

Precept Outline

- Review of Lectures 3 and 4:
 - Stacks and Queues
 - Advanced Java
- Stacks: resizing arrays + linked lists

Relevant Book Sections

- 1.3 (Queues and Stacks)

A. Review: Stacks, Queues and Iterators/Iterables

Your preceptor will briefly review key points of this week's lectures.

Here are some code snippets that your instructor might refer to as examples:

```
1 Stack<String> stack = new Stack<String>();
2
3 stack.push("One");
4 stack.push("Two");
5 stack.push("Three");
6 stack.push("Four");
7 stack.push("Five");
8
9 for (i = 0; i < 5; i++)
0     StdOut.println(stack.pop());
```

```
1 Queue<String> queue = new Queue<String>();
2
3 queue.enqueue("One");
4 queue.enqueue("Two");
5 queue.enqueue("Three");
6 queue.enqueue("Four");
7 queue.enqueue("Five");
8
9 for (i = 0; i < 5; i++)
0     StdOut.println(queue.dequeue());
```

```
1 public class YourClass<Item> implements Iterable<Item> {
2
3     public Iterator<Item> iterator() {
4         return new YourClassIterator();
5     }
6
7     private class YourClassIterator implements Iterator<Item> {
8         // instance variable(s) to keep track of where iterator is
9
10        public boolean hasNext() {
11            // condition to end iteration
12        }
13
14        public Item next() {
15            // returns next item and updates instance variable(s)
16        }
17    }
18 }
```

```
1 Stack<String> stack = new Stack<String>();
2 // initialize stack
3
4 Iterator<String> iter = stack.iterator();
5
6 while (iter.hasNext()) {
7     String s = iter.next();
8     // do something with s
9 }
```

```
1 Stack<String> stack = new Stack<String>();
2 // initialize stack
3
4
5
6 for (String s : stack) {
7
8     // do something with s
9 }
```

B. Stacks and Queues

Part 1: Resizing arrays

In lecture, you saw how the *repeated doubling* strategy solves the problem of resizing arrays too often. There was a caveat, however: we resize up at 100% capacity but resize down at 25% (rather than 50%) capacity.

(Warm-up) Recall what goes wrong if we resize down at 50%: give an example of a sequence of `push()` and `pop()` operations with $\Theta(n)$ amortized cost. The cost of a sequence of operations (as in lecture) is the total number of array accesses made throughout their execution.

Consider the following “resizing policies”:

1. Double at 100% capacity, halve at 25%;
2. Triple at 100% capacity, multiply by 1/3 at 1/3;
3. Triple at 100% capacity, multiply by 2/3 at 1/3;
4. Double at 75% capacity, halve at 25%.

Identify which policies have worst-case $\Theta(n)$ amortized cost for n operations and which have $\Theta(1)$.

Part 2: Linked Lists

Recall that in a singly linked list, each node stores an item (of generic type) and a reference to the next node in the list. Describe a method that produces a new linked list with the same elements of `list` but in *reverse order*.

Assume that the input list is given as its `first` node. You can create extra nodes or linked lists, but not modify the input list. Feel free to write code or pseudocode.

(c) Assume that the Stack data type is implemented as a *resizable array*. How many times would the array *shrink* when calling `parseUndos(str)`, where `str` is a string that consists of $16n$ non-undo characters followed by $13n$ undo characters?

Assume that the stack is 100% full after `parseUndos` processes the first $16n$ non-undo characters, and recall that `pop()` resizes the array (to half of its size) when it reaches 25% capacity.

- 0.
 - 1.
 - 2.
 - Constant strictly greater than 2.
 - $\sim \frac{1}{4} \log n$.
 - $\sim \log n$.
- (d) When the stack is implemented as a linked list, the *worst-case* running time of `parseUndos()` on a string of length n is $\Theta(n)$.
- True.
 - False.
- (e) When the stack is implemented as a resizable array, the *worst-case* running time of `parseUndos()` on a string of length n is $\Theta(n)$.
- True.
 - False.

C. Assignment Overview: Queues

Your preceptor will introduce and give an overview of your [second assignment](#). Please don't hesitate to ask questions!

Summary of the assignment.

- Implement a *deque* data structure, which is a queue that supports insertion and removal of items from both ends. This will be done in `Deque.java`.
- Implement a randomized queue, which is a queue that differs from a typical queue in that items are removed uniformly at random, not based on the sequence they were added. This will be done in `RandomizedQueue.java`.
- Implement an application of the randomized queue to read a sequence of strings and print a subset of them uniformly at random. This will be done in `Permutation.java`.
- Both deque and randomized queue require iterators that support operations in constant worst-case time and use space efficiently.

D. Optional Bonus Problems

Part 1: Obstructed Skyline (Challenge)

Suppose you are given the shape of a city's skyline in the form of a length- n array $h = [h_0, h_1, \dots, h_{n-1}]$ (where the height of building i is h_i). In other words, the skyline is a $1 \times n$ grid where the i th column/building has height $h_i \leq k$.

Design an algorithm to find the rectangle with the largest area that is blocked by the skyline. (E.g., your algorithm should output 2 when $h = [2, 1]$, 4 when $h = [2, 2]$ and 12 with the input drawn below.) It should run in $\Theta(n)$ time and space.

