## Computer Science


7. Digital Circuits

- boolean algebra
- logic gates
- adder circuit


## Context

Q. How are computers built?
A. Not nearly as complicated as you might think.

This lecture. Introduction to digital circuits.

- Digital = all signals are either 0 or 1 .
- Analog = signals vary continuously.
- Advantages of digital: accurate, reliable, fast, cheap, scalable, ...


Applications. Laptop, smartphone, gaming console, pacemaker, microprocessor, ...



## 7. Digital Circuits

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## Boolean algebra

Boolean algebra. Developed by George Boole in 1840s to study logic problems.

- Values of variables are true (1) or false (0).
- Primitive operations are NOT, AND, and OR.
- Widely used in mathematics, logic, computer science, ...

| operation | logic <br> notation | Java <br> notation | circuit <br> notation | precedence |
| :---: | :---: | :---: | :---: | :---: |
| NOT | $\neg \mathrm{x}$ | $!\mathrm{x}$ | $x^{\prime}$ | highest |
| OR | $\mathrm{x} \wedge \mathrm{y}$ | $\mathrm{x} \& \& \mathrm{y}$ | $x \cdot y$ | middle |
| $\mathrm{x} \vee \mathrm{y}$ | $\mathrm{x} \\| \mathrm{y}$ | $x+y$ | lowest |  |
|  |  |  |  |  |
|  |  |  |  |  |

Relevance to circuits. Provides the mathematical foundation.


George Boole is Coole


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## Truth tables

Boolean function. A function whose arguments and result assume the values 0 and 1 .

Truth table. A systematic way to define a boolean function.

- One row for each possible assignment of arguments.
- Each row gives the function value for the specified arguments.
- The truth table of a boolean function of $n$ variables has $2^{n}$ rows.



## Boolean algebra properties

Boolean algebra shares many properties with elementary algebra. $\qquad$ justifies use of • and + for AND and OR


## Proving a theorem in Boolean algebra

Q. How to prove a theorem, such as De Morgan's law?

A1. Apply sequence of axioms or known theorems.

A2. For each possible assignment of truth values to variables, evaluate the purported theorem; confirm that it is true.

Ex. De Morgan's law: $(x \cdot y)^{\prime}=\left(x^{\prime}+y^{\prime}\right)$.

| $x$ | $y$ | $x \cdot y$ | $(x \cdot y)^{\prime}$ |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 1 |
| 0 | 1 | 0 | 1 |
| 1 | 0 | 0 | 1 |
| 1 | 1 | 1 | 0 |

truth table for LHS

| $x$ | $y$ | $x^{\prime}$ | $y^{\prime}$ | $x^{\prime}+y^{\prime}$ |
| :--- | :--- | :--- | :--- | :--- |
| 0 | 0 | 1 | 1 | 1 |
| 0 | 1 | 1 | 0 | 1 |
| 1 | 0 | 0 | 1 | 1 |
| 1 | 1 | 0 | 0 | 0 |

## Boolean functions of two variables

Boolean function. A function whose arguments and result assume the values 0 and 1 .

| $x$ | $y$ | $A N D$ | $O R$ | NAND | NOR | XOR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 1 | 1 | 0 |
| 0 | 1 | 0 | 1 | 1 | 0 | 1 |
| 1 | 0 | 0 | 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 | 0 | 0 | 0 |

commonly used boolean functions of 2 variables


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## Boolean functions of three (and more) variables

Boolean function. A function whose arguments and result assume the values 0 and 1 .

| $x$ | $y$ | $z$ | AND | OR | MAJ | $O D D$ |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | function | shorthand | description |
| 0 | 0 | 1 | 0 | 1 | 0 | 1 |  |  |  |
| 0 | 1 | 0 | 0 | 1 | 0 | 1 | logical AND | AND | all inputs are 1 |
| 0 | 1 | 1 | 0 | 1 | 1 | 0 | logical OR | OR | any input is 1 |
| 1 | 0 | 0 | 0 | 1 | 0 | 1 | majority | MAJ | more inputs are 1 than 0 |
| 1 | 0 | 1 | 0 | 1 | 1 | 0 | odd parity | $O D D$ | odd number of inputs are 1 |
| 1 | 1 | 0 | 0 | 1 | 1 | 0 |  |  |  |
| 1 | 1 | 1 | 1 | 1 | 1 | 1 | these functions all extends to $n$ variab |  |  |

## Sum-of-products

Sum-of-products. Every boolean function can be represented as a sum of products.

- Products: form an AND term for each 1 in truth table.
- Sum: combine the terms with the $O R$ function.
also known as
"disjunctive normal form"

| $x$ | $y$ | $z$ | $M A J$ | $x^{\prime} \cdot y \cdot z$ | $x \cdot y^{\prime} \cdot z$ | $x \cdot y \cdot z^{\prime}$ | $x \cdot y \cdot z$ | $\quad$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 1 | 1 | 1 | 1 | 0 | 0 | 0 | 1 |
| 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 |
| 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 | 0 | 0 | 0 | 1 | 1 |

Digital circuits: quiz 1

## Which of the following does NOT represent majority function?

A. $(x \cdot y)+(y \cdot z)+(x \cdot z)$
B. $z\left(x^{\prime} y+x y^{\prime}\right)+x y$
C. $(x \cdot y)+(y \cdot z)$
D.

```
public static boolean majority(boolean x, boolean y, boolean z) {
    int count = 0;
    if (x) count++;
    if (y) count++;
    if (z) count++;
    return count >= 2;
}
```


## Universality

Def. A set of operations is universal if every boolean function can be expressed using just those operations.

Proposition. $\{A N D, O R, N O T\}$ is a universal set of operations.
Pf. Sum-of-products construction on previous slide.

Proposition. $\{N A N D$ \} is a universal set of operations.
Pf. $\{A N D, O R, N O T\}$ can be constructed from NAND.

| $x$ | $y$ | NAND |
| :---: | :---: | :---: |
| 0 | 0 | 1 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 0 |

NAND


## 7. Digital Circuits

- boolean a algèbra
- logic gates
- adder circuit


## A basis for digital devices

Claude Shannon. Identified the deep connection between Boolean algebra and circuits.

- Demonstrated how circuits could be analyzed using Boolean algebra.
- Designed circuits to perform mathematical operations on binary numbers. $\qquad$ add, subtract, multiply, factor, ...


Claude Shannon's master's thesis at MIT (1937)

Impact. Every electronic device we use today is based upon Shannon's foundational work.


## Primitive logic gates: AND, OR, and NOT

Logic gate. Physical device that implement a boolean function with one output.


## Digital circuits

Digital circuit. A network of logic gates connected by wires.

- Every wire is either on (1) or off (0).
- Can connect output of one gate to input of another gate.
- Any wire connected to a wire that is on is also on (and same for off).



## Digital circuits

Digital circuit. A network of logic gates connected by wires.

- Every wire is either on (1) or off (0).
- Can connect output of one gate to input of another gate.
- Any wire connected to a wire that is on is also on (and same for off).


| $x$ | $y$ | XOR |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 0 |

## Digital circuits: quiz 2

For which values of $x$ and $y$ does the following circuit output 1 ?
A. $x=0, y=0$
B. $x=0, y=1$
C. $x=1, y=0$
D. $x=1, y=1$
E. None of the above.


Multiway AND gates

Multiway AND gate.

- 1 if all inputs are 1 .
- 0 if any input is 0 .



## Multiway $O R$ gates

Multiway $O R$ gate.

- 1 if any input is 1 .
- 0 if all inputs are 0 .


8-way OR gate implementation (tree of 2-way OR gates)

## Generalized AND gates

## Generalized AND gate.

- 1 for exactly one set of input values.
- 0 for all other sets of input values.


4-way generalized AND gate implementation (tree of 2-way AND gates, plus NOT gates)

Majority function

Sum-of-products construction.

- Identify rows of truth table where the function is 1 .
- Use a generalized $A N D$ gate for each term.
- Combine the terms using an $O R$ gate.


Ex 1. Majority function.

| $x$ | $y$ | $z$ | $M A J$ |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | $1 \longleftarrow x^{\prime} y z$ |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | $1 \longleftarrow x y^{\prime} z$ |
| 1 | 1 | 0 | $1 \longleftarrow x y z^{\prime}$ |
| 1 | 1 | 1 | $1 \longleftarrow x y z$ |

$\operatorname{MAJ}(x, y, z)=x^{\prime} y z+x y^{\prime} z+x y z^{\prime}+x y z$


## Odd-parity function

## Sum-of-products construction.

- Identify rows of truth table where the function is 1 .
- Use a generalized $A N D$ gate for each term.
- Combine the terms using an $O R$ gate.

3-way odd parity circuit

[^0]

## Sum-of-products construction (summary)

Goal. Design a digital circuit that computes a given boolean function.

## Recipe.

- Step 1: Represent input and output with boolean variables.
- Step 2: Construct truth table to define the function.
- Step 3: Identify rows where the function is 1 .
- Step 4: Use a generalized $A N D$ gate for each row, and $O R$ the results.

Profound consequence. Can design a digital circuit for ANY boolean function.

## Optimized digital circuits

Caveat. Sum-of-products construction is not optimal in terms of:

- Space $=$ number of gates.
- Time = depth of circuit.


## Ex. Majority function.



Digital circuits: quiz 3

How many 3-way generalized AND gates are needed to build the sum-of-products circuit for the following truth table?
A. 1
B. 2
C. 3
D. 4

| $x$ | $y$ | $z$ | $E Q$ |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 1 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 0 |
| 1 | 1 | 1 | 1 |



## 7. Digital Circuits

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- logic gáatés
- adder circuit


## Let's make an adder circuit!

Adder circuit. Compute $z=x+y$ for 4-bit binary integers. ignore integer overflow

First step. Represent inputs and outputs in binary.


## Let's make an adder circuit!

Adder circuit. Compute $z=x+y$ for 4-bit binary integers. ignore integer overflow

First step. Represent inputs and outputs in binary.

$+$| 0 | 1 | 0 | 1 |
| :--- | :--- | :--- | :--- |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 1 | 1 |



Digital circuits: quiz 4

What is the binary sum $1011+0110 ?$
A. 0001
B. 1001
C. 1101

D. 1121
E. 10001

## Let's make an adder circuit!

Adder circuit. Compute $z=x+y$ for 4-bit binary integers.
exceeds number of
electrons in universe (!)
Straw-person solution. Build a truth table for each output bit.
Approach is not scalable! Truth table for 128 -bit adder would have $2^{256}$ rows.

truth table for 4-bit adder

## Let's make an adder circuit!

Adder circuit. Compute $z=x+y$ for 4-bit binary integers.

$$
+\begin{array}{llll} 
& & & c_{3} \\
c_{2} & c_{1} & c_{0} \\
x_{3} & x_{2} & x_{1} & x_{0} \\
y_{3} & y_{2} & y_{1} & y_{0} \\
\hline z_{3} & z_{2} & z_{1} & z_{0}
\end{array}
$$

Efficient solution. Do one bit at a time.

- Build truth table for each carry bit. $\qquad$ majority function (!)
- Build truth table for each sum bit.

| $x_{i}$ | $y_{i}$ | $c_{i}$ | $c_{i+1}$ | $M A J$ |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 0 |
| 0 | 1 | 0 | 0 | 0 |
| 0 | 1 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 | 1 |
| 1 | 1 | 0 | 1 | 1 |
| 1 | 1 | 1 | 1 | 1 |
| truth table for carry bit | $c_{i+1}=\operatorname{MAJ}\left(x_{i}, y_{i}, c_{i}\right)$ |  |  |  |

## Let's make an adder circuit!

Adder circuit. Compute $z=x+y$ for 4-bit binary integers.

$$
+\begin{array}{cccc}
c_{3} & c_{2} & c_{1} & c_{0} \\
x_{3} & x_{2} & x_{1} & x_{0} \\
y_{3} & y_{2} & y_{1} & y_{0} \\
\hline z_{3} & z_{2} & z_{1} & z_{0}
\end{array}
$$

Efficient solution. Do one bit at a time.

- Build truth table for each carry bit $\qquad$
- Build truth table for each sum bit. $\longleftarrow$ odd-parity function (!)

| $x_{i}$ | $y_{i}$ | $c_{i}$ | $z_{i}$ | $O D D$ |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 1 | 1 |
| 0 | 1 | 0 | 1 | 1 |
| 0 | 1 | 1 | 0 | 0 |
| 1 | 0 | 0 | 1 | 1 |
| 1 | 0 | 1 | 0 | 0 |
| 1 | 1 | 0 | 0 | 0 |
| 1 | 1 | 1 | 1 | 1 |

## Let's make an adder circuit!

Adder circuit. Compute $z=x+y$ for 4-bit binary integers.


Efficient solution. Do one bit at a time.

- Carry bit is MAJ.
- Sum bit is $O D D$.
- Chain 1-bit adders to "ripple" carries.

Size of circuit. $\Theta(n)$ gates for $n$-bit adder.


## Adder circuit trace

Circuit trace. Trace the execution of the adder circuit on a given input.


## Encapsulation

Encapsulation in circuit design mirrors familiar software design principle.

- API describes behavior (input and outputs) of circuit.
- Implementation gives details of how to build it from wires and gates.
- Client uses circuit as a black box.


Bottom line. We manage complexity by encapsulating circuits.

## Layers of abstraction

Layers of abstraction apply with a vengeance.

- On/off.
- Switch.
- Primitive gates (AND, OR, NOT).
- Composite gates (multiway AND/OR, MAJ, ODD).
- Adder circuit.
- Memory.
- Arithmetic logic unit (ALU).
- Central processing unit (CPU).
- Input and output.
- Your computer.

Want to learn more? See ECE 206 and ECE 365.


## Credits

Co-instructors, course admin, and graduate student preceptors.


Alan Kaplan


Sebastian Caldas


Kobi Kaplan

Undergrad graders and lab TAs. Apply to be one next semester!

## A final thought

## Credits

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[^0]:    $O D D(x, y, z)=x^{\prime} y^{\prime} z+x^{\prime} y z^{\prime}+x y^{\prime} z^{\prime}+x y z$

