

### **Basic architecture**

Clients submit applications to the cluster manager Cluster manager assigns cluster resources to applications Each Worker launches containers for each application Driver containers run main method of user program Executor containers run actual computation

*Examples* of cluster manager: YARN, Mesos *Examples* of computing frameworks: Hadoop MapReduce, Spark

### Two levels of scheduling

Cluster-level: Cluster manager assigns resources to applications Application-level: Driver assigns *tasks* to run on executors

A task is a unit of execution that operates on one *partition* 

#### Some advantages:

Applications need not be concerned with resource fairness Cluster manager need not be concerned with individual tasks Easy to implement priorities and preemption

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# Application: Word count

Hello my love. I love you, my dear. Goodbye.

hello: 1, my: 2, love: 2, i: 1, dear: 1, goodbye: 1

# Case Study: MapReduce

#### (Data-parallel programming at scale)

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# Application: Word count

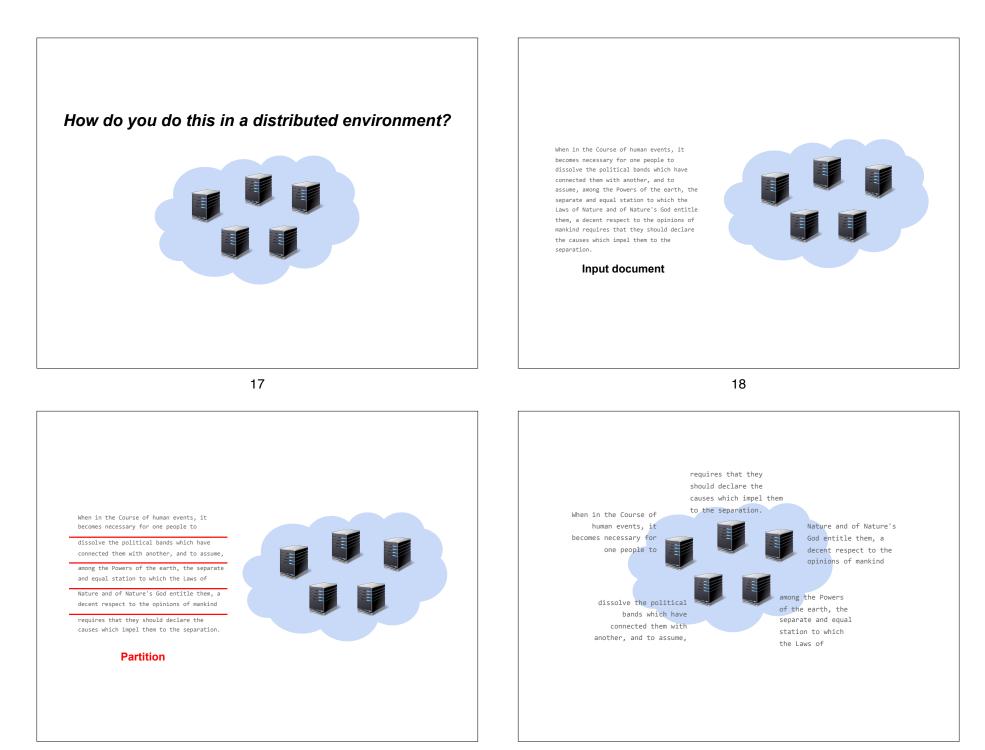
Locally: tokenize and put words in a hash map

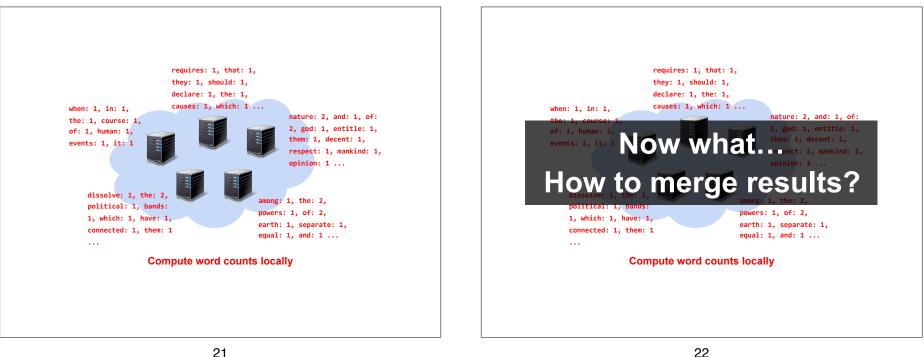
#### How do you parallelize this?

Split document by half

Build two hash maps, one for each half

Merge the two hash maps (by key)





### Merging results computed locally

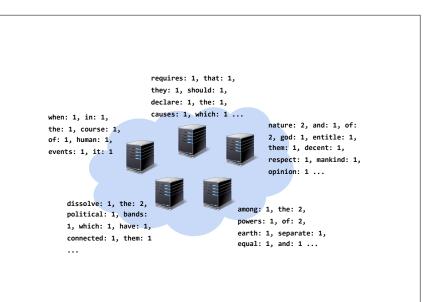
Several options

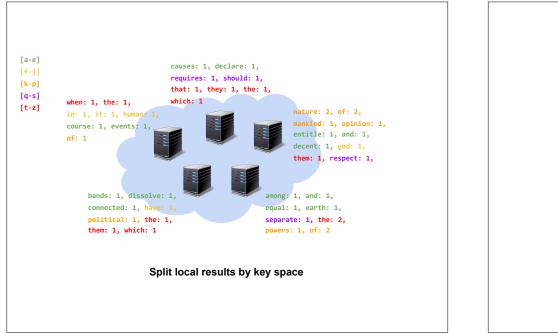
Don't merge — requires additional computation for correct results

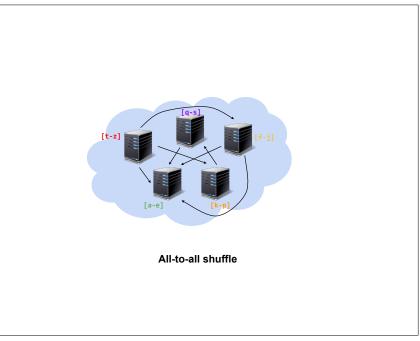
Send everything to one node - what if data is too big? Too slow...

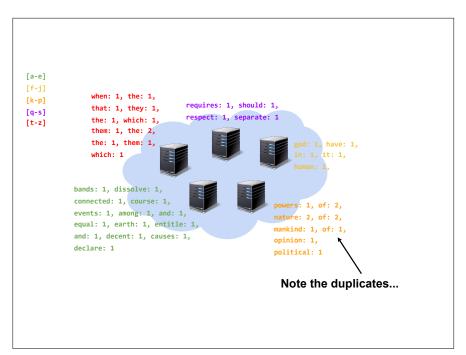
Partition key space among nodes in cluster (e.g. [a-e], [f-j], [k-p] ...)

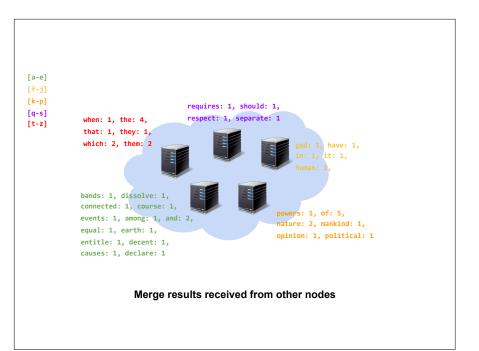
- 1. Assign a key space to each node
- 2. Partition local results by the key spaces
- 3. Fetch and merge results that correspond to the node's key space











#### MapReduce

Partition dataset into many chunks

Map stage: Each node processes one or more chunks locally

**Reduce stage:** Each node fetches and merges partial results from all other nodes

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### MapReduce: Word count

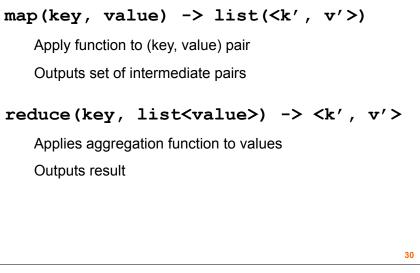
#### map(key, value):

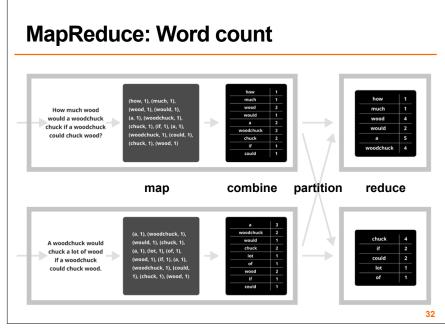
// key = document name
// value = document contents
for each word w in value:
 emit (w, 1)

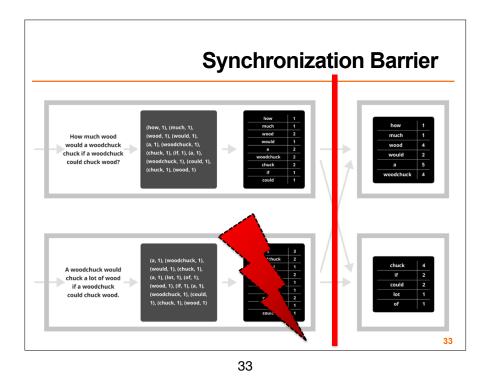
#### reduce(key, values):

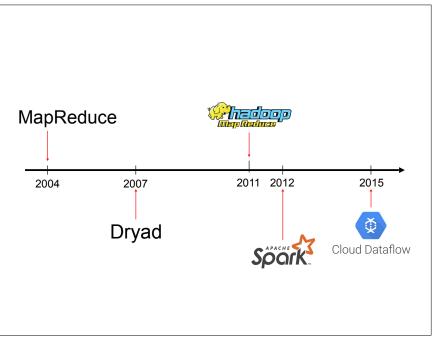
// key = the word
// values = number of occurrences of that word
count = sum(values)
emit (key, count)

### **MapReduce Interface**









### **Brainstorm: Top K**

Top K is the problem of finding the largest K values from a set of numbers

How would you express this as a distributed application?

In particular, what would the map and reduce phases look like?

Hint: use a heap...

# Brainstorm: Top K

Assuming that a set of K integers fit in memory...

#### Key idea...

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Map phase: everyone maintains a heap of K elements Reduce phase: merge the heaps until you're left with one

### Brainstorm: Top K

Problem: What are the keys and values here?

No notion of key here, just assign the same key to all the values (e.g. key = 1) Map task 1: [10, 5, 3, 700, 18, 4]  $\rightarrow$  (1, heap(700, 18, 10)) Map task 2: [16, 4, 523, 100, 88]  $\rightarrow$  (1, heap(523, 100, 88)) Map task 3: [3, 3, 3, 3, 300, 3]  $\rightarrow$  (1, heap(300, 3, 3)) Map task 4: [8, 15, 20015, 89]  $\rightarrow$  (1, heap(20015, 89, 15)) Then all the heaps will go to a single reducer responsible for the key 1 This works, but clearly not scalable...

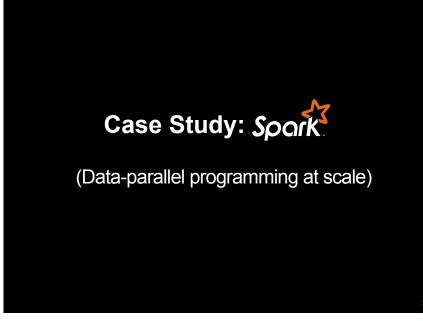
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### **Brainstorm: Top K**

Idea: Use X different keys to balance load (e.g. X = 2 here) Map task 1: [10, 5, 3, 700, 18, 4]  $\rightarrow$  (1, heap(700, 18, 10)) Map task 2: [16, 4, 523, 100, 88]  $\rightarrow$  (1, heap(523, 100, 88)) Map task 3: [3, 3, 3, 3, 300, 3]  $\rightarrow$  (2, heap(300, 3, 3)) Map task 4: [8, 15, 20015, 89]  $\rightarrow$  (2, heap(20015, 89, 15))

Then all the heaps will (hopefully) go to X different reducers Rinse and repeat (*what's the runtime complexity?*)

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# What is **Spork**?

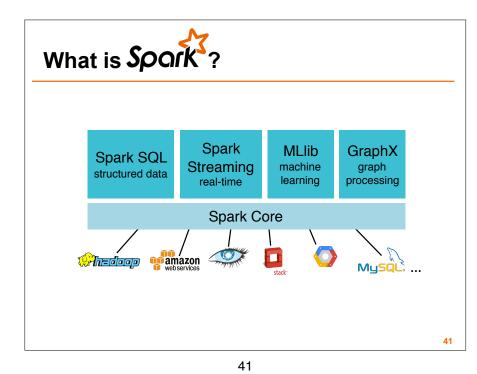
General distributed data execution engine

#### Key optimizations

General computation graphs (pipelining, lazy execution) In-memory data sharing (caching) Fine-grained fault tolerance

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#1 most active big data project

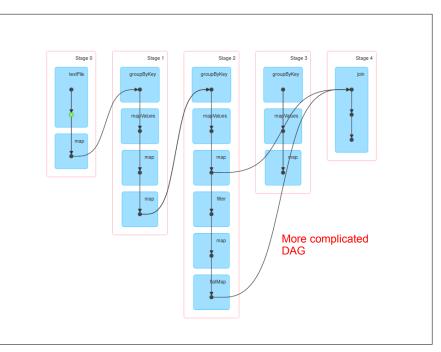


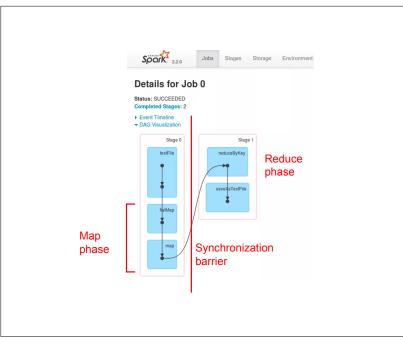
### Spark computational model

Most computation can be expressed in terms of two phases Map phase defines how each machine processes its individual partition Reduce phase defines how to merge map outputs from previous phase

Spark expresses computation as a DAG of maps and reduces

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### **Spark: Word count**

Transformations express how to process a dataset

Actions express how to turn a transformed dataset into results

```
sc.textFile("declaration-of-independence.txt")
```

```
.flatMap { line => line.split(" ") }
.map { word => (word, 1) }
.reduceByKey { case (counts1, counts2) => counts1 + counts2 }
.collect()
```

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### **In-memory caching**

Store intermediate results in memory to bypass disk access

Important for iterative workloads (e.g. machine learning!)

#### Example:

```
val cached = data.map { ... }.filter { ... }.cache()
(1 to 100).foreach { i =>
      cached.reduceByKey { ... }.saveAsTextFile(...)
}
```

# **Pipelining + Lazy execution**

Transformations can be pipelined until we hit

A synchronization barrier (e.g. reduce), or

An action

#### Example:

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data.map { ... }.filter { ... }.flatMap { ... }.groupByKey().count()

These three operations can all be run in the same task

This allows lazy execution; we don't need to eagerly execute map

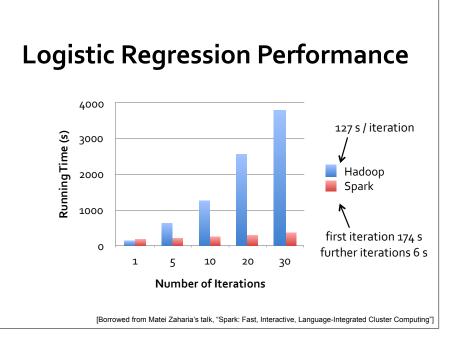
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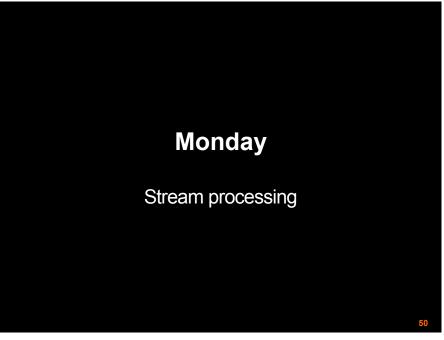
# **Reusing map outputs**

Reusing map outputs allows Spark to avoid redoing map stages Along with caching, this makes iterative workloads much faster

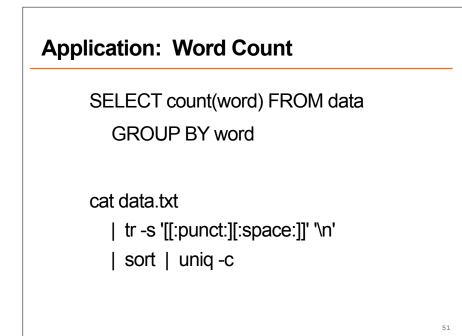
#### Example:

```
val transformed = data.map { ... }.reduceByKey { ... }
transformed.collect()
transformed.collect() // does not run map phase again
```





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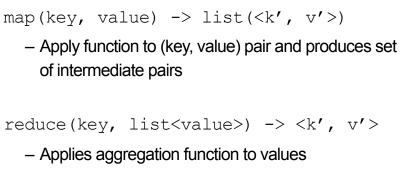
# Using partial aggregation

- 1. Compute word counts from individual files
- 2. Then merge intermediate output
- 3. Compute word count on merged outputs

### Using partial aggregation

- 1. In parallel, send to worker:
  - Compute word counts from individual files
  - Collect result, wait until all finished
- 2. Then merge intermediate output
- 3. Compute word count on merged intermediates

### MapReduce: Programming Interface



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Outputs result

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### MapReduce: Programming Interface

```
map(key, value):
   for each word w in value:
      EmitIntermediate(w, "1");
reduce(key, list(values):
   int result = 0;
```

for each v in values:

```
result += ParseInt(v);
```

```
Emit(AsString(result));
```

MapReduce: Optimizations

combine(list<key, value>) -> list<k,v>

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Perform partial aggregation on mapper node:

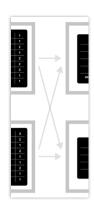
<the, 1>, <the, 1>, <the, 1>  $\rightarrow$  <the, 3>

- reduce() should be commutative and associative

partition(key, int) -> int

- Need to aggregate intermediate vals with same key
- Given n partitions, map key to partition  $0 \le i < n$
- Typically via hash(key) mod n

### Fault Tolerance in MapReduce



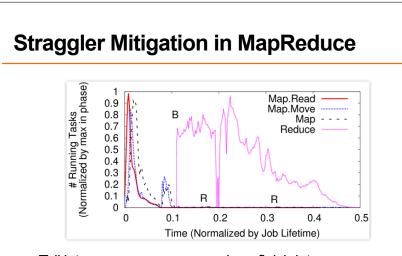
- Map worker writes intermediate output to local disk, separated by partitioning. Once completed, tells master node.
- Reduce worker told of location of map task outputs, pulls their partition's data from each mapper, execute function across data
- Note:
  - "All-to-all" shuffle b/w mappers and reducers
  - Written to disk ("materialized") b/w each stage

### Fault Tolerance in MapReduce

- Master node monitors state of system
  - If master failures, job aborts and client notified
- · Map worker failure
  - Both in-progress/completed tasks marked as idle
  - Reduce workers notified when map task is re-executed on another map worker
- Reducer worker failure
  - In-progress tasks are reset to idle (and re-executed)
  - Completed tasks had been written to global file system

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- Tail latency means some workers finish late
- For slow map tasks, execute in parallel on second map worker as "backup", race to complete task

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