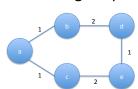


Suppose network operator Olivia decides to bring down the link c – e for maintenance. Olivia figures she can issue a series of link weight changes in the network to shift traffic away from c – e such that no temporary forwarding loops occur.

She's right; what series of changes to c – e's weight would achieve this?

Routing Loops



Change c – e weight to 4.

a to e moves to a - b.

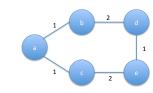
c to e remains on c - e.

c to d?

d to c moves to d - b.

e to c remains on c – e.

Routing Loops



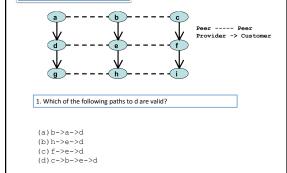
Change c – e weight to 6.

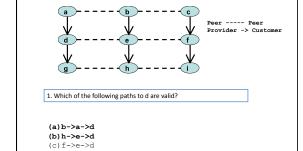
c to e moves to c – a.

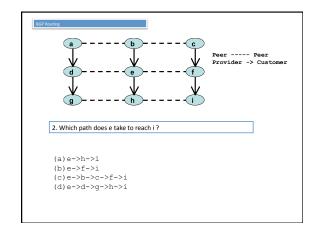
e to c moves to e - d.

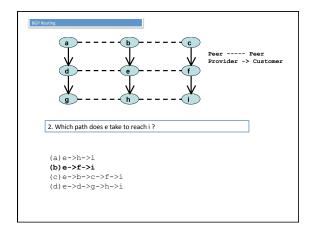
(d)c->b->e->d

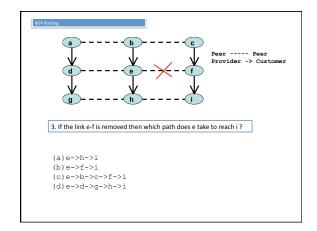
No traffic left on c – e. Olivia can bring down the link.

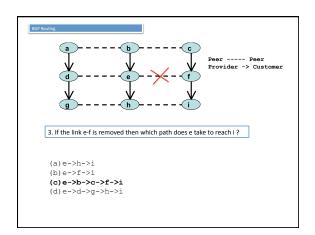


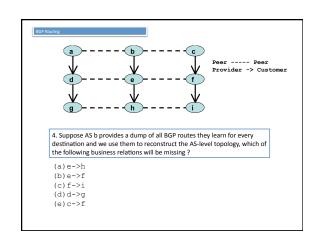


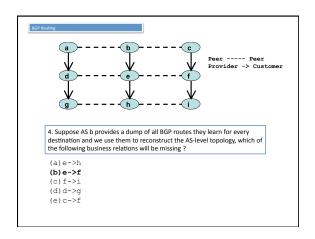


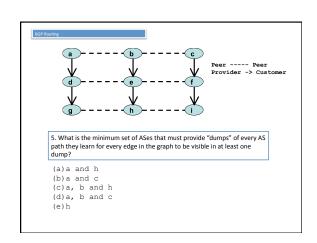


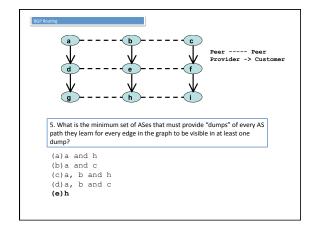


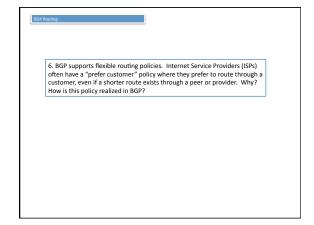












6. BGP supports flexible routing policies. Internet Service Providers (ISPs) often have a "prefer customer" policy where they prefer to route through a customer, even if a shorter route exists through a peer or provider. Why? How is this policy realized in BGP?

Directing traffic through a customer generates revenue, whereas sending through a peer or provider is (at best) revenue neutral and may, in fact, cost money.

The policy is realized in BGP by having an import policy that assigns a higher local-preference value to routes learned from customer ASes.