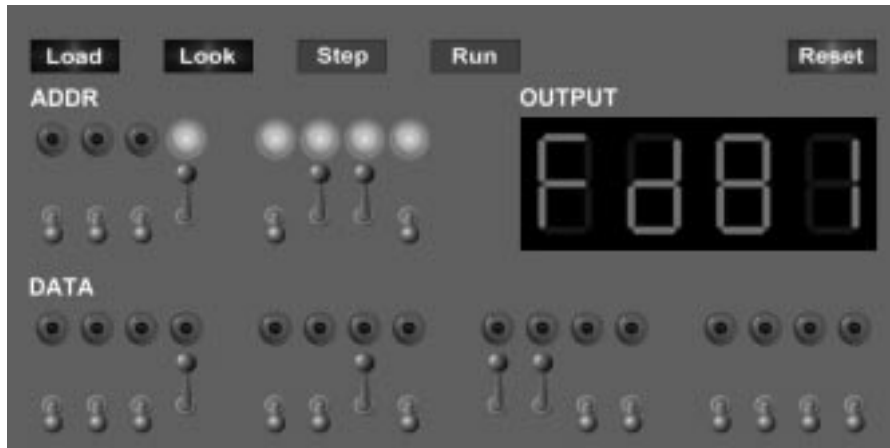


# Lecture A1: The X-TOY Machine



## What is X-TOY?

An imaginary machine similar to:

- Ancient computers.
- Today's microprocessors.

Why study?

- Machine language programming.
  - how do C programs relate to computer?
  - still (a few) situations today where it is really necessary
- Computer architecture.
  - how is a computer put together?
  - how does it work?
- Simplified machine.
  - captures essence of real computers

## Inside the Box

Switches.

- Input data and programs.

Lights.

- View data.

Registers.

- Fastest form of storage.
  - each stores 16 bits
- Use as scratch space during computation.
- 16 registers.
  - each stores 16 bits
- Register 0 is always 0.

Program counter (PC).

- An extra 8-bit register.
- Keeps track of next instruction to be executed.

Memory.

- Store data and programs.
- 256 "words."
  - each word stores 16 bits
- FF for stdin / stdout.

ALU (arithmetic-logic unit).

- Manipulate data.

## Data and Programs Encoded in Binary

Each bit consists of two states:

- Switch is ON or OFF.
- High voltage or low voltage.
- 1 or 0.
- True or false.

| Dec | Bin  | Dec | Bin  |
|-----|------|-----|------|
| 0   | 0000 | 8   | 1000 |
| 1   | 0001 | 9   | 1001 |
| 2   | 0010 | 10  | 1010 |
| 3   | 0011 | 11  | 1011 |
| 4   | 0100 | 12  | 1100 |
| 5   | 0101 | 13  | 1101 |
| 6   | 0110 | 14  | 1110 |
| 7   | 0111 | 15  | 1111 |

How to represent integers?

- Use binary encoding.
- Ex:  $6375_{10} = 0001100011100111_2$

| 15 | 14 | 13 | 12 | 11 | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
|----|----|----|----|----|----|---|---|---|---|---|---|---|---|---|---|
| 0  | 0  | 0  | 1  | 1  | 0  | 0 | 0 | 1 | 1 | 1 | 0 | 0 | 1 | 1 | 1 |

$$\begin{aligned}
 6375_{10} &= +2^{12} + 2^{11} && +2^7 + 2^6 + 2^5 && +2^2 + 2^1 + 2^0 \\
 &= 4096 + 2048 && +128 + 64 + 32 && +4 + 2 + 1
 \end{aligned}$$

## Shorthand Notation

Use hexadecimal (base 16) representation.

- Binary code, four bits at a time.
- Ex:  $6375_{10} = 0001100011100111_2 = 18E7_{16}$

| Dec | Bin  | Hex | Dec | Bin  | Hex |
|-----|------|-----|-----|------|-----|
| 0   | 0000 | 0   | 8   | 1000 | 8   |
| 1   | 0001 | 1   | 9   | 1001 | 9   |
| 2   | 0010 | 2   | 10  | 1010 | A   |
| 3   | 0011 | 3   | 11  | 1011 | B   |
| 4   | 0100 | 4   | 12  | 1100 | C   |
| 5   | 0101 | 5   | 13  | 1101 | D   |
| 6   | 0110 | 6   | 14  | 1110 | E   |
| 7   | 0111 | 7   | 15  | 1111 | F   |

| 15 | 14 | 13 | 12 | 11 | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
|----|----|----|----|----|----|---|---|---|---|---|---|---|---|---|---|
| 0  | 0  | 0  | 1  | 1  | 0  | 0 | 0 | 1 | 1 | 1 | 0 | 0 | 1 | 1 | 1 |
| 1  |    |    |    | 8  |    |   |   | E |   |   |   | 7 |   |   |   |

$$6375_{10} = 1 \times 16^3 + 8 \times 16^2 + 14 \times 16^1 + 7 \times 16^0$$

$$= 4096 + 2048 + 224 + 7$$

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## Machine "Core" Dump

EVERYTHING is encoded in binary (hex).

- Integers.
- Machine instructions.
- Text.
- Reals.
- ...

Machine contents at a particular place and time.

- Record of what program has done.
- Completely determines what program will do.

| Registers |      |      |      |      |      |      |      | PC |
|-----------|------|------|------|------|------|------|------|----|
| 0         | 1    | 2    | 3    | 4    | 5    | 6    | 7    | 10 |
| 0000      | 0788 | B700 | 0010 | 0401 | 0002 | 0003 | 00A0 | 10 |
| 8         | 9    | A    | B    | C    | D    | E    | F    |    |
| 0000      | 0788 | B700 | 0010 | 0401 | 0002 | 0003 | 00A0 |    |

| Main Memory |      |      |      |      |      |      |      |      |
|-------------|------|------|------|------|------|------|------|------|
| 00:         | 0000 | 0000 | 0000 | 0000 | 0000 | 0000 | 0000 | 0000 |
| 08:         | 0000 | 0000 | 0000 | 0000 | 0000 | 0000 | 0000 | 0000 |
| 10:         | 9222 | 9120 | 1121 | A120 | 1121 | A121 | 7211 | 0000 |
| 18:         | 0000 | 0001 | 0002 | 0003 | 0004 | 0005 | 0006 | 0007 |
| 20:         | 0008 | 0009 | 000A | 000B | 000C | 000D | 000E | 000F |
| 28:         | 0000 | 0000 | 0000 | FE10 | FACE | CAFE | ACED | CEDE |
| .           |      |      |      |      |      |      |      |      |
| .           |      |      |      |      |      |      |      |      |
| E8:         | 1234 | 5678 | 9ABC | DEF0 | 0000 | 0000 | F00D | 0000 |
| F0:         | 0000 | 0000 | EEEE | 1111 | EEEE | 1111 | 0000 | 0000 |
| F8:         | B1B2 | F1F5 | 0000 | 0000 | 0000 | 0000 | 0000 | 0000 |

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## Program and Data

Program:

- Sequence of instructions.

16 instruction types:

- 16-bit word (interpreted one way).
- Changes contents of registers, memory, and PC in specified, well-defined ways.

Data:

- 16-bit word (interpreted other way).

Program counter (PC):

- Stores memory address of "next instruction."

| Instructions |                 |
|--------------|-----------------|
| 0:           | halt            |
| 1:           | add             |
| 2:           | subtract        |
| 3:           | and             |
| 4:           | xor             |
| 5:           | shift left      |
| 6:           | shift right     |
| 7:           | load address    |
| 8:           | load            |
| 9:           | store           |
| A:           | load indirect   |
| B:           | store indirect  |
| C:           | branch zero     |
| D:           | branch positive |
| E:           | jump register   |
| F:           | jump and link   |

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## Using the X-TOY Machine: Input, Output

To enter a program or data:

- Set 8 memory address switches.
- Set 16 data switches.
- Press LOAD.
  - data written into addressed word of memory

To view the results of a program:

- Set 8 memory address switches.
- Press LOOK.
  - contents of addressed word of memory appears in lights



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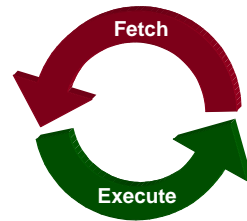
## Using the X-TOY Machine: Run

### To run the program:

- Set 8 memory address switches to address of first instruction.
- Press the RUN or STEP button.
  - loads PC from address switches
  - initiates fetch-execute cycle
  - machine repeats fetch-execute cycle until halt instruction

### Fetch-execute cycle.

- FETCH.**
  - get instruction from memory
- EXECUTE.**
  - update PC
  - move data to or from memory, registers
  - perform calculations



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## X-TOY Instruction Set Architecture

### X-TOY instruction set architecture (ISA).

- Interface that specifies behavior of machine.
- 16 register, 256 words of main memory, 16-bit words.
- 16 instructions.

### Each instruction consists of 16 bits.

- Bits 12-15 encode one of 16 instruction types or opcodes.
- Bits 8-11 encode destination register d.
- Bits 0-7 encode:
  - Format 1: source registers s and t
  - Format 2: 8-bit memory address or constant

|                 |        |    |    |    |        |    |   |   |          |   |   |   |          |   |   |   |
|-----------------|--------|----|----|----|--------|----|---|---|----------|---|---|---|----------|---|---|---|
|                 | 15     | 14 | 13 | 12 | 11     | 10 | 9 | 8 | 7        | 6 | 5 | 4 | 3        | 2 | 1 | 0 |
|                 | 1      | 0  | 1  | 1  | 1      | 0  | 1 | 0 | 0        | 0 | 0 | 0 | 0        | 1 | 0 | 0 |
| <b>Format 1</b> | opcode |    |    |    | dest d |    |   |   | source s |   |   |   | source t |   |   |   |
| <b>Format 2</b> | opcode |    |    |    | dest d |    |   |   | addr     |   |   |   |          |   |   |   |

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## Load Address

### Load address. (opcode 7)

- Loads an 8-bit integer into a register.
- Format 2.
- 7A04 means:
  - load the value 0004 into register A
  - $R[A] \leftarrow 0004$

| C code |
|--------|
| a = 4; |

|                 |    |    |    |                 |    |   |   |                 |   |   |   |                 |   |   |   |
|-----------------|----|----|----|-----------------|----|---|---|-----------------|---|---|---|-----------------|---|---|---|
| 15              | 14 | 13 | 12 | 11              | 10 | 9 | 8 | 7               | 6 | 5 | 4 | 3               | 2 | 1 | 0 |
| 0               | 1  | 1  | 1  | 1               | 0  | 1 | 0 | 0               | 0 | 0 | 0 | 0               | 1 | 0 | 0 |
| 7 <sub>16</sub> |    |    |    | A <sub>16</sub> |    |   |   | 0 <sub>16</sub> |   |   |   | 4 <sub>16</sub> |   |   |   |
| opcode          |    |    |    | dest d          |    |   |   | addr            |   |   |   |                 |   |   |   |

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## Load, Store

### Load. (opcode 8)

- Loads the contents of some memory location into a register.
- 8A04 means:
  - load the contents of memory location 04 into register A
  - $R[A] \leftarrow \text{mem}[04]$

### Store. (opcode 9)

- Opposite of load.
- Store the contents of a register into main memory.

| load-store.toy |                     |
|----------------|---------------------|
| 04:            | 1111                |
| 05:            | 0000                |
| 10:            | 8A04 R[A] ← mem[04] |
| 11:            | 9A05 mem[05] ← R[A] |
| 12:            | 0000 halt           |

|                 |    |    |    |                 |    |   |   |                 |   |   |   |                 |   |   |   |
|-----------------|----|----|----|-----------------|----|---|---|-----------------|---|---|---|-----------------|---|---|---|
| 15              | 14 | 13 | 12 | 11              | 10 | 9 | 8 | 7               | 6 | 5 | 4 | 3               | 2 | 1 | 0 |
| 1               | 0  | 0  | 0  | 1               | 0  | 1 | 0 | 0               | 0 | 0 | 0 | 0               | 1 | 0 | 0 |
| 8 <sub>16</sub> |    |    |    | A <sub>16</sub> |    |   |   | 0 <sub>16</sub> |   |   |   | 4 <sub>16</sub> |   |   |   |
| opcode          |    |    |    | dest d          |    |   |   | addr            |   |   |   |                 |   |   |   |

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## Add, Subtract

### Add. (opcode 1)

- Add contents of two registers and store sum in a third.
- 1CAB means:

- add contents of registers A and B
- put result in register C
- $R[C] \leftarrow R[A] + R[B]$

| add.toy  |                                  |
|----------|----------------------------------|
| 00: 0005 | 5                                |
| 01: 0008 | 8                                |
| 10: 8A00 | $R[A] \leftarrow \text{mem}[00]$ |
| 11: 8B01 | $R[B] \leftarrow \text{mem}[01]$ |
| 12: 1CAB | $R[C] \leftarrow R[A] + R[B]$    |
| 13: 9C02 | $\text{mem}[02] \leftarrow R[C]$ |
| 14: 0000 | halt                             |

### Subtract. (opcode 2)

- Analogous to add.

| 15       | 14 | 13 | 12 | 11       | 10 | 9 | 8 | 7        | 6 | 5 | 4 | 3        | 2 | 1 | 0 |
|----------|----|----|----|----------|----|---|---|----------|---|---|---|----------|---|---|---|
| 0        | 0  | 0  | 1  | 1        | 1  | 0 | 0 | 1        | 0 | 1 | 0 | 1        | 0 | 1 | 1 |
| $1_{16}$ |    |    |    | $C_{16}$ |    |   |   | $A_{16}$ |   |   |   | $B_{16}$ |   |   |   |
| opcode   |    |    |    | dest d   |    |   |   | source s |   |   |   | source t |   |   |   |

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## Branch Zero, Branch Positive

### Branch if zero. (opcode C)

- Changes PC depending of value of some register.
- Used to implement loops, if-else.

### Multiply.

- Load in integers a and b, and store  $c = a \times b$ .
- Brute-force algorithm:
  - initialize  $c = 0$
  - add b to c, a times
- Problems.

| multiply.c               |         |
|--------------------------|---------|
| int a = 7, b = 8;        |         |
| int c;                   |         |
| for (c = 0; a != 0; a--) | c += b; |

### Branch if positive. (opcode D)

- Analogous.

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## Branch Zero, Branch Positive

### Branch if zero. (opcode C)

- Changes PC depending of value of some register.
- Used to implement loops, if-else.

### Multiply.

- Load in integers a and b, and store  $c = a \times b$ .
- Brute-force algorithm:
  - initialize  $c = 0$
  - add b to c, a times

| multiply.toy |                                  |
|--------------|----------------------------------|
| 0A: 0007     | 7                                |
| 0B: 0008     | 8                                |
| 10: 8A0A     | $R[A] \leftarrow \text{mem}[0A]$ |
| 11: 8B0B     | $R[B] \leftarrow \text{mem}[0B]$ |
| 12: 7C00     | $R[C] \leftarrow 00$             |
| 13: 7101     | $R[1] \leftarrow 01$             |
| 14: CA18     | if (R[A] == 0) pc ← 18           |
| 15: 1CCB     | $R[C] += R[B]$                   |
| 16: 2AA1     | $R[A]--$                         |
| 17: C014     | pc ← 14                          |
| 18: 9C0C     | $\text{mem}[0C] \leftarrow R[C]$ |
| 19: 0000     | halt                             |

Annotations: "result" points to R[C] in line 12, "always 1" points to R[1] in line 13, "loop" points to the branch instruction in line 14, and "opcode C branch if zero" points to the branch instruction in line 14.

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## Step-By-Step Trace of multiply.toy

|          |                                  | R[1] | R[A] | R[B] | R[C] |
|----------|----------------------------------|------|------|------|------|
| 10: 8A0A | $R[A] \leftarrow \text{mem}[0A]$ |      | 0003 |      |      |
| 11: 8B0B | $R[B] \leftarrow \text{mem}[0B]$ |      |      | 0007 |      |
| 12: 7C00 | $R[C] \leftarrow 00$             |      |      |      | 0000 |
| 13: 7101 | $R[1] \leftarrow 01$             | 0001 |      |      |      |
| 14: CA18 | if (R[A] == 0) goto 18           |      |      |      |      |
| 15: 1CCB | $R[C] += R[B]$                   |      |      |      | 0007 |
| 16: 2AA1 | $R[A]--$                         |      | 0002 |      |      |
| 17: C014 | goto 14                          |      |      |      |      |
| 14: CA18 | if (R[A] == 0) goto 18           |      |      |      |      |
| 15: 1CCB | $R[C] += R[B]$                   |      |      |      | 000E |
| 16: 2AA1 | $R[A]--$                         |      | 0001 |      |      |
| 17: C014 | goto 14                          |      |      |      |      |
| 14: CA18 | if (R[A] == 0) goto 18           |      |      |      |      |
| 15: 1CCB | $R[C] += R[B]$                   |      |      |      | 0015 |
| 16: 2AA1 | $R[A]--$                         |      | 0000 |      |      |
| 17: C014 | goto 14                          |      |      |      |      |
| 14: CA18 | if (R[A] == 0) goto 18           |      |      |      |      |
| 18: 9C0C | $\text{mem}[0C] \leftarrow R[C]$ |      |      |      |      |
| 19: 0000 | halt                             |      |      |      |      |

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## A Little History

### ENIAC. (Eckert and Mauchly, 1946)

- First general purpose electronic computer.
- 30 x 50 x 8.5 ft.
- 17,468 vacuum tubes.
- 300 multiplication per second.
- Conditional jumps, programmable.
  - code: set switches
  - data: punch cards
  - used to compute artillery firing tables

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## TOY Cheat Sheet

|          | 15     | 14 | 13 | 12 | 11     | 10 | 9 | 8 | 7        | 6 | 5 | 4 | 3        | 2 | 1 | 0 |
|----------|--------|----|----|----|--------|----|---|---|----------|---|---|---|----------|---|---|---|
| Format 1 | opcode |    |    |    | dest d |    |   |   | source s |   |   |   | source t |   |   |   |
| Format 2 | opcode |    |    |    | dest d |    |   |   | addr     |   |   |   |          |   |   |   |

| #  | Operation       | Fmt | Pseudocode  |
|----|-----------------|-----|---|
| 0: | halt            | 1   | exit(0)   |
| 1: | add             | 1   | $R[d] \leftarrow R[s] + R[t]$                                 |
| 2: | subtract        | 1   | $R[d] \leftarrow R[s] - R[t]$                                 |
| 3: | and             | 1   | $R[d] \leftarrow R[s] \& R[t]$                                |
| 4: | xor             | 1   | $R[d] \leftarrow R[s] \wedge R[t]$                            |
| 5: | shift left      | 1   | $R[d] \leftarrow R[s] \ll R[t]$                               |
| 6: | shift right     | 1   | $R[d] \leftarrow R[s] \gg R[t]$                               |
| 7: | load addr       | 2   | $R[d] \leftarrow \text{addr}$                                 |
| 8: | load            | 2   | $R[d] \leftarrow \text{mem}[\text{addr}]$                     |
| 9: | store           | 2   | $\text{mem}[\text{addr}] \leftarrow R[d]$                     |
| A: | load indirect   | 1   | $R[d] \leftarrow \text{mem}[R[t]]$                            |
| B: | store indirect  | 1   | $\text{mem}[R[t]] \leftarrow R[d]$                            |
| C: | branch zero     | 2   | if $(R[d] == 0)$ $\text{pc} \leftarrow \text{addr}$           |
| D: | branch positive | 2   | if $(R[d] > 0)$ $\text{pc} \leftarrow \text{addr}$            |
| E: | jump register   | 2   | $\text{pc} \leftarrow R[d]$                                   |
| F: | jump and link   | 2   | $R[d] \leftarrow \text{pc}; \text{pc} \leftarrow \text{addr}$ |

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