

What is the computational cost of automating brilliance or serendipity?

(Computational complexity and P vs NP
question)

COS 116: 4/15/2008

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Combination lock

Why is it secure?
(Assume it cannot be picked)



Ans: Combination has 3 numbers 0-39...
thief must try $39^3 = 59,319$ combinations

Exponential running time

2^n time to solve instances of “size” n

Increase n by 1 \rightarrow running time doubles!

Main fact to remember:

For $n = 300$,

$2^n >$ number of atoms in the visible universe.

Boolean satisfiability

$$(A + B + C) \cdot (\bar{D} + F + G) \cdot (\bar{A} + G + K) \cdot (\bar{B} + P + Z) \cdot (C + \bar{U} + \bar{X})$$

- Does it have a satisfying assignment?
- What if instead we had 100 variables?
- 1000 variables?
- How long will it take to determine the assignment?

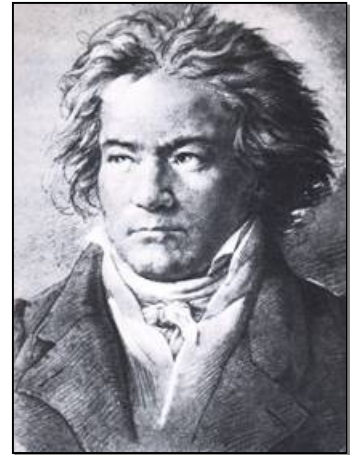
Discussion

Is there an inherent difference between

being creative / brilliant

and

being able to appreciate creativity / brilliance?



What is a computational analogue of this phenomenon?


A Proposal

Brilliance = Ability to find
“needle in a haystack”

Beethoven found
“satisfying assignments”
to our neural circuits
for music appreciation

Comments??

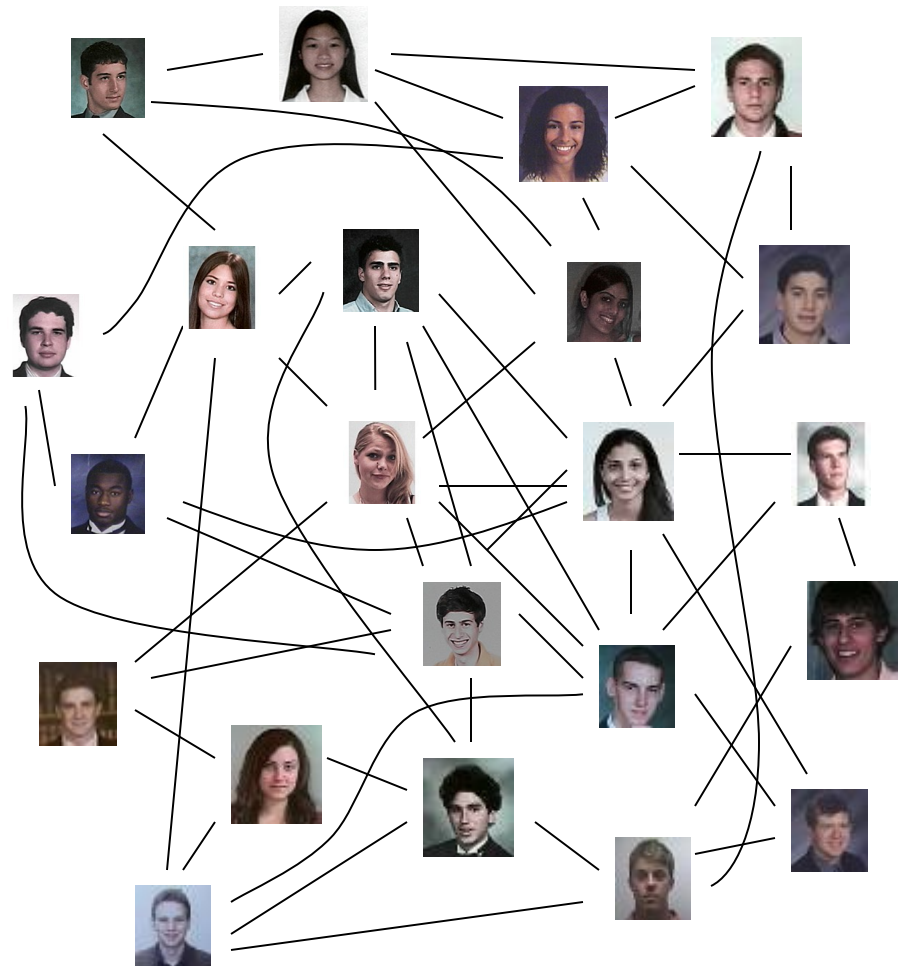




There are many computational problems where finding a solution involves “finding a needle in a haystack”....

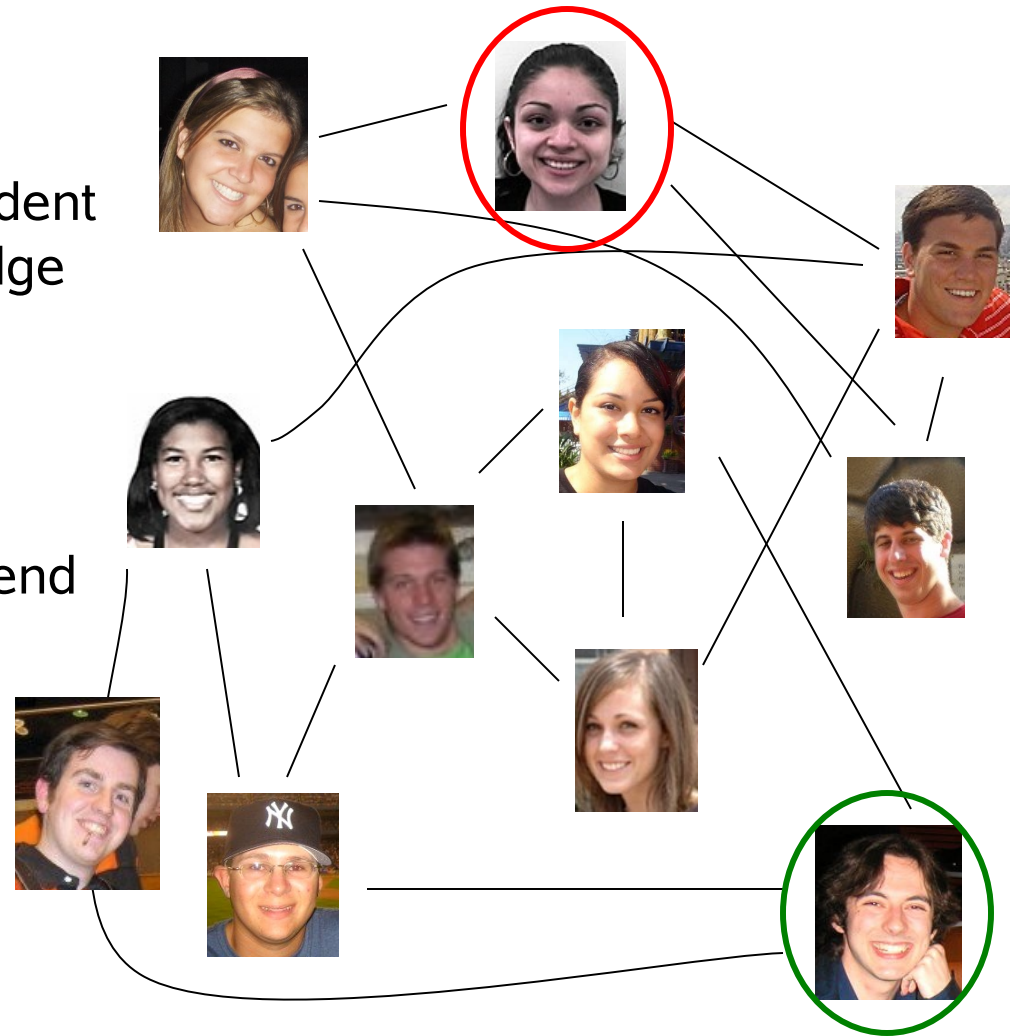
CLIQUE Problem

- In this social network, is there a CLIQUE with 5 or more students?
- CLIQUE: Group of students, every pair of whom are friends
- What is a good algorithm for detecting cliques?
- How does efficiency depend on network size and desired clique size?



Rumor mill problem

- Social network for COS 116
- Each node represents a student
- Two nodes connected by edge if those students are friends
- Johanna starts a rumor
- Will it reach Kieran?
- Suggest an algorithm
- How does running time depend on network size?
- Internet servers solve this problem all the time (“traceroute” in Lab 9).



Exhaustive Search / Combinatorial Explosion

Naïve algorithms for many “needle in a haystack” tasks involve checking all possible answers → exponential running time.

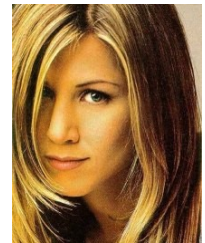
- Ubiquitous in the computational universe
- Can we design smarter algorithms (as for “Rumor Mill”)? Say, n^2 running time.

Harmonious Dorm Floor

Given: Social network involving n students.

Edges correspond to pairs of students who don't get along.

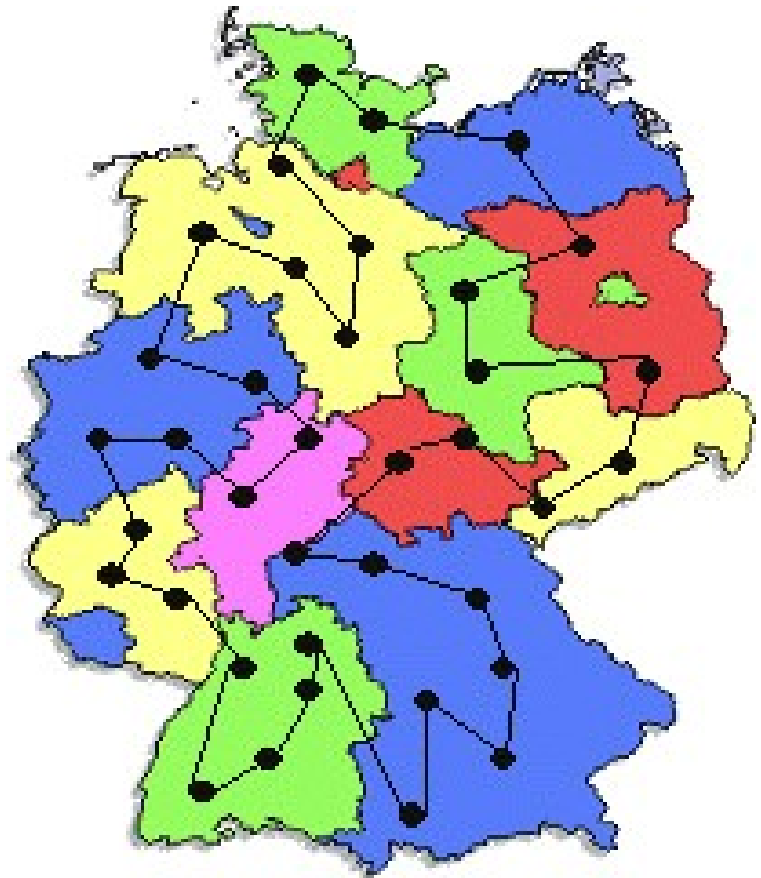
Decide if there is a set of k students who would make a harmonious group (everybody gets along).



Just the Clique problem in **disguise!**

Traveling Salesman Problem (aka UPS Truck problem)

- Input: n points and all pairwise inter-point distances, and a distance k
- Decide: is there a path that visits all the points (“salesman tour”) whose total length is at most k ?



Finals scheduling



- Input: n students, k classes, enrollment lists, m time slots in which to schedule finals
- Define “conflict”: a student is in two classes that have finals in the same time slot
- Decide: if schedule with at most 100 conflicts exists?

The P vs NP Question



- P: problems for which solutions can be found in polynomial time (n^c where c is a fixed integer and n is “input size”). Example: Rumor Mill
- NP: problems where a *good solution* can be checked in n^c time. Examples: Boolean Satisfiability, Traveling Salesman, Clique, Finals Scheduling
- Question: **Is P = NP?**
“Can we automate brilliance?”

(Note: Choice of computational model --- Turing-Post, pseudocode, C++ etc. --- irrelevant.)

NP-complete Problems

Problems in NP that are “the hardest”

- If they are in P then so is **every** NP problem.

Examples: Boolean Satisfiability, Traveling Salesman, Clique,
Finals Scheduling, 1000s of others

How could we possibly prove these problems are “the hardest”?



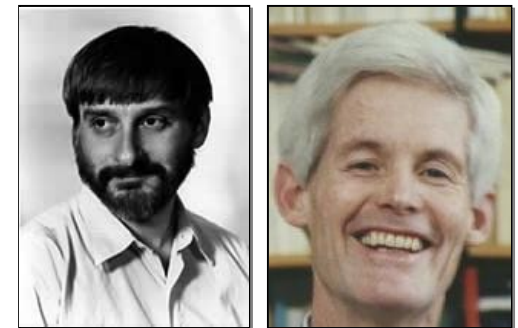
“Reduction”

“If you give me a place to stand, I will move the earth.”
– Archimedes (~ 250BC)



“If you give me a polynomial-time algorithm for Boolean Satisfiability, I will give you a polynomial-time algorithm for every NP problem.” --- Cook, Levin (1971)

“Every NP problem is a satisfiability problem in disguise.”





Dealing with NP-complete problems

1. **Heuristics** (algorithms that produce reasonable solutions in practice)
2. **Approximation algorithms** (compute provably near-optimal solutions)

Computational Complexity Theory:

Study of Computationally Difficult problems.

A new lens on the world?



- Study matter → look at mass, charge, etc.
- Study processes → look at computational difficulty

Example 1: Economics

General equilibrium theory:

- Input: n agents, each has some initial endowment (goods, money, etc.) and preference function
- General equilibrium: system of prices such that for every good, demand = supply.
- Equilibrium exists [Arrow-Debreu, 1954]. Economists assume markets find it (“invisible hand”)
- But, *no known* efficient algorithm to compute it. How does the market compute it?





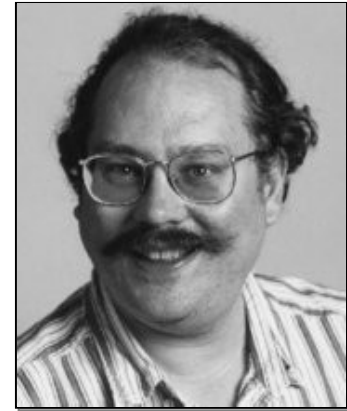
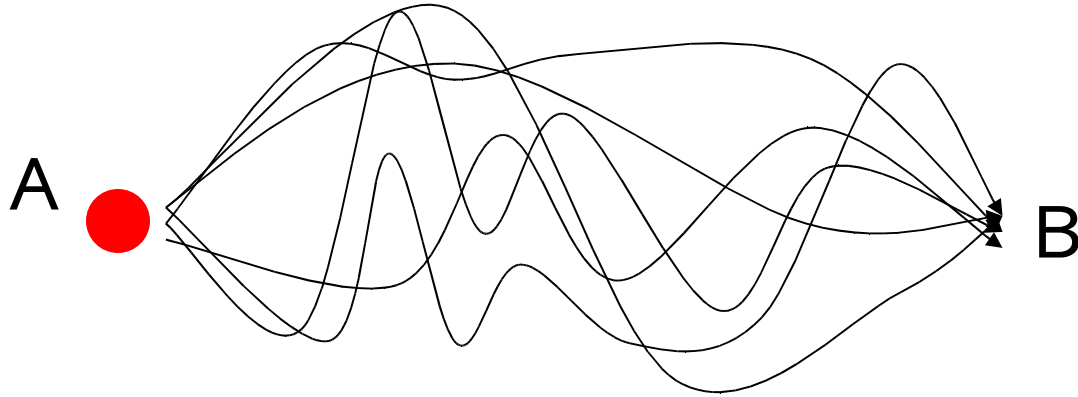
Example 2: Factoring problem

Given a number n , find two numbers p , q (neither of which is 1) such that $n = p \times q$.

Any suggestions how to solve it?

Fact: This problem is believed to be hard.
It is the basis of much of cryptography.
(More next time.)

Example 3: Quantum Computation



Peter Shor

- Central tenet of quantum mechanics: when a particle goes from A to B, it takes all possible paths all at the same time
- [Shor'97] Can use quantum behavior to efficiently factor integers (and break cryptosystems!)
- Can quantum computers be built, or is quantum mechanics not a correct description of the world?

Example 4: Artificial Intelligence

What is computational complexity of language recognition?

Chess playing?

Etc. etc.



Potential way to show the brain is not a computer:
Show it routinely solves some problem that provably takes exponential time on computers.

(Will talk more about AI in a couple weeks)

Why is P vs NP a Million-dollar open problem?

- If $P = NP$ then Brilliance becomes routine (best schedule, best route, best design, best math proof, etc...)
- If $P \neq NP$ then we know something *new and fundamental* not just about computers but about the world (akin to “Nothing travels faster than light”).

Next time: Cryptography (practical use of computational complexity)

