

## Midterm Solutions

1. **8 sorting algorithms.**

0 6 5 2 4 9 3 8 7 1

2. **Sorting equal keys.**

Insertion	$A_0$	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$	$A_6$
Selection	$A_0$	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$	$A_6$
Shellsort	$A_0$	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$	$A_6$
Mergesort	$A_0$	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$	$A_6$
Quicksort	$A_4$	$A_3$	$A_5$	$A_6$	$A_0$	$A_2$	$A_1$
Heapsort	$A_1$	$A_2$	$A_3$	$A_4$	$A_5$	$A_6$	$A_0$

3. **Analysis of algorithms.**

(a) I and II only

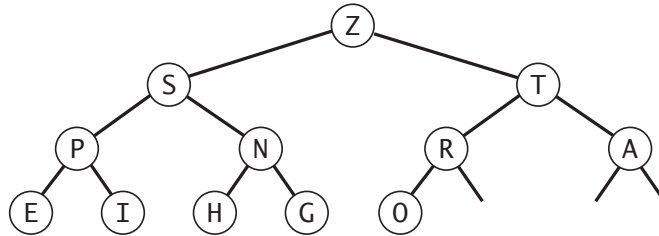
*Big-Oh notation and tilde notation both suppress lower order terms.*

(b) I only

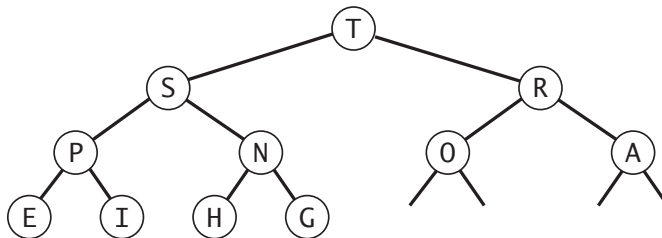
*Amortized analysis provides a worst-case guarantee on any sequence of operations starting from an empty data structure.*

4. Binary heaps.

(a)



(b)

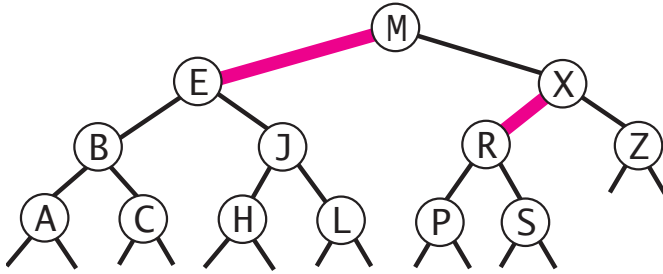


(c) True.

5. Ordered-array implementation of a set.

<code>add(key)</code>	<i>add the key to the set</i>	$N$
<code>contains(key)</code>	<i>is the key in the set?</i>	$\log N$
<code>ceiling(key)</code>	<i>smallest key in set <math>\geq</math> given key</i>	$\log N$
<code>rank(key)</code>	<i>number of keys in set <math>&lt;</math> given key</i>	$\log N$
<code>select(i)</code>	<i><math>i</math>th largest key in the set</i>	1
<code>min()</code>	<i>minimum key in the set</i>	1
<code>delete(key)</code>	<i>delete the given key from the set</i>	$N$
<code>iterator()</code>	<i>iterate over all <math>N</math> keys in the set in order</i>	$N$

6. Red-black trees.



7. Line intersection.

(a) There are two cases:

- If the two lines have the same slope ( $a_0 = a_1$ ), then return no intersection.
- Otherwise, the point  $(x, y)$  of intersection is given by:

$$x = -\frac{b_1 - b_0}{a_1 - a_0}, \quad y = a_0x + b_0$$

(b) To determine whether the  $i$ th line is involved in an intersection with 3 (or more) lines:

- Create a symbol table with key = point, value = list (say, a queue) of lines.
- For each line  $j \neq i$  in order:
  - Compute the intersection point  $p$  between line  $i$  and line  $j$ .
  - If they don't intersect, continue.
  - If the key  $p$  is not already in the symbol table, add an entry to the symbol table with key =  $p$  and value = empty list.
  - Add line  $j$  to the end of the list associated with  $p$ .
- For each key in the symbol table, if it's list contains 2 (or more) lines, they correspond to 3 (or more) lines intersecting at a single point (line  $i$ , plus the lines in the list).

Implement the symbol table using a separate-chaining (or linear-probing) hash table so that each insert/search takes  $O(1)$  time. Thus, the overall subroutine takes  $O(N)$  time.

To determine whether any 3 (or more lines) intersect at a point, run the previous subroutine  $N$  times, once for each line  $i$ . The total running time is  $O(N^2)$ .

(c) Only print out a set of lines in the last step of the subroutine if the index of the first line in the list is greater than  $i$ . This guarantees we only find a set of lines once, when using the line with the smallest index as the base line.