

# COS 461: Computer Networks

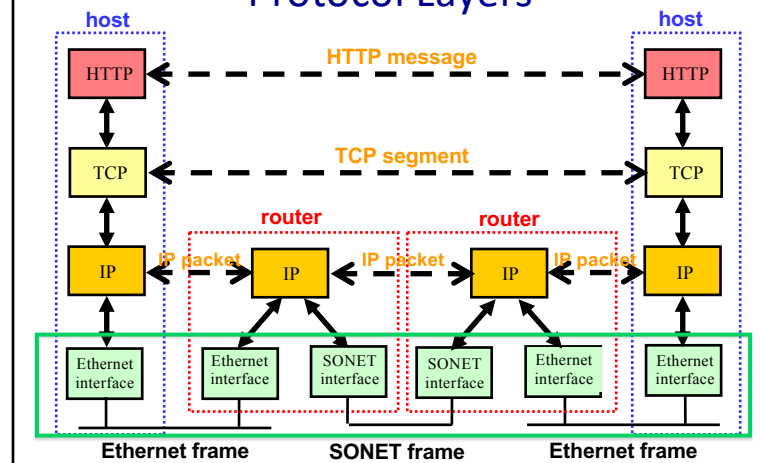
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Spring 2020

Lectures: MW 1:30-2:50pm in Richardson 002

<http://www.cs.princeton.edu/courses/archive/spring20/cos461/>

## Protocol Layers



Link = Medium + Adapters

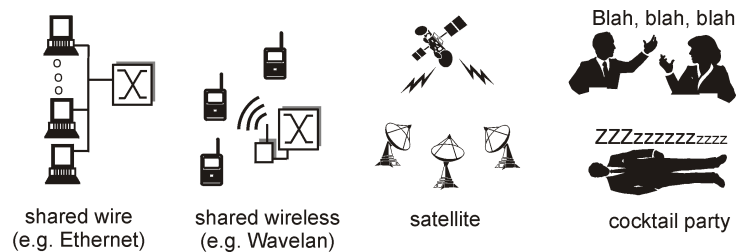
## What is a Link?

Communication Medium

Network Adapter

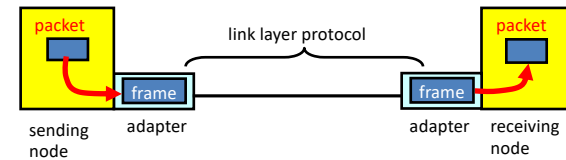


## Broadcast Links: Shared Media



5

## Adaptors Communicating



- **Sending side**
  - Encapsulates packet in a frame
  - Adds error checking bits, flow control, etc.
- **Receiving side**
  - Looks for errors, flow control, etc.
  - Extracts datagram and passes to receiving node

6

## Link-Layer Services

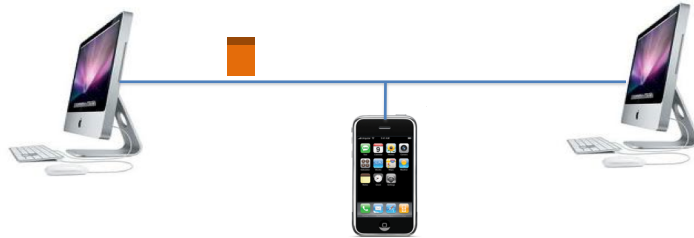
- **Encoding**
  - Represent the 0s and 1s
- **Framing**
  - Encapsulate packet into frame, adding header/trailer
- **Error detection**
  - Receiver detecting errors with checksums
- **Error correction**
  - Receiver optionally correcting errors
- **Flow control**
  - Pacing between sending and receiving nodes

7

## Addresses

8

## Medium Access Control Address



- **Identify the sending and receiving adapter**
  - Unique identifier for each network adapter
  - Identifies the intended receiver(s) of the frame
  - ... and the sender who sent the frame

9

## Medium Access Control Address

- **MAC address (e.g., 00-15-C5-49-04-A9)**
  - Numerical address used within a link
  - Unique, hard-coded in the adapter when it is built
  - Flat name space of 48 bits
- **Hierarchical allocation: Global uniqueness!**
  - **Blocks**: assigned to vendors (e.g., Dell) by the IEEE
  - **Adapters**: assigned by the vendor from its block
- **Broadcast address (i.e., FF-FF-FF-FF-FF-FF)**
  - Send the frame to *all* adapters

10

## As an Aside: Promiscuous Mode

- **Normal adapter: receives frames sent to**
  - The local MAC address
  - Broadcast address FF-FF-FF-FF-FF-FF
- **Promiscuous mode**
  - Receive *everything*, independent of destination MAC
- **Useful for packet sniffing**
  - Network monitoring
  - E.g., wireshark, tcpdump



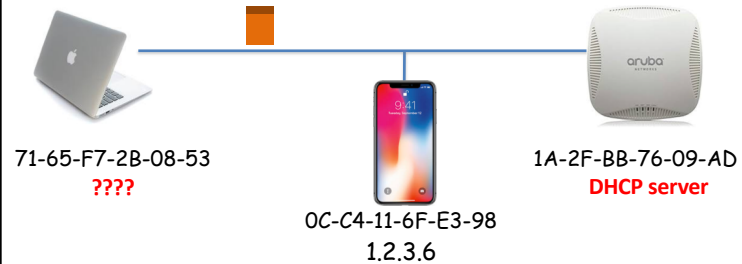
11

## Why Not Just Use IP Addresses?

- **Links can support *any* network protocol**
  - Not just for IP (e.g., IPX, Appletalk, X.25, ...)
  - Different addresses on different kinds of links
- **An adapter may move to a new location**
  - So, cannot simply assign a static IP address
  - Instead, must reconfigure the adapter's IP address
- **Must identify the adapter during bootstrap**
  - Need to talk to the adapter to assign it an IP address

12

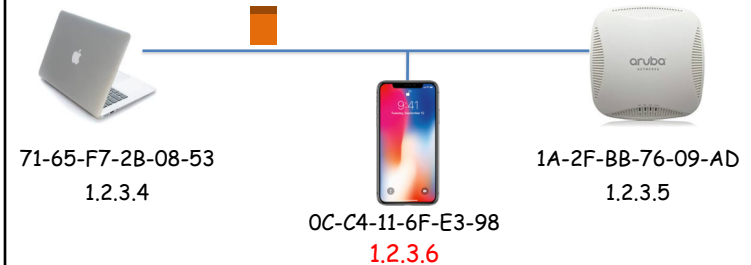
## Who Am I: Acquiring an IP Address



- **Dynamic Host Configuration Protocol (DHCP)**
  - Broadcast “I need an IP address, please!”
  - Response “You can have IP address 1.2.3.4.”

13

## Who Are You: Discovering the Receiver



- **Address Resolution Protocol (ARP)**
  - Broadcast “who has IP address 1.2.3.6?”
  - Response “OC-C4-11-6F-E3-98 has 1.2.3.6!”

14

## Sharing the Medium

15

## Collisions

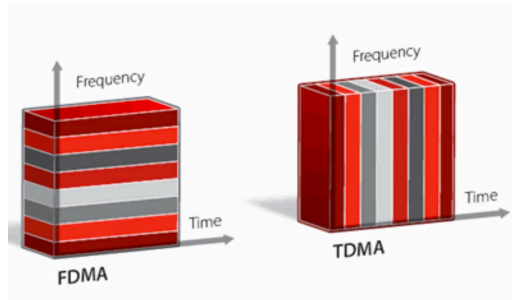


- **Single shared broadcast channel**
  - Avoid having multiple nodes speaking at once
  - Otherwise, collisions lead to garbled data

16

## Multi-Access Protocol

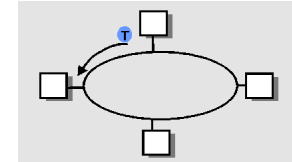
- **Divide the channel into pieces**
  - in frequency vs. in time



17

## Multi-Access Protocol

- **Take turns**
  - Do not transmit w/o token
  - With token, transmit for up to some max time/length
  - Pass token to right
- **Punt**
  - Let collisions happen
  - ... and detect and recover from them



18

## Like Human Conversation...

- **Carrier sense**
  - Listen before speaking
  - ...and don't interrupt!
- **Collision detection**
  - Detect simultaneous talking
  - ... and shut up!
- **Random access**
  - Wait for a random period of time
  - ... before trying to talk again!

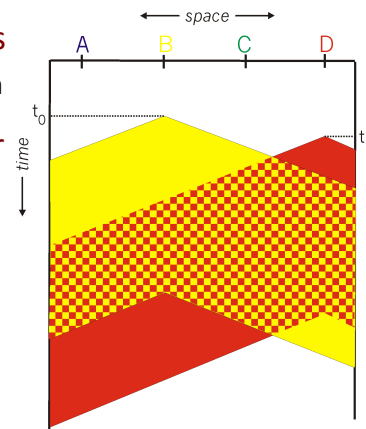


*Please  
Wait...*

19

## Carrier Sense Multiple Access

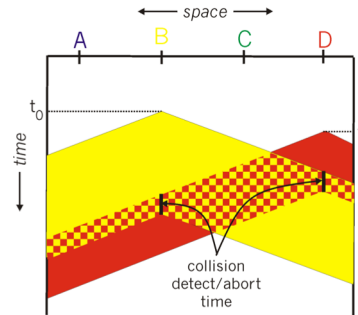
- **Listen for other senders**
  - Then transmit your data
- **Collisions can still occur**
  - Propagation delay
  - Wasted transmission



20

## CSMA/CD Collision Detection

- **Detect collision**
  - Abort transmission
  - Jam the link
- **Wait random time**
  - Transmit again
- **Hard in wireless**
  - Must receive data while transmitting



21

## Comparing the Three Approaches

- **Channel partitioning is**
  - (a) Efficient/fair at high load, inefficient at low load
  - (b) Inefficient at high load, efficient/fair at low load
- **“Taking turns”**
  - (a) Inefficient at high load
  - (b) Efficient at all loads
  - (c) Robust to failures
- **Random access**
  - (a) Inefficient at low load
  - (b) Efficient at all load
  - (c) Robust to failures

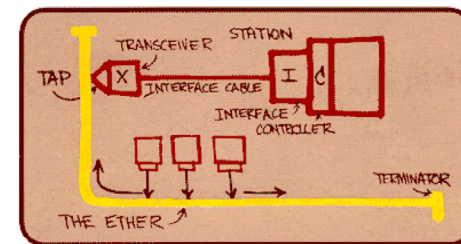
22

## Ethernet

24

## Ethernet

- Dominant wired LAN technology
- First widely used LAN technology
- Kept up with speed race: 10 Mbps – 40 Gbps



Metcalfe's  
Ethernet  
sketch

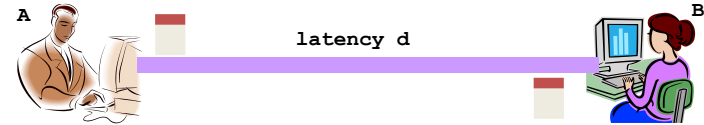
25

## Ethernet Uses CSMA/CD

- **Carrier Sense: wait for link to be idle**
  - Channel idle: start transmitting
  - Channel busy: wait until idle
- **Collision Detection: listen while transmitting**
  - No collision: transmission is complete
  - Collision: abort transmission, and send jam signal
- **Random Access: exponential back-off**
  - After collision, wait random time before trying again
  - After  $m^{\text{th}}$  collision, choose  $K$  randomly from  $\{0, \dots, 2^m - 1\}$
  - ... and wait for  $K * 512$  bit times before trying again

26

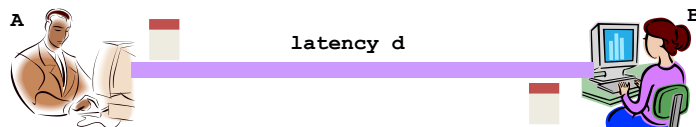
## Limitations on Ethernet Length



- **Latency depends on physical length of link**
  - Time to propagate a packet from one end to other
- **Suppose A sends a packet at time  $t$** 
  - And B sees an idle line at a time just before  $t+d$
  - ... so B happily starts transmitting a packet
- **B detects a collision, and sends jamming signal**
  - But A doesn't see collision till  $t+2d$

27

## Limitations on Ethernet Length

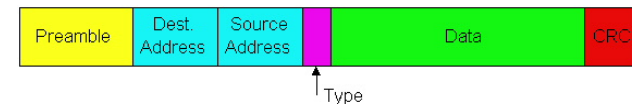


- **A needs to wait for time  $2d$  to detect collision**
  - So, A should keep transmitting during this period
  - ... and keep an eye out for a possible collision
- **Imposes restrictions on Ethernet**
  - Maximum length of the wire: 2500 meters
  - Minimum length of the packet: 512 bits (64 bytes)

28

## Ethernet Frame Structure

- **Sending adapter encapsulates packet in frame**

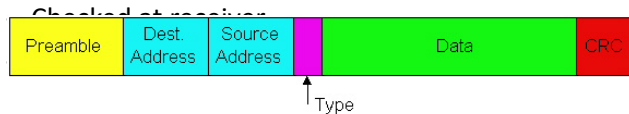


- **Preamble: synchronization**
  - Seven bytes with pattern 10101010, followed by one byte with pattern 10101011
  - Used to synchronize receiver, sender clock rates

29

## Ethernet Frame Structure

- **Addresses: source and destination MAC addresses**
  - Adaptor passes frame to network-level protocol
    - If destination is local MAC address or broadcast address
  - Otherwise, adaptor discards frame
- **Type: indicates the higher layer protocol**
  - Usually IP
  - But also Novell IPX, AppleTalk, ...
- **CRC: cyclic redundancy check**



30

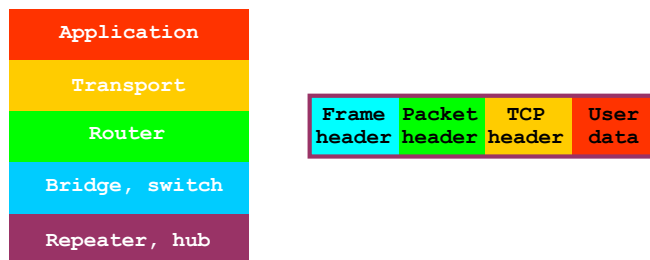
## Unreliable, Connectionless Service

- **Connectionless**
  - No handshaking between send and receive adapter
- **Unreliable**
  - Receiving adapter doesn't send ACKs or NACKs
  - Packets passed to network layer can have gaps
  - Gaps can be filled by transport protocol (e.g., TCP)
  - Otherwise, the application will see the gaps

31

## Summary: Multiple Layers

- **Different devices switch different things**
  - Network layer: packets (routers)
  - Link layer: frames (bridges and switches)
  - Physical layer: electrical signals (repeaters and hubs)



32

## Conclusion

- **Links**
  - Connect two or more network adapters
  - ... each with a unique address
  - ... over a shared communication medium
- **Coming next**
  - Network layer (IP)

33