Using and storing the index

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Review: Model

- Document: sequence of {terms + attributes}
  - equivalently, set of (term, attributes) pairs
    - positions of a term are attributes
- Query: sequence of terms
  - Can make more complicated: Advanced search
- Satisfying: most common now: AND model
  - for Web, terms "contained" in doc. includes:
    - in anchor text of pointers to this doc from other docs
    - in URL
- Ranking: wide open function of document and terms

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Review: Inverted Index

- For each term, keep list of document entries, one for each document in which it appears: a postings list
  - Document entry is list of positions at which term occurs and attributes for each occurrence: a posting
- Keep summary term information
- Keep summary document information

meta-data

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Index structure

term$_1$:(doc ID (position, attributes)
(position, attributes),
…
(position, attributes))
(doc ID (position, attributes)
(position, attributes),
…
(position, attributes))
…
term$_2$:(doc ID (position, attributes)
(position, attributes),
…
(position, attributes))
…

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Consider "advanced search" queries

To know if satisfied need:

Content	Meta-data
• Phrases	•Language
• OR	•Geographic region
• NOT	•File format
• Numeric range	•Date published
• Where in page	•From specific domain
	•Specific licensing rights
	•Filtered by "safe search"

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Retrieval of satisfying documents

- Inverted index will allow retrieval for content queries
- Keep meta-data on docs for meta-data queries
- Issue of efficient retrieval

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Basic retrieval algorithms

- One term
- AND of several terms
- OR of several terms
- NOT term
- proximity

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Basic postings list processing: Merging posting lists

- Have two lists must coordinate
  - Find shared entries and do "something"
  - "something" changes for different operations
    - Set operations UNION? INTERSECTION? DIFFERENCE? …

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Basic retrieval algorithms

- One term:
  - look up posting list in (inverted) index
- AND of several terms:
  - Intersect posting lists of the terms: a list merge
- OR of several terms:
  - Union posting lists of the terms
  - eliminate duplicates: a list merge
- NOT term
  - If *terms* AND NOT(other *terms*), take a difference
  - a list merge (similar to AND)
- Proximity
  - a list merge (similar to AND)

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Algorithms for Merging Postings Lists: unsorted lists

- X Read 2nd list over and over - once for each entry on 1st list
  - computationally expensive
    time $O(|L_1|*|L_2|)$ where |L| length list L
- Build hash table on entry values;
  insert entries of one list, then other;
  look for collisions
  - must have good hash table
  - unwanted collisions expensive
- Sort lists; use algorithm for sorted lists
  - often lists on disk: external sort
  - can sort in O(|L| log |L|) operations

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Algorithms for Merging Postings Lists: sorted lists

- Lists sorted by some key
  - same key both lists
- Read both lists in "parallel"
  - Classic list merge:
    (sorted list_1 , sorted list_2) ⇒ sorted set union
  - General merge: if no duplicates, get time $|L_1|+|L_2|$
- Build lists so sorted
  - pay cost at most once
  - maybe get sorted order "naturally"
- If only one list sorted, can do binary search of sorted list for entries of other list
  - Must be able to binary search! - rare!
    - can't binary search disk

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Keys for documents

For posting lists, entries are documents
What value is used to sort?

- Unique document IDs
  - can still be duplicate documents
  - consider for Web when consider crawling
- document scoring function that is independent of query
  - PageRank, HITS authority
  - sort on document IDs as secondary key
  - allows for approximate "highest k" retrieval
    - approx. k highest ranking doc.s for a query

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Keys within document list

Processing within document posting

- Proximity of terms
  - merge lists of terms occurrences within 1 doc.
- Sort on term position

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Computing document score

- "On fly"- as find each satisfying document
- Separate phase after build list of satisfying documents
- For either, must sort doc.s by score

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Web query processing: limiting size

- For Web-scale collections, may not process complete posting list for each term in query
  - at least not initially
- Need docs sorted first on global (static) quantity
  - why not by term frequency for doc?
- Only take first k doc.s on each term list
  - k depends on query - how?
  - k depends on how many want to be able to return
    - Google: 1000 max returns
  - Flaws w/ partial retrieval from each list?
- Other limits? query size
  - Google: 32 words max query size

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Limiting size with term-based sorting

- Can sort doc.s on postings list by score of term
  - term frequency + …
- Lose linear merge - salvage any?
- Tiered index:
  - tier 1: docs with highest term-based scores, sorted by ID or global quantity
  - tier 2: docs in next bracket of score quality, sorted
  - etc.
  - need to decide size or range of brackets
- If give up AND of query terms, can use idf too
  - only consider terms with high idf = rarer terms

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Data structure for inverted index?

How access individual terms and each associated postings list?

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Data structure for inverted index?

- Sorted array:
  - binary search IF can keep in memory
  - High overhead for additions
- Hashing
  - Fast look-up
  - Collisions
- Search trees: B+-trees
  - Maintain balance - always log look-up time
  - Can insert and delete

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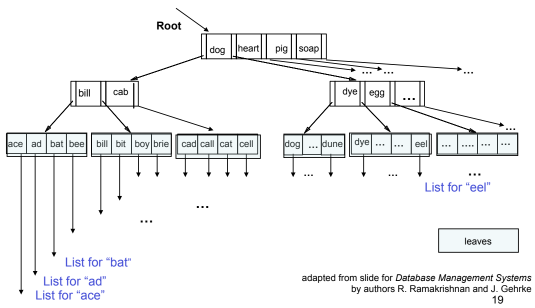
Example B+ Tree

order = 2: 2 to 4 search keys per interior node



adapted from slide for *Database Management Systems* by authors R. Ramakrishnan and J. Gehrke

B+- trees

- All index entries are at leaves
- Order *m* B+ tree has *m* to *2m* children for each interior node
- Look up: follow root to leaf by keys in interior nodes
- Insert:
  - find leaf in which belongs
  - If leaf full, split
  - Split can propagate up tree
- Delete:
  - Merge or redistribute from too-empty leaf
  - Merge can propagate up tree

Disk-based B+ trees for large data sets

- Each leaf is file page (block) on disk
- Each interior node is file page on disk
- Keep top of tree in buffer (RAM)
- Typical sizes:
  - m ~ 200;
  - average fanout ~ 267
    - Height 4 gives ~ 5 billion entries

prefix key B+ trees

- Save space
- Each interior node key is shortest prefix of word needed to distinguish which child pointer to follow
  - Allows more keys per interior node
  - higher fanout
    - fanout determined by what can fit
    - keep at least 1/2 full