Dynamo / Bayou

Feb 23rd & 24th, 2022

[Adapted from Andrew Or’s]
Some context...

Dynamo and Bayou both offer high availability and weak consistency

Most traditional databases offer strong consistency and low availability
Not suitable for modern applications with super high demands

What are some example applications of each?
Flight ticket booking (HA)
Amazon shopping carts (HA)
Offline edits (HA)
Billing services (SC)
Bank accounts (SC)
Availability is important

Tens of millions of customers at peak times

Tens of millions of shopping cart requests, 3 million checkouts per day

Hundreds of thousands of concurrently active sessions

Strict Service-Level Agreements (SLAs) translate to business value
Dynamo

Fully decentralized, highly available key-value store

Always writeable, resolve conflicts during reads --- Eventual Consistency

API for clients to specify requirements (99.9th percentile)

Departure from RDBMS: simpler functionality, fewer guarantees, runs on commodity hardware (low-end, broadly compatible, non-specialized machines)

(Read the original paper, especially Section 4)
Techniques for achieving availability

**Consistent hashing** for partitioning key space

**Vector clocks** for reconciling conflicts during reads

**Sloppy quorums** for handling temporary failures

**Anti-entropy using Merkle trees** for syncing key-value pairs

**Gossip-based protocol** for membership notifications
Techniques for achieving availability

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Consistent Hashing

Assign each node a random position on the ring

Node owns the preceding key range

For fault tolerance, replicate each key at N successor nodes in the ring

**Virtual nodes**: each physical node gets assigned multiple nodes on the ring (e.g. B, D, F)
Consistent Hashing

*Desirable properties?*

- Uniform distribution of load
- Minimum object movements when nodes join or leave the ring
- Number of virtual nodes can be adjusted for device heterogeneity
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Conflict resolution

Two machines write different values to the same key

*Vector clocks*: list of (node, count) pairs where count is incremented on write

If one vector clock subsumes another, discard older value

Else, return all conflicting values to client
Context contains vector clocks

Dynamo client API is simple:

get(key) (value, context)

put(key, value, context)

Common pattern: put after get
<table>
<thead>
<tr>
<th>Product Name</th>
<th>Description</th>
<th>Stock Status</th>
<th>Availability</th>
<th>Gift Option</th>
<th>Delete</th>
<th>Save for later</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Cat-Opoly</strong> by LatefortheSky</td>
<td>In Stock</td>
<td><strong>Prime</strong></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td><strong>Fancy Feast Wet Cat Food, Grilled, Seafood Feast Variety Pack, 3-Ounce Can, Pack of 24</strong> by Purina Fancy Feast</td>
<td>In stock. Usually ships within 3 to 4 days.</td>
<td>Shipped from: Connect Buy</td>
<td></td>
<td>Gift options not available. Learn more</td>
<td></td>
<td></td>
</tr>
<tr>
<td><strong>Furhaven Orthopedic Mattress Pet Bed, Large, Chocolate, for Dogs and Cats</strong> by Furhaven Pet</td>
<td>In Stock</td>
<td><strong>Prime</strong></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Product</td>
<td>by</td>
<td>Status</td>
<td>Shipping Info</td>
<td>Gift Options</td>
<td>Actions</td>
<td></td>
</tr>
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<td>----------------------------------------------</td>
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<td>-------------------------------------------------------------------------------</td>
<td>--------------------</td>
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<td></td>
</tr>
<tr>
<td>Cat-Opoly</td>
<td>LatefortheSky</td>
<td>In Stock</td>
<td>- Prime</td>
<td>- This is a gift</td>
<td>Delete, Save for later</td>
<td></td>
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<td></td>
</tr>
<tr>
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<td>by Furhaven Pet</td>
<td>In Stock</td>
<td>Prime</td>
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<td></td>
</tr>
<tr>
<td>Chocolate, for Dogs and Cats</td>
<td></td>
<td></td>
<td></td>
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Sloppy Quorums

Write to N nodes, return success when W < N nodes respond

Read from N nodes, return value(s) from R < N nodes

Typically, W+R > N means at least one writer and one reader overlap, so values are consistent

Sloppy here means skip nodes that have failed, such that even if W+R > N, the readers and writers may not overlap = not consistent!
Sloppy Quorums

Example:

Typical values are $N = 3$, $W = R = 2$

Nodes C and D have failed, so key $k$ is written to E and F instead

Nodes C and D recover, and now client tries to read from C and D = stale value
Hinted Handoff

“Hint” refers to the node the data originally belongs to

Example:

Nodes E and F remember they are writing on behalf of C and D

As soon as C and D recovers, E and F transfer their values for $k$ to C and D
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Anti-entropy using Merkle trees

Goal: minimize durability loss from above techniques

Nodes responsible for the same key spaces exchange Merkle trees

Find differences quickly while exchanging little information
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Membership notification

Gossip-based protocol to propagate membership changes

Each node learns the key spaces handled by all other nodes

**Result:** zero-hop distributed hash table (DHT)

*Clearly not infinitely scalable*, but storage requirement not a problem in practice
Bayou

What is it?
- Weakly consistent, replicated storage system

Goals:
- Maximize availability, *support offline collaboration*
- Minimize network communication
- Agree on all values (eventually)
Bayou Writes

Legend
Commit Timestamp:Write Timestamp:Write Server

Primary

Versions

P: 0
A: 0
B: 0

W(X, 4)  Client 1

A

Versions

P: 0
A: 0
B: 0

B

Versions

P: 0
A: 0
B: 0
Bayou Writes

Legend
Commit Timestamp: Write Timestamp: Write Server

Client 1

Primary Versions

∞:1:P  w(x,4)
P: 1
A: 0
B: 0

A Versions

P: 0
A: 0
B: 0

B Versions

P: 0
A: 0
B: 0
Bayou Writes

Legend
Commit Timestamp:Write Timestamp:Write Server

Client 1
W(Y, 8)

Client 2
W(X, 3)
Bayou Writes

Legend
Commit Timestamp:Write Timestamp:Write Server

Client 1
W(Z, 8)

Client 2
W(Y, 4)

A
\[\infty:7: A \text{ W}(X, 3)\]
\begin{align*}
P &: 0 \\
A &: 7 \\
B &: 0
\end{align*}

Versions

Primary
\begin{align*}
\infty:1: P & \text{ W}(X, 4) \\
\infty:7: P & \text{ W}(Y, 8)
\end{align*}

Versions

B
\begin{align*}
P &: 0 \\
A &: 0 \\
B &: 0
\end{align*}
Bayou Writes

Legend
Commit Timestamp:Write Timestamp:Write Server

Primary Versions
∞:1:P \text{ w}(X,4)
∞:7:P \text{ w}(Y,8)

A Versions
∞:7:A \text{ w}(X,3)
∞:12:A \text{ w}(Y,4)

B Versions
∞:5:B \text{ w}(Z,8)

P: 7
A: 0
B: 0

P: 0
A: 12
B: 0

P: 0
A: 0
B: 5
Bayou Anti-Entropy

Anti-Entropy Session
A & B
Bayou Anti-Entropy
Bayou Commit

Primary commits its entries
Bayou Write

Write after anti-entropy session
Write timestamp = max(clock, max(TS)+1)
Bayou Anti-Entropy

Anti-Entropy Session
P & B
Bayou Anti-Entropy

Anti-Entropy Session
P & B
Primary respects causality
Bayou Commit

Primary commits Its entries
Bayou

After a number of commits and anti-entropy sessions (without further writes)
Bayou and Dynamo similarities

Anti-entropy to achieve eventual consistency

Exchange vector clocks to determine order of operations

Expose conflict resolution to application

High availability!