

### Cryptography

COS 461: Computer Networks
Princeton University

#### Overview

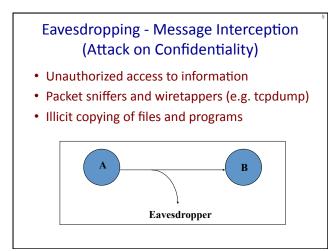
- Network security and definitions
- Brief introduction to cryptography
  - Cryptographic hash functions
  - Symmetric-key crypto
  - Public-key crypto
  - Hybrid crypto

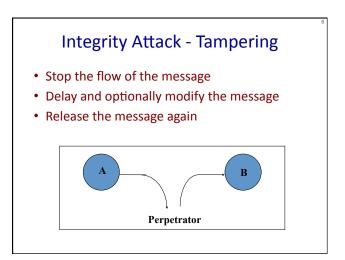
## Internet's Design: Insecure

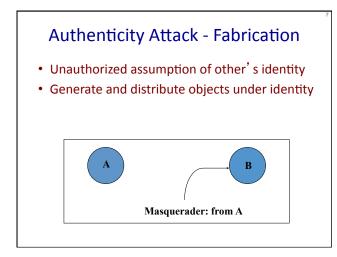
- · Designed for simplicity
- "On by default" design
- Readily available zombie machines
- Attacks look like normal traffic
- Internet's federated operation obstructs cooperation for diagnosis/mitigation

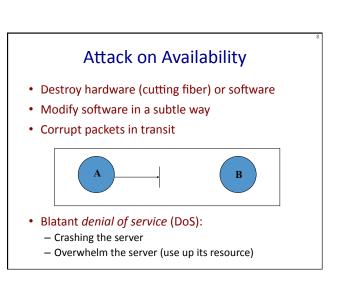
# **Basic Components**

- Confidentiality: Concealment of information or resources
- Authenticity: Identification and assurance of origin of info
- Integrity: Trustworthiness of data or resources in terms of preventing improper and unauthorized changes
- Availability: Ability to use desired info or resource
- Non-repudiation: Offer of evidence that a party indeed is sender or a receiver of certain information
- Access control: Facilities to determine and enforce who is allowed access to what resources (host, software, network, ...)









## Introduction to Cryptography

# What is Cryptography?

- Comes from Greek meaning "secret writing"
  - Primitives also can provide integrity, authentication
- Cryptographers invent secret codes to attempt to hide messages from unauthorized observers

	encryption		decryption	
plaintext	$\stackrel{\hspace{1cm}\longrightarrow}$	ciphertext	${-\!\!\!-\!\!\!\!-}$	plaintext

- Modern encryption:
  - Algorithm public, key secret and provides security
  - May be symmetric (secret) or asymmetric (public)

### Cryptographic Algorithms: Goal

- One-way functions: cryptographic hash
  - Easy to compute hash
  - Hard to invert
- "Trapdoor" functions: encryption/signatures
  - Given ciphertext alone, hard to compute plaintext (invert)
  - Given ciphertext and key (the "trapdoor"), relatively easy to compute plaintext
  - "Level" of security often based on "length" of key

### Cryptographic hash functions

### **Cryptography Hash Functions**

- Take message, m, of arbitrary length and produces a smaller (short) number, h(m)
- One-way function
  - Easy to compute h(m)
  - Hard to find an m, given h(m)
  - Often assumed: output of hash fn's "looks" random
- · Collision resistance:
  - Strong: Find any m!= m' such that h(m) == h(m')
  - Weak: Given m, find m' such that h(m) == h(m')
  - For 160-bit hash (SHA-1)
    - Strong collision are birthday paradox:  $2^{160/2} = 2^80$
    - Weak collision requires 2^160

### Example use #1: Passwords

- Can't store passwords in a file that could be read
   Concerned with insider attacks!
- Must compare typed passwords to stored passwords
   Does hash (typed) == hash (password)?
- Actually, a "salt" is often used: hash (input | | salt)

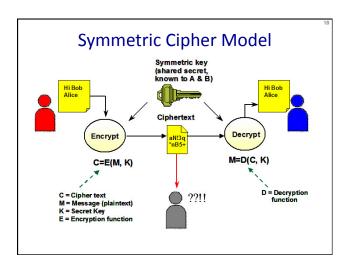
#### Example use #2: Self-certifying naming

- File-sharing software (LimeWire, BitTorrent)
  - File named by F<sub>name</sub> = hash (data)
  - Participants verify that hash (downloaded) ==  $F_{name}$ 
    - If check fails, reject data
    - To successfully "forge" data', must find weak collision
- · Recursively applied...
  - BitTorrent file has many chunks
  - Control file downloaded from tracker includes:
    - \forall chunks, F<sub>chunk name</sub> = hash (chunk)
  - BitTorrent client verifies each individual chunk

Symmetric (Secret) Key Cryptography

### **Symmetric Encryption**

- Also: "conventional / private-key / single-key"
  - Sender and recipient share a common key
  - All classical encryption algorithms are private-key
  - Dual use: confidentiality or authentication/integrity
    - Encryption vs. msg authentication code (MAC)
- Was only type of encryption prior to invention of public-key in 1970's
  - Most widely used
  - (Much) more computationally efficient than "public key"



## **Distribution of Symmetric Keys**

- Options: (between A and B).
  - A selects a key and physically delivers it to B.
  - If A and B already have have a viable key, it can be used to distribute a new key.
  - A trusted third party key distribution center (KDC) selects a key and physically delivers it to A and B (through a secure channel).

## **Distribution of Symmetric Keys**

- Manual delivery is challenging...
- The number of keys grows quadratically with the number of endpoints (n\*(n-1)/2)
  - Further complexity for application/user level encryption
- Key distribution center (KDC) a good alternative
  - Only *n* master keys required
  - KDC generate session key for Alice and Bob

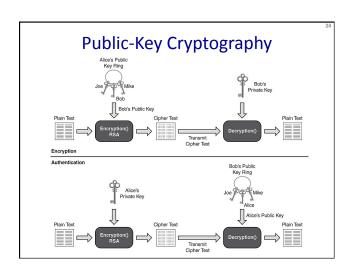
# **Public-Key Cryptography**

# Why Public-Key Cryptography?

- Developed to address two key issues:
  - Key distribution: Secure communication w/o having to trust a key distribution center with your key
  - Digital signatures: Verify msg comes intact from claimed sender (w/o prior establishment)
- Public invention due to Whitfield Diffie & Martin Hellman in 1976
  - Known earlier in classified community

# **Public-Key Cryptography**

- **Public-key:** Known by anybody, and can be used to encrypt messages and verify signatures
- **Private-key:** Known only to recipient, used to decrypt messages and sign (create) signatures
- Can encrypt messages or verify signatures w/o ability to decrypt messages or create signatures



# **Security of Public Key Schemes**

- Public-key encryption is a "trap-door" function:
  - Easy to compute  $c \leftarrow F(m, k_p)$
  - Hard to compute  $m \leftarrow F^{-1}(c)$  without knowing  $k_s$
  - Easy to compute  $m \leftarrow F^{-1}(c, k_s)$  by knowing  $k_s$
- Like private key schemes, brute force search possible
  - But keys used are too large (e.g., >= 2048 bits)
  - Hence is slow compared to private key schemes

#### (Simple) RSA Algorithm

- Security due to cost of factoring large numbers
  - Factorization takes O(e log n log log n) operations (hard)
  - Exponentiation takes O((log n)<sup>3</sup>) operations (easy)
- To encrypt a message M the sender:
  - Obtain public key {e,n}; compute C = Me mod n
- To decrypt the ciphertext C the owner:
  - Use private key  $\{d, n\}$ ; computes  $M = C^d \mod n$
- Note that msg M must be smaller than the modulus n
  - Otherwise, hybrid encryption:
    - ullet Generate random symmetric key r
    - ullet Use public key encryption to encrypt r
    - Use symmetric key encryption under r to encrypt M

## Symmetric vs. Asymmetric

- Symmetric Pros and Cons
  - Simple and really very fast (order of 1000 to 10000 faster than asymmetric mechanisms)
  - Must agree/distribute the key beforehand
  - AES/CBC (256-bit)  $\rightarrow$  80 MB / s
  - (for 2048 bits, that's .003 ms)
- · Public Key Pros and Cons
  - Easier predistribution for public keys
    - Public Key Infrastructure (in textbook)
  - Much slower
  - 2048-RSA → 6.1ms Decrypt, 0.16ms Encrypt