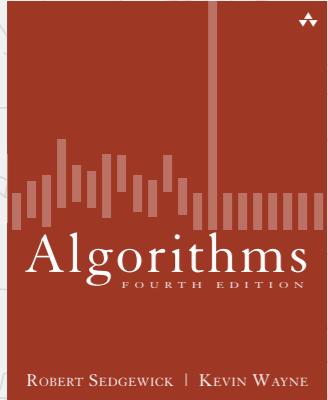


Algorithms

ROBERT SEDGEWICK | KEVIN WAYNE



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<http://algs4.cs.princeton.edu>

2.1 ELEMENTARY SORTS

- ▶ rules of the game
- ▶ selection sort
- ▶ insertion sort
- ▶ shuffling
- ▶ comparators

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- ▶ comparators

Sorting problem

Ex. Student records in a university.

Chen	3	A	(991) 878-4944	308 Blair
Rohde	2	A	(232) 343-5555	343 Forbes
Gazsi	4	B	(800) 867-5309	101 Brown
Furia	1	A	(766) 093-9873	101 Brown
Kanaga	3	B	(898) 122-9643	22 Brown
Andrews	3	A	(664) 480-0023	097 Little
Battle	4	C	(874) 088-1212	121 Whitman

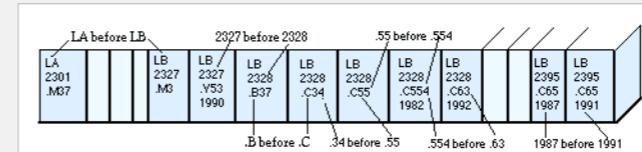
item →

key →

Sort. Rearrange array of N items in ascending order by key.

Andrews	3	A	(664) 480-0023	097 Little
Battle	4	C	(874) 088-1212	121 Whitman
Chen	3	A	(991) 878-4944	308 Blair
Furia	1	A	(766) 093-9873	101 Brown
Gazsi	4	B	(800) 867-5309	101 Brown
Kanaga	3	B	(898) 122-9643	22 Brown
Rohde	2	A	(232) 343-5555	343 Forbes

Sorting applications



Library of Congress numbers



FedEx packages



contacts



playing cards



Hogwarts houses

Sample sort client 1

Goal. Sort **any** type of data.

Ex 1. Sort random real numbers in ascending order.

seems artificial (stay tuned for an application)

```
public class Experiment
{
    public static void main(String[] args)
    {
        int N = Integer.parseInt(args[0]);
        Double[] a = new Double[N];
        for (int i = 0; i < N; i++)
            a[i] = StdRandom.uniform();
        Insertion.sort(a);
        for (int i = 0; i < N; i++)
            StdOut.println(a[i]);
    }
}
```

```
% java Experiment 10
0.08614716385210452
0.09054270895414829
0.10708746304898642
0.21166190071646818
0.363292849257276
0.460954145685913
0.5340026311350087
0.7216129793703496
0.9003500354411443
0.9293994908845686
```

5

Sample sort client 2

Goal. Sort **any** type of data.

Ex 2. Sort strings in alphabetical order.

```
public class StringSorter
{
    public static void main(String[] args)
    {
        String[] a = StdIn.readAllStrings();
        Insertion.sort(a);
        for (int i = 0; i < a.length; i++)
            StdOut.println(a[i]);
    }
}

% more words3.txt
bed bug dad yet zoo ... all bad yes

% java StringSorter < words3.txt
all bad bed bug dad ... yes yet zoo
[suppressing newlines]
```

6

Sample sort client 3

Goal. Sort **any** type of data.
Ex 3. Sort the files in a given directory by filename.

```
import java.io.File;

public class FileSorter
{
    public static void main(String[] args)
    {
        File directory = new File(args[0]);
        File[] files = directory.listFiles();
        Insertion.sort(files);
        for (int i = 0; i < files.length; i++)
            StdOut.println(files[i].getName());
    }
}
```

```
% java FileSorter .
Insertion.class
Insertion.java
InsertionX.class
InsertionX.java
Selection.class
Selection.java
Shell.class
Shell.java
ShellX.class
ShellX.java
```

7

Total order

Goal. Sort **any** type of data (for which sorting is well defined).

A **total order** is a binary relation \leq that satisfies:

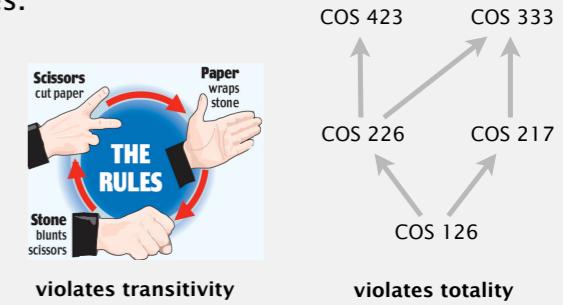
- Antisymmetry: if both $v \leq w$ and $w \leq v$, then $v = w$.
- Transitivity: if both $v \leq w$ and $w \leq x$, then $v \leq x$.
- Totality: either $v \leq w$ or $w \leq v$ or both.

Ex.

- Standard order for natural and real numbers.
- Chronological order for dates or times.
- Lexicographic order for strings.

Not transitive. Ro-sham-bo.

Not total. PU course prerequisites.



8

Callbacks

Goal. Sort **any** type of data (for which sorting is well defined).

Q. How can sort() compare data of type Double, String, and java.io.File without hardwiring in type-specific information.

Callback = reference to executable code.

- Client passes array of objects to sort() function.
- The sort() function calls object's compareTo() method as needed.

Implementing callbacks.

- Java: interfaces.
- C: function pointers.
- C++: class-type functors.
- C#: delegates.
- Python, Perl, ML, Javascript: first-class functions.

Callbacks: Java interfaces

Interface. Specifies a set of methods that a concrete class can provide.

```
public interface Comparable<Item>
{
    public int compareTo(Item that);
}
```

contract: one method with this signature and prescribed behavior

Concrete class. Can provide the set of methods in the interface.

```
public class String implements Comparable<String>
{
    ...
    public int compareTo(String that)
    {
        ...
    }
}
```

class promises to honor the contract

class honors the contract

"polymorphism"

Impact.

- You can treat any String object as an object of type Comparable.
- On a Comparable object, you can invoke (only) the compareTo() method.
- Enables callbacks.

Callbacks: roadmap

client (StringSorter.java)

```
public class StringSorter
{
    public static void main(String[] args)
    {
        String[] a = StdIn.readAllStrings();
        Insertion.sort(a);
        for (int i = 0; i < a.length; i++)
            StdOut.println(a[i]);
    }
}
```

java.lang.Comparable interface

```
public interface Comparable<Item>
{
    public int compareTo(Item that);
}
```

sort implementation (Insertion.java)

```
public static void sort(Comparable[] a)
{
    int N = a.length;
    for (int i = 0; i < N; i++)
        for (int j = i; j > 0; j--)
            if (a[j].compareTo(a[j-1]) < 0)
                exch(a, j, j-1);
            else break;
}
```

data type implementation (String.java)

```
public class String
implements Comparable<String>
{
    ...
    public int compareTo(String that)
    {
        ...
    }
}
```

key point: no dependence on String data type

Elementary sorts: quiz 1

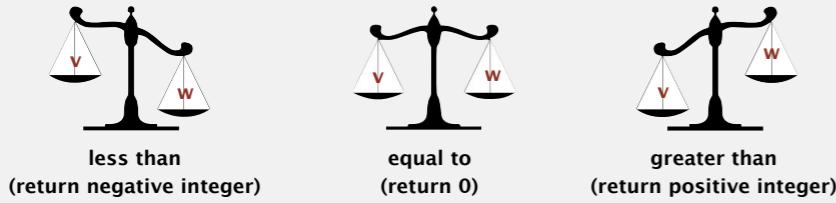
Suppose that the Java architects leave out `implements Comparable<String>` in the class declaration for String. What would be the effect?

- String.java won't compile.
- StringSorter.java won't compile.
- Insertion.java won't compile.
- Insertion.java will throw a run-time exception.
- I don't know.

java.lang.Comparable API

Implement `compareTo()` so that `v.compareTo(w)`

- Defines a total order.
- Returns a negative integer, zero, or positive integer if v is less than, equal to, or greater than w , respectively.
- Throws an exception if incompatible types (or either is null).



Built-in comparable types. Integer, Double, String, Date, File, ...

User-defined comparable types. Implement the Comparable interface.

Implementing the Comparable interface

Date data type. Simplified version of `java.util.Date`.

```
public class Date implements Comparable<Date>
{
    private final int month, day, year;

    public Date(int m, int d, int y)
    {
        month = m;
        day   = d;
        year  = y;
    }

    public int compareTo(Date that)
    {
        if (this.year < that.year) return -1;
        if (this.year > that.year) return +1;
        if (this.month < that.month) return -1;
        if (this.month > that.month) return +1;
        if (this.day < that.day) return -1;
        if (this.day > that.day) return +1;
        return 0;
    }
}
```

only compare dates to other dates

<http://algs4.cs.princeton.edu/12oop/Date.java.html>

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Algorithms

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<http://algs4.cs.princeton.edu>

rules of the game

selection sort

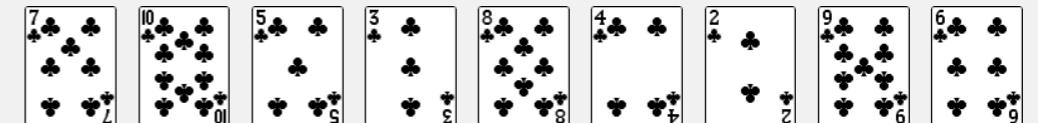
insertion sort

shuffling

comparators

Selection sort demo

- In iteration i , find index \min of smallest remaining entry.
- Swap $a[i]$ and $a[\min]$.



initial



16

Selection sort

Algorithm. ↑ scans from left to right.

Invariants.

- Entries the left of ↑ (including ↑) fixed and in ascending order.
- No entry to right of ↑ is smaller than any entry to the left of ↑.



Selection sort inner loop

To maintain algorithm invariants:

- Move the pointer to the right.

```
i++;
```



- Identify index of minimum entry on right.

```
int min = i;
for (int j = i+1; j < N; j++)
    if (less(a[j], a[min]))
        min = j;
```



- Exchange into position.

```
exch(a, i, min);
```



Two useful sorting abstractions

Helper functions. Refer to data only through compares and exchanges.

Less. Is item v less than w?

```
private static boolean less(Comparable v, Comparable w)
{   return v.compareTo(w) < 0; }
```

Exchange. Swap item in array a[] at index i with the one at index j.

```
private static void exch(Object[] a, int i, int j)
{
    Object swap = a[i];
    a[i] = a[j];
    a[j] = swap;
}
```

Selection sort: Java implementation

```
public class Selection
{
    public static void sort(Comparable[] a)
    {
        int N = a.length;
        for (int i = 0; i < N; i++)
        {
            int min = i;
            for (int j = i+1; j < N; j++)
                if (less(a[j], a[min]))
                    min = j;
            exch(a, i, min);
        }
    }

    private static boolean less(Comparable v, Comparable w)
    { /* see previous slide */ }

    private static void exch(Object[] a, int i, int j)
    { /* see previous slide */ }
}
```

Generic methods

Oops. The compiler complains.

```
% javac Selection.java
Note: Selection.java uses unchecked or unsafe operations.
Note: Recompile with -Xlint:unchecked for details.
```

```
% javac -Xlint:unchecked Selection.java
Selection.java:83: warning: [unchecked] unchecked call to
compareTo(T) as a member of the raw type java.lang.Comparable
    return (v.compareTo(w) < 0);
           ^
1 warning
```

Q. How to fix?

Generic methods

Pedantic (type-safe) version. Compiles cleanly.

```
generic type variable
(type inferred from argument; must be Comparable)

public class SelectionPedantic
{
    public static <Key extends Comparable<Key>> void sort(Key[] a)
    { /* as before */ }

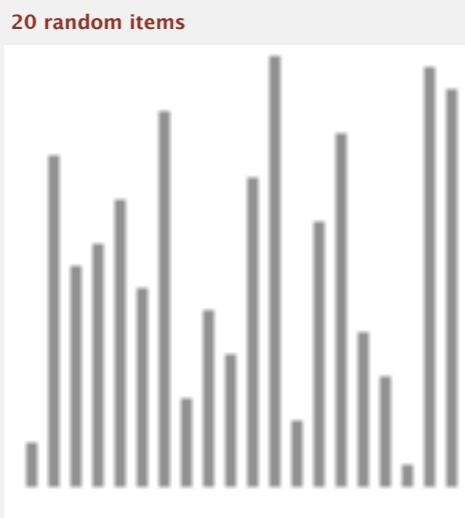
    private static <Key extends Comparable<Key>> boolean less(Key v, Key w)
    { /* as before */ }

    private static Object void exch(Object[] a, int i, int j)
    { /* as before */ }
}
```

<http://algs4.cs.princeton.edu/21elementary/SelectionPedantic.java.html>

Remark. Use type-safe version in system code (but not in lecture).

Selection sort: animations



▲ algorithm position
— in final order
— not in final order

<http://www.sorting-algorithms.com/selection-sort>

Selection sort: animations



▲ algorithm position
— in final order
— not in final order

<http://www.sorting-algorithms.com/selection-sort>

Elementary sorts: quiz 1

How many compares does selection sort make to sort an array of N keys?

- A. $\sim N$
- B. $\sim 1/4 N^2$
- C. $\sim 1/2 N^2$
- D. $\sim N^2$
- E. I don't know.

Selection sort: mathematical analysis

Proposition. Selection sort uses $(N-1) + (N-2) + \dots + 1 + 0 \sim N^2/2$ compares and N exchanges to sort any array of N items.

i	min	0	1	2	3	4	5	6	7	8	9	10
0	6	S	O	R	T	E	X	A	M	P	L	E
1	4	S	O	R	T	E	X	A	M	P	L	E
2	10	A	O	R	T	E	X	S	M	P	L	E
3	9	A	E	E	T	O	X	S	M	P	L	R
4	7	A	E	E	L	O	X	S	M	P	T	R
5	7	A	E	E	L	M	X	S	O	P	T	R
6	8	A	E	E	L	M	O	S	X	P	T	R
7	10	A	E	E	L	M	O	P	X	S	T	R
8	8	A	E	E	L	M	O	P	R	S	T	X
9	9	A	E	E	L	M	O	P	R	S	T	X
10	10	A	E	E	L	M	O	P	R	S	T	X
		A	E	E	L	M	O	P	R	S	T	X

Trace of selection sort (array contents just after each exchange)

Running time insensitive to input. Quadratic time, even if input is sorted.
Data movement is minimal. Linear number of exchanges.

25

26

The image shows the front cover of the book 'Algorithms' by Robert Sedgewick and Kevin Wayne. The title 'Algorithms' is at the bottom left. The main title '2.1 ELEMENTARY SORTS' is centered above a list of topics. The topics include: rules of the game, selection sort, insertion sort, shuffling, and comparators. The background of the cover features a light gray abstract geometric pattern.

2.1 ELEMENTARY SORTS

- ▶ rules of the game
- ▶ selection sort
- ▶ insertion sort
- ▶ shuffling
- ▶ comparators

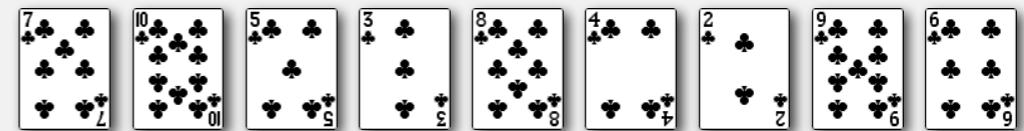
Algorithms

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Insertion sort demo

- In iteration i , swap $a[i]$ with each larger entry to its left.



28

Insertion sort

Algorithm. ↑ scans from left to right.

Invariants.

- Entries to the left of ↑ (including ↑) are in ascending order.
- Entries to the right of ↑ have not yet been seen.



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Insertion sort: inner loop

To maintain algorithm invariants:

- Move the pointer to the right.

i++;



in order not yet seen

- Moving from right to left, exchange a[i] with each larger entry to its left.

```
for (int j = i; j > 0; j--)  
    if (less(a[j], a[j-1]))  
        exch(a, j, j-1);  
    else break;
```



in order not yet seen

30

Insertion sort: Java implementation

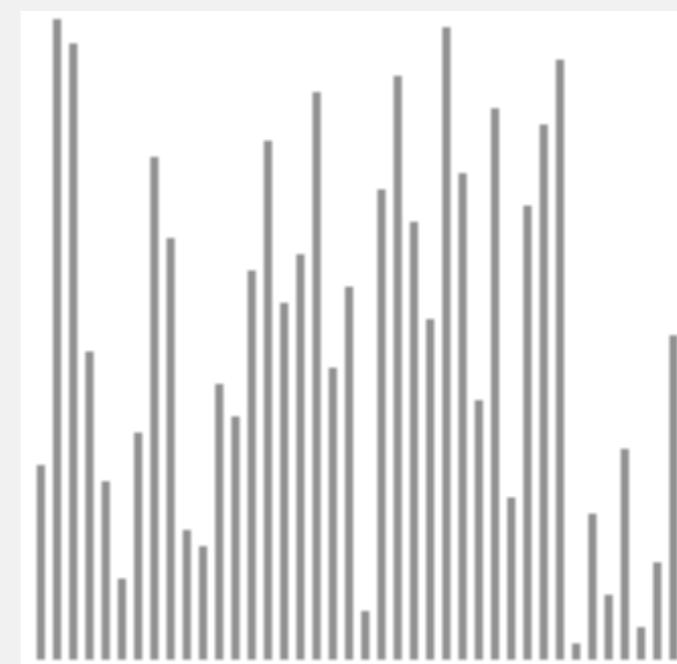
```
public class Insertion  
{  
    public static void sort(Comparable[] a)  
    {  
        int N = a.length;  
        for (int i = 0; i < N; i++)  
            for (int j = i; j > 0; j--)  
                if (less(a[j], a[j-1]))  
                    exch(a, j, j-1);  
                else break;  
    }  
  
    private static boolean less(Comparable v, Comparable w)  
    { /* as before */ }  
  
    private static void exch(Object[] a, int i, int j)  
    { /* as before */ }  
}
```

<http://algs4.cs.princeton.edu/21elementary/Insertion.java.html>

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Insertion sort: animation

40 random items



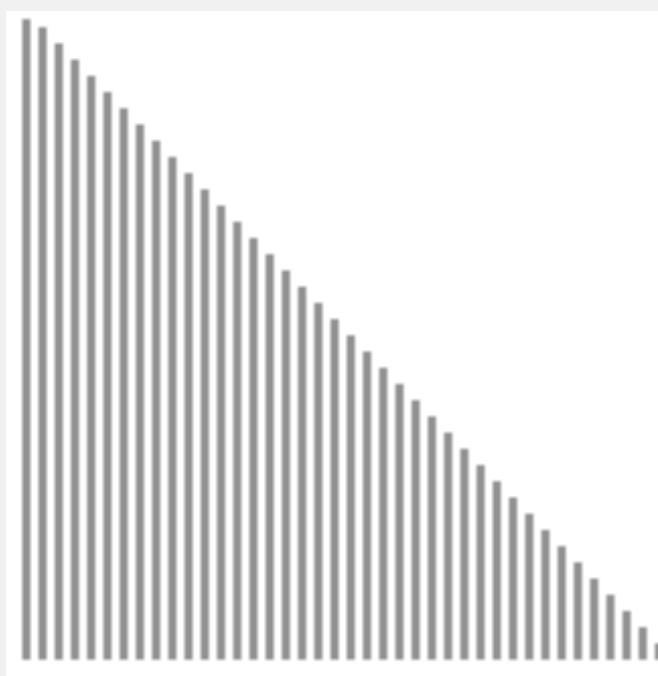
▲ algorithm position
— in order
■ not yet seen

<http://www.sorting-algorithms.com/insertion-sort>

32

Insertion sort: animation

40 reverse-sorted items



<http://www.sorting-algorithms.com/insertion-sort>

Insertion sort: mathematical analysis

Proposition. To sort a randomly-ordered array with distinct keys, insertion sort uses $\sim \frac{1}{4} N^2$ compares and $\sim \frac{1}{4} N^2$ exchanges on average.

Pf. Expect each entry to move halfway back.

		a[]										
i	j	0	1	2	3	4	5	6	7	8	9	10
1	0	O	S	R	T	E	X	A	M	P	L	E
2	1	O	R	S	T	E	X	A	M	P	L	E
3	3	O	R	S	T	E	X	A	M	P	L	E
4	0	E	O	R	S	T	X	A	M	P	L	E
5	5	E	O	R	S	T	X	A	M	P	L	E
6	0	A	E	O	R	S	T	X	M	P	L	E
7	2	A	E	M	O	R	S	T	X	P	L	E
8	4	A	E	M	O	P	R	S	T	X	L	E
9	2	A	E	L	M	O	P	R	S	T	X	E
10	2	A	E	E	L	M	O	P	R	S	T	X

Trace of insertion sort (array contents just after each insertion)

Elementary sorts: quiz 2

How many compares does insertion sort make to sort an array of N distinct keys in reverse order?

- A. $\sim N$
- B. $\sim \frac{1}{4} N^2$
- C. $\sim \frac{1}{2} N^2$
- D. $\sim N^2$
- E. I don't know.

Insertion sort: analysis

Worst case. If the array is in descending order (and no duplicates), insertion sort makes $\sim \frac{1}{2} N^2$ compares and $\sim \frac{1}{2} N^2$ exchanges.

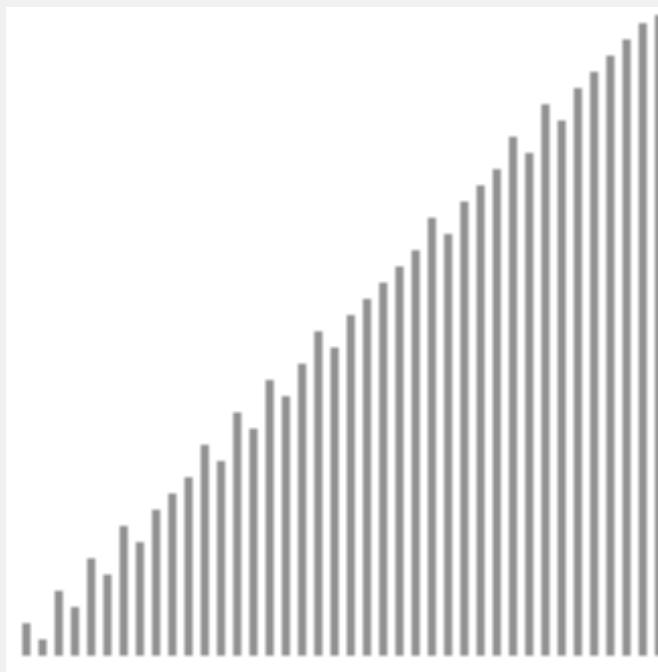
X T S R P O M L F E A

Best case. If the array is in ascending order, insertion sort makes $N-1$ compares and 0 exchanges.

A E E L M O P R S T X

Insertion sort: animation

40 partially-sorted items



<http://www.sorting-algorithms.com/insertion-sort>

Insertion sort: partially-sorted arrays

Def. An **inversion** is a pair of keys that are out of order.

A E E L M O T R X P S

T-R T-P T-S R-P X-P X-S
(6 inversions)

Def. An array is **partially sorted** if the number of inversions is $\leq c N$.

- Ex 1. A sorted array has 0 inversions.
- Ex 2. A subarray of size 10 appended to a sorted subarray of size N .

Proposition. For partially-sorted arrays, insertion sort runs in linear time.

Pf. Number of exchanges equals the number of inversions.

↑
number of compares \leq exchanges + $(N - 1)$

Insertion sort: practical improvements

Half exchanges. Shift items over (instead of exchanging).

- Eliminates unnecessary data movement.
- No longer uses only `less()` and `exch()` to access data.

A C H H I M N N P Q X Y K B I N A R Y

Binary insertion sort. Use binary search to find insertion point.

- Number of compares $\sim N \lg N$.
- But still a quadratic number of array accesses.

A C H H I M N N P Q X Y K B I N A R Y

binary search for first key > K

Elementary sorts: quiz 3

Which is faster in practice, selection sort or insertion sort?

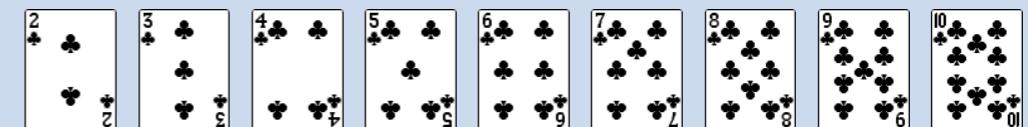
- A. Selection sort.
- B. Insertion sort.
- C. No significant difference.
- D. *I don't know.*



Interview question: shuffle an array

Goal. Rearrange array so that result is a uniformly random permutation.

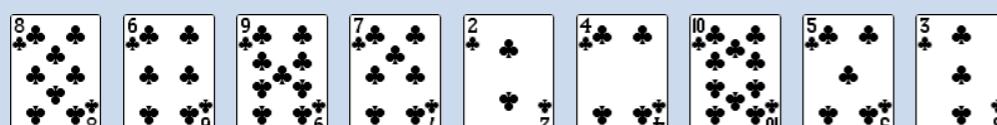
all $N!$ permutations
equally likely



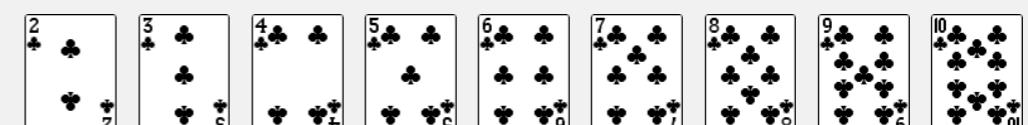
42

Shuffle sort

- Generate a random real number for each array entry.
- Sort the array.



all $N!$ permutations
equally likely

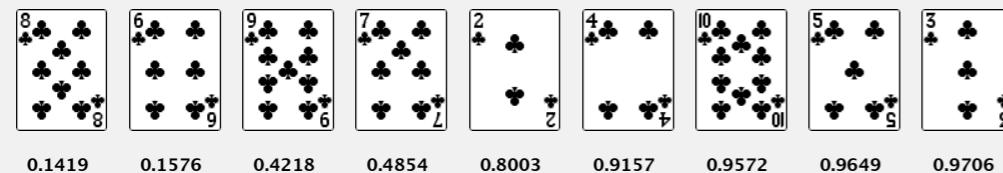


43

44

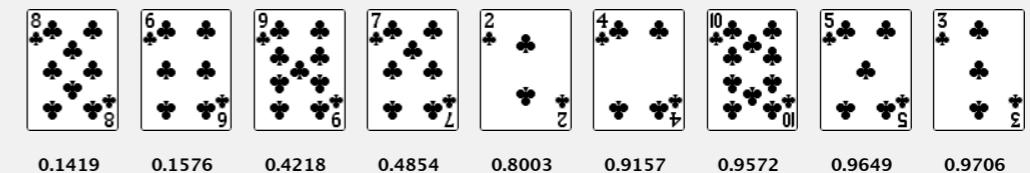
Shuffle sort

- Generate a random real number for each array entry.
- Sort the array.



Shuffle sort

- Generate a random real number for each array entry.
- Sort the array.



Proposition. Shuffle sort produces a uniformly random permutation.

Application. Shuffle columns in a spreadsheet.

assuming real numbers are uniformly random (and no ties)

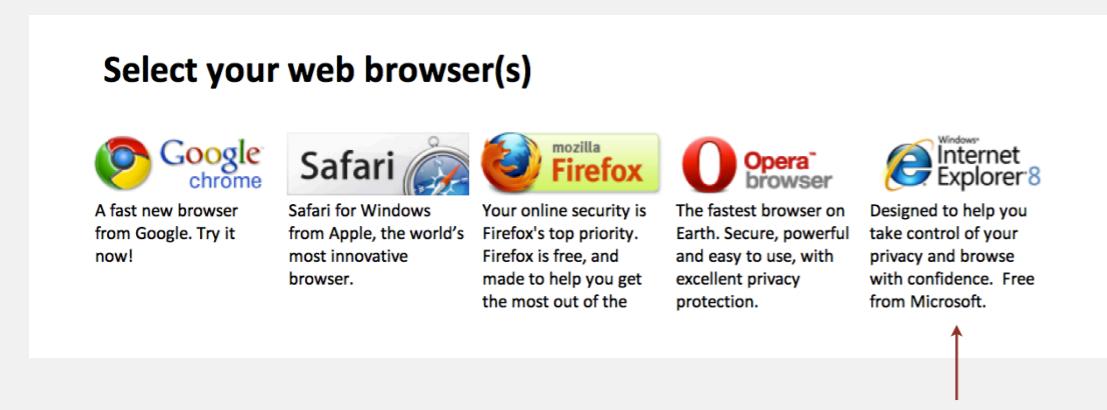
45

46

War story (Microsoft)

Microsoft antitrust probe by EU. Microsoft agreed to provide a randomized ballot screen for users to select browser in Windows 7.

<http://www.browserchoice.eu>



War story (Microsoft)

Microsoft antitrust probe by EU. Microsoft agreed to provide a randomized ballot screen for users to select browser in Windows 7.

Solution? Implement shuffle sort by making comparator always return a random answer.

```
public int compareTo(Browser that)
{
    double r = Math.random();
    if (r < 0.5) return -1;
    if (r > 0.5) return +1;
    return 0;
}
```

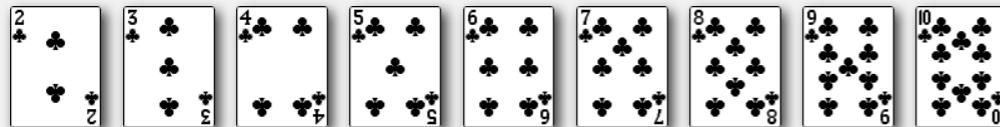
← browser comparator
(should implement a total order)

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Knuth shuffle demo

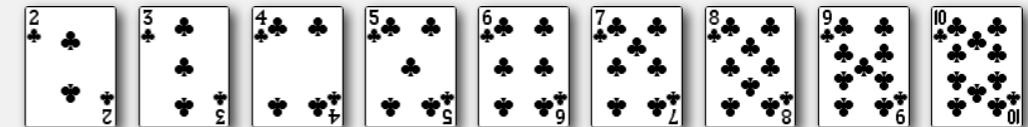
- In iteration i , pick integer r between 0 and i uniformly at random.
- Swap $a[i]$ and $a[r]$.



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Knuth shuffle

- In iteration i , pick integer r between 0 and i uniformly at random.
- Swap $a[i]$ and $a[r]$.



Proposition. [Fisher-Yates 1938] Knuth shuffling algorithm produces a uniformly random permutation of the input array in linear time.

assuming integers
uniformly at random

50

Knuth shuffle

- In iteration i , pick integer r between 0 and i uniformly at random.
- Swap $a[i]$ and $a[r]$.

common bug: between 0 and $N - 1$
correct variant: between i and $N - 1$

```
public class Knuth
{
    public static void shuffle(Object[] a)
    {
        int N = a.length;
        for (int i = 0; i < N; i++)
        {
            int r = StdRandom.uniform(i + 1);           ← between 0 and i
            exch(a, i, r);
        }
    }
}
```

<http://algs4.cs.princeton.edu/11model/Knuth.java.html>

Broken Knuth shuffle

- Q. What happens if integer is chosen between 0 and $N - 1$?
A. Not uniformly random!

instead of
between 0 and i

permutation	Knuth shuffle	broken shuffle
A B C	1/6	4/27
A C B	1/6	5/27
B A C	1/6	5/27
B C A	1/6	5/27
C A B	1/6	4/27
C B A	1/6	4/27

probability of each permutation when shuffling { A, B, C }

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War story (online poker)

Texas hold'em poker. Software must shuffle electronic cards.



How We Learned to Cheat at Online Poker: A Study in Software Security
http://www.digital.com/papers/download/developer_gambling.php

War story (online poker)

Shuffling algorithm in FAQ at www.planetpoker.com

```
for i := 1 to 52 do begin
    r := random(51) + 1; ← between 1 and 51
    swap := card[r];
    card[r] := card[i];
    card[i] := swap;
end;
```

- Bug 1. Random number r never 52 \Rightarrow 52nd card can't end up in 52nd place.
- Bug 2. Shuffle not uniform (should be between 1 and i).
- Bug 3. `random()` uses 32-bit seed $\Rightarrow 2^{32}$ possible shuffles.
- Bug 4. Seed = milliseconds since midnight \Rightarrow 86.4 million shuffles.

“The generation of random numbers is too important to be left to chance.”

— Robert R. Coveyou

War story (online poker)

Best practices for shuffling (if your business depends on it).

- Use a hardware random-number generator that has passed both the FIPS 140-2 and the NIST statistical test suites.
- Continuously monitor statistic properties: hardware random-number generators are fragile and fail silently.
- Use an unbiased shuffling algorithm.



RANDOM.ORG

Bottom line. Shuffling a deck of cards is hard!

2.1 ELEMENTARY SORTS



ROBERT SEDGEWICK | KEVIN WAYNE

<http://algs4.cs.princeton.edu>

- ▶ rules of the game
- ▶ selection sort
- ▶ insertion sort
- ▶ shuffling
- ▶ comparators

Sort music library by artist



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Sort music library by song name



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Comparable interface: review

Comparable interface: sort using a type's natural order.

```
public class Date implements Comparable<Date>
{
    private final int month, day, year;

    public Date(int m, int d, int y)
    {
        month = m;
        day   = d;
        year  = y;
    }

    ...
    public int compareTo(Date that)
    {
        if (this.year < that.year) return -1;
        if (this.year > that.year) return +1;
        if (this.month < that.month) return -1;
        if (this.month > that.month) return +1;
        if (this.day   < that.day)  return -1;
        if (this.day   > that.day)  return +1;
        return 0;
    }
}
```

natural order

Comparator interface

Comparator interface: sort using an alternate order.

```
public interface Comparator<Item>
{
    public int compare(Item v, Item w);
}
```

Required property. Must be a total order.

string order	example
natural order	Now is the time
case insensitive	is Now the time
Spanish language	café cafetero cuarto churro nube ñoño
British phone book	McKinley Mackintosh

pre-1994 order for
digraphs ch and ll and rr

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Comparator interface: system sort

To use with Java system sort:

- Create Comparator object.
- Pass as second argument to `Arrays.sort()`.

```
String[] a;           uses natural order
...
Arrays.sort(a);
...
Arrays.sort(a, String.CASE_INSENSITIVE_ORDER);
...
Arrays.sort(a, Collator.getInstance(new Locale("es")));
...
Arrays.sort(a, new BritishPhoneBookOrder());
...
```

Bottom line. Decouples the definition of the data type from the definition of what it means to compare two objects of that type.

Comparator interface: using with our sorting libraries

To support comparators in our sort implementations:

- Pass Comparator to both `sort()` and `less()`, and use it in `less()`.
- Use Object instead of Comparable.

```
import java.util.Comparator;

public class Insertion
{
    ...
    public static void sort(Object[] a, Comparator comparator)
    {
        int N = a.length;
        for (int i = 0; i < N; i++)
            for (int j = i; j > 0 && less(comparator, a[j], a[j-1]); j--)
                exch(a, j, j-1);
    }

    private static boolean less(Comparator comparator, Object v, Object w)
    {   return comparator.compare(v, w) < 0; }
}
```

<http://algs4.cs.princeton.edu/21elementary/Insertion.java.html>

Comparator interface: implementing

To implement a comparator:

- Define a (nested) class that implements the Comparator interface.
- Implement the `compare()` method.
- Provide client access to Comparator.

```
import java.util.Comparator;

public class Student
{
    private final String name;
    private final int section;
    ...
    one Comparator for the class
    public static Comparator<Student> byNameOrder()
    {   return new NameOrder(); }

    private static class NameOrder implements Comparator<Student>
    {
        public int compare(Student v, Student w)
        {   return v.name.compareTo(w.name); }
    }
    ...
}
```

Comparator interface: implementing

To implement a comparator:

- Define a (nested) class that implements the Comparator interface.
- Implement the `compare()` method.
- Provide client access to Comparator.

```
import java.util.Comparator;

public class Student
{
    private final String name;
    private final int section;
    ...
    public static Comparator<Student> bySectionOrder()
    {   return new SectionOrder(); }

    private static class SectionOrder implements Comparator<Student>
    {
        public int compare(Student v, Student w)
        {   return v.section - w.section; }
    }
    ...
}
```

this trick works here
since no danger of overflow

Comparator interface: implementing

To implement a comparator:

- Define a (nested) class that implements the Comparator interface.
- Implement the compare() method.
- Provide client access to Comparator.

Insertion.sort(a, Student.byNameOrder());

Andrews	3	A	(664) 480-0023	097 Little
Battle	4	C	(874) 088-1212	121 Whitman
Chen	3	A	(991) 878-4944	308 Blair
Fox	3	A	(884) 232-5341	11 Dickinson
Furia	1	A	(766) 093-9873	101 Brown
Gazsi	4	B	(800) 867-5309	101 Brown
Kanaga	3	B	(898) 122-9643	22 Brown
Rohde	2	A	(232) 343-5555	343 Forbes

Insertion.sort(a, Student.bySectionOrder());

Furia	1	A	(766) 093-9873	101 Brown
Rohde	2	A	(232) 343-5555	343 Forbes
Andrews	3	A	(664) 480-0023	097 Little
Chen	3	A	(991) 878-4944	308 Blair
Fox	3	A	(884) 232-5341	11 Dickinson
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Kanaga	3	B	(898) 122-9643	22 Brown
Battle	4	C	(874) 088-1212	121 Whitman
Gazsi	4	B	(800) 867-5309	101 Brown

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Stability

A typical application. First, sort by name; then sort by section.

Selection.sort(a, Student.byNameOrder());

Andrews	3	A	(664) 480-0023	097 Little
Battle	4	C	(874) 088-1212	121 Whitman
Chen	3	A	(991) 878-4944	308 Blair
Fox	3	A	(884) 232-5341	11 Dickinson
Furia	1	A	(766) 093-9873	101 Brown
Gazsi	4	B	(800) 867-5309	101 Brown
Kanaga	3	B	(898) 122-9643	22 Brown
Rohde	2	A	(232) 343-5555	343 Forbes

Selection.sort(a, Student.bySectionOrder());

Furia	1	A	(766) 093-9873	101 Brown
Rohde	2	A	(232) 343-5555	343 Forbes
Chen	3	A	(991) 878-4944	308 Blair
Fox	3	A	(884) 232-5341	11 Dickinson
Andrews	3	A	(664) 480-0023	097 Little
Kanaga	3	B	(898) 122-9643	22 Brown
Gazsi	4	B	(800) 867-5309	101 Brown
Battle	4	C	(874) 088-1212	121 Whitman

@#%&@! Students in section 3 no longer sorted by name.

A **stable** sort preserves the relative order of items with equal keys.

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Elementary sorts: quiz 4

Which sorting algorithms are stable?

- A. Selection sort.
- B. Insertion sort.
- C. Both A and B.
- D. Neither A nor B.
- E. I don't know.

Stability: insertion sort

Proposition. Insertion sort is **stable**.

```
public class Insertion
{
    public static void sort(Comparable[] a)
    {
        int N = a.length;
        for (int i = 0; i < N; i++)
            for (int j = i; j > 0 && less(a[j], a[j-1]); j--)
                exch(a, j, j-1);
    }
}
```

i	j	0	1	2	3	4
0	0	B ₁	A ₁	A ₂	A ₃	B ₂
1	0	A ₁	B ₁	A ₂	A ₃	B ₂
2	1	A ₁	A ₂	B ₁	A ₃	B ₂
3	2	A ₁	A ₂	A ₃	B ₁	B ₂
4	4	A ₁	A ₂	A ₃	B ₁	B ₂
		A ₁	A ₂	A ₃	B ₁	B ₂

Pf. Equal items never move past each other.

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Stability: selection sort

Proposition. Selection sort is **not stable**.

```
public class Selection
{
    public static void sort(Comparable[] a)
    {
        int N = a.length;
        for (int i = 0; i < N; i++)
        {
            int min = i;
            for (int j = i+1; j < N; j++)
                if (less(a[j], a[min]))
                    min = j;
            exch(a, i, min);
        }
    }
}
```

i	min	0	1	2
0	2	B ₁	B ₂	A
1	1	A	B ₂	B ₁
2	2	A	B ₂	B ₁

Pf by counterexample. Long-distance exchange can move one equal item past another one.