

### 2.3 RECURSION

- foundations
- a classic example
- recursive graphics
- exponential waste
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## Overview

Recursion is when something is specified in terms of itself.

Why learn recursion?

- Represents a new mode of thinking.
- Provides a powerful programming paradigm.
- Reveals insight into the nature of computation.


Many computational artifacts are naturally self-referential.

- File system with folders containing folders.
- Binary trees.
- Fractal patterns.
- Depth-first search.
- Divide-and-conquer algorithms.

- ...


## Recursive functions (in Java)

Recursive function. A function that calls itself.

- Base case: If the result can be computed directly, do so.
- Reduction step: Otherwise, simplify by calling the function with one (or more) other arguments.

Ex. Factorial function: $n!=n \times(n-1) \times \cdots \times 3 \times 2 \times 1$.

- Base case: 1 ! = 1
- Reduction step:
$n!=n \times \frac{(n-1)!}{\uparrow}$
same function
with simpler argument
public class Factorial \{
public static int factorial(int n) \{

\}
public static void main(String[] args) \{
int $\mathrm{n}=$ Integer.parseInt(args[0]);
int result = factorial(n);
Std0ut.println(resu7t);
\}
\}


## Review: mechanics of a function call

1. Evaluate argument expressions and assign values to corresponding parameter variables.
2. Save environment (values of all local variables and call location).
3. Transfer control to the function.
4. Restore environment (with function-call expression evaluating to return value).
5. Transfer control back to the calling code.


## Factorial function demo

```
public static int factorial(int n) {
    if (n == 1) return 1;
    return n * factorial(n-1);
}
```

factorial(1)
factorial(2)
factorial (3)
main()

## Function-call trace

Function-call trace.

- Print name and arguments when each function is called.
- Print function's return value just before returning.
- Add indentation on function calls and subtract on returns.

```
factorial(5)
    factorial(4)
        factorial(3)
            factorial(2)
                    factorial(1)
                    return 1
            return 2 * 1 = 2
        return 3 * 2 = 6
        return 4 * 6 = 24
    return 5 * 24 = 120
```

function-call trace for factorial(5)

```
public static int factorial(int n) {
    if (n == 1) return 1;
    return n * factorial(n-1);
}
```


## Stack overflow errors



Problems with recursion?

https://www.smbc-comics.com

https://www.safelyendangered.com/comic/oh-bother

## Recursion: quiz 1

What is printed by a call to collatz(6)?
A. 63105168421
B. 12481651036
C. 2481651036
D. 631
E. stack overflow error

```
public static void collatz(int n) {
    StdOut.print(n + " ");
    if (n == 1) return;
    if (n % 2 == 0) collatz(n / 2);
    else collatz(3*n + 1);
}

\section*{Collatz sequence}

Famous unsolved problem. Does collatz(n) terminate for all \(n \geq 1\) ?
Partial answer. Yes, for all \(1 \leq n \leq 2^{68}\).


\section*{Saying the digits of a number}

Goal. Say the decimal digits in a positive integer \(n\).
- Base case: say nothing when \(n\) is zero.
- Reduction step: otherwise,
- recursively say most significant digits
- then, say the least significant digit

sayDigits(126)
sayDigits(12)
sayDigits(1)
sayDigits(0)
play "1.wav"
play "2.wav"
play "6.wav"
function-call trace of sayDigits(126)
Q. How to say digits in binary (instead of decimal)?
A. Replace constant 10 with constant 2.

\section*{Recursion: quiz 2}

What does sayDigits(126) do with this version of sayDigits() ?
```

public static void sayDigits(int n)
if (n == 0) return;
StdAudio.play((n % 10) + ".wav"); | «. the order of these two
sayDigits(n / 10);
}

```
A. Speaks "1 2 6."
B. Speaks "6 2 1."
C. Stack overflow error.


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\section*{Warmup: ruler function (revisited)}

Goal. Function ruler( \(n\) ) that returns first \(2^{n}-1\) values of ruler function.
- Base case: empty for \(n=0\).
- Reduction step: sandwich \(n\) between two copies of ruler \((n-1)\).

```

public class Ruler {
public static String ruler(int n) {
if (n == 0) return " "; \longleftarrow
return ruler(n-1) + n + ruler(n-1); \longleftarrow « reduction step
}
public static void main(String[] args) {
int n = Integer.parseInt(args[0]);
String result = ruler(n);
StdOut.println(result);
}
}

```
```

~/cos126/recursion> java-introcs Ruler 1
1
~/cos126/recursion> java-introcs Ruler 2
121
~/cos126/recursion> java-introcs Ruler 3
1213121
~/cos126/recursion> java-introcs Ruler 4
121312141213121

```

\section*{Tracing a recursive program}

Draw the function-call tree.
- One node for each function call.
- Label node with return value after children are labeled.

function-call tree for ruler(4)

\section*{Recursion: quiz 2}

Which string does ruler(3) return for this version of ruler()?
A. "1 1221123 "
B. "1 2131121 "
C. "3 211211 "
D. "3 2211111 "
```

public static String ruler(int n) {
if (n == 0) return "";
return n + " " + ruler(n-1) + ruler(n-1);
}

```

\section*{Towers of Hanoi puzzle}

\section*{A legend of uncertain origin.}
- \(n=64\) disks of differing size; 3 poles; disks on middle pole, from largest to smallest.
- An ancient prophecy has commanded monks to move the disks to another pole.
- When the task is completed, the world will end.

\section*{Rules.}
- Can move only one disk at a time.
- Cannot put a larger disk on top of a smaller disk.

Q1. How to generate a list of instruction for monks.
Q2. When might the world end?
start

\(\mathrm{n}=10\)
goal

\section*{Towers of Hanoi solution}

For instructions, use cyclic wraparound.
- Move right means 1 to 2 , 2 to 3 , or 3 to 1 .
- Move left means 1 to 3,3 to 2 , or 2 to 1 .


A recursive solution. [to move stack of \(n\) disks to the right]
- Base case: if \(n=0\) disks, do nothing.
- Reduction step: otherwise,
- move \(n-1\) smallest disks to the left (recursively)
- move largest disk to the right
- move \(n-1\) smallest disks to the left (recursively)


\section*{Towers of Hanoi solution \((\mathrm{n}=3)\)}

Notation. Label disks from smallest (1) to largest (n).


\section*{Towers of Hanoi: mutually recursive solution}

Goal. Function hanoiRight(n) that returns instructions for \(n\) disk puzzle. \(\qquad\) and also a similar function hanoi Left ( n )
- Base case: if \(n=0\) disks, do nothing.
- Reduction step: otherwise, sandwich moving disk \(n\) right
between two calls to hanoiLeft(n-1)
```

public class Hanoi {
public static String hanoiRight(int n) {
if (n == 0) return " ";
return hanoiLeft(n-1) + n + "R" + hanoiLeft(n-1);
}
return hanoiLeft(n-1) + $\mathrm{n}+\mathrm{"R}^{\mathrm{R}}+\mathrm{hanoiLeft(n-1)} \mathrm{;}$

```
    public static String hanoiLeft(int n) \{

move stack of move stack of
 \(n\) disks right
    if ( \(n==0\) ) return " ";
    return hanoiRight(n-1) + \(\mathrm{n}+\mathrm{"L"}+\) hanoiRight(n-1);
    \}
    public static void main(String[] args) \{
        int \(\mathrm{n}=\) Integer. parseInt(args[0]);
        StdOut. println(hanoiRight(n));
    \}
\}
concise but tricky code; read carefully!
                    concise but tricky code; read carefully!

\section*{Function-call tree for towers of Hanoi}

\section*{Properties.}
- Each disk always moves in the same direction.
- Moving smallest disk always alternates with (unique legal) move not involving smallest disk.
- Solution to puzzle with \(n\) disks makes \(2^{n}-1\) moves.


\section*{Answers for towers of Hanoi}
Q. How to generate instructions for monks?

A1. [long form] 1L 2R 1L 3L 1L 2R 1L 4R 1L 2R 1L 3L 1L 2R 1L 5L 1L 2R 1L 3L 1L 2R 1L 4R ...
A2. [short form] Alternate 1 L with the only legal move not involving disk 1.
Q. When might the world end?
A. Not soon. Takes \(2^{64}-1\) moves.
recursive solution
provably uses fewest moves


\section*{Recursion vs. iteration}

Fact 1. Any recursive program can be rewritten with loops (and no recursion).
Fact 2. Any program with loops can be rewritten with recursion (and no loops).
\begin{tabular}{cc} 
loops & recursion \\
\hline \begin{tabular}{c} 
more memory efficient \\
(no function-call stack)
\end{tabular} & concise and elegant code \\
easier to trace code & easier to reason about code \\
(fewer variables) & (fewer mutable variables)
\end{tabular}
Q. When should I use recursion?

A1. The problem is naturally recursive (e.g., towers of Hanoi).
A2. The data is naturally recursive (e.g., filesystem with folders).


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\section*{Recursive graphics in the wild}

"Hello, World" of recursive graphics: H-trees

\section*{H-tree of order \(n\).}
- Base case: if \(n\) is 0 , draw nothing.
- Reduction step:
- draw an H
- draw four H -trees of order \(n-1\) and half the size, centered at the tips of the H

order 1

order 3

\section*{H－trees}

\section*{Application．Connect a large set of regularly spaced sites to a single source．}

保安
的山幽织
近 4
 Y
约山幽
 HGHyyyyyy 隹出







宸

为


\section*{Recursive H -tree implementation}
```

public class Htree {
if (n == 0) return;
| double x0 = x - size/2, x1 = x + size/2;

```

```

    StdDraw.line(x1, y0, x1, y1);
    draw(n-1, size/2, x0, y0); // Tower 7eft
    draw(n-1, size/2, x0, y1); // upper left
    draw(n-1, size/2, x1, y0); // lower right
    draw(n-1, size/2, x1, y1); // upper right
    }
public static void main(String[] args) {
StdDraw.setPenRadius(0.005);
int n = Integer.parseInt(args[0]);
draw(n, 0.5, 0.5, 0.5);
}
}

```
public static void draw(int \(n\), double size, double \(x\), double \(y\) ) \{


\section*{Recursion: quiz 3}

Suppose that Htree (with \(\mathrm{n}=4\) ) is stopped after drawing the \(30^{\text {th }} \mathrm{H}\).
Which drawing will result?
A.

B.

C.

D.

Q. What will happen if we add the following statements to draw(), just before the recursive calls?
```

double freq = Synth.midiToFrequency(n + 45);
double duration = 0.25* n;
doub7e[] a = Synth.supersaw(freq, duration);
StdAudio.play(a);

```


\section*{Every semester, Princeton University's COS 126 invites students to use their newly acquired programming skills to create some amazing pieces of recursive art!}

Here is what the Fall 2023 class has come up with!
( PRINCETON UNIVERSITY


Interdisciplingry Approach
ROBERT SEDGEWICK
https://introcs.cs.princeton.edu

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Fibonacci numbers

Fibonacci numbers. \(0,1,1,2,3,5,8,13,21,34,55,89, \ldots\)



3



5


34


8


55


13


89

Fibonacci numbers: recursive approach

Fibonacci numbers. \(0,1,1,2,3,5,8,13,21,34,55,89, \ldots\)
\[
F_{n}= \begin{cases}0 & \text { if } n=0 \\ 1 & \text { if } n=1 \\ F_{n-1}+F_{n-2} & \text { if } n>1\end{cases}
\]

Goal. Given \(n\), compute \(F_{n}\).

Recursive approach.
- Base cases: \(\quad F_{0}=0, F_{1}=1\).
- Reduction step: \(F_{n}=F_{n-1}+F_{n-2}\).
```

public static long fib(int n) {
if (n == 0) return 0;
if (n == 1) return 1;
return fib(n-1) + fib(n-2);
}

```

Recursion: quiz 4

How long dose it take to compute fib(80)?
A. Much less than 1 second.
B. About 1 second.
C. About 1 minute.
D. About 1 hour.
E. More than 1 hour.
```

public static long fib(int n) {
if (n == 0) return 0;
if (n == 1) return 1;
return fib(n-1) + fib(n-2);
}

```

\section*{Recursion tree for Fibonacci numbers}

Recursion tree.
- One node for each recursive call.
- Label node with return value after children are labelled.


\section*{Exponential waste}

Exponential waste. Same overlapping subproblems are solved repeatedly.
- fib(5) is called 1 time.
- fib(4) is called 2 times.
- fib(3) is called 3 times.
- fib(2) is called 5 times.
- fib(1) is called 8 times.

> number of recursive calls \(\longleftarrow \quad\) are Fibonacci numbers (and grow exponentially)


\section*{Exponential waste dwarfs progress in technology}

Lesson. If you engage in exponential waste, you will not be able to solve a large problem.
\begin{tabular}{|c|c|c|c|c|}
\hline n & recursive calls & \[
\begin{aligned}
& \text { VAX-11 } \\
& (1970 s)
\end{aligned}
\] & \[
\begin{aligned}
& \text { MacBook Pro } \\
& (2020 \mathrm{~s})
\end{aligned}
\] & \\
\hline 30 & 2,692,536 & minute & & \[
=\mathrm{E}
\] \\
\hline 40 & 331,160,280 & hours & & \\
\hline 50 & 40,730,022,146 & weeks & minute & \\
\hline 60 & 5,009,461,563,920 & years & hours & -11/780 \\
\hline 70 & 616,123,042,340,256 & centuries & weeks & \\
\hline 80 & 75,778,124,746,287,810 & millenia & years & \\
\hline 90 & 9,320,093,220,751,060,616 & & centuries &  \\
\hline 100 & 1,146,295,688,027,634,168,200 & & millenia & \\
\hline ! & exponential growth (!) & & : & Macbook Pro (10,000 \(\times\) faster) \\
\hline
\end{tabular}

\footnotetext{
time to compute \(f i b(n)\) using recursive code
}

\section*{Avoiding exponential waste with memoization}

\section*{Memoization.}
- Maintain an array to remember all computed values.
- If value to compute is known, just return it; otherwise, compute it; remember it; and return it.

Impact. Calls fibR(i) at most twice for each i.
```

~/cos126/recursion> java-introcs FibonacciMemo 6
8
~/cos126/recursion> java-introcs FibonacciMemo }8
23416728348467685
instantaneous (!)

```

Design paradigm. This is a simple example of memoization (top-down dynamic programming).

\section*{Summary}

Recursive function. A function that calls itself.

\section*{Why learn recursion?}
- Represents a new mode of thinking.
- Provides a powerful programming paradigm.
- Reveals insight into the nature of computation.

Dynamic programming. A powerful technique to avoid exponential waste. \(\qquad\)

```

// Ackermann function
public static long ack(long m, long n) {
if (m == 0) return n+1;
if (n == 0) return ack(m-1, 1);
return ack(m-1, ack(m, n-1));
}

```
challenge for bored: compute ack(5, 2)

\section*{Credits}
\begin{tabular}{ccc} 
media & source & \multicolumn{1}{c}{ license } \\
\hline Painting Hands & \(\underline{\text { Adobe Stock }}\) & \(\underline{\text { education license }}\) \\
Bugs & \(\underline{\text { Adobe Stock }}\) & \(\underline{\text { education license }}\) \\
Stack Overflow Logo & \(\underline{\text { Stack Overflow }}\) & \\
Problems with Recursion & \(\underline{\text { Zach Weinersmith }}\) & \\
You're Eating Recursion & \(\underline{\text { Safely Endangered }}\) & \\
Collatz Game & \(\underline{\text { Quanta magazine }}\) & \\
File System with Folders & \(\underline{\text { Adobe Stock }}\) & \(\underline{\text { education license }}\) \\
Wooden Towers of Hanoi & Adobe Stock & education license \\
Towers of Hanoi Visualization & \(\underline{\text { Imaginative Animations }}\) & \\
\hline
\end{tabular}

\section*{Credits}
\begin{tabular}{ccc} 
media & source & license \\
\hline Droste Cocoa & \(\underline{\text { Droste }}\) & \\
Recursive Giraffe & \(\underline{\text { Farley Katz }}\) & \\
Circle Limit IV & \(\underline{\text { M.C. Escher }}\) & \\
Recursive Mona Lisa & \(\underline{\text { Mr. Rallentando }}\) & \\
Recursive New York Times & \(\underline{\text { Serkan Ozkaya }}\) & \\
Leonardo Fibonacci & \(\underline{\text { Wikimedia }}\) & public domain \\
VAX 11/780 & \(\underline{\text { Digital Equipment Corporation }}\) & \\
Macbook Pro M1 & \(\underline{\text { Apple }}\) & \\
Menger Sponge & \(\underline{\text { Niabot }}\) & \(\underline{\text { CC BY 3.0 }}\)
\end{tabular}```

