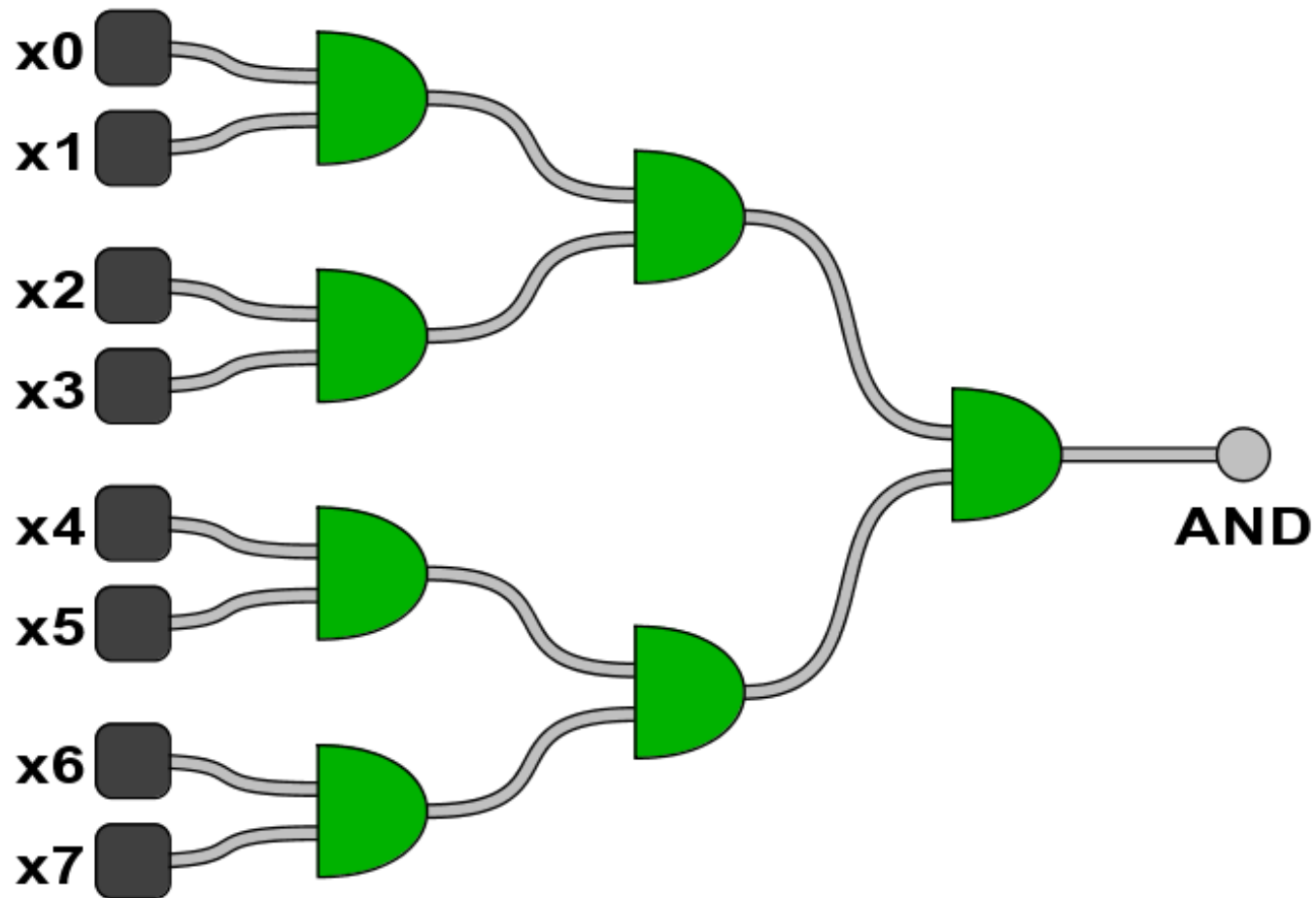


OR

Multiway AND Gates

$AND(x_0, x_1, x_2, x_3, x_4, x_5, x_6, x_7)$.

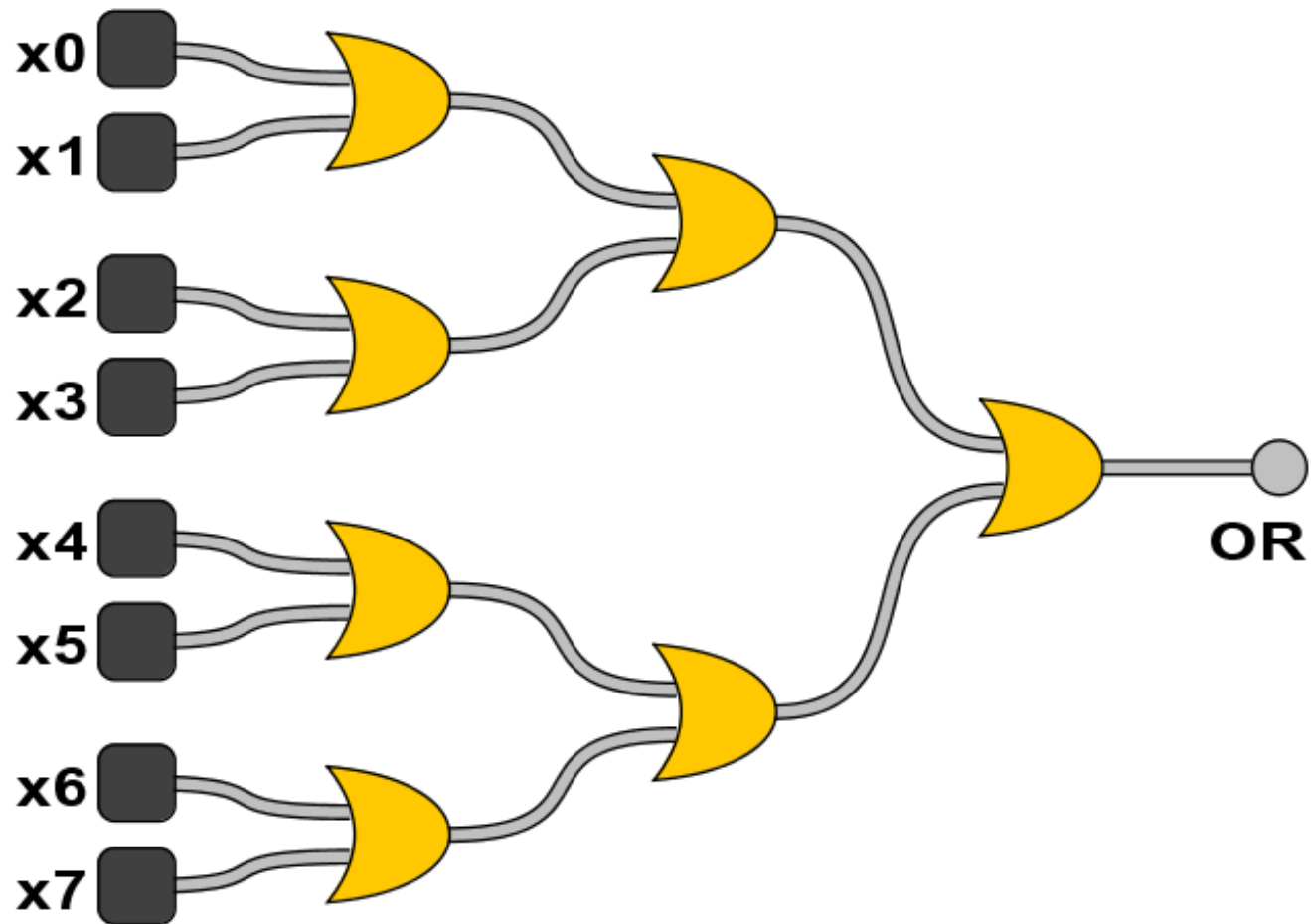
- 1 if all inputs are 1.
- 0 otherwise.



Multiway OR Gates

$OR(x_0, x_1, x_2, x_3, x_4, x_5, x_6, x_7)$.

- 1 if at least one input is 1.
- 0 otherwise.



Boolean Algebra

History.

- Developed by Boole to solve mathematical logic problems (1847).
- Shannon master's thesis applied it to digital circuits (1937).

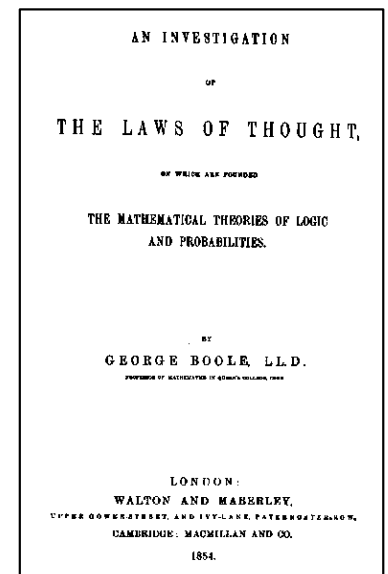
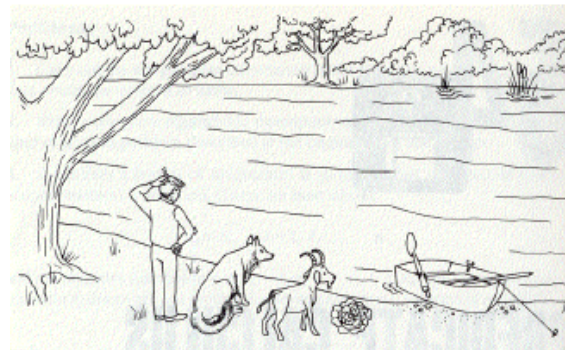
"possibly the most important, and also the most famous, master's thesis of the [20th] century" --Howard Gardner

Basics.

- Boolean variable: value is 0 or 1.
- Boolean function: function whose inputs and outputs are 0, 1.

Relationship to circuits.

- Boolean variables: signals.
- Boolean functions: circuits.

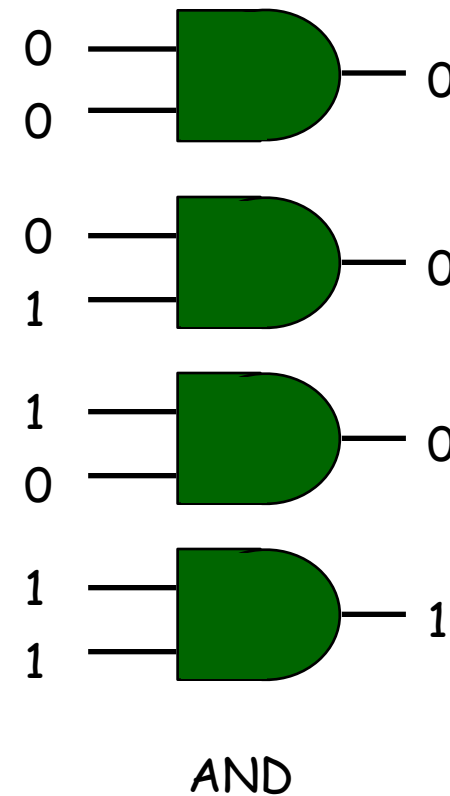


Truth Table

Truth table.

- Systematic method to describe Boolean function.
- One row for each possible input combination.
- N inputs $\Rightarrow 2^N$ rows.

AND Truth Table		
x	y	AND(x, y)
0	0	0
0	1	0
1	0	0
1	1	1



Truth Table for Functions of 2 Variables

Truth table.

- 16 Boolean functions of 2 variables.
 - every 4-bit value represents one

Truth Table for All Boolean Functions of 2 Variables									
x	y	ZERO	AND		x		y	XOR	OR
0	0	0	0	0	0	0	0	0	0
0	1	0	0	0	0	1	1	1	1
1	0	0	0	1	1	0	0	1	1
1	1	0	1	0	1	0	1	0	1

Truth Table for All Boolean Functions of 2 Variables									
x	y	NOR	EQ	y'		x'		NAND	ONE
0	0	1	1	1	1	1	1	1	1
0	1	0	0	0	0	1	1	1	1
1	0	0	0	1	1	0	0	1	1
1	1	0	1	0	1	0	1	0	1

Truth Table for Functions of 3 Variables

Truth table.

- 16 Boolean functions of 2 variables.
 - every 4-bit value represents one
- 256 Boolean functions of 3 variables.
 - every 8-bit value represents one
- $2^{(2^N)}$ Boolean functions of N variables!

Some Functions of 3 Variables						
x	y	z	AND	OR	MAJ	ODD
0	0	0	0	0	0	0
0	0	1	0	1	0	1
0	1	0	0	1	0	1
0	1	1	0	1	1	0
1	0	0	0	1	0	1
1	0	1	0	1	1	0
1	1	0	0	1	1	0
1	1	1	1	1	1	1



Universality of AND, OR, NOT

Any Boolean function can be expressed using AND, OR, NOT.

- "Universal."
- $XOR(x,y) = xy' + x'y$

Notation	Meaning
x'	NOT x
$x y$	x AND y
$x + y$	x OR y

Expressing XOR Using AND, OR, NOT							
x	y	x'	y'	x'y	xy'	x'y + xy'	XOR
0	0	1	1	0	0	0	0
0	1	1	0	1	0	1	1
1	0	0	1	0	1	1	1
1	1	0	0	0	0	0	0

Exercise. Show {AND, NOT}, {OR, NOT}, {NAND}, {AND, XOR} are universal.

Hint. Use DeMorgan's Law: $(xy)' = (x' + y')$ and $(x + y)' = (x'y')$

Sum-of-Products

Any Boolean function can be expressed using AND, OR, NOT.

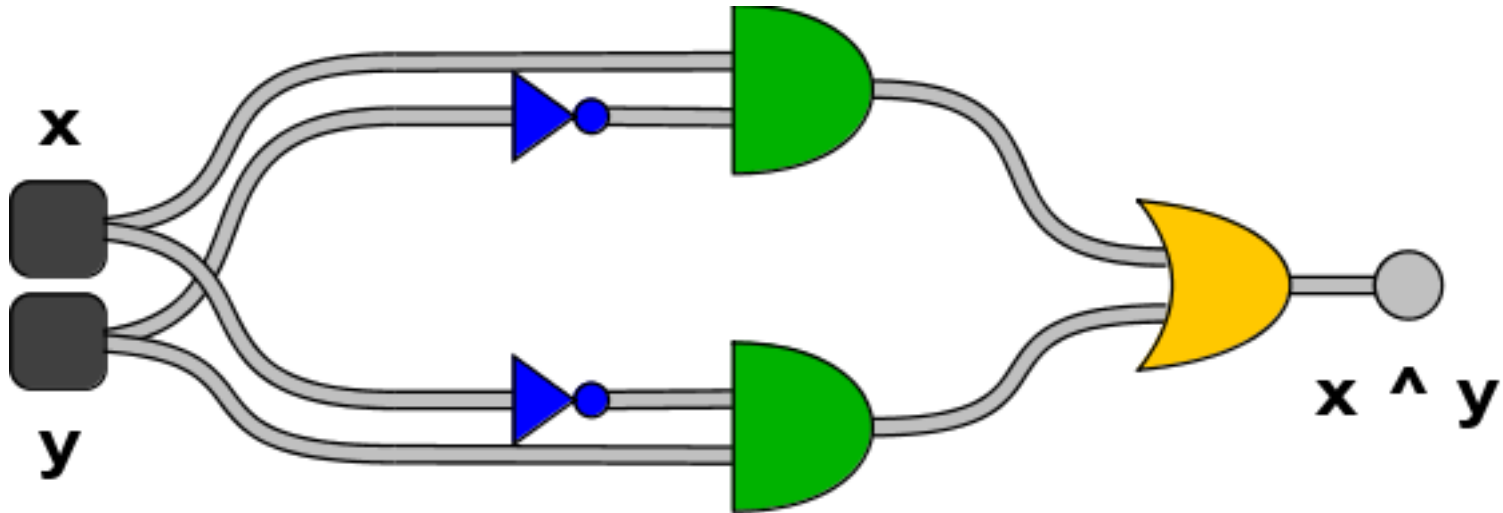
- Sum-of-products is systematic procedure.
 - form AND term for each 1 in truth table of Boolean function
 - OR terms together

Expressing MAJ Using Sum-of-Products								
x	y	z	MAJ	$x'yz$	$xy'z$	xyz'	xyz	$x'yz + xy'z + xyz' + xyz$
0	0	0	0	0	0	0	0	0
0	0	1	0	0	0	0	0	0
0	1	0	0	0	0	0	0	0
0	1	1	1	1	0	0	0	1
1	0	0	0	0	0	0	0	0
1	0	1	1	0	1	0	0	1
1	1	0	1	0	0	1	0	1
1	1	1	1	0	0	0	1	1

Translate Boolean Formula to Boolean Circuit

Use sum-of-products form.

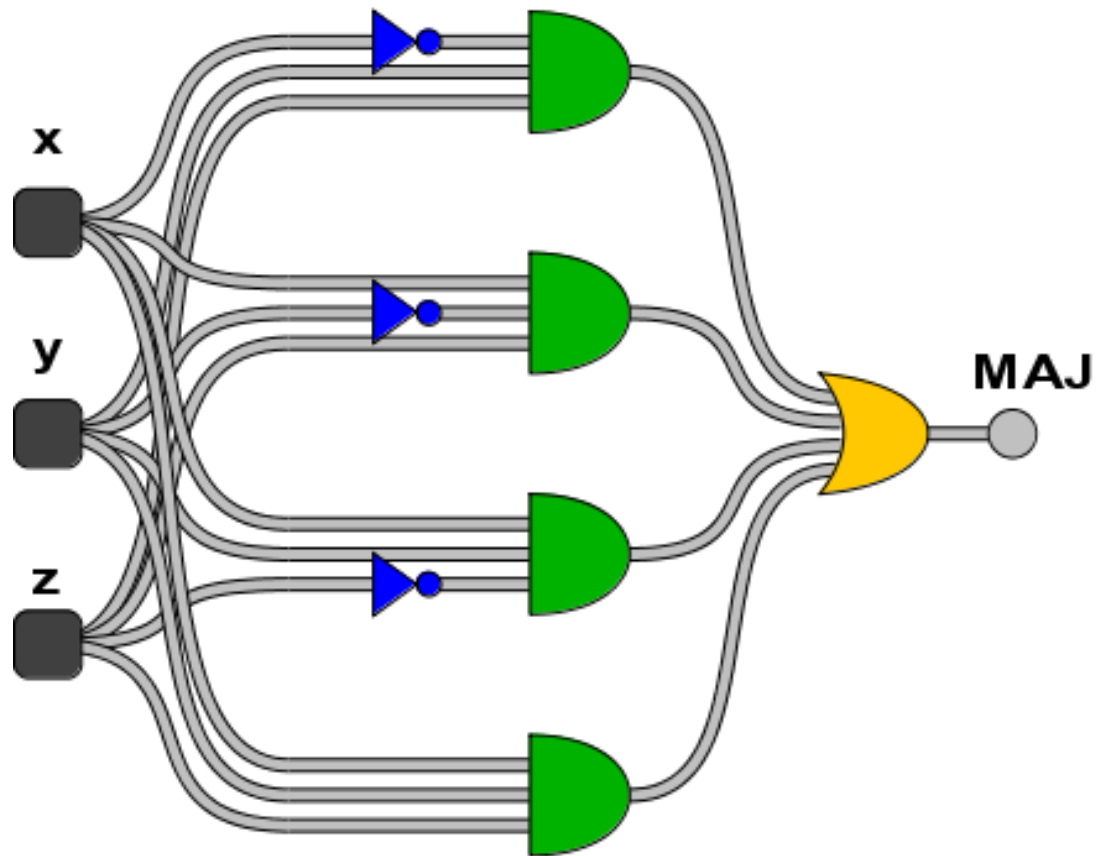
- $XOR(x, y) = xy' + x'y$.



Translate Boolean Formula to Boolean Circuit

Use sum-of-products form.

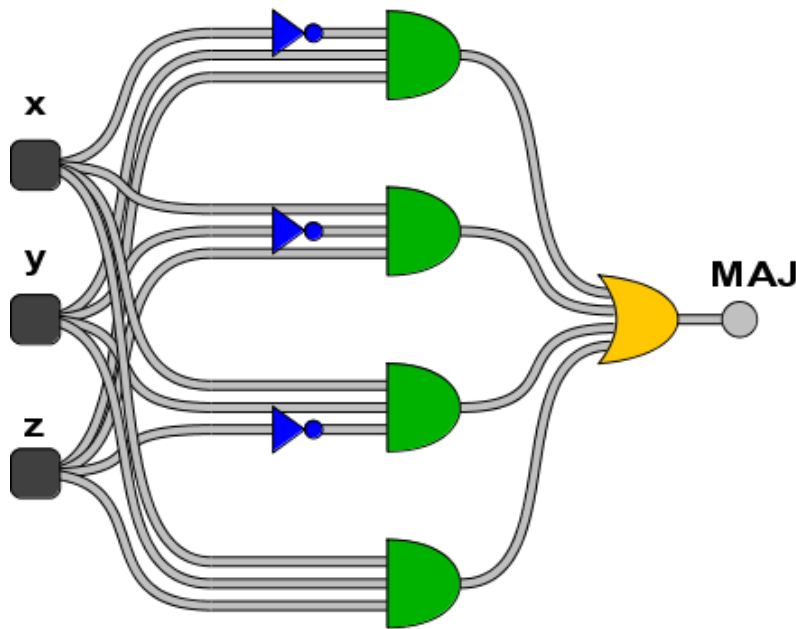
- $MAJ(x, y, z) = x'yz + xy'z + xyz' + xyz.$



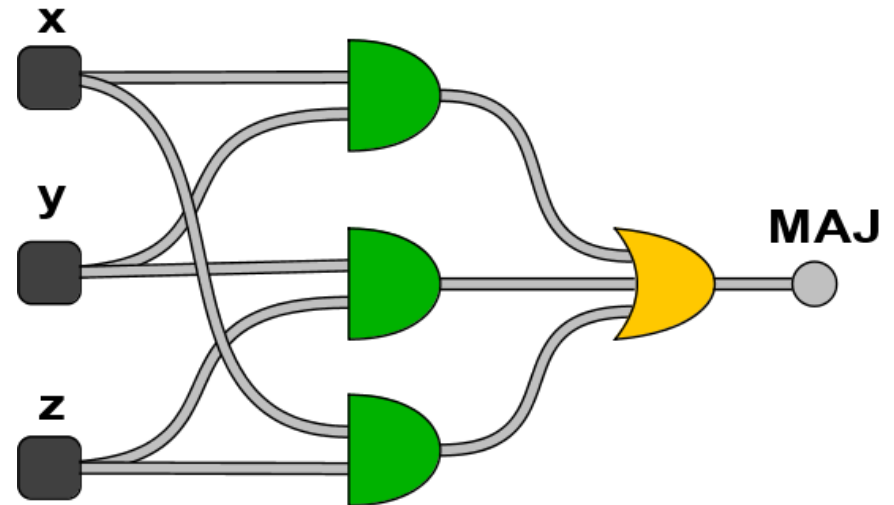
Simplification Using Boolean Algebra

Many possible circuits for each Boolean function.

- Sum-of-products not necessarily optimal in:
 - number of gates (space)
 - depth of circuit (time)
- $MAJ(x, y, z) = x'yz + xy'z + xyz' + xyz = xy + yz + xz.$



size = 8, depth = 3



size = 4, depth = 2

Expressing a Boolean Function Using AND, OR, NOT

Ingredients.

- AND gates.
- OR gates.
- NOT gates.
- Wire.

Instructions.

- Step 1: represent input and output signals with Boolean variables.
- Step 2: construct truth table to carry out computation.
- Step 3: derive (simplified) Boolean expression using sum-of products.
- Step 4: transform Boolean expression into circuit.

ODD Parity Circuit

$ODD(x, y, z)$.

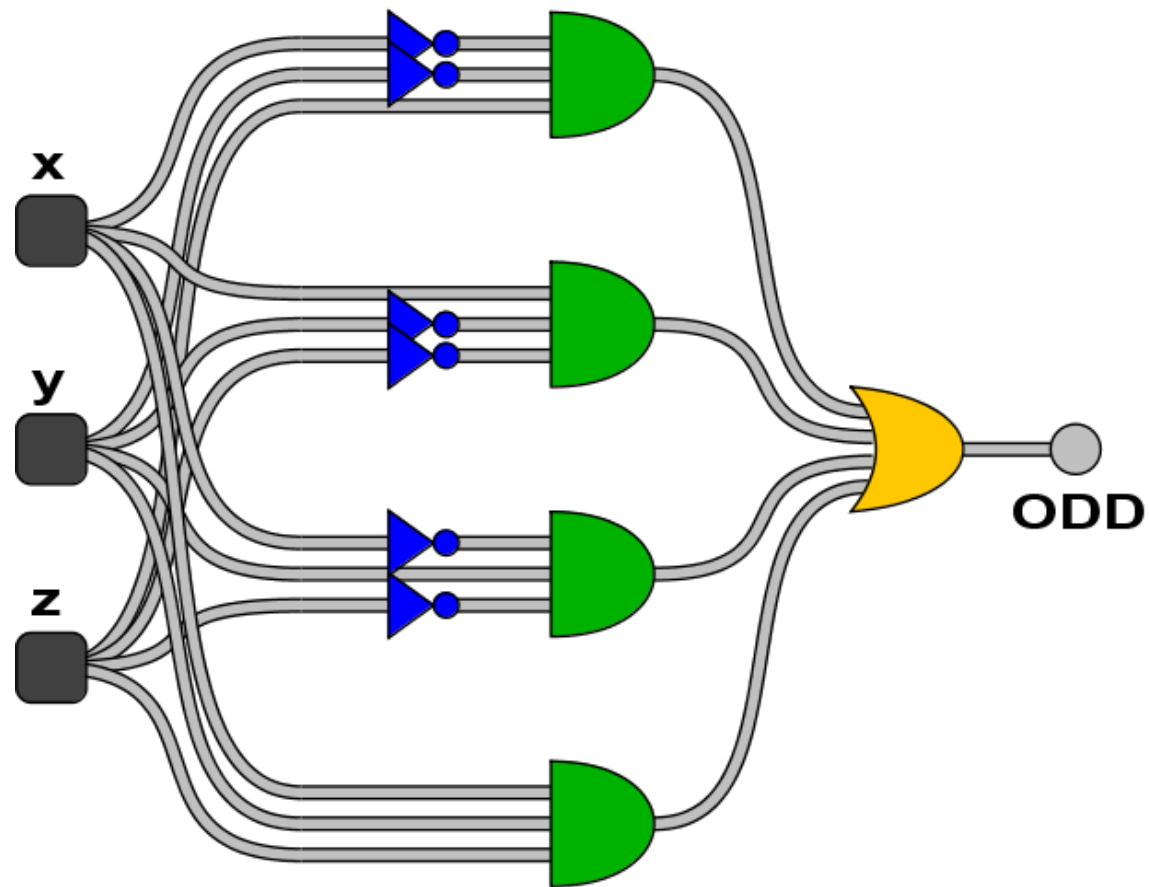
- 1 if odd number of inputs are 1.
- 0 otherwise.

Expressing ODD Using Sum-of-Products								
x	y	z	ODD	$x'y'z$	$x'yz'$	$xy'z'$	xyz	$x'y'z + x'yz' + xy'z' + xyz$
0	0	0	0	0	0	0	0	0
0	0	1	1	1	0	0	0	1
0	1	0	1	0	1	0	0	1
0	1	1	0	0	0	0	0	0
1	0	0	1	0	0	1	0	1
1	0	1	0	0	0	0	0	0
1	1	0	0	0	0	0	0	0
1	1	1	1	0	0	0	1	1

ODD Parity Circuit

$ODD(x, y, z)$.

- 1 if odd number of inputs are 1.
- 0 otherwise.



Let's Make an Adder Circuit

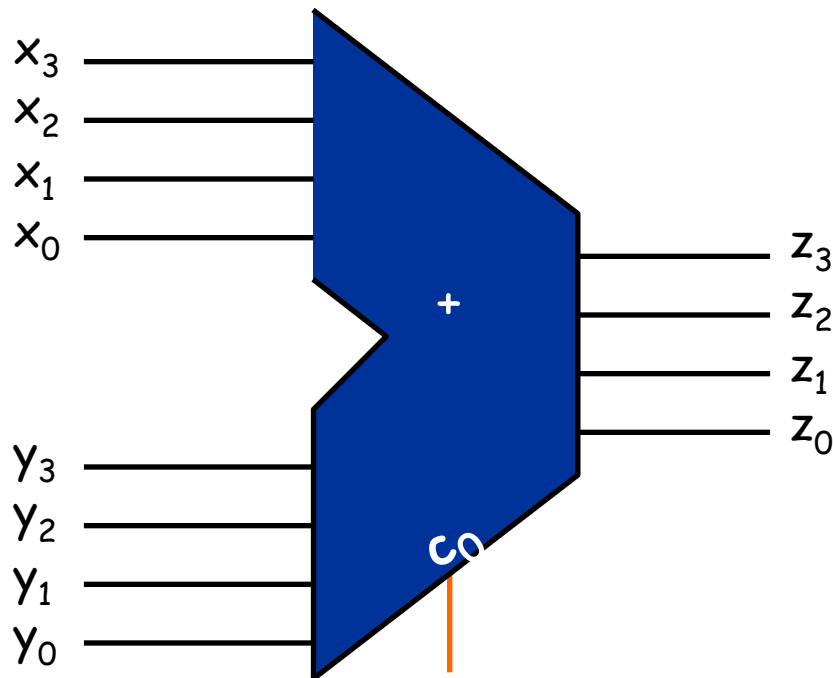
Goal: $x + y = z$ for 4-bit integers.

- We build 4-bit adder: 9 inputs, 4 outputs.
- Same idea scales to 128-bit adder.
- Key computer component.

	1	1	1	0
	2	4	8	7
+	3	5	7	9
<hr/>				
	6	0	6	6

Step 1.

- Represent input and output in binary.



	1	1	0	0
	0	0	1	0
+	0	1	1	1
<hr/>				
	1	0	0	1

	c_3	c_2	c_1	c_0
	x_3	x_2	x_1	x_0
+	y_3	y_2	y_1	y_0
<hr/>				
	z_3	z_2	z_1	z_0

Let's Make an Adder Circuit

Goal: $x + y = z$ for 4-bit integers.

Step 2. (first attempt)

- Build truth table.
- Why is this a bad idea?
 - 128-bit adder: 2^{256+1} rows > # electrons in universe!

				c_0
	x_3	x_2	x_1	x_0
+	y_3	y_2	y_1	y_0
	z_3	z_2	z_1	z_0

4-Bit Adder Truth Table

4-Bit Adder Truth Table												
c_0	x_3	x_2	x_1	x_0	y_3	y_2	y_1	y_0	z_3	z_2	z_1	z_0
0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	0	0	1	0	0	0	1	0
0	0	0	0	0	0	0	1	1	0	0	1	1
0	0	0	0	0	0	1	0	0	0	1	0	0
0	0	0	0	0	0	1	0	1	0	1	0	1
.
1	1	1	1	1	1	1	1	1	1	1	1	1

} $2^{8+1} = 512$ rows!

Let's Make an Adder Circuit

Goal: $x + y = z$ for 4-bit integers.

Step 2. (do one bit at a time)

- Build truth table for carry bit.
- Build truth table for summand bit.

	c_3	c_2	c_1	$c_0 = 0$
	x_3	x_2	x_1	x_0
+	y_3	y_2	y_1	y_0
	z_3	z_2	z_1	z_0

Carry Bit			
x_i	y_i	c_i	c_{i+1}
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

Summand Bit			
x_i	y_i	c_i	z_i
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

Let's Make an Adder Circuit

Goal: $x + y = z$ for 4-bit integers.

Step 3.

- Derive (simplified) Boolean expression.

	c_3	c_2	c_1	$c_0 = 0$
	x_3	x_2	x_1	x_0
+	y_3	y_2	y_1	y_0
	z_3	z_2	z_1	z_0

Carry Bit				
x_i	y_i	c_i	c_{i+1}	MAJ
0	0	0	0	0
0	0	1	0	0
0	1	0	0	0
0	1	1	1	1
1	0	0	0	0
1	0	1	1	1
1	1	0	1	1
1	1	1	1	1

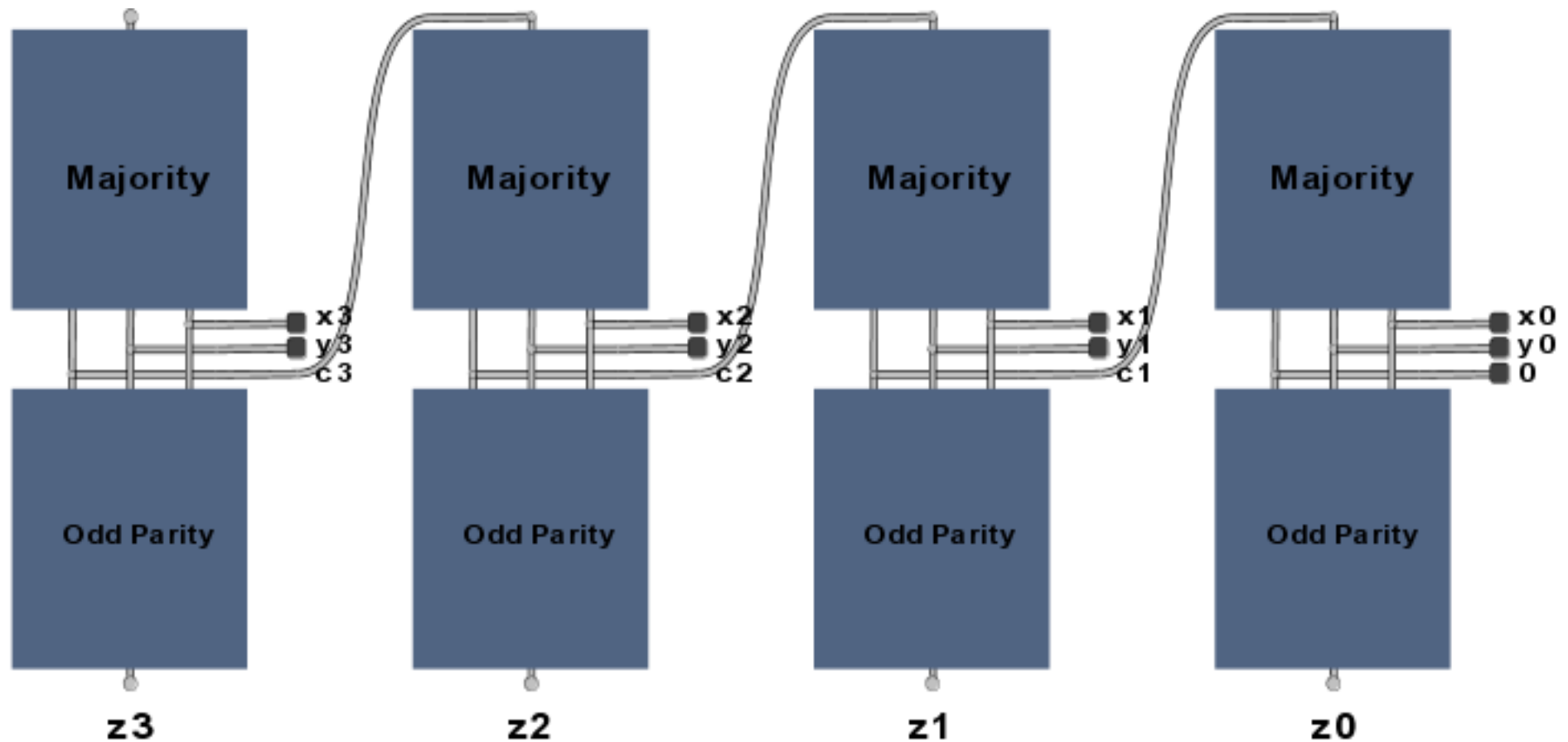
Summand Bit				
x_i	y_i	c_i	z_i	ODD
0	0	0	0	0
0	0	1	1	1
0	1	0	1	1
0	1	1	0	0
1	0	0	1	1
1	0	1	0	0
1	1	0	0	0
1	1	1	1	1

Let's Make an Adder Circuit

Goal: $x + y = z$ for 4-bit integers.

Step 4.

- Transform Boolean expression into circuit.
- Chain together 1-bit adders.

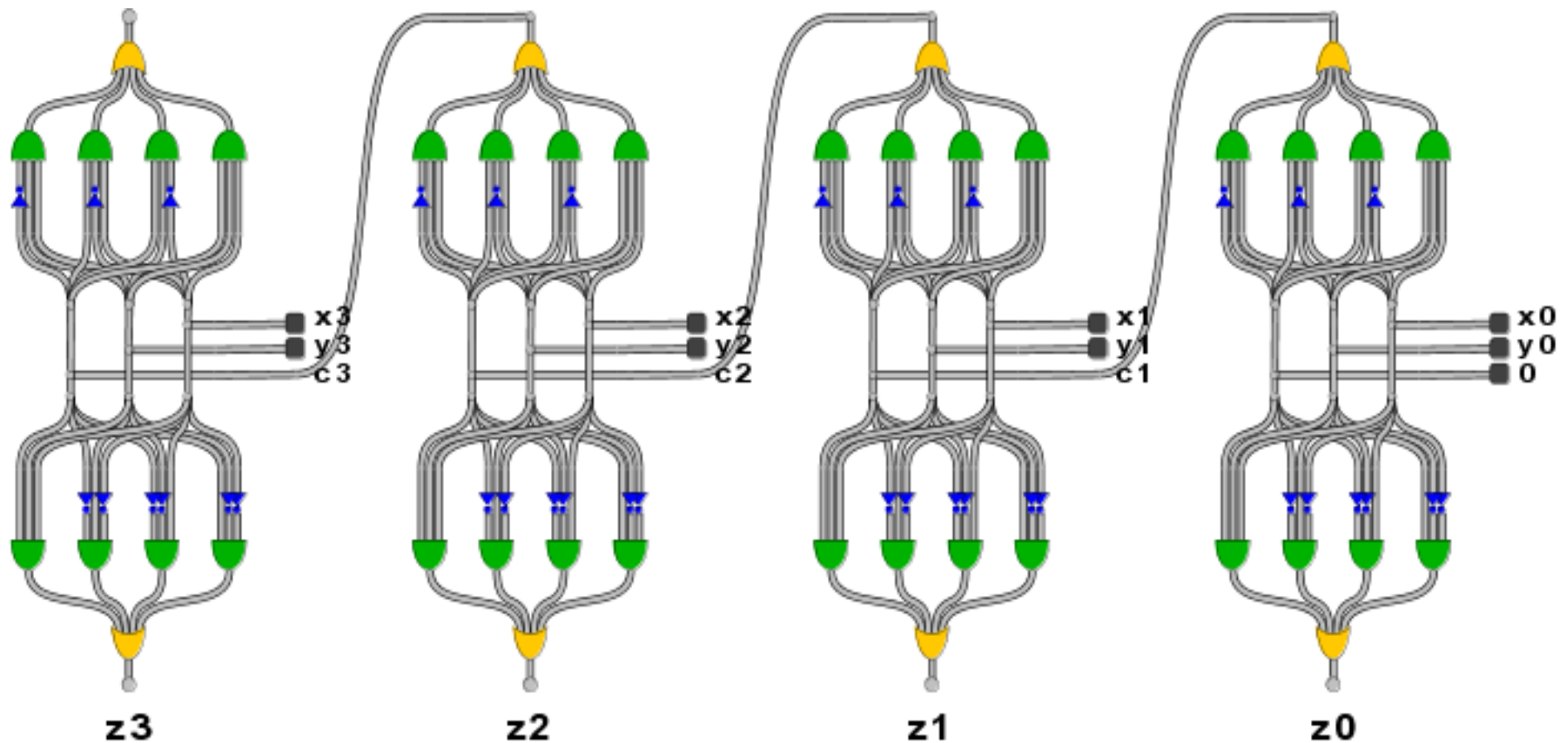


Let's Make an Adder Circuit

Goal: $x + y = z$ for 4-bit integers.

Step 4.

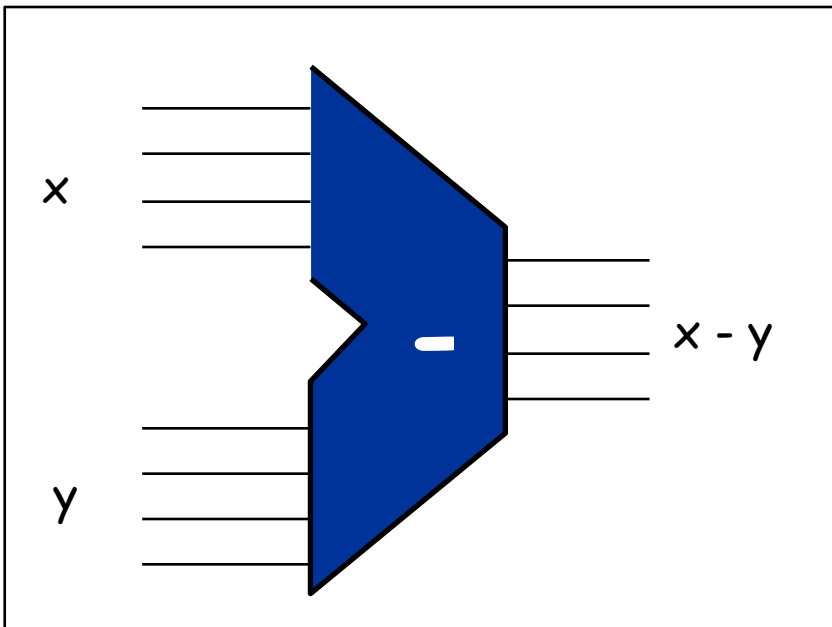
- Transform Boolean expression into circuit.
- Chain together 1-bit adders.



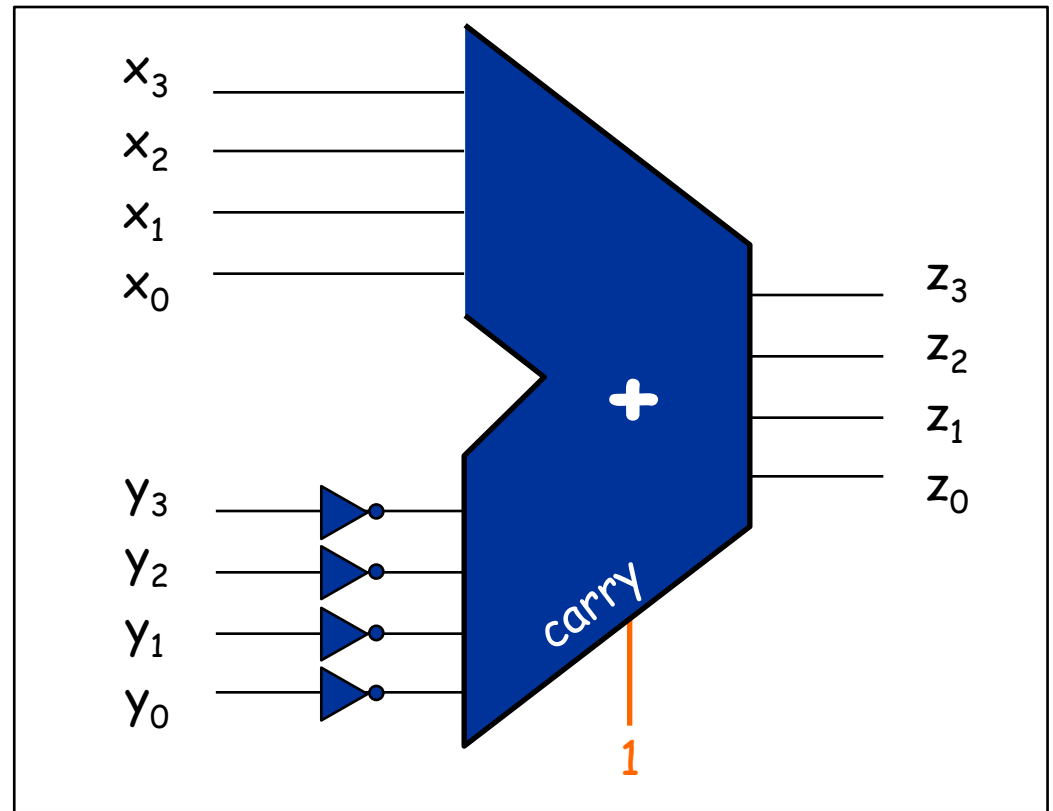
Subtractor

Subtractor circuit: $z = x - y$.

- One approach: new design, like adder circuit.
- Better idea: **reuse** adder circuit.
 - 2's complement: to negate an integer, flip bits, then add 1



4-Bit Subtractor Interface



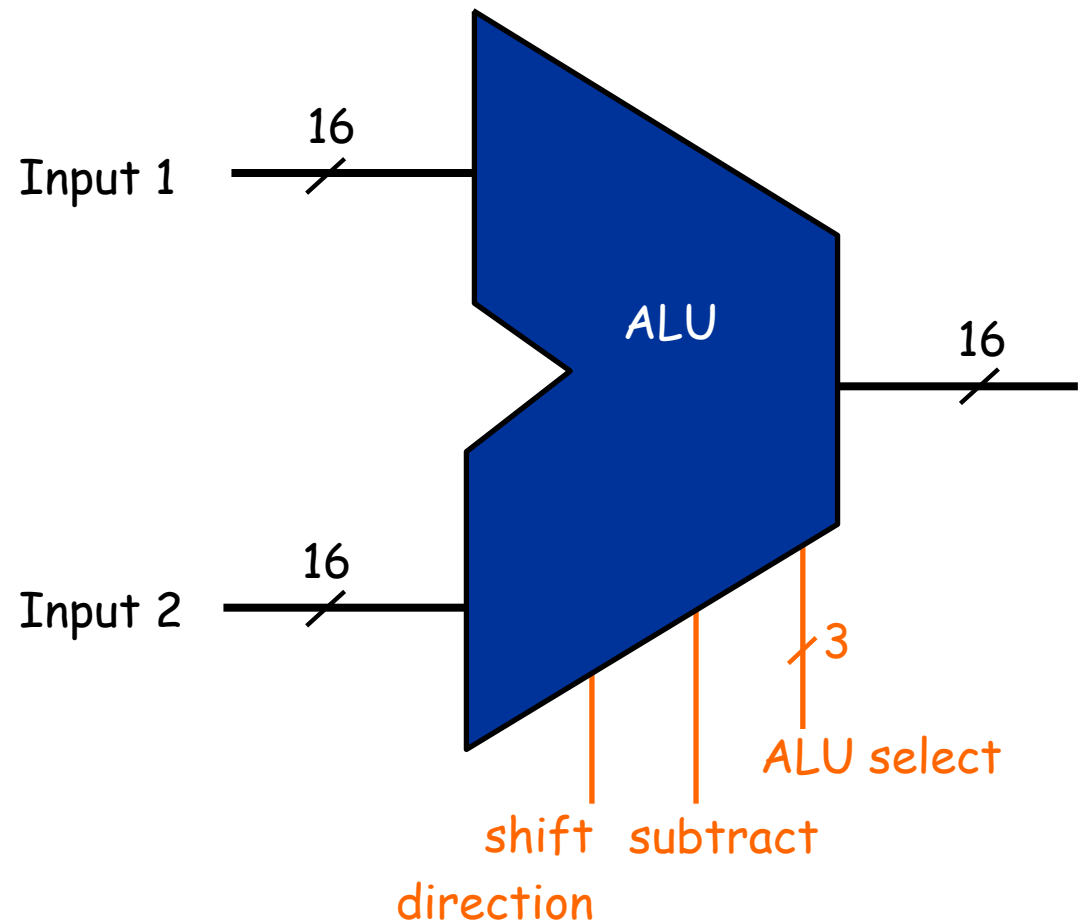
4-Bit Subtractor Implementation

TOY Arithmetic Logic Unit: Interface

ALU Interface.

- Add, subtract, bitwise and, bitwise xor, shift left, shift right, copy.
- Associate 3-bit integer with 5 primary ALU operations.
 - ALU performs operations in parallel
 - control wires select which result ALU outputs

op	2	1	0
+, -	0	0	0
&	0	0	1
^	0	1	0
<<, >>	0	1	1
input 2	1	0	0

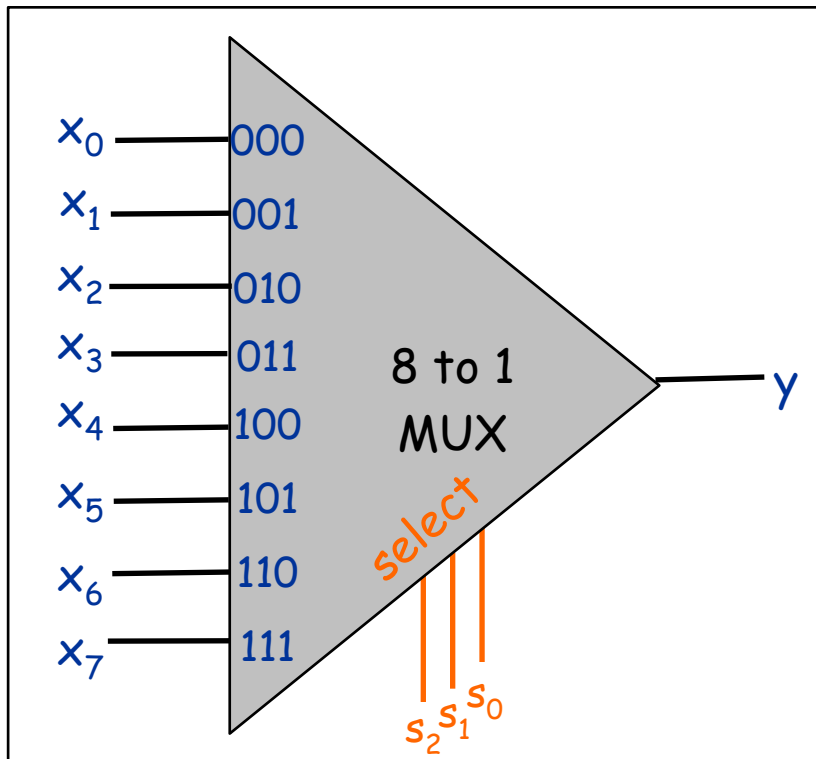


2ⁿ-to-1 Multiplexer

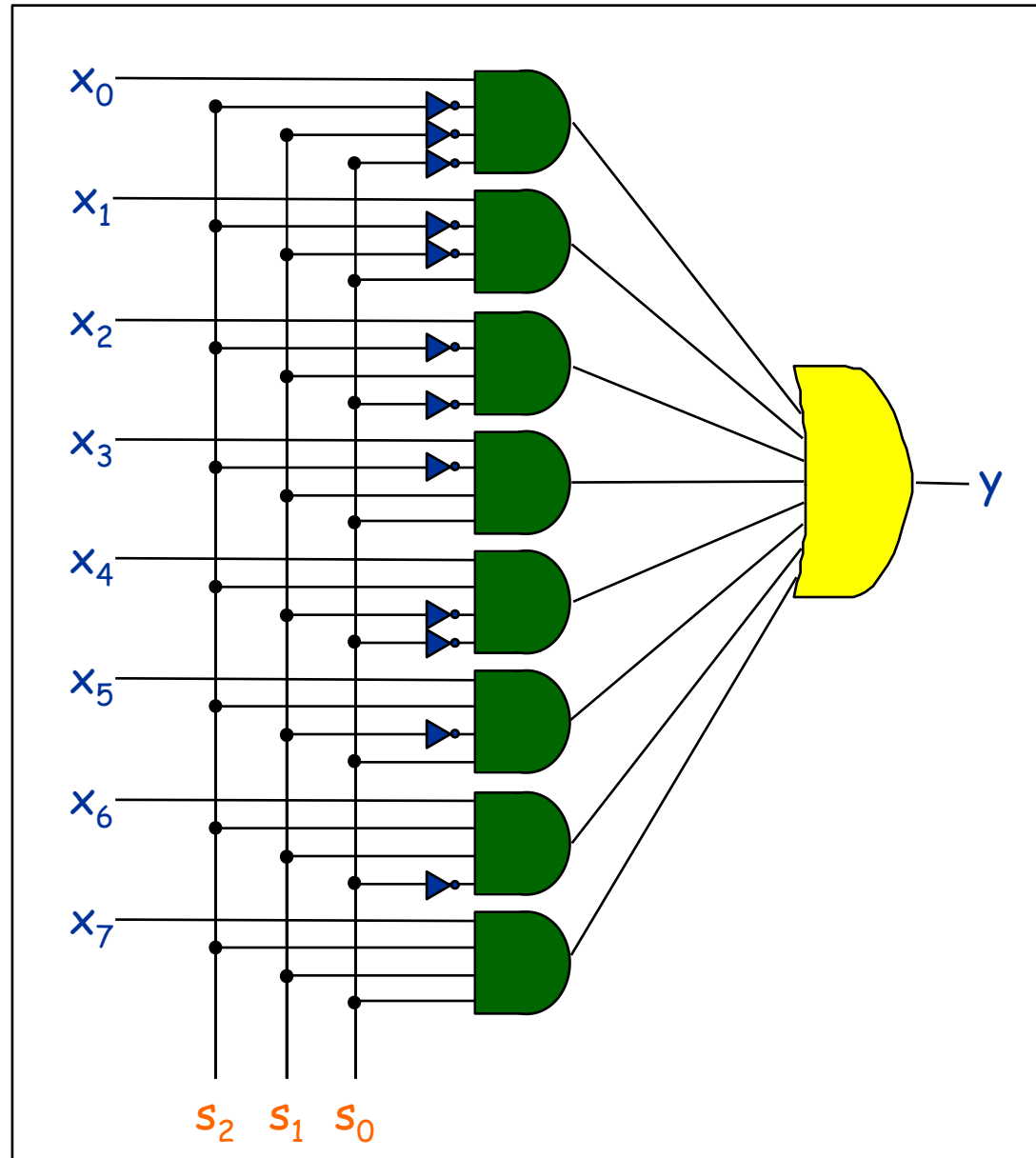
↙ n = 8 for main memory

2ⁿ-to-1 multiplexer.

- n select inputs, 2ⁿ data inputs, 1 output.
- Copies "selected" data input bit to output.

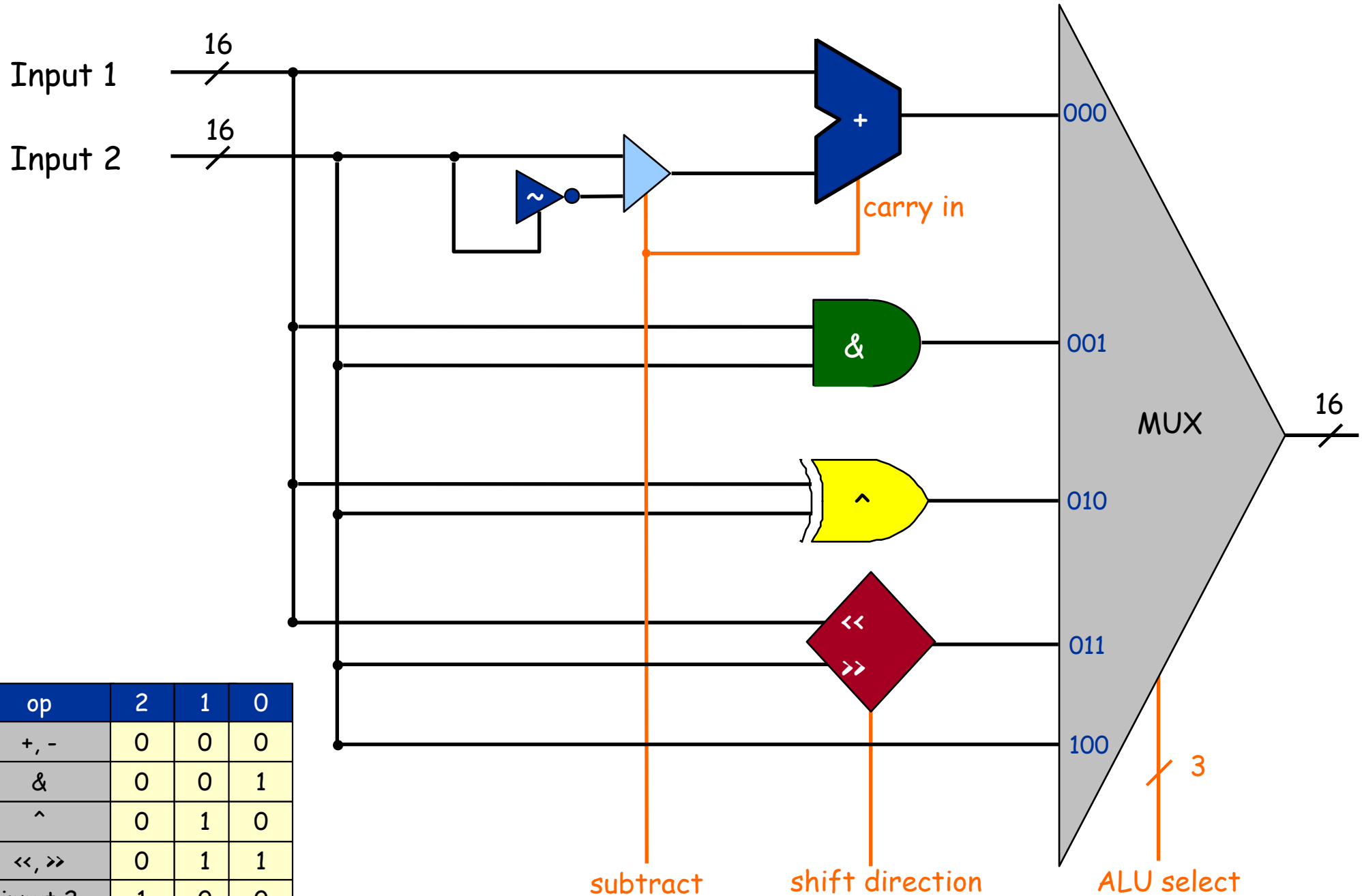


8-to-1 Mux Interface

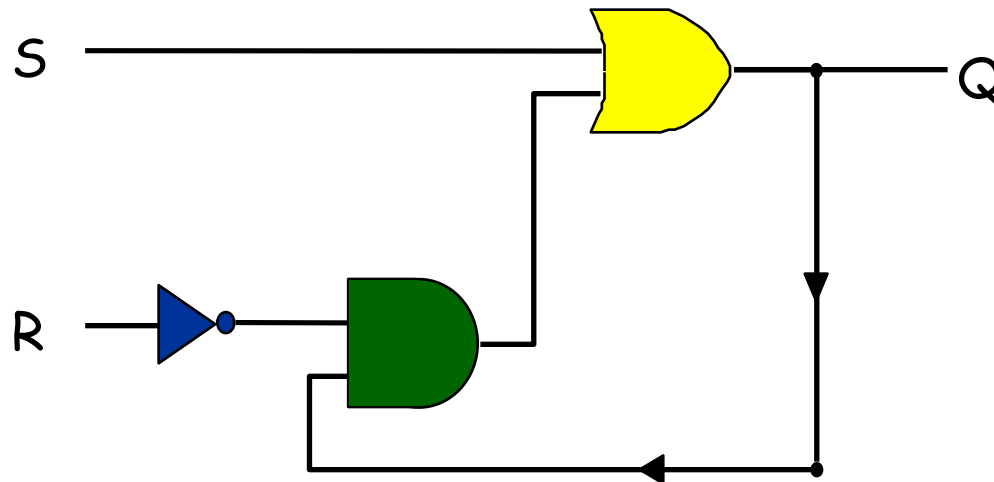


8-to-1 Mux Implementation

TOY Arithmetic Logic Unit: Implementation



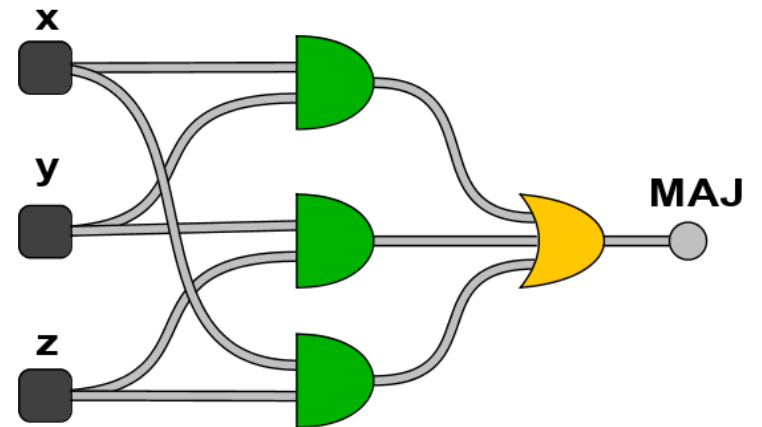
6.2: Sequential Circuits



Sequential vs. Combinational Circuits

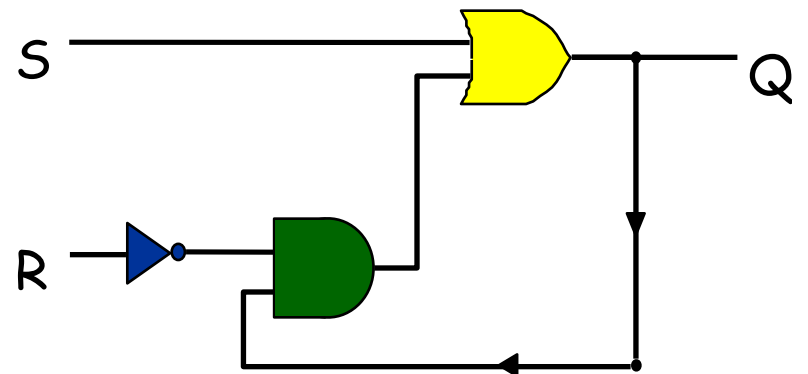
Combinational circuits.

- Output determined solely by inputs.
- Can draw solely with left-to-right signal paths.



Sequential circuits.

- Output determined by inputs AND previous outputs.
- Feedback loop.



Flip-Flop

Flip-flop.

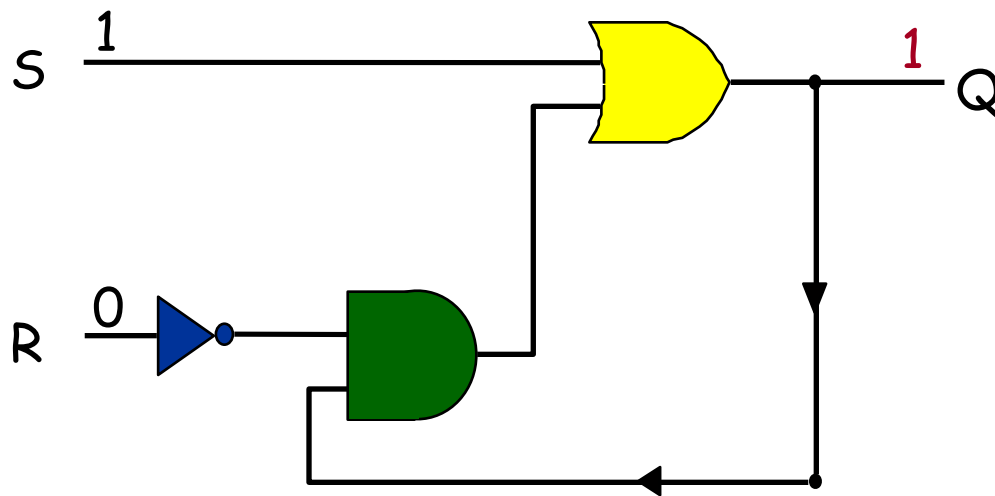
- A small and useful sequential circuit.
- Abstraction that "remembers" one bit.
- Basis of important computer components:
 - memory
 - counter

We will consider several flavors.

SR Flip-Flop

What is the value of Q if:

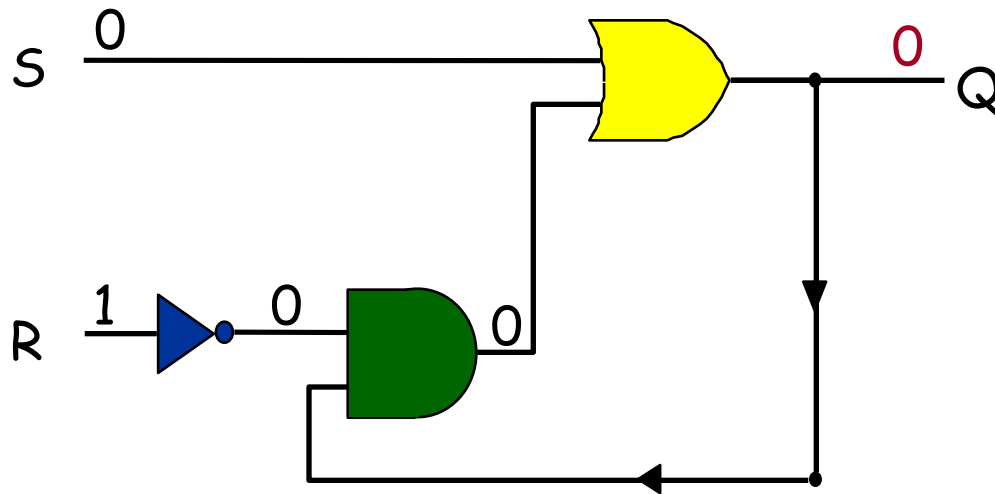
- $S = 1$ and $R = 0$? \Rightarrow **Q is surely 1**



SR Flip-Flop

What is the value of Q if:

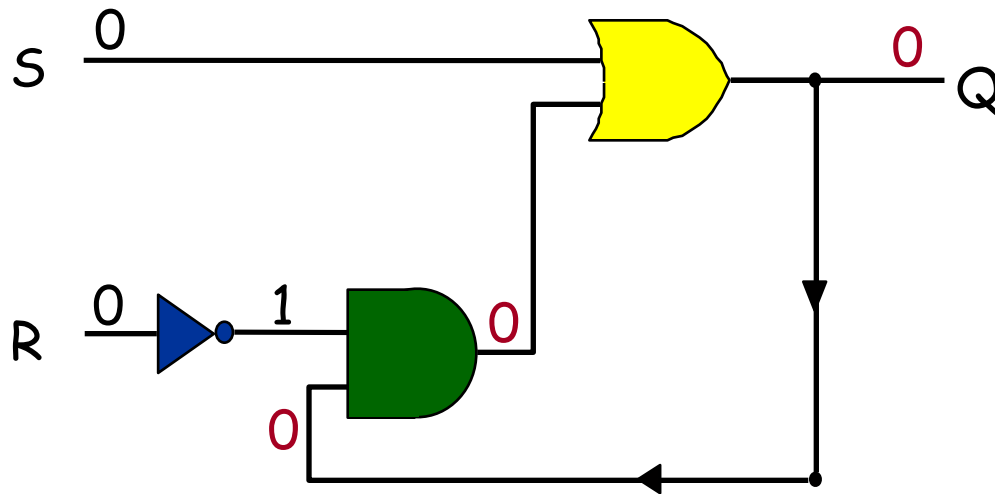
- $S = 1$ and $R = 0$? \Rightarrow Q is surely 1
- $S = 0$ and $R = 1$? \Rightarrow Q is surely 0



SR Flip-Flop

What is the value of Q if:

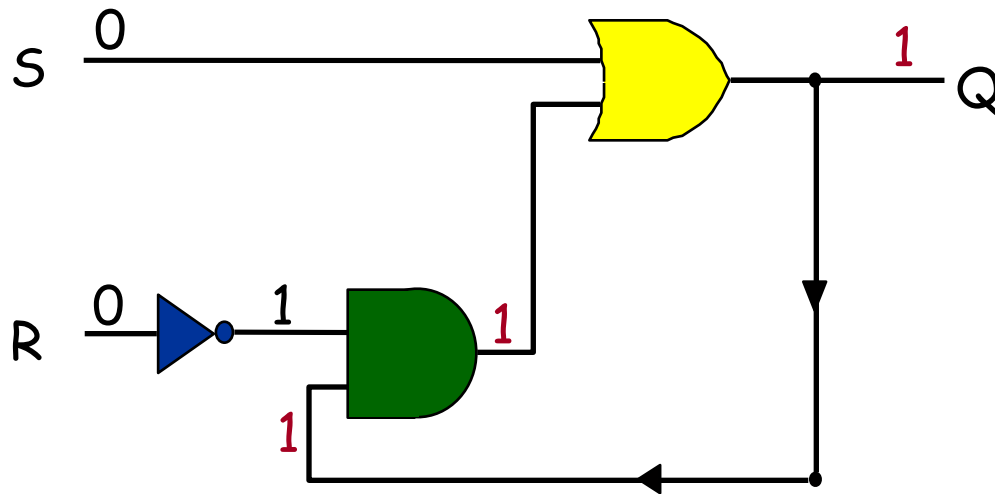
- $S = 1$ and $R = 0$? \Rightarrow Q is surely 1
- $S = 0$ and $R = 1$? \Rightarrow Q is surely 0
- $S = 0$ and $R = 0$? \Rightarrow Q is possibly 0



SR Flip-Flop

What is the value of Q if:

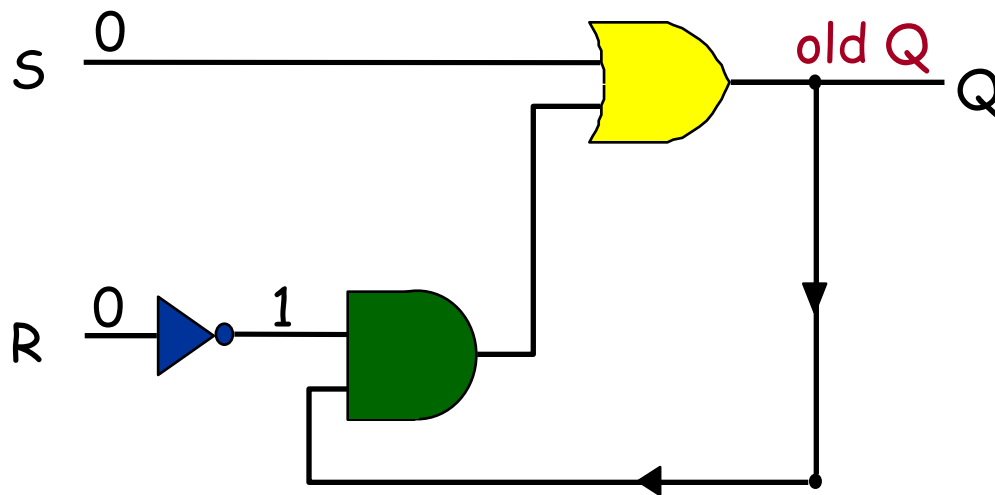
- $S = 1$ and $R = 0$? \Rightarrow Q is surely 1
- $S = 0$ and $R = 1$? \Rightarrow Q is surely 0
- $S = 0$ and $R = 0$? \Rightarrow Q is possibly 0 . . . or possibly 1 !



SR Flip-Flop

What is the value of Q if:

- $S = 1$ and $R = 0$? \Rightarrow Q is surely 1
- $S = 0$ and $R = 1$? \Rightarrow Q is surely 0
- $S = 0$ and $R = 0$? \Rightarrow Q is possibly 0 . . . or possibly 1 !

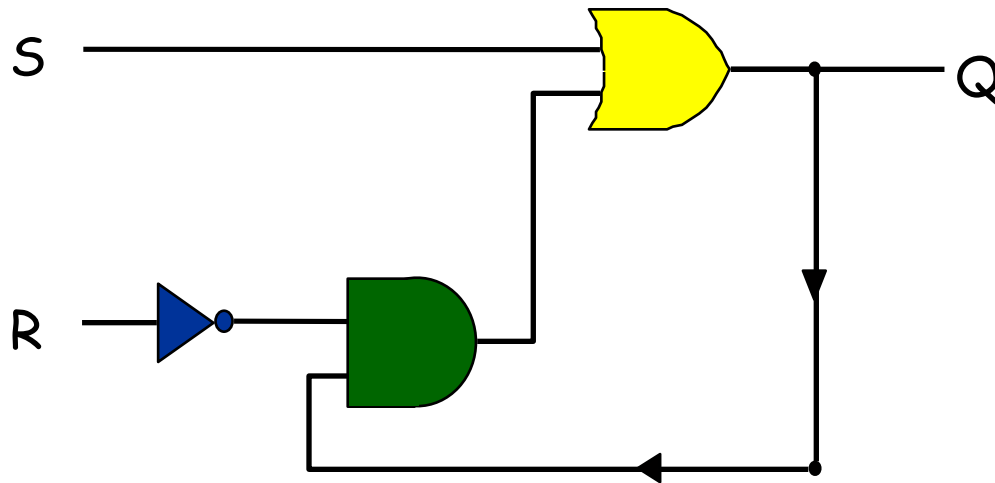


While $S = R = 0$, Q remembers what it was the last time S or R was 1.

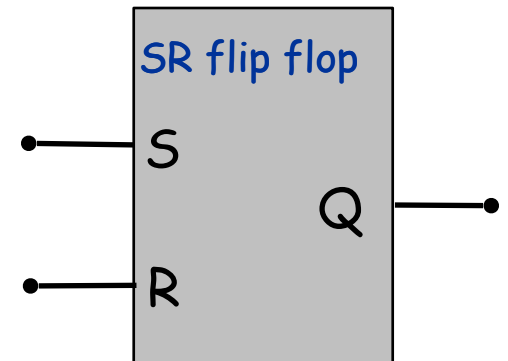
SR Flip-Flop

SR Flip-Flop.

- $S = 1, R = 0$ (**S**et) \Rightarrow "Flips" bit on.
- $S = 0, R = 1$ (**R**eset) \Rightarrow "Flops" bit off.
- $S = R = 0$ \Rightarrow Status quo.
- $S = R = 1$ \Rightarrow Not allowed.



Implementation



Interface

Clock

Clock.

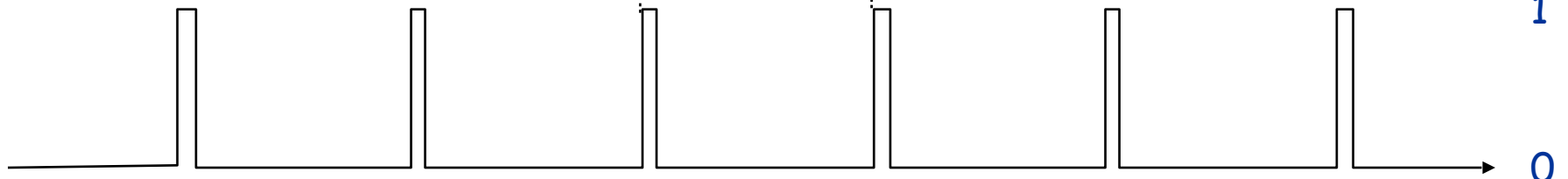
- Fundamental abstraction.
 - regular on-off pulse
- From some oscillating device, possibly external.
- Synchronizes operations of different circuit elements.
- 1 GHz clock means 1 billion pulses per second.



cycle time



Clock



How much does it Hertz?

Frequency is inverse of cycle time.

- Expressed in hertz.
- Frequency of 1 Hz means that there is 1 cycle per second.
- Hence:
 - 1 kilohertz (kHz) means 1000 cycles/sec.
 - 1 megahertz (MHz) means 1 million cycles/sec.
 - 1 gigahertz (GHz) means 1 billion cycles/sec.
 - 1 terahertz (THz) means 1 trillion cycles/sec.

By the way, no such thing as 1 "hert" !

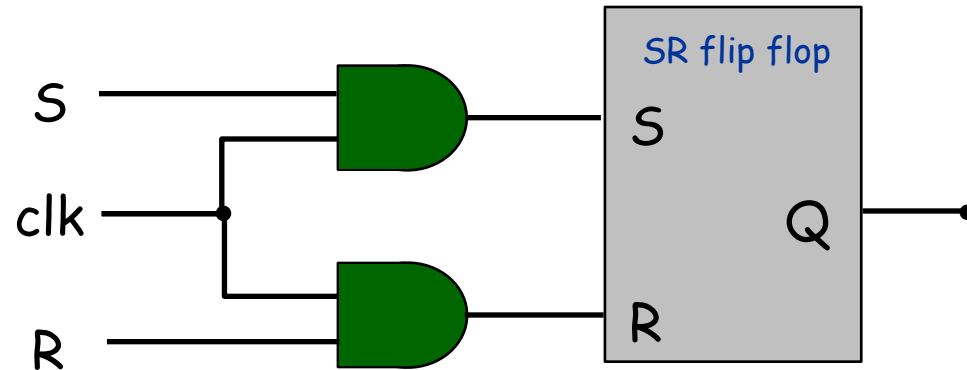


Heinrich Rudolf Hertz
(1857-1894)

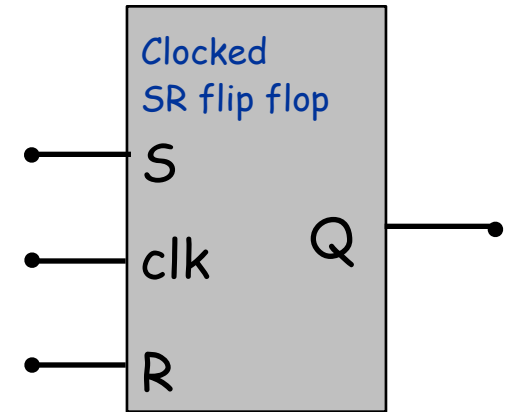
Clocked SR Flip-Flop

Clocked SR Flip-Flop.

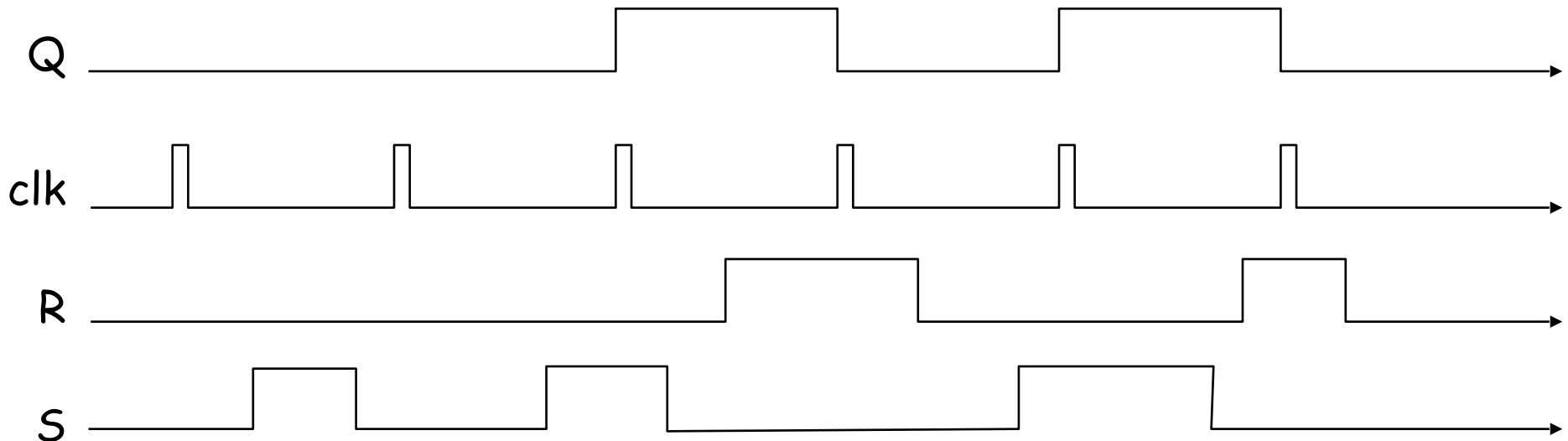
- Same as SR flip-flop except S and R only active when clock is 1.



Implementation



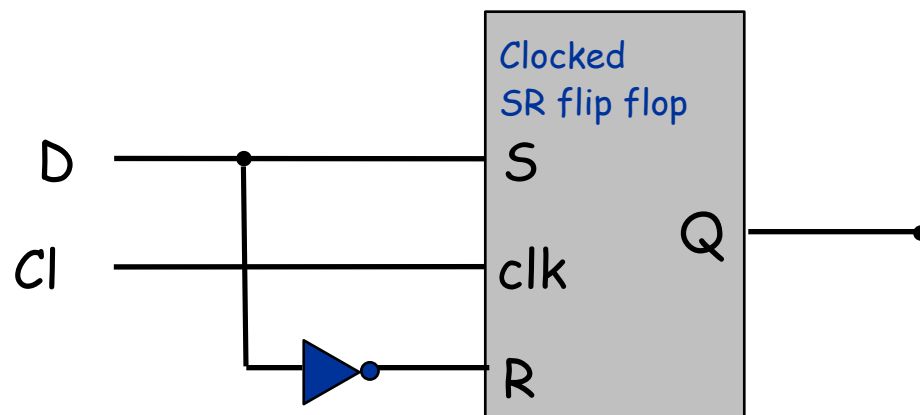
Interface



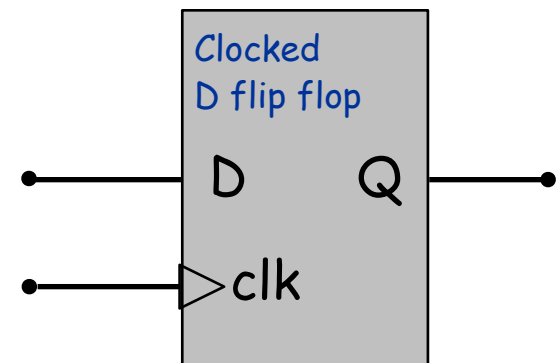
Clocked D Flip-Flop

Clocked D Flip-Flop.

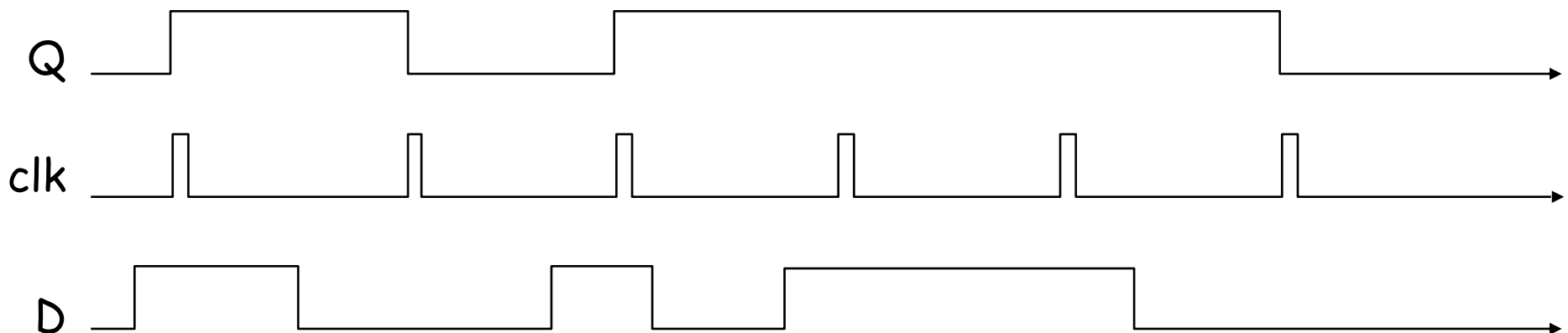
- Output follows D input while clock is 1.
- Output is remembered while clock is 0.



Implementation



Interface



Summary

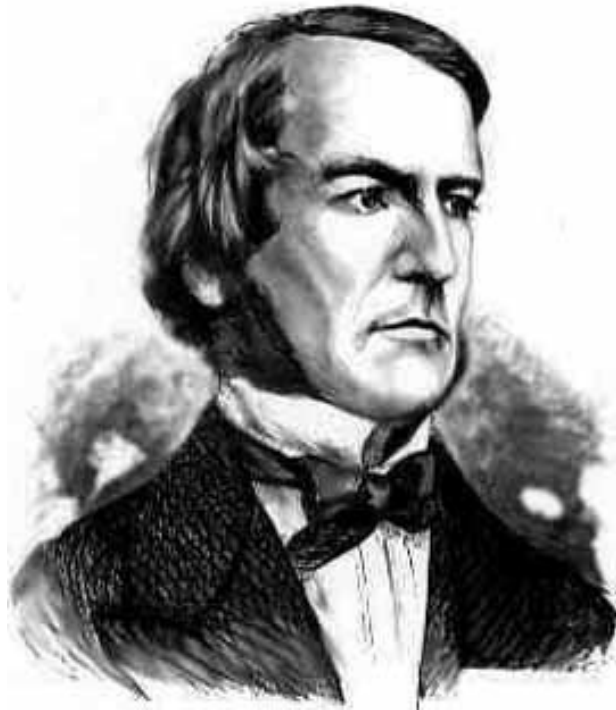
Combinational circuits implement Boolean functions

- Gates and wires Fundamental building blocks.
- Truth tables. Describe Boolean functions.
- Sum-of-products. Systematic method to implement functions.

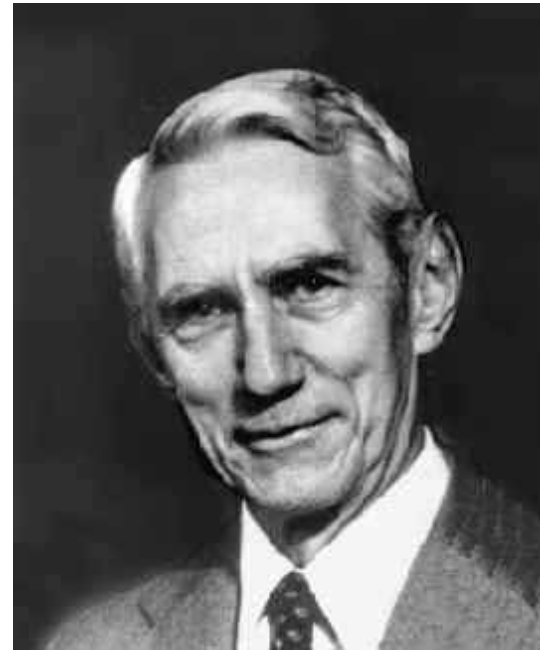
Sequential circuits add "state" to digital hardware.

- Flip-flop. Represents 1 bit.
- TOY register. 16 D flip-flops.
- TOY main memory. 256 registers.

Next time: we build a complete TOY computer (oh yes).



George Boole (1815 - 1864)



Claude Shannon (1916 - 2001)