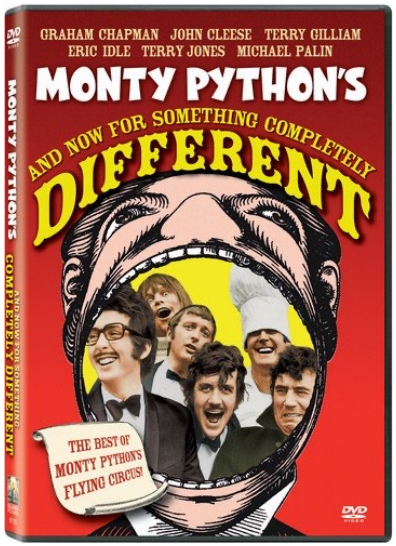
COS 217: Introduction to Programming Systems

Assembly Language

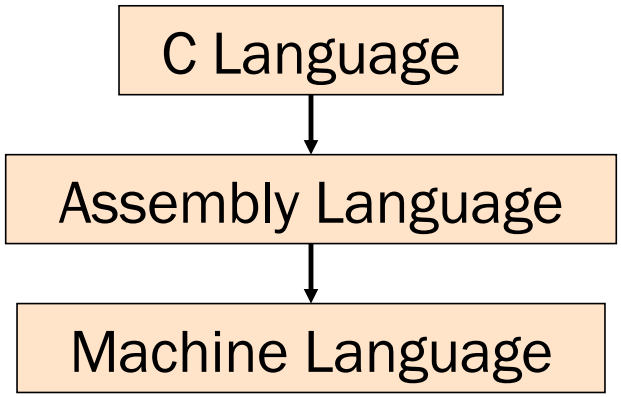
Part 1



PRINCETON UNIVERSITY

Context of this Lecture

"Under the hood"



Agenda

Language Levels

Architecture

Assembly Language: Performing Arithmetic

Assembly Language: Load/Store and Defining Global Data

High-Level Languages

Characteristics

- Portable (to varying degrees)
- Complex
  - One statement can do a lot of work – good ratio of functionality to code size
- Human readable
  - Structured: if(), for(), while(), etc.
  - Variable names can hide details of where data is stored (stack, heap, etc.)
  - Type system allows compiler to check usage details without burdening reader

```
int collatz(int n)
{
    int count = 0;
    while (n > 1) {
        count++;
        if (n & 1)
            n = 3 * n + 1;
        else
            n /= 2;
    }
    return count;
}
```

Machine Languages

Characteristics

- Not portable (hardware-specific)
- Simple
  - Every instruction does a simple task – low expressivity (ratio of functionality to code size)
- Not human readable
  - Not structured
  - Requires a lot of effort
  - Requires tool support

```
0000 0000 0000 0000 0000 0000 0000 0000
0000 0000 0000 0000 0000 0000 0000 0000
9222 9120 1121 A120 1121 A121 7211 0000
0000 0001 0002 0003 0004 0005 0006 0007
0008 0009 000A 000B 000C 000D 000E 000F
0000 0000 0000 FE10 FACE CAFE ACED CEDE

1234 5678 9ABC DEF0 0000 0000 F00D 0000
0000 0000 EEEE 1111 EEEE 1111 0000 0000
B1B2 F1F5 0000 0000 0000 0000 0000 0000
```

Assembly Languages

Characteristics

- Not portable
  - Every assembly language instruction maps to one machine instruction
- Simple
  - Every instruction does a simple task
- Human readable …

```
        mov     w1, 0
loop:
        cmp     w0, 1
        ble     endloop
        add     w1, w1, #1
        ands    wzr, w0, #1
        beq     else
        add     w2, w0, w0
        add     w0, w0, w2
        add     w0, w0, 1
        b       endif
else:
        asr     w0, w0, 1
endif:
        b       loop
endloop:
```

Why Learn Assembly Language?

Knowing assembly language helps you:

- Write faster code
  - In assembly language
  - Potentially even in a high-level language
- Write safer code
  - Understanding mechanism of potential security problems helps you avoid them – even in high-level languages
- Understand what's happening under the hood
  - Someone needs to develop future computer systems
- Become more comfortable with levels of abstraction
  - Become a better programmer at all language levels

Why Learn ARMv8 Assembly Language?

Pros of learning ARMv8 (a.k.a. AARCH64 or A64) assembly

- ARM is the most widely used processor architecture in the world (in your phone, in your Mac, in your Chromebook, in Armlab, IoT devices)
- ARM has a modern and (relatively) elegant instruction set ("RISC" – Reduced Instruction Set Computer) with each instruction being the same size (4 bytes). C.f., the expansive but overwhelming x86-64 instruction set

Cons

- x86-64 still has a huge presence in desktop/laptop/cloud

Lectures vs. Precepts

Approach to studying assembly language:

Lectures	Precepts
Study partial programs	Study complete programs
Begin with simple constructs; proceed to complex ones	Begin with small programs; proceed to large ones
Emphasis on reading code	Emphasis on writing code

Agenda

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Assembly Language: Load/Store and Defining Global Data

John von Neumann (1903-1957)

In computing

- Stored program computers
- Cellular automata, self-replication,
- Game theory
- mergesort

Other interests

- Mathematics, statistics
- Nuclear physics

Princeton connection

- Princeton University & IAS, 1930-1957
- https://paw.princeton.edu/article/early-history-computing-princeton

Known for the "Von Neumann architecture"

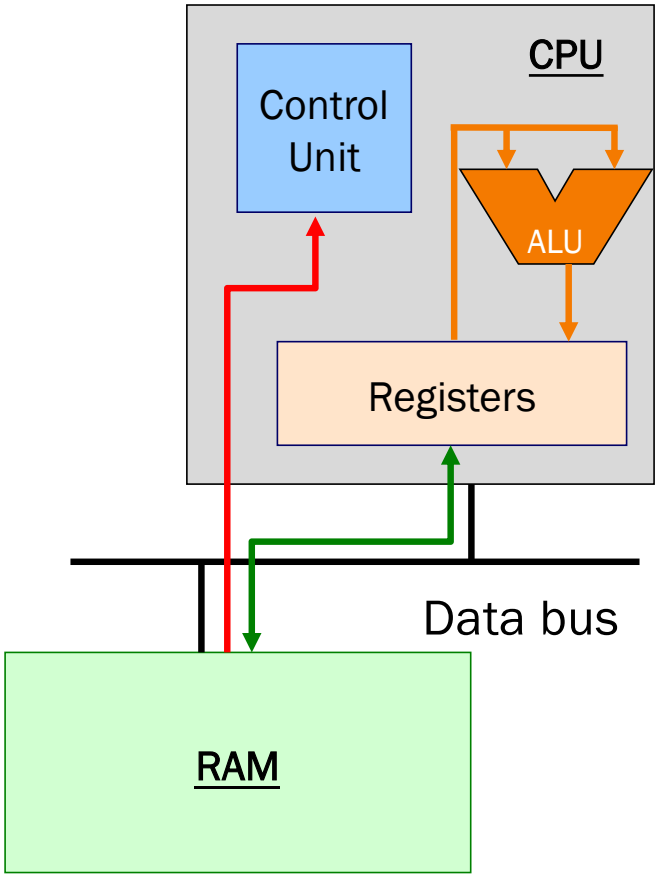
- In which (machine-language) programs are just data in memory
- a.k.a. "Princeton architecture" – contrast to the now-mostly-obsolete "Harvard architecture"

Von Neumann Architecture

Instructions (encoded within words) are fetched from RAM

Control unit interprets instructions:

- to shuffle data between registers and RAM
- to move data from registers to ALU (arithmetic+logic unit) where operations are performed



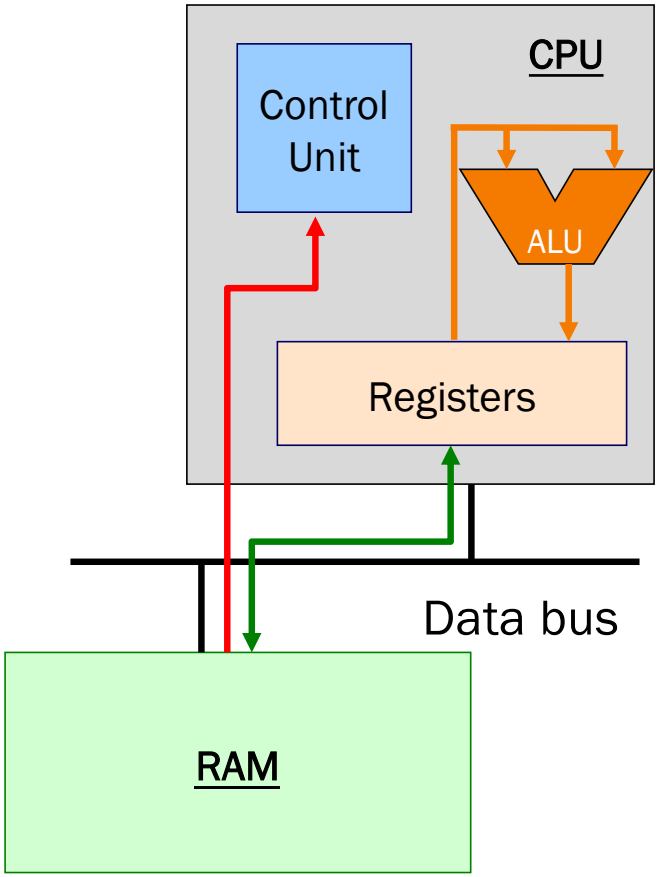
Von Neumann Architecture

Registers

Small amount of storage on the CPU

- Top of the “storage hierarchy”
- Very {small, expensive, fast}

ALU instructions operate on registers



Von Neumann Architecture

RAM (Random Access Memory)

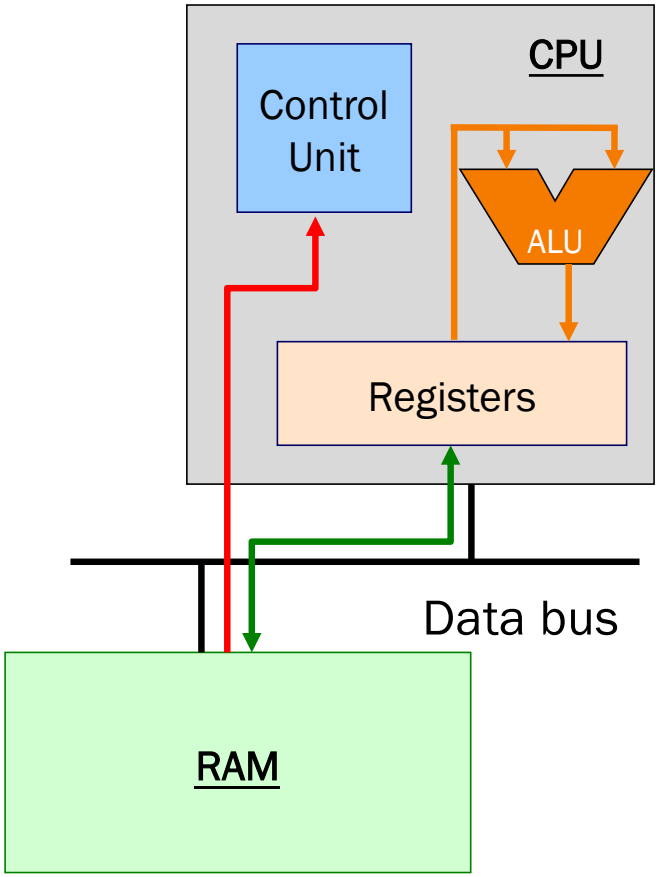
Conceptually: large array of bytes

(gigabytes+ in modern machines)

(registers are ~512 bytes)

- Contains data (variables, data structures)
- And the program itself, in machine code

Instructions are fetched from RAM



Time to reminisce about old TOYs

Thinking back to COS 126,
how did you feel about TOY?

A. Loved it

B. Wasn't a fan.

C. I took ECE115, so I have no idea what you're talking about

D. I placed out, so I have no idea what you're talking about

Time to reminisce about old TOYs

```
                TOY REFERENCE CARD


INSTRUCTION FORMATS

              | . . . . | . . . . | . . . . | . . . .|
  Format RR:  | opcode  |    d    |    s    |    t    |  (0-6, A-B)
  Format A:   | opcode  |    d    |       addr        |  (7-9, C-F)


ARITHMETIC and LOGICAL o
    1: add
    2: subtract
    3: and
    4: xor
    5: shift left
    6: shift right

TRANSFER between registe
    7: load address
    8: load
    9: store
    A: load indirect
    B: store indirect

CONTROL
    0: halt
    C: branch zero
    D: branch positive
    E: jump register
    F: jump and link


Register 0 always reads 0.
Loads from M[FF] come from stdin.
Stores to M[FF] go to stdout.

16-bit registers (two's complement)
16-bit memory locations
 8-bit program counter
```

Word size. The TOY machine has two types of storage: main memory and registers. Each entity stores one *word* of information. On the TOY machine, a word is a sequence of 16 bits. Typically, we interpret these 16 bits as a hexadecimal integer in the range `0000` through `FFFF`. Using *two's complement notation*, we can also interpret it as a decimal integer in the range -32,768 to +32,767. See Section 5.1 for a refresher on number representations and two's complement integers.

Main memory. The TOY machine has 256 words of *main memory*. Each memory location is labeled with a unique *memory address*. By convention, we use the 256 hexadecimal integers in the range `00` through `FF`. Think of a memory location as a mailbox, and a memory address as a postal address. Main memory is used to store instructions and data.

Registers. The TOY machine has 16 *registers*, indexed from `0` through `F`. Registers are much like main memory: each register stores one 16-bit word. However, registers provide a faster form of storage than main memory. Registers are used as scratch space during computation and play the role of variables in the TOY language. Register 0 is a special register whose output value is always 0.

Program counter. The *program counter* or `pc` is an extra register that keeps track of the next instruction to be executed. It stores 8 bits, corresponding to a hexadecimal integer in the range `00` through `FF`. This integer stores the memory address of the next instruction to execute.

https://introcs.cs.princeton.edu/java/62toy/

Registers and RAM

Typical pattern:

- **Load** data from RAM to registers
- **Manipulate** data in registers
- **Store** data from registers to RAM

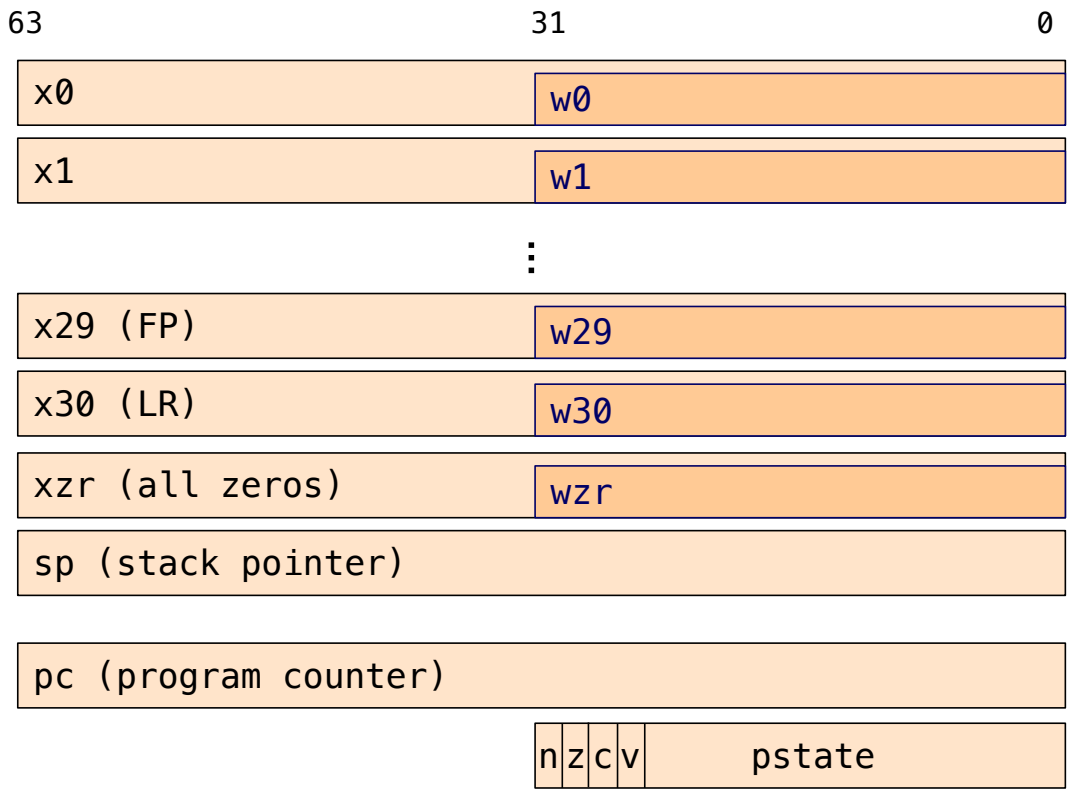
On AARCH64, this pattern is enforced

- "Manipulation" instructions can *only* access registers
- This is known as a **load-store architecture** (as opposed to "register-memory" architectures)
- Characteristic of RISC architectures (vs."CISC" or Complex Instruction Set Computer, e.g. x86)

Even RISC compilers don't use all their instructions

- We'll use this fact to "beat the compiler" in Assignment A5. CISC is much worse

Registers (ARM-64 architecture)



General-Purpose 64-bit Registers

X0 … X30

- Scratch space for instructions, parameter passing to/from functions, return address for function calls, etc.
- Some have special roles defined *in hardware* (e.g. X30) or defined *by software convention* (e.g. X29)
- Also available as 32-bit versions: W0 … W30

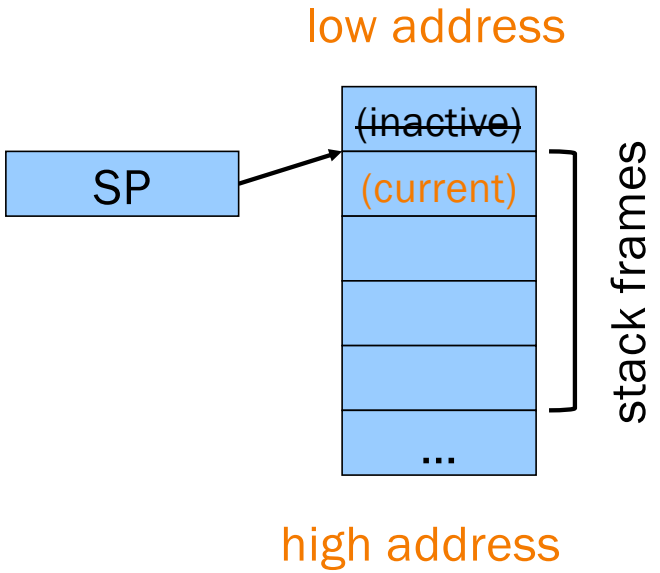
XZR

- On read: all zeros
- On write: data thrown away
- Also available as 32-bit version: WZR

SP Register

Special-purpose register…

- **SP (Stack Pointer):**
  Contains address of top (low memory address) of current function's stack frame



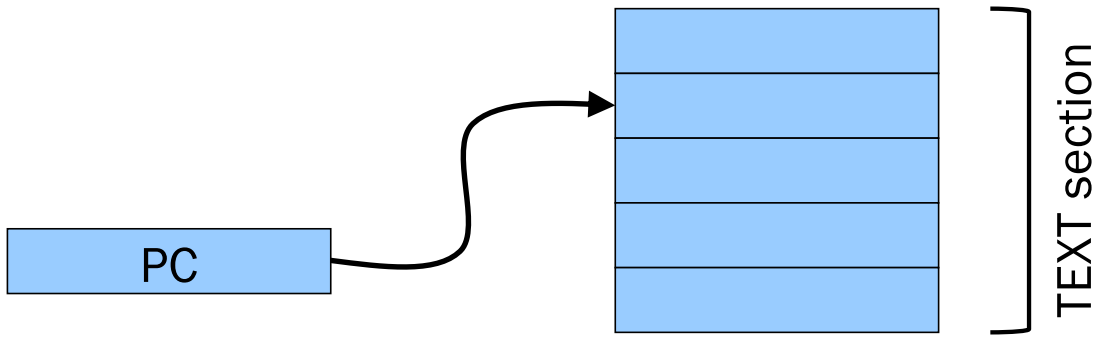
Allows use of the STACK section of memory

(See **Assembly Language: Function Calls** lecture later)

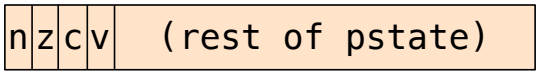
PC Register

Special-purpose register…

- **PC (Program Counter)**
- Stores the location of the next instruction
  - Address (in TEXT section) of machine-language instruction to be executed next
- Its value is changed either:
  - Automatically to implement sequential control flow (increment by 4 bytes)
  - By branch instructions to implement selection, repetition



PSTATE Register



Special-purpose register…

- Contains **condition flags:**

  n (**N**egative), z (**Z**ero), c (**C**arry), v (o**V**erflow)

- Affected by compare (`cmp`) instruction
  - And many others, if requested (e.g. with s suffix on arithmetic instructions)
- Used by conditional branch instructions
  - `beq`, `bne`, `blo`, `bhi`, `ble`, `bge`, …
  - (See **Assembly Language: Part 2** lecture)

Agenda

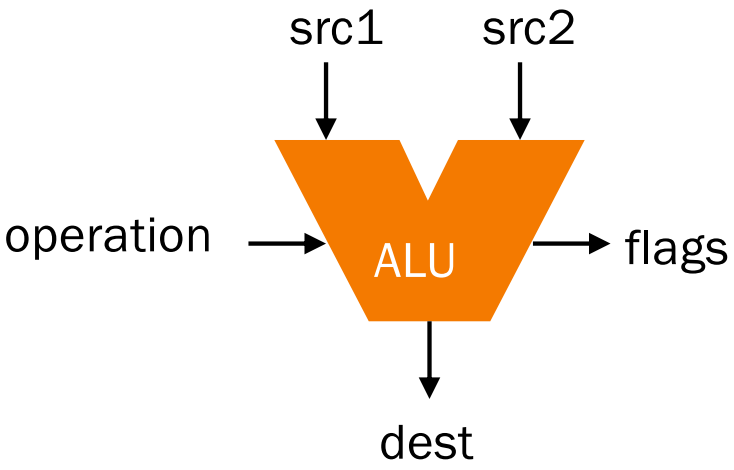
Language Levels

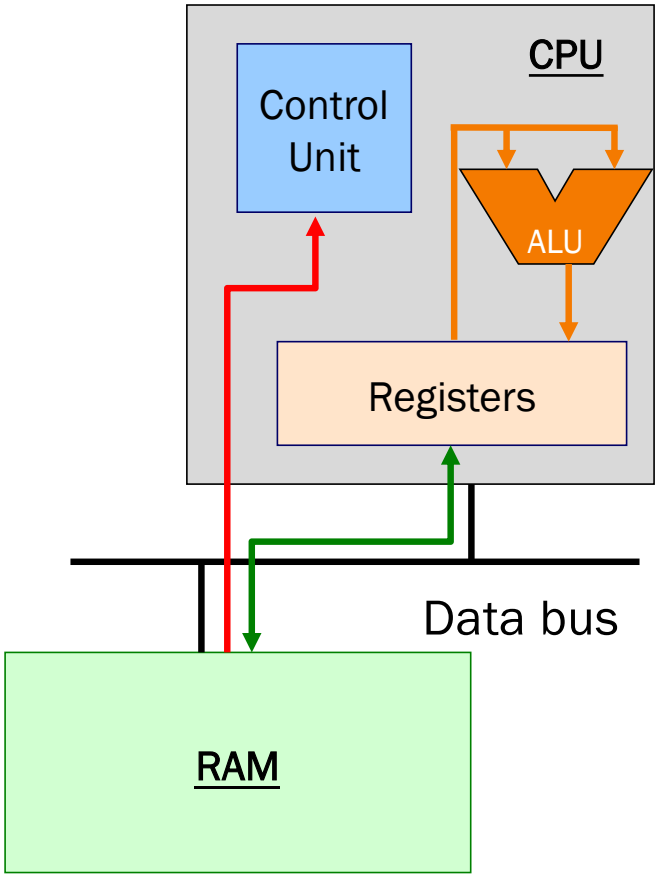
Architecture

Assembly Language: Performing Arithmetic

Assembly Language: Load/Store and Defining Global Data

ALU Arithmetic Example



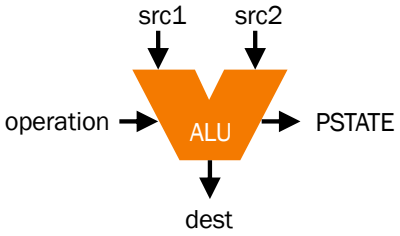



```
add dest, src1, src2
```

Instruction Format

Many instructions have this format:

```
name{,s} dest, src1, src2
name{,s} dest, src1, immed
```



- **name:** mnemonic name of the instruction (`add`, `sub`, `mul`, `and`, etc.)
- **s:** if present, specifies that condition flags should be **S**et

- dest and src1,src2 are **x** registers: 64-bit operation
- dest and src1,src2 are **w** registers: 32-bit operation
  - No mixing and matching between x and w registers

- src2 may be a constant ("immediate" value) instead of a register

64-bit Arithmetic

C code:

```
static long length;
static long width;
static long perim;
...
perim =
  (length + width) * 2;
```

Assume that…

- length held in x1
- width held in x2
- perim held in x3
- they're not local variables (we'll see those later)

Assembly code:

```
add x3, x1, x2
lsl x3, x3, 1
```

Recall use of left shift by 1 bit to multiply by 2

You'll see #1 instead of 1 in book. Also works, but not necessary.

More Arithmetic

```
static long x;
static long y;
static long z;
...
z = x - y;
z = x * y;
z = x / y;
z = x & y;
z = x | y;
z = x ^ y;
z = x >> y;
```

Assume that…

- x held in x1
- y held in x2
- z held in x3

Assembly code:

```
sub  x3, x1, x2
mul  x3, x1, x2
sdiv x3, x1, x2
and  x3, x1, x2
orr  x3, x1, x2
eor  x3, x1, x2
asr  x3, x1, x2
```

Note: unlike C, separate instructions for arithmetic and logical shifts

More Arithmetic: Shortcuts

```
static long x;
static long y;
static long z;
...
z = x;
z = -x;
```

Assume that…

- x held in x1
- y held in x2
- z held in x3

Assembly code:

```
mov  x3, x1
neg  x3, x1
```

```
orr  x3, xzr, x1
sub  x3, xzr, x1
```

Abstraction

These are actually assembler shortcuts for instructions that use XZR.

Precepts will cover signed/unsigned arithmetic and handling smaller operand sizes

Signed vs Unsigned?

```
static long x;
static unsigned long y;
...
x++;
y--;
```

Assume that…

- x held in x1
- y held in x2

Assembly code:

```
add x1, x1, 1
sub x2, x2, 1
```

Mostly the same algorithms, same instructions

- Can set different condition flags in PSTATE
- But some exceptions…

Signed vs Unsigned: Exceptions

```
static long x;
static unsigned long y;
...
x /= 17;
y /= 42;
x >>= 1;
y >>= 2;
```

Assume that…

- x held in x1
- y held in x2

Assembly code:

```
sdiv x1, x1, 17
udiv x2, x2, 42
asr x1, x1, 1
lsr x2, x2, 2
```

Division needs different algorithms for signed and unsigned, and hence different instructions

"Arithmetic" right shift (shift in sign bit on left) vs. "logical" right shift (shift in zeros on left)

32-bit Arithmetic using “w” registers

C code:

```
static int length;
static int width;
static int perim;
...
perim =
  (length + width) * 2;
```

Assume that…

- length held in w1
- width held in w2
- perim held in w3

Assembly code:

```
add w3, w1, w2
lsl w3, w3, 1
```

8- and 16-bit Arithmetic?

```
static char x;
static short y;
...
x++;
y--;
```

No specialized arithmetic instructions

- Just use the “w” registers
- Corresponds to C language semantics: all arithmetic is implicitly done on (at least) ints, values are implicitly promoted to 32 bit
- Specialized “load” and “store” instructions for transfer of shorter data types from / to memory – we’ll see these next

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Loads and Stores

Most basic way to load (from RAM) and store (to RAM):

```
ldr dest, [src]
str src, [dest]
```

- `dest`, `src` are registers. The memory operand always comes second
- **Memory addressing must be via registers (not literal addresses as in TOY)**
  - Every memory access effectively dereferences a 'pointer' (the register)
- Contents of registers in [brackets] must be memory addresses
  - Note how this represents memory as a C array with the addresses being its indices

- The memory addresses (`src` for `ldr`, `dest` for `str`) must be x-flavored
- Other (register) operands (`dest` for `ldr`, `src` for `str`) can be x-flavored or w-flavored

Signed vs Unsigned, 8- and 16-bit

```
ldrb  dest, [src]
ldrh  dest, [src]
strb  src, [dest]
strh  src, [dest]

ldrsb dest, [src]
ldrsh dest, [src]
ldrsw dest, [src]
```

Special instructions for reading/writing **B**ytes (8 bit) and shorts ("**H**alf-words": 16 bit)

- See appendix of these slides for ordering information: little-endian vs. big-endian

Special instructions for signed loads (tell how to fill in leftmost bits in register)

- "**S**ign-extend" byte, half-word, or word to 32 or 64 bits

A Note on Loads and Stores

```
ldr dest, [src]
str src, [dest]
```

- How to get correct memory address into register?
  - Depends on whether data are on stack (local variables), on heap (dynamically-allocated memory), or global / static
  - For today, we'll look only at the global / static data case
  - `adr` instruction puts the address of a given label into a given register
    - The load/store then uses the register as a pointer (its value as a memory address)

Making a "Pointer"

```
static int length = 1;
static int width = 2;
static int perim = 0;

int main()
{
  perim =
   (length + width) * 2;
   return 0;
}
```

Generating addresses

adr: puts address of a label in a register

```
  .section .data
length: .word 1
width:  .word 2
perim:  .word 0
  .section .text
  .global main
main:
adr     x0, length
ldr     w1, [x0]
adr     x0, width
ldr     w2, [x0]
add     w1, w1, w2
lsl     w1, w1, 1
adr     x0, perim
str     w1, [x0]
mov     w0, 0
ret
```

Our First Full Program (Perimeter)*

```
static int length = 1;
static int width = 2;
static int perim = 0;

int main()
{
  perim =
   (length + width) * 2;
   return 0;
}
```

```
  .section .data
length: .word 1
width:  .word 2
perim:  .word 0
  .section .text
  .global main
main:
adr     x0, length
ldr     w1, [x0]
adr     x0, width
ldr     w2, [x0]
add     w1, w1, w2
lsl     w1, w1, 1
adr     x0, perim
str     w1, [x0]
mov     w0, 0
ret
```


* You'll see the more typical, "Hello, World!" program in precept …

Memory Sections

```
static int length = 1;
static int width = 2;
static int perim = 0;

int main()
{
  perim =
   (length + width) * 2;
   return 0;
}
```

```
  .section .data
length: .word 1
width:  .word 2
perim:  .word 0
  .section .text
  .global main
main:
adr     x0, length
ldr     w1, [x0]
adr     x0, width
ldr     w2, [x0]
add     w1, w1, w2
lsl     w1, w1, 1
adr     x0, perim
str     w1, [x0]
mov     w0, 0
ret
```

Sections (Stack/heap are different!)

- .rodata: read-only
- .data: read-write
- .bss: read-write (uninit or init to 0)
- .text: read-only, program code

Variable Definitions and Usage

```
static int length = 1;
static int width = 2;
static int perim = 0;

int main()
{
  perim =
   (length + width) * 2;
   return 0;
}
```

```
  .section .data
length: .word 1
width:  .word 2
perim:  .word 0
  .section .text
  .global main
main:
adr     x0, length
ldr     w1, [x0]
adr     x0, width
ldr     w2, [x0]
add     w1, w1, w2
lsl     w1, w1, 1
adr     x0, perim
str     w1, [x0]
mov     w0, 0
ret
```

We have variables in assembly language

- Abstraction

Declaring data

"Labels" for locations in memory

.word: 32-bit int and initial value

See appendix for variables in other sections, with other types.

main()

```
static int length = 1;
static int width = 2;
static int perim = 0;

int main()
{
  perim =
   (length + width) * 2;
   return 0;
}
```

Global visibility

.global: Declare “main” to be a globally-visible label

```
  .section .data
length: .word 1
width:  .word 2
perim:  .word 0
  .section .text
  .global main
main:
adr     x0, length
ldr     w1, [x0]
adr     x0, width
ldr     w2, [x0]
add     w1, w1, w2
lsl     w1, w1, 1
adr     x0, perim
str     w1, [x0]
mov     w0, 0
ret
```

In C external linkage is default.
In ARM, internal linkage is default

main function is called by _start, so must be globally visible

Make a “Pointer”

```
static int length = 1;
static int width = 2;
static int perim = 0;

int main()
{
  perim =
   (length + width) * 2;
   return 0;
}
```

Generating addresses

adr: puts address of a label in a register

```
  .section .data
length: .word 1
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  .section .text
  .global main
main:
adr     x0, length
ldr     w1, [x0]
adr     x0, width
ldr     w2, [x0]
add     w1, w1, w2
lsl     w1, w1, 1
adr     x0, perim
str     w1, [x0]
mov     w0, 0
ret
```

Loads and Stores

```
static int length = 1;
static int width = 2;
static int perim = 0;

int main()
{
  perim =
   (length + width) * 2;
   return 0;
}
```

Load and store

Use x0 as a “pointer” to load from and store to memory

```
  .section .data
length: .word 1
width:  .word 2
perim:  .word 0
  .section .text
  .global main
main:
adr     x0, length
ldr     w1, [x0]
adr     x0, width
ldr     w2, [x0]
add     w1, w1, w2
lsl     w1, w1, 1
adr     x0, perim
str     w1, [x0]
mov     w0, 0
ret
```

Return

```
static int length = 1;
static int width = 2;
static int perim = 0;

int main()
{
  perim =
   (length + width) * 2;
   return 0;
}
```

```
  .section .data
length: .word 1
width:  .word 2
perim:  .word 0
  .section .text
  .global main
main:
adr     x0, length
ldr     w1, [x0]
adr     x0, width
ldr     w2, [x0]
add     w1, w1, w2
lsl     w1, w1, 1
adr     x0, perim
str     w1, [x0]
mov     w0, 0
ret
```

Return a value

ret: return to the caller, with register 0* holding the return value

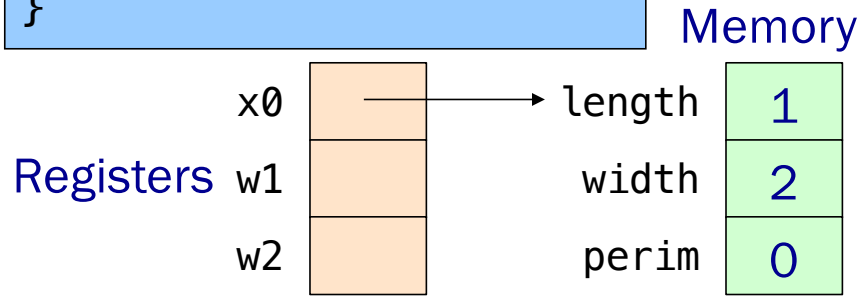
* w0 for int, x0 for long

Tracing Our Program

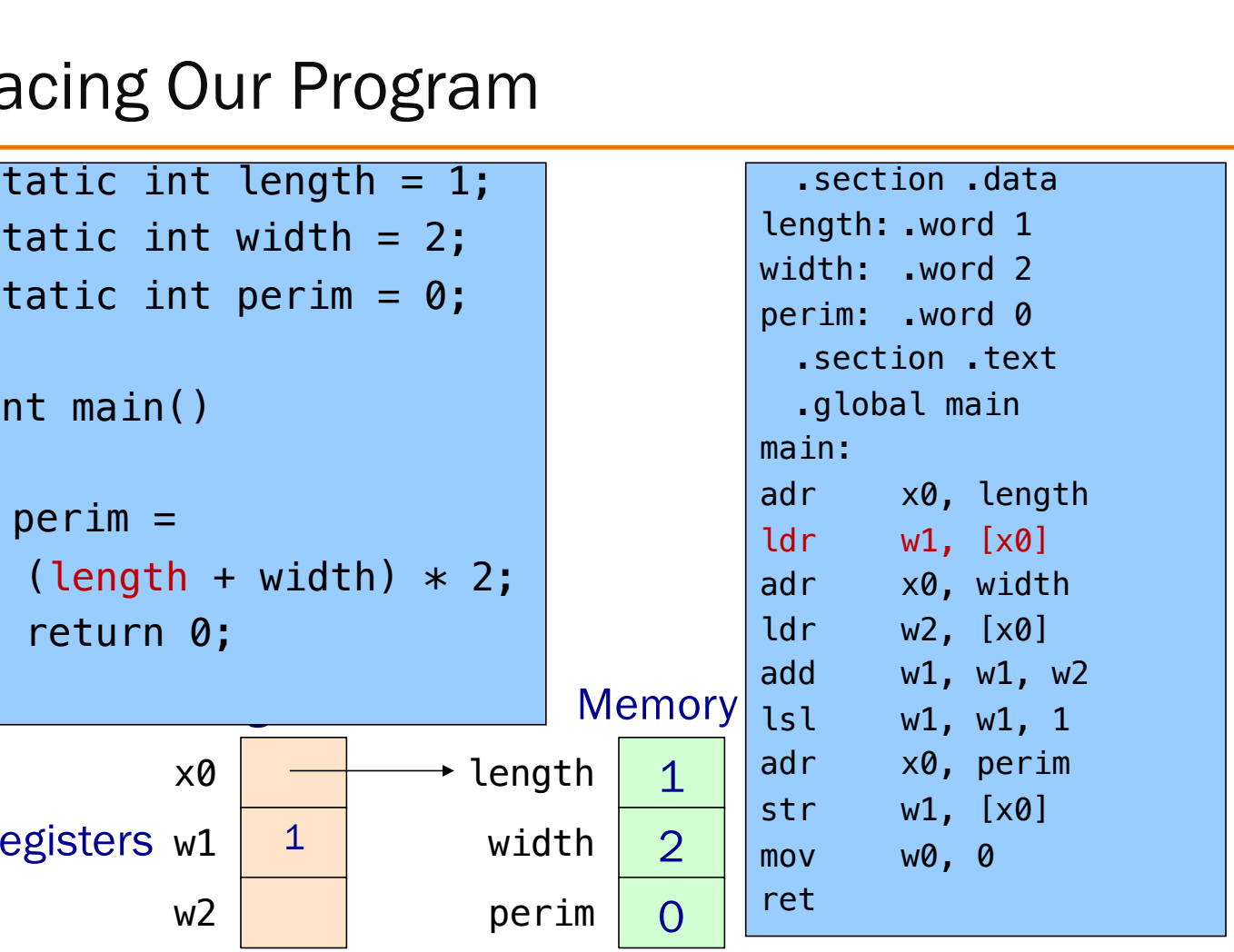
```
static int length = 1;
static int width = 2;
static int perim = 0;

int main()
{
  perim =
   (length + width) * 2;
   return 0;
}
```

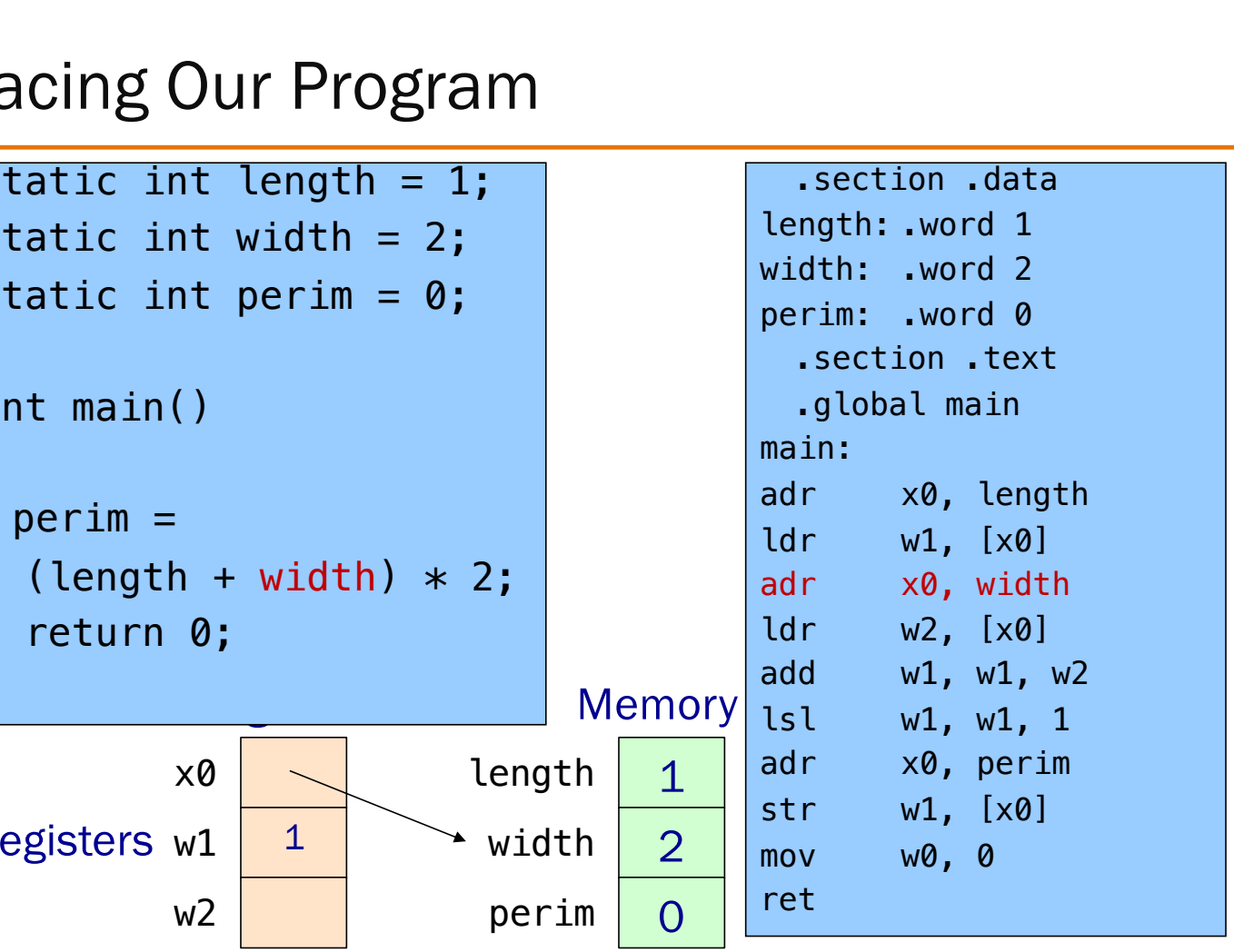
```
  .section .data
length: .word 1
width:  .word 2
perim:  .word 0
  .section .text
  .global main
main:
adr     x0, length
ldr     w1, [x0]
adr     x0, width
ldr     w2, [x0]
add     w1, w1, w2
lsl     w1, w1, 1
adr     x0, perim
str     w1, [x0]
mov     w0, 0
ret
```



Tracing Our Program



Tracing Our Program



Tracing Our Program

```
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static int perim = 0;

int main()
{
  perim =
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   return 0;
}
```

```
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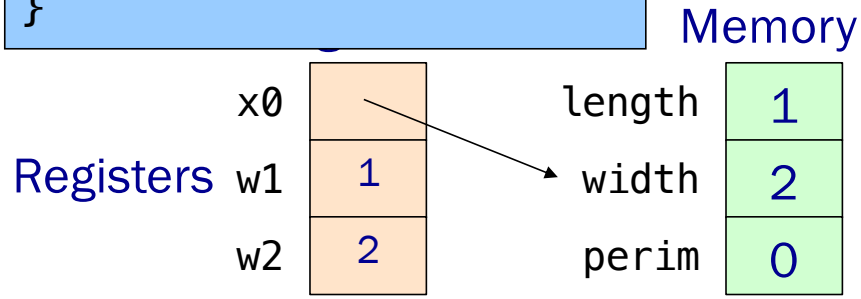
width:  .word 2
perim:  .word 0
  .section .text
  .global main
main:
adr     x0, length
ldr     w1, [x0]
adr     x0, width
ldr     w2, [x0]
add     w1, w1, w2
lsl     w1, w1, 1
adr     x0, perim
str     w1, [x0]
mov     w0, 0
ret
```



Tracing Our Program

```
static int length = 1;
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static int perim = 0;

int main()
{
  perim =
   (length + width) * 2;
   return 0;
}
```

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  .section .data
length: .word 1
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perim:  .word 0
  .section .text
  .global main
main:
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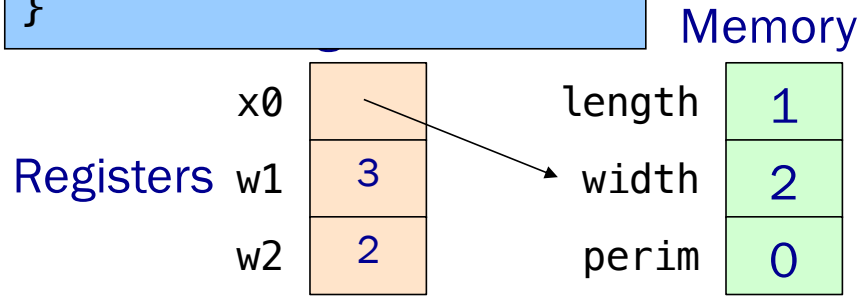
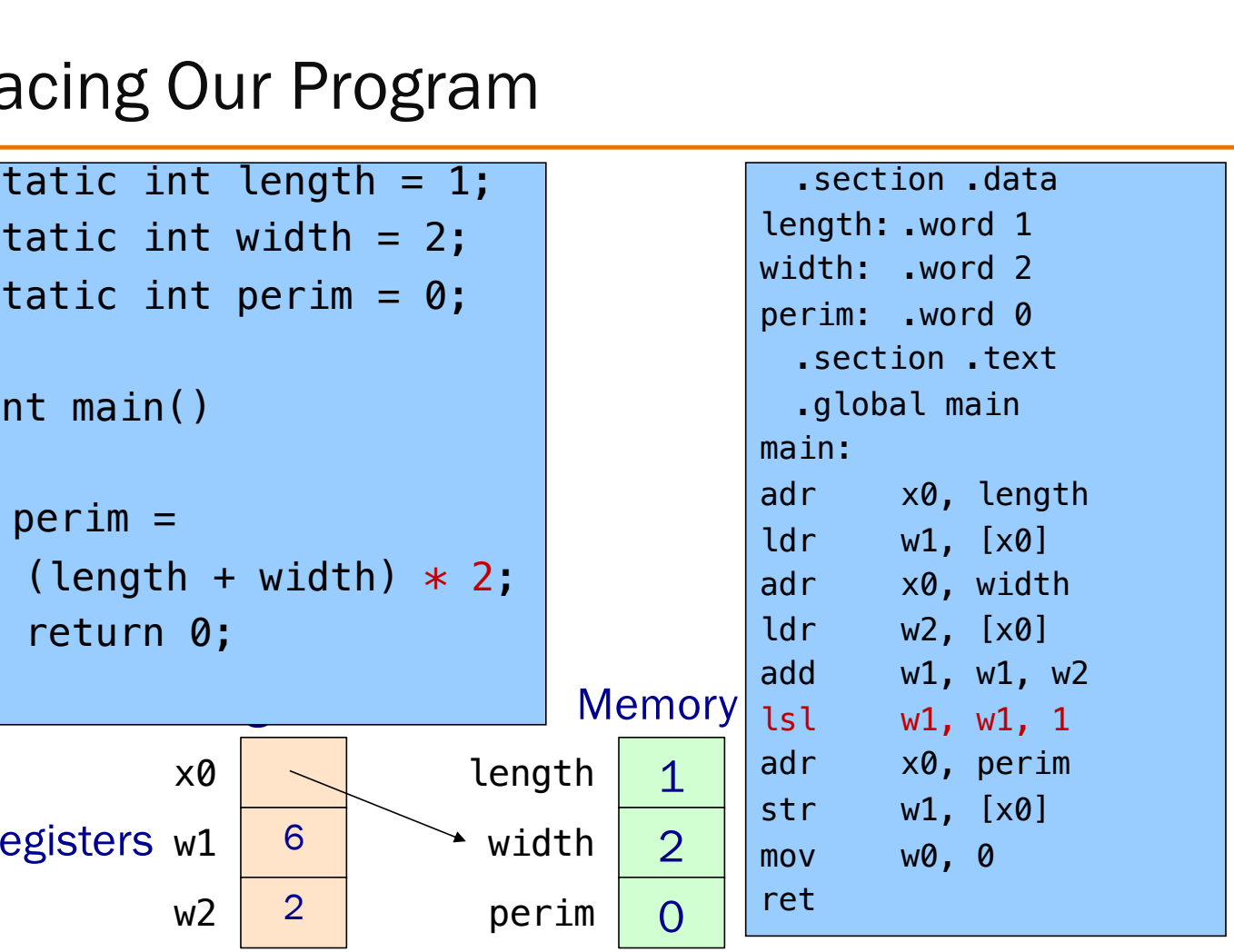
ldr     w1, [x0]
adr     x0, width
ldr     w2, [x0]
add     w1, w1, w2
lsl     w1, w1, 1
adr     x0, perim
str     w1, [x0]
mov     w0, 0
ret
```



Tracing Our Program

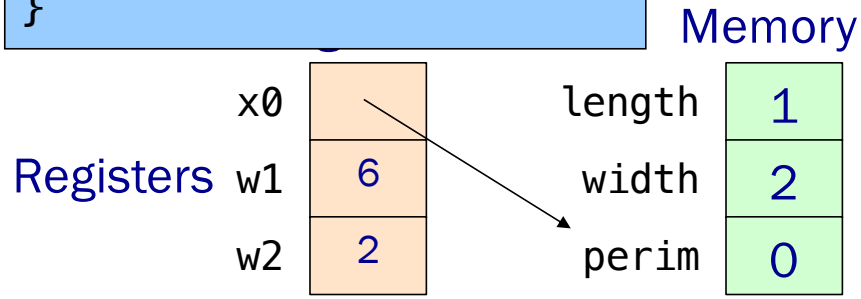


Tracing Our Program

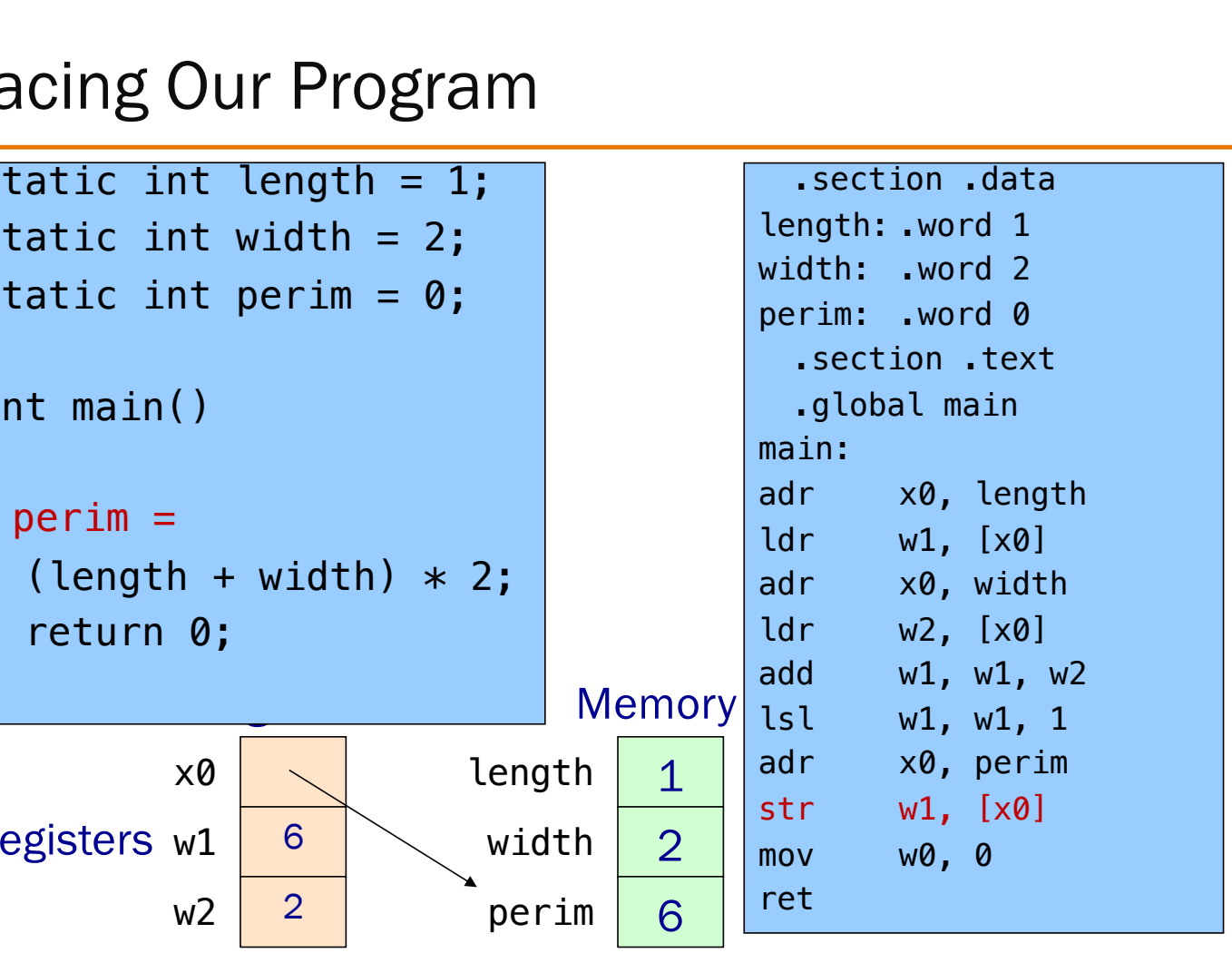
```
static int length = 1;
static int width = 2;
static int perim = 0;

int main()
{
  perim =
   (length + width) * 2;
   return 0;
}
```

```
  .section .data
length: .word 1
width:  .word 2
perim:  .word 0
  .section .text
  .global main
main:
adr     x0, length
ldr     w1, [x0]
adr     x0, width
ldr     w2, [x0]
add     w1, w1, w2
lsl     w1, w1, 1
adr     x0, perim
str     w1, [x0]
mov     w0, 0
ret
```



Tracing Our Program



Tracing Our Program

```
static int length = 1;
static int width = 2;
static int perim = 0;

int main()
{
  perim =
   (length + width) * 2;
   return 0;
}
```

```
  .section .data
length: .word 1
width:  .word 2
perim:  .word 0
  .section .text
  .global main
main:
adr     x0, length
ldr     w1, [x0]
adr     x0, width
ldr     w2, [x0]
add     w1, w1, w2
lsl     w1, w1, 1
adr     x0, perim
str     w1, [x0]
mov     w0, 0
ret
```

Return value

Passed back in register w0

Tracing Our Program

```
static int length = 1;
static int width = 2;
static int perim = 0;

int main()
{
  perim =
   (length + width) * 2;
   return 0;
}
```

Return to caller

ret instruction

```
  .section .data
length: .word 1
width:  .word 2
perim:  .word 0
  .section .text
  .global main
main:
adr     x0, length
ldr     w1, [x0]
adr     x0, width
ldr     w2, [x0]
add     w1, w1, w2
lsl     w1, w1, 1
adr     x0, perim
str     w1, [x0]
mov     w0, 0
ret
```

Summary

Language levels: High-level language, assembly language, machine language

The basics of computer architecture

- Enough to understand AARCH64 assembly language

The basics of AARCH64 assembly language

- Instructions to perform arithmetic
- Instructions to define global data and perform data transfer

To learn more

- Study more assembly language examples
  - Chapters 2-5 of Pyeatt and Ughetta book
- Study compiler-generated assembly language code (though it will be challenging!)
  - `gcc217 –S somefile.c`
- We will have three more lectures and four precepts on AARCH64

Appendix 1

DEFINING DATA: OTHER SECTIONS AND SIZES

Defining Data: DATA Section 1

```
static char c = 'a';
static short s = 12;
static int i = 345;
static long l = 6789;
```

```
    .section ".data"
c:
    .byte 'a'
s:
    .short 12
i:
    .word 345
l:
    .quad 6789
```

Notes:

- `.section` directive
  (to announce DATA section)
- label definition
  (marks a spot in RAM)
- `.byte` directive (1 byte)
- `.short` directive (2 bytes)
- `.word` directive (4 bytes)
- `.quad` directive (8 bytes)

Defining Data: DATA Section 2

```
char c = 'a';
short s = 12;
int i = 345;
long l = 6789;
```

```
    .section ".data"
    .global c
c:  .byte 'a'
    .global s
s:  .short 12
    .global i
i:  .word 345
    .global l
l:  .quad 6789
```

Notes:

Can place label on same line as next instruction or directive

`.global` directive can also apply to variables, not just functions

Defining Data: BSS Section

```
static char c;
static short s;
static int i;
static long l;
```

```
   .section ".bss"
c:
   .skip 1
s:
   .skip 2
i:
   .skip 4
l:
   .skip 8
```

Notes:

- `.section` directive
  (to announce BSS section)
- `.skip` directive
  (to specify number of bytes)

Defining Data: RODATA Section

```
…
…"hello\n"…;
…
```

```
   .section ".rodata"
helloLabel:
   .string "hello\n"
```

Notes:

- `.section` directive (to announce RODATA section)
- `.string` directive

Appendix 2

BYTE ORDER: BIG-ENDIAN VS LITTLE-ENDIAN

Byte Order

AARCH64 is a **little endian** architecture

- **Least** significant byte of multi-byte entity is stored at lowest memory address
- "Little end goes first"

The int (four bytes) 5 at address 1000 :

1000	00000101
1001	00000000
1002	00000000
1003	00000000

Some other systems use **big endian**

- **Most** significant byte of multi-byte entity is stored at lowest memory address
- "Big end goes first"

The int (four bytes) 5 at address 1000:

1000	00000000
1001	00000000
1002	00000000
1003	00000101

Byte Order Example 1

```
#include <stdio.h>
int main(void)
{  unsigned int i = 0x003377ff;
   unsigned char *p;
   int j;
   p = (unsigned char *)&i;
   for (j = 0; j < 4; j++)
      printf("Byte %d: %2x\n", j, p[j]);
}
```

Output on a little-endian machine

```
Byte 0: ff
Byte 1: 77
Byte 2: 33
Byte 3: 00
```

Output on a big-endian machine

```
Byte 0: 00
Byte 1: 33
Byte 2: 77
Byte 3: ff
```

Byte Order Example 2

Note:

Flawed code; uses “b” instructions to load from a four-byte memory area

```
    .section ".data"
foo: .word 7
    .section ".text"
    .global “main”
main:
adr   x0, foo
ldrb  w0, [x0]
ret
```

AARCH64 is little endian, so what will be the value returned from w0?

What would be the value returned from w0 if AARCH64 were big endian?