



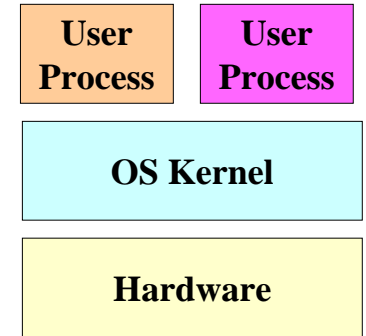
# Processes

CS 217



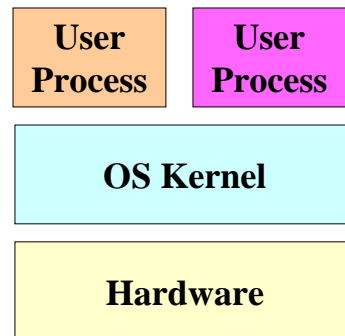
# Operating System

- Supports virtual machines
  - Promises each process the illusion of having whole machine to itself
- Provides services:
  - Protection
  - Scheduling
  - Memory management
  - File systems
  - Synchronization
  - etc.



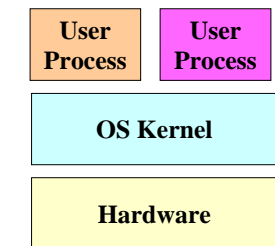
# What is a Process?

- A process is a running program with its own ...
  - Processor state
    - EIP, EFLAGS, registers
  - Address space (memory)
    - Text, bss, data, heap, stack



# Operating System

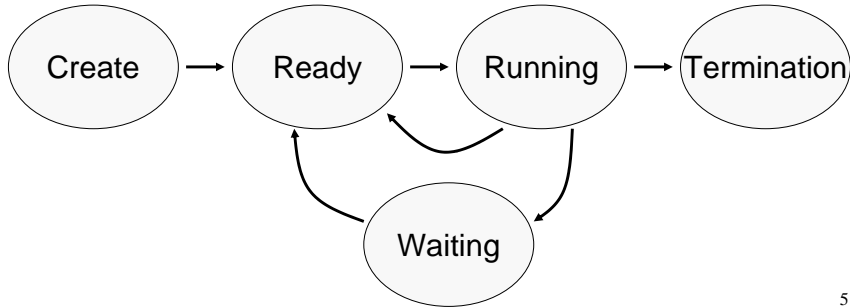
- Resource allocation
  - Sharing
  - Protection
  - Fairness
  - Higher-level abstractions
- Common strategies
  - Chop up resources into small pieces and allocate small pieces at fine-grain level
  - Introduce level of indirection and provide mapping from virtual resources to physical ones
  - Use past history to predict future behavior



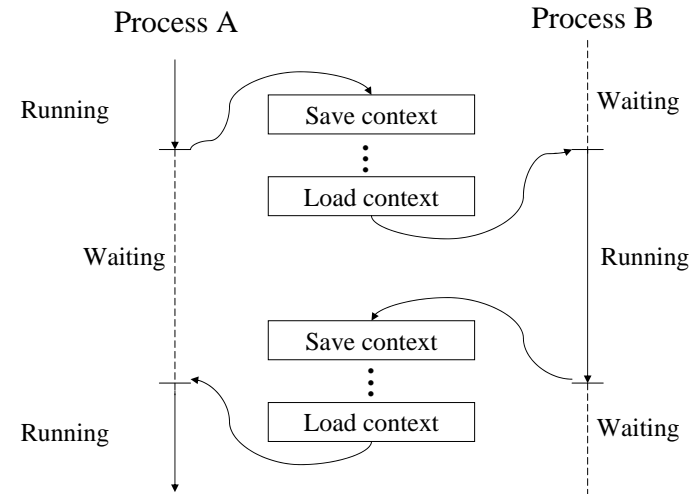


# Life Cycle of a Process

- Running: instructions are being executed
- Waiting: waiting for some event (e.g., i/o finish)
- Ready: ready to be assigned to a processor

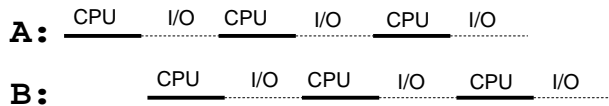


# Context Switch



# Overlap CPU with I/O operations

- Schedule CPU for process B while process A is waiting for I/O
  - Better utilize CPU



# Process Control Block

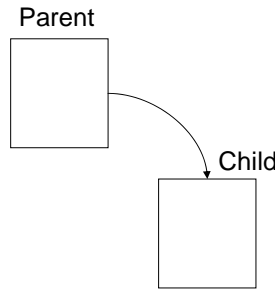
- For each process, the kernel keeps track of ...
  - Process state (new, ready, waiting, halted)
  - CPU registers (EIP, EFLAGS, EAX, EBX, ...)
  - CPU scheduling information (priority, queues, ...)
  - Memory management information (page tables, ...)
  - Accounting information (time limits, group ID, ...)
  - I/O status information (open files, I/O requests, ...)



# Fork

- Create a new process (system call)
  - child process inherits state from parent process
  - parent and child have separate copies of that state
  - parent and child share access to any open files

```
pid = fork();
if (pid != 0) {
    /* in parent */
    ...
}
/* in child */
...
```



# Wait

- Parent waits for a child (system call)
  - blocks until a child terminates
  - returns pid of the child process
  - returns -1 if no children exist (already exited)
  - status

```
#include <sys/types.h>
#include <sys/wait.h>

pid_t wait(int *status);
```

- Parent waits for a specific child to terminate

```
#include <sys/types.h>
#include <sys/wait.h>

pid_t waitpid(pid_t pid, int *status, int options);
```



# Fork

- |  |   |
|--|---|
| <ul style="list-style-type: none"> <li>• Inherited:           <ul style="list-style-type: none"> <li>◦ user and group IDs</li> <li>◦ environment</li> <li>◦ close-on-exec flag</li> <li>◦ signal handling settings</li> <li>◦ supplementary group IDs</li> <li>◦ set-user-ID mode bit</li> <li>◦ set-group-ID mode bit</li> <li>◦ profiling on/off/mode status</li> <li>◦ debugger tracing status</li> <li>◦ nice value</li> <li>◦ stdin</li> <li>◦ scheduler class</li> <li>◦ all shared memory segments</li> <li>◦ all mapped files</li> <li>◦ file pointers</li> <li>◦ non-degrading priority</li> <li>◦ process group ID</li> <li>◦ session ID</li> <li>◦ current working directory</li> <li>◦ root directory</li> <li>◦ file mode creation mask</li> <li>◦ resource limits</li> <li>◦ controlling terminal</li> <li>◦ all machine register states</li> <li>◦ control register(s)</li> </ul> </li> </ul> | <ul style="list-style-type: none"> <li>• Separate in child           <ul style="list-style-type: none"> <li>◦ process ID</li> <li>◦ address space (memory)</li> <li>◦ file descriptors</li> <li>◦ active process group ID.</li> <li>◦ parent process ID</li> <li>◦ process locks, file locks, page locks, text locks and data locks</li> <li>◦ pending signals</li> <li>◦ timer signal reset times</li> <li>◦ share mask</li> </ul> </li> </ul> |
|--|---|



# Exec

- Overlay current process image with a specified image file (system call)
  - affects process memory and registers
  - has no affect on file table

- Example:
 

```
execlp("ls", "ls", "-l", NULL);
fprintf(stderr, "exec failed\n");
exit(1);
```



## Exec (cont)

- Many variations of `exec`

```
int execlp(const char *file,
           const char *arg, ...)
```

```
int execl(const char *path,
           const char *arg, ...)
```

```
int execv(const char *path,
           char * const argv[])
```

```
int execl(const char *path,
           const char *arg, ...,
           char * const envp[])
```
- Also `execve` and `execvp`

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## Fork/Exec

- Commonly used together by the shell
 

```
... parse command line ...
pid = fork()
if (pid == -1)
    fprintf(stderr, "fork failed\n");
else if (pid == 0) {
    /* in child */
    execvp(file, argv);
    fprintf(stderr, "exec failed\n");
}
else {
    /* in parent */
    pid = wait(&status);
}
... return to top of loop ...
```

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## System

- Convenient way to invoke `fork/exec/wait`
  - Forks new process
  - Execs command
  - Waits until it is complete

```
int system(const char *cmd);
```

- Example:

```
int main()
{
    system("echo Hello world");
}
```

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## Summary

- Operating systems manage resources
  - Divide up resources (e.g., quantum time slices)
  - Allocate them (e.g., process scheduling)
- A processes is a running program with its own ...
  - Processor state
  - Address space (memory)
- Create and manage processes with ...
  - `fork`
  - `exec`
  - `wait`
  - `system`

} Used in shell

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