



Exceptions and Processes

The material for this lecture is drawn from
Computer Systems: A Programmer's Perspective (Bryant & O'Hallaron) Chapter 8

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Goals of this Lecture

- Help you learn about:
 - **Exceptions**
 - The **process** concept ... and thereby...
 - How operating systems work
 - How application programs interact with operating systems and hardware

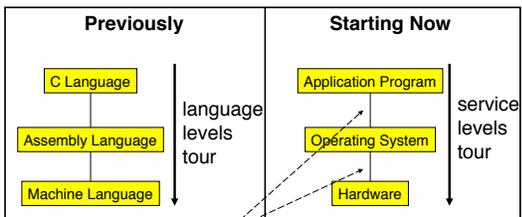
The **process** concept is one of the most important concepts in systems programming

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Context of this Lecture

Second half of the course



Application programs, OS, and hardware interact via **exceptions**

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Motivation



Question:

- Executing program thinks it has exclusive control of the CPU
- But multiple executing programs must share one CPU (or a few CPUs)
- How is that illusion implemented?

Question:

- Executing program thinks it has exclusive use of all of memory
- But multiple executing programs must share one memory
- How is that illusion implemented?

Answers: Exceptions...

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Exceptions



• Exception

- An abrupt change in control flow in response to a change in processor state

• Examples:

- Application program:
 - Requests I/O
 - Requests more heap memory
 - Attempts integer division by 0
 - Attempts to access privileged memory
 - Accesses variable that is not in real memory (see upcoming "Virtual Memory" lecture)
- User presses key on keyboard
- Disk controller finishes reading data

Synchronous

Asynchronous

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Exceptions Note



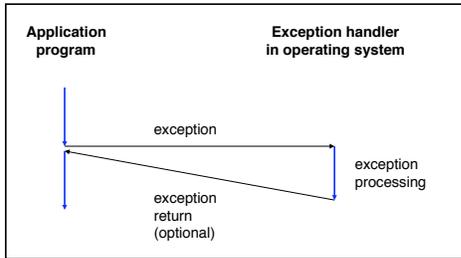
• Note:

Exceptions in OS \neq exceptions in Java

Implemented using
try/catch
and **throw** statements

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Exceptional Control Flow



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Exceptions vs. Function Calls



- Exceptions are **similar to function calls**
 - Control transfers from original code to other code
 - Other code executes
 - Control returns to original code
- Exceptions are **different from function calls**
 - Processor pushes **additional state** onto stack
 - E.g. values of all registers
 - Processor pushes data onto **OS's stack**, not application pgm's stack
 - Handler runs in **privileged mode, not in user mode**
 - Handler can execute all instructions and access all memory
 - Control **might return to next instruction**
 - Control sometimes returns to **current** instruction
 - Control sometimes does not return at all!

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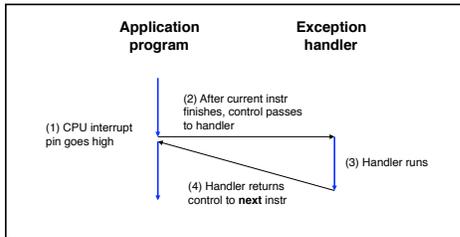
Classes of Exceptions



- There are 4 classes of exceptions...

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(1) Interrupts



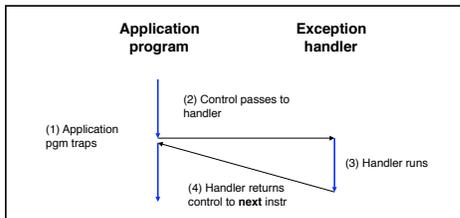
Cause: Signal from I/O device

Examples:

- User presses key
- Disk controller finishes reading/writing data

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(2) Traps



Cause: Intentional (application pgm requests OS service)

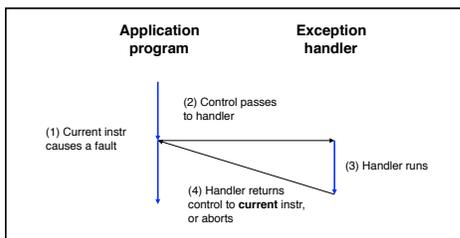
Examples:

- Application pgm requests more heap memory
- Application pgm requests I/O

Traps provide a function-call-like interface between application pgm and OS

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(3) Faults



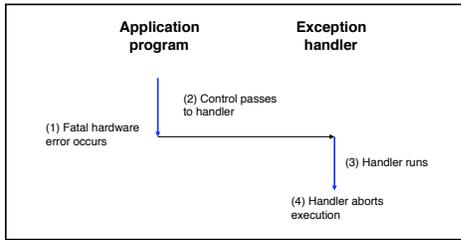
Cause: Application pgm causes (possibly) recoverable error

Examples:

- Application pgm accesses privileged memory (seg fault)
- Application pgm accesses data that is not in real memory (page fault)

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(4) Aborts



Cause: Non-recoverable error

Example:

Parity check indicates corruption of memory bit (overheating, cosmic ray!, etc.)

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Summary of Exception Classes



Class	Cause	Asynch/Synch	Return Behavior
Interrupt	Signal from I/O device	Asynch	Return to next instr
Trap	Intentional	Sync	Return to next instr
Fault	(Maybe) recoverable error	Sync	(Maybe) return to current instr
Abort	Non-recoverable error	Sync	Do not return

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Exceptions in Intel Processors



Each exception has a number
Some exceptions in Intel processors:

Exception #	Exception
0	Fault: Divide error
13	Fault: Segmentation fault
14	Fault: Page fault (see "Virtual Memory" lecture)
18	Abort: Machine check
32-127	Interrupt or trap (OS-defined)
128	Trap
129-255	Interrupt or trap (OS-defined)

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Traps in Intel Processors



- To execute a trap, application program should:
 - Place number in EAX register indicating desired functionality
 - Place parameters in EBX, ECX, EDX registers
 - Execute assembly language instruction "int 128"
- Example: To request more heap memory...

```
movl $45, %eax
movl $1024, %ebx
int $128
```

In Linux, 45 indicates request for more heap memory

Causes trap

Request is for 1024 bytes

System-Level Functions



- For convenience, traps are wrapped in **system-level functions**
- Example: To request more heap memory...

```
/* unistd.h */
void *sbrk(intptr_t increment);
...
```

sbrk () is a system-level function

```
/* unistd.s */
Defines sbrk() in assembly lang
Executes int instruction
...
```

```
/* client.c */
...
sbrk(1024);
...
```

A call of a system-level function, that is, a **system call**

See Appendix for list of some Linux system-level functions¹⁷

Processes



- **Program**
 - Executable code
- **Process**
 - An instance of a program in execution
- Each program runs in the **context** of some process
- **Context** consists of:
 - Process ID
 - Address space
 - TEXT, RODATA, DATA, BSS, HEAP, and STACK
 - Processor state
 - EIP, EFLAGS, EAX, EBX, etc. registers
 - Etc.

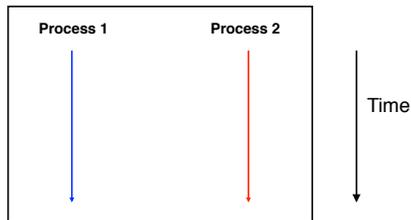
Significance of Processes



- **Process** is a profound abstraction in computer science
- The process abstraction provides application pgms with two key illusions:
 - Private control flow
 - Private address space

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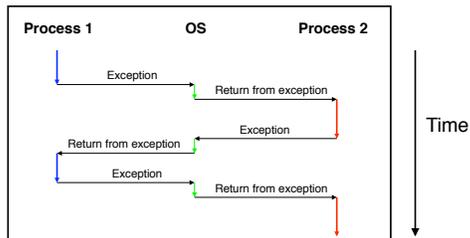
Private Control Flow: Illusion



Hardware and OS give each application process the illusion that it is the only process running on the CPU

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Private Control Flow: Reality



All application processes -- and the OS process -- share the same CPU(s)

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Context Switches



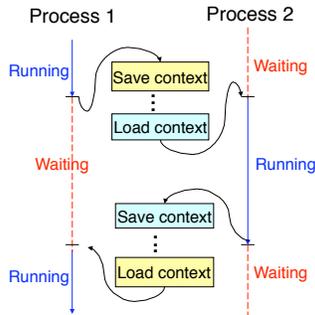
- **Context switch**
 - The activity whereby the OS assigns the CPU to a different process
 - Occurs during exception handling, at discretion of OS
- Exceptions can be caused:
 - Synchronously, by application pgm (trap, fault, abort)
 - Asynchronously, by external event (interrupt)
 - **Asynchronously, by hardware timer**
 - So no process can dominate the CPUs
- Exceptions are the mechanism that enables the illusion of private control flow

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Context Switch Details



- **Context**
 - State the OS needs to restart a preempted process
- **Context switch**
 - Save the context of current process
 - Restore the saved context of some previously preempted process
 - Pass control to this newly restored process

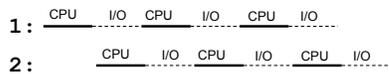


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When Should OS Do Context Switch?



- **When a process is stalled waiting for I/O**
 - Better utilize the CPU, e.g., while waiting for disk access



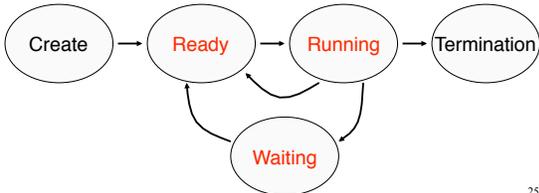
- **When a process has been running for a while**
 - Sharing on a fine time scale to give each process the illusion of running on its own machine
 - Trade-off efficiency for a finer granularity of fairness

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Life Cycle of a Process



- **Running**: instructions are being executed
- **Waiting**: waiting for some event (e.g., I/O finish)
- **Ready**: ready to be assigned to a processor



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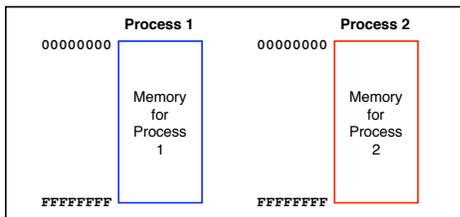
Context Details



- What does the OS need to save/restore during a context switch?
 - Process state
 - New, ready, waiting, terminated
 - CPU registers
 - EIP, EFLAGS, EAX, EBX, ...
 - I/O status information
 - Open files, I/O requests, ...
 - Memory management information
 - Page tables (see "Virtual Memory" lecture)
 - Accounting information
 - Time limits, group ID, ...
 - CPU scheduling information
 - Priority, queues

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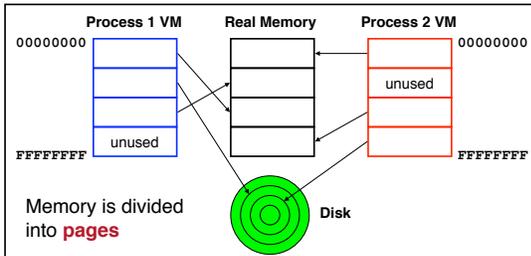
Private Address Space: Illusion



Hardware and OS give each application process the illusion that it is the only process using memory

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Private Address Space: Reality



All processes use the same real memory
Hardware and OS provide application pgms with a **virtual** view of memory, i.e. **virtual memory (VM)**

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Private Address Space Details



- Exceptions (specifically, page faults) are the mechanism that enables the illusion of private address spaces
- See the **Virtual Memory** lecture for details

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Summary



- **Exception**: an abrupt change in control flow
 - **Interrupts**: asynchronous; e.g. I/O completion, hardware timer
 - **Traps**: synchronous; e.g. app pgm requests more heap memory, I/O
 - **Faults**: synchronous; e.g. seg fault
 - **Aborts**: synchronous; e.g. parity error
- **Process**: An instance of a program in execution
 - Hardware and OS use exceptions to give each process the illusion of:
 - Private control flow (reality: **context switches**)
 - Private address space (reality: **virtual memory**)

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Appendix: System-Level Functions

Linux system-level functions for **I/O management**

Number	Function	Description
3	<code>read()</code>	Read data from file descriptor Called by <code>getchar()</code> , <code>scanf()</code> , etc.
4	<code>write()</code>	Write data to file descriptor Called by <code>putchar()</code> , <code>printf()</code> , etc.
5	<code>open()</code>	Open file or device Called by <code>fopen()</code>
6	<code>close()</code>	Close file descriptor Called by <code>fclose()</code>
8	<code>creat()</code>	Open file or device for writing Called by <code>fopen(..., "w")</code>

Described in **I/O Management** lecture

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Appendix: System-Level Functions

Linux system-level functions for **process management**

Number	Function	Description
1	<code>exit()</code>	Terminate the process
2	<code>fork()</code>	Create a child process
7	<code>waitpid()</code>	Wait for process termination
7	<code>wait()</code>	(Variant of previous)
11	<code>exec()</code>	Execute a program in current process
20	<code>getpid()</code>	Get process id

Described in **Process Management** lecture

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Appendix: System-Level Functions

Linux system-level functions for **I/O redirection and inter-process communication**

Number	Function	Description
41	<code>dup()</code>	Duplicate an open file descriptor
42	<code>pipe()</code>	Create a channel of communication between processes
63	<code>dup2()</code>	Close an open file descriptor, and duplicate an open file descriptor

Described in **Process Management** lecture

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Appendix: System-Level Functions

Linux system-level functions for **dynamic memory management**

Number	Function	Description
45	<code>brk()</code>	Move the program break, thus changing the amount of memory allocated to the HEAP
45	<code>sbrk()</code>	(Variant of previous)
90	<code>mmap()</code>	Map a virtual memory page
91	<code>munmap()</code>	Unmap a virtual memory page

Described in **Dynamic Memory Management** lectures

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Appendix: System-Level Functions

Linux system-level functions for **signal handling**

Number	Function	Description
27	<code>alarm()</code>	Deliver a signal to a process after a specified amount of wall-clock time
37	<code>kill()</code>	Send signal to a process
67	<code>sigaction()</code>	Install a signal handler
104	<code>setitimer()</code>	Deliver a signal to a process after a specified amount of CPU time
126	<code>sigprocmask()</code>	Block/unblock signals

Described in **Signals** lecture

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