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1.4 ANALYSIS OF ALGORITHMS

- *introduction*
- *running time (experimental analysis)*
- *running time (mathematical models)*
- *memory usage*

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1.4 Analysis of Algorithms

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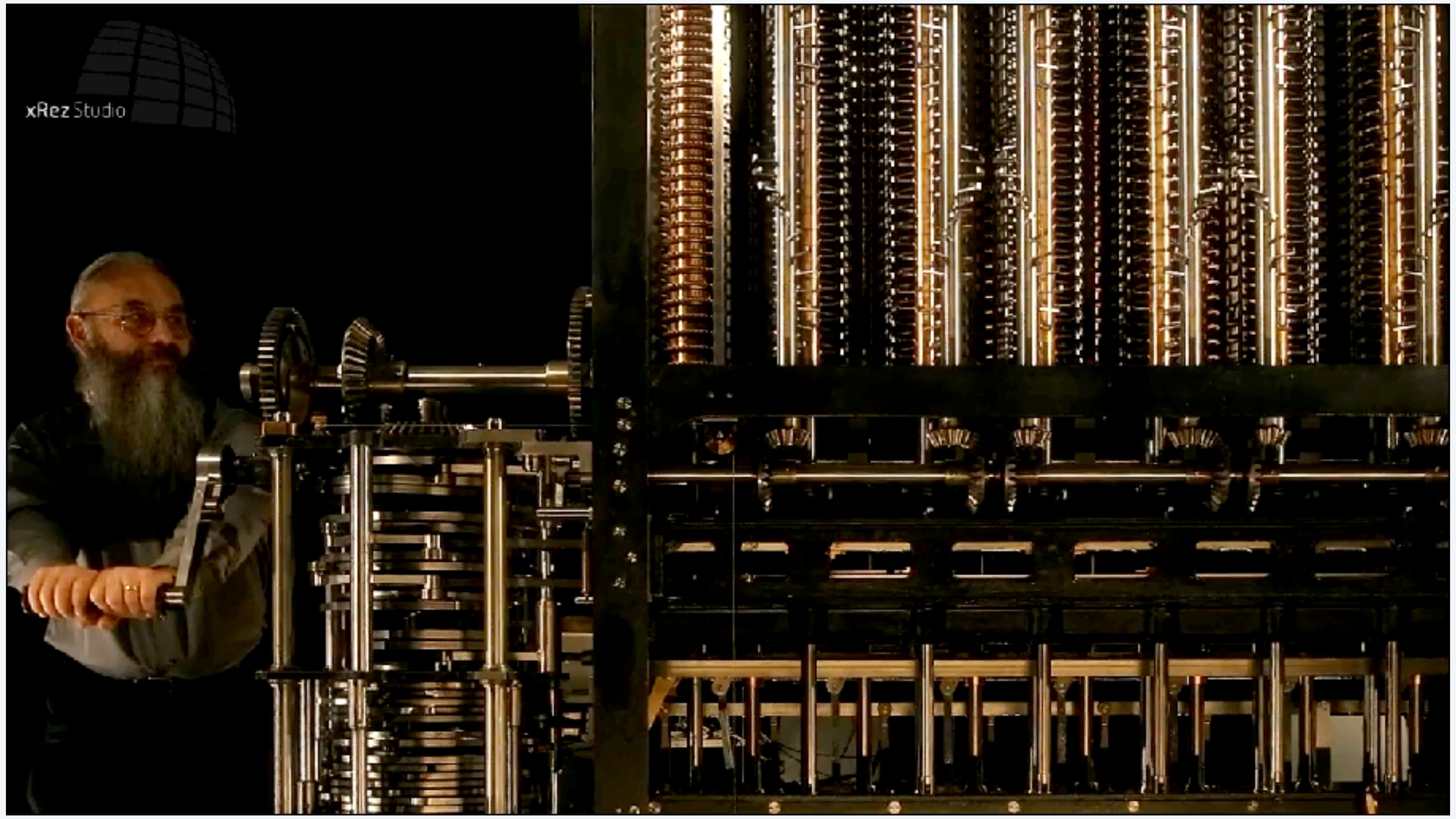
Running time

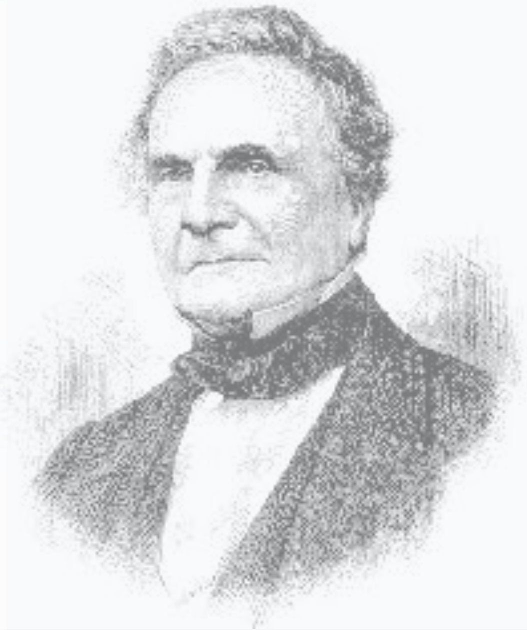
" *As soon as an Analytical Engine exists, it will necessarily guide the future course of the science. Whenever any result is sought by its aid, the question will then arise—By what course of calculation can these results be arrived at by the machine in the shortest time?* " — Charles Babbage (1864)

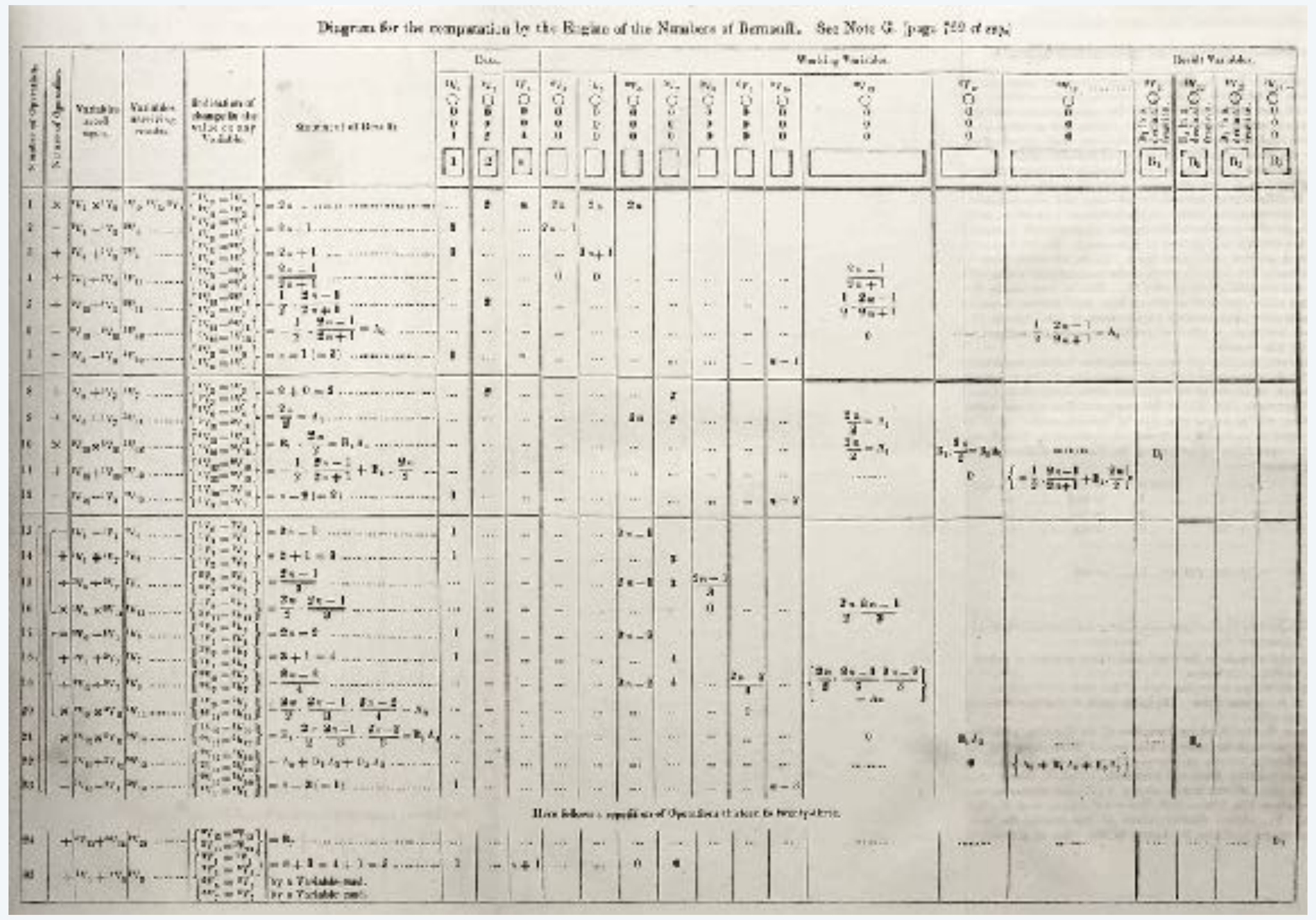
how many times do you have to turn the crank?



https://vimeo.com/49080293

“ *As soon as an Analytical Engine exists, it will necessarily guide the future course of the science. Whenever any result is sought by its aid, the question will then arise—By what course of calculation can these results be arrived at by the machine in the shortest time* ? ” — Charles Babbage (1864)





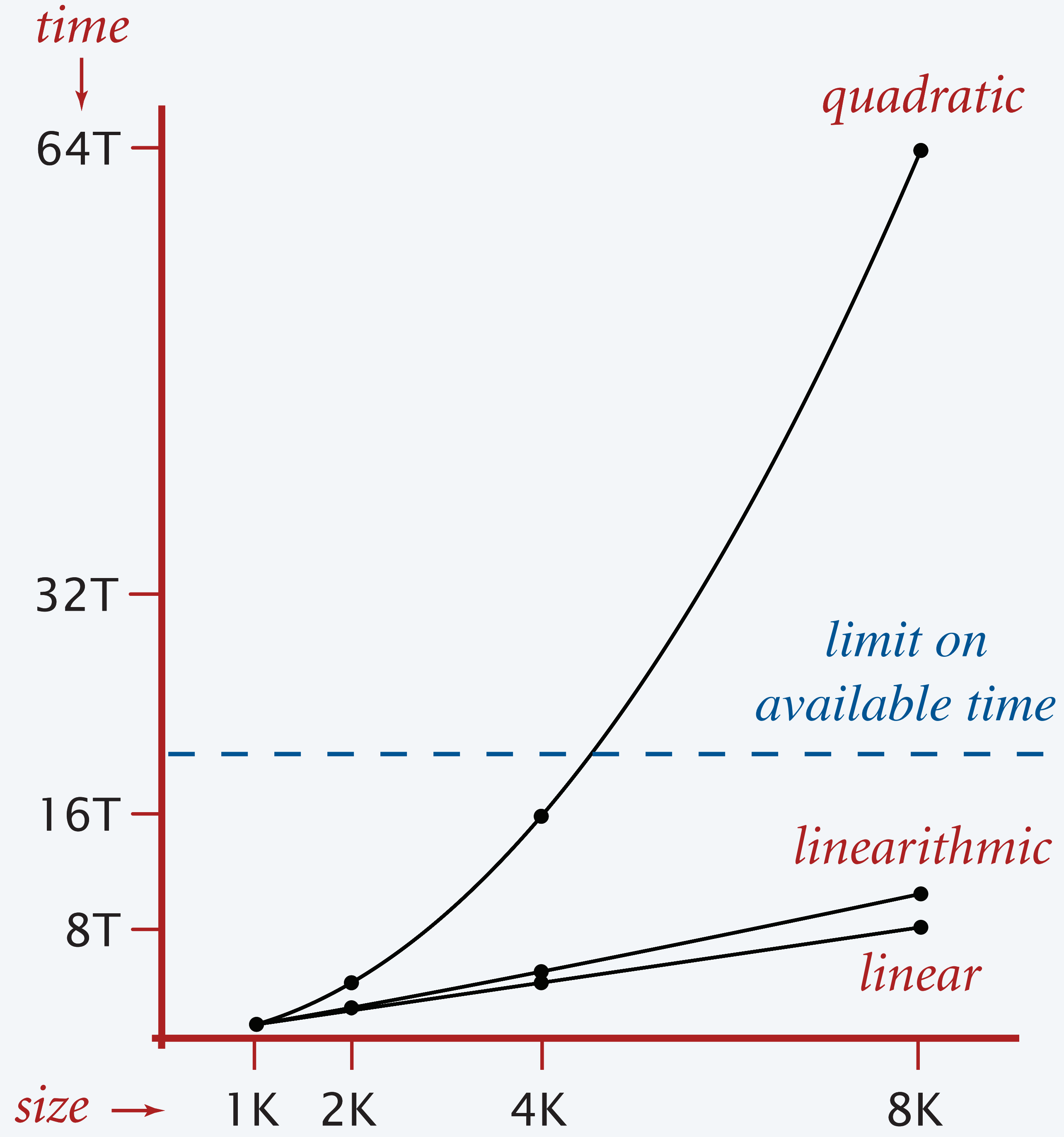
Ada Lovelace’s algorithm
to compute Bernoulli numbers
on Analytic Engine (1843)

An algorithmic success story

N-body simulation.

- Simulate gravitational interactions among n bodies.
- Applications: cosmology, fluid dynamics, semiconductors, …
- Brute force: $\Theta(n^2)$ steps.
- Barnes-Hut algorithm: $\Theta(n \log n)$ steps, enables new research.

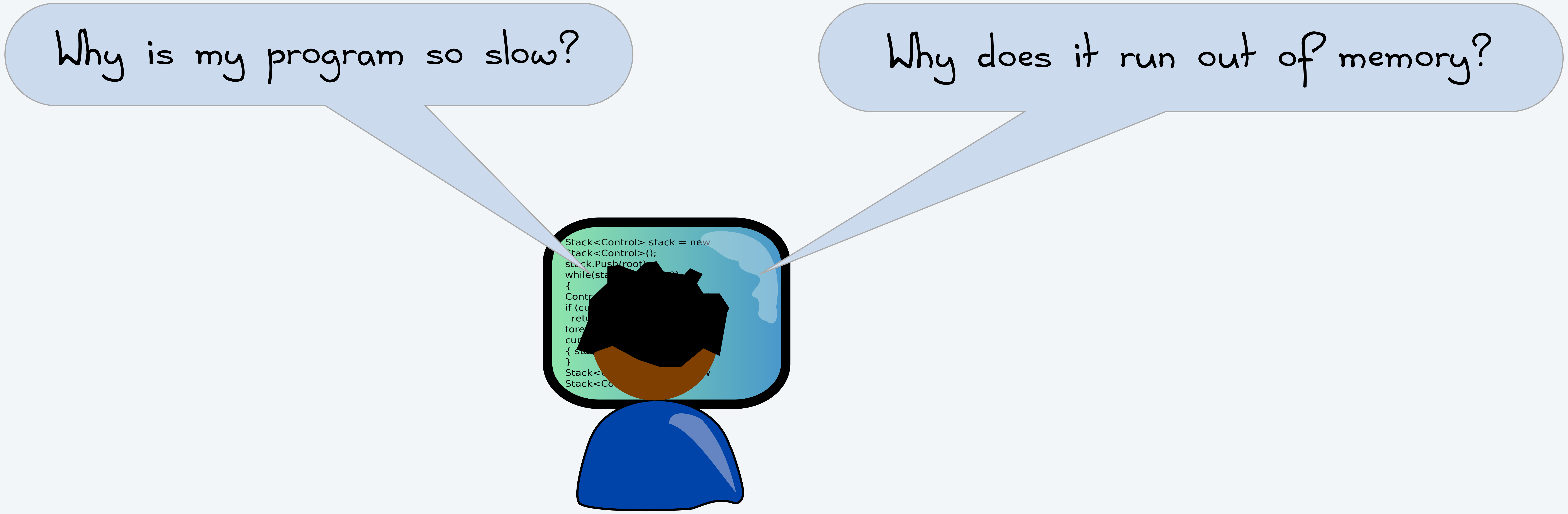
Andrew Appel
PU ’81



The challenge

Q1. Will my program be able to solve a large practical input?

Q2. If not, how might I understand its performance characteristics so as to improve it?



Our approach. Combination of experiments and mathematical modeling.

Example: 3-Sum

3-Sum. Given n distinct integers, how many triples sum to exactly zero?

```
~/cos226/3sum> more 8ints.txt
8
30 -40 -20 -10 40 0 10 5

~/cos226/3sum> java ThreeSum 8ints.txt
4
```

	a[i]	a[j]	a[k]	sum	
1	30	–40	10	0	✔
2	30	–20	–10	0	✔
3	–40	40	0	0	✔
4	–10	0	10	0	✔

Context. Connected with problems in computational geometry (computer games!)

Open! What is the running time of the optimal algorithm for 3-Sum ?

3-SUM: brute-force algorithm

```
public class ThreeSum {

   public static int count(int[] a) {
      int n = a.length;
      int count = 0;
      for (int i = 0; i < n; i++)
         for (int j = i+1; j < n; j++)
            for (int k = j+1; k < n; k++)
               if (a[i] + a[j] + a[k] == 0)
                  count++;
      return count;
   }

   public static void main(String[] args) {
      In in = new In(args[0]);
      int[] a = in.readAllInts();
      StdOut.println(count(a));
   }
}
```

check distinct triples

for simplicity,
ignore integer overflow

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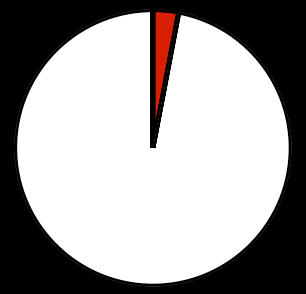
Measuring the running time

Running time. Run the program for inputs of varying size; measure running time.

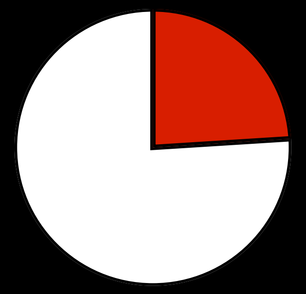
Observation. The running time $T(n)$ grows as a function of the input size n.

```
~/cos226/analysis> java ThreeSum 1Kints.txt
70
```



tick tick

```
~/cos226/analysis> java ThreeSum 2Kints.txt
528
```



tick tick tick tick tick tick tick tick
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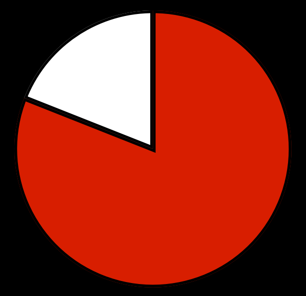
```
~/cos226/analysis> java ThreeSum 3Kints.txt
1670
```



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tick tick tick tick tick tick

Measuring the running time

Running time. Run the program for inputs of varying size; measure running time.

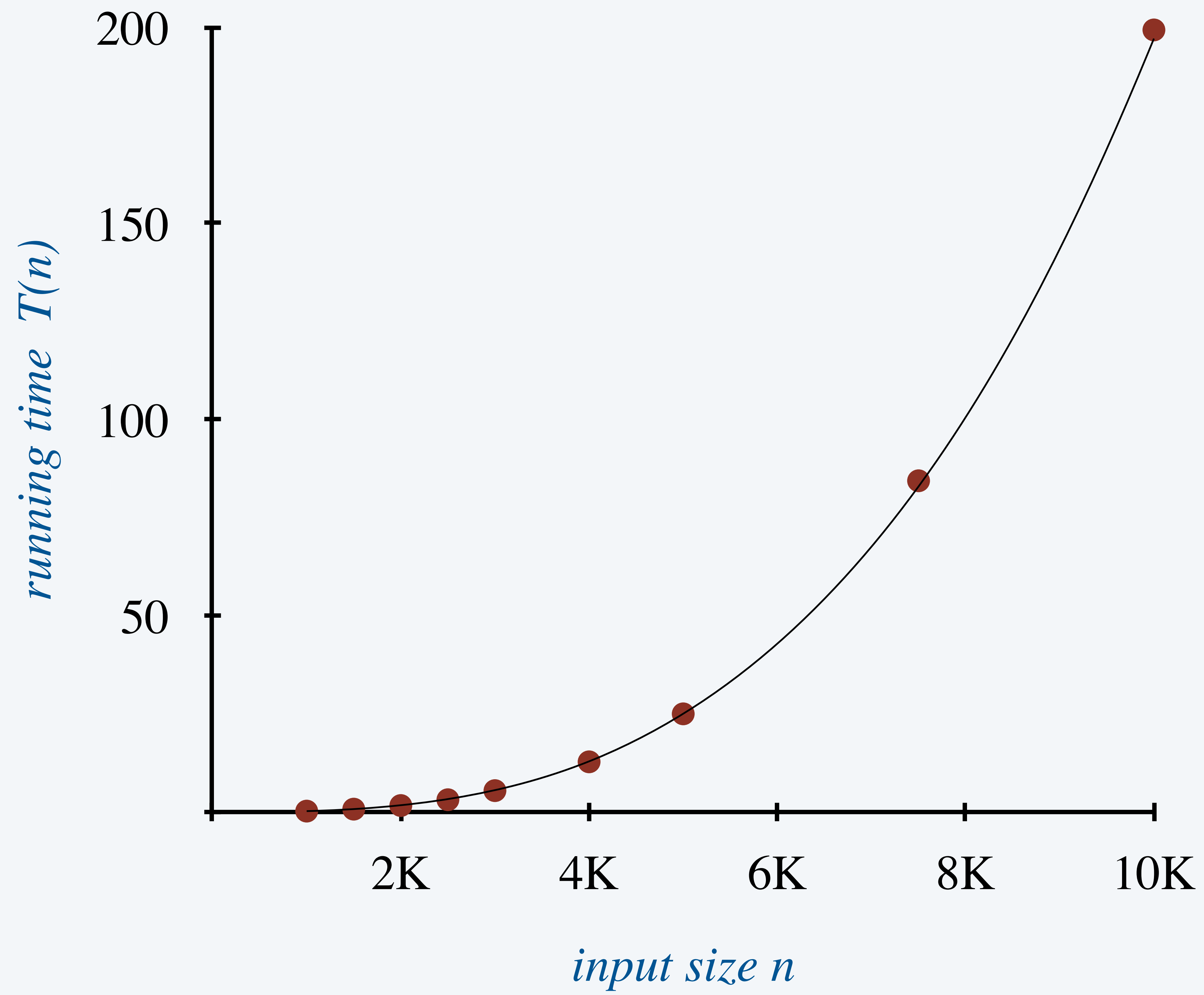
n	time (seconds) †
1,000	0.21
1,500	0.71
2,000	1.63
2,500	3.11
3,000	5.43
4,000	12.8
5,000	25.0
7,500	84.4
10,000	199.3

† *Apple M2 Pro with 32 GB memory running OpenJDK 11 on MacOS Ventura*

Data analysis: standard plot

Standard plot. Plot running time $T(n)$ vs. input size n.

n	time (seconds) †
1,000	0.21
1,500	0.71
2,000	1.63
2,500	3.11
3,000	5.43
4,000	12.8
5,000	25.0
7,500	84.4
10,000	199.3



Hypothesis. The running time obeys a power law: $T(n) = a \times n^b$ seconds.

Questions. How to validate hypothesis? How to estimate constants a and b ?

Doubling test: estimating the exponent b

Doubling test. Run program, doubling the size of the input.

- Assume running time satisfies the "power law" $T(n) = a \times n^b$.
- Estimate $b = \log_2$ ratio.

n	time (seconds) †	ratio	$\log_2$ ratio
500	0.05	–	–
1,000	0.21	4.20	2.07
2,000	1.63	7.76	2.96
4,000	12.8	7.85	2.97
8,000	103.1	8.05	3.01
16,000	819.0	7.94	2.99

$\log_2 (103.1 / 12.8) = 3.01$

seems to converge to a constant $b \approx 3.0$

$$\frac{T(n)}{T(n/2)} = \frac{an^b}{a(n/2)^b} = 2^b$$

$$\implies b = \log_2 \frac{T(n)}{T(n/2)}$$

why the $\log_2$ ratio works

Doubling test: estimating the leading coefficient a

Doubling test. Run program, doubling the size of the input.

- Assume running time satisfies $T(n) = a \times n^b$.
- Estimate $b = \log_2$ ratio.
- Estimate a by solving $T(n) = a \times n^b$ for a sufficiently large value of n.

n	time (seconds) †	ratio	$\log_2$ ratio	
500	0.05	–	–	
1,000	0.21	4.20	2.07	
2,000	1.63	7.76	2.96	
4,000	12.8	7.85	2.97	
8,000	103.1	8.05	3.01	
16,000	819.0	7.94	2.99	$819.0 = a \times 16{,}000^3 \Rightarrow a = 2.00 \times 10^{-10}$

Hypothesis. Running time is about $2.00 \times 10^{-10} \times n^3$ seconds.

Estimate the running time to solve a problem of size $n = 96{,}000$.

A. 39 seconds

B. 52 seconds

C. 117 seconds

D. 350 seconds

n	time (seconds)
1,000	0.02
2,000	0.05
4,000	0.20
8,000	0.81
16,000	3.25
32,000	13.01

Experimental algorithmics

System independent effects.

- Algorithm.
- Input data.

determines exponent b in power law $T(n) = a \times n^b$

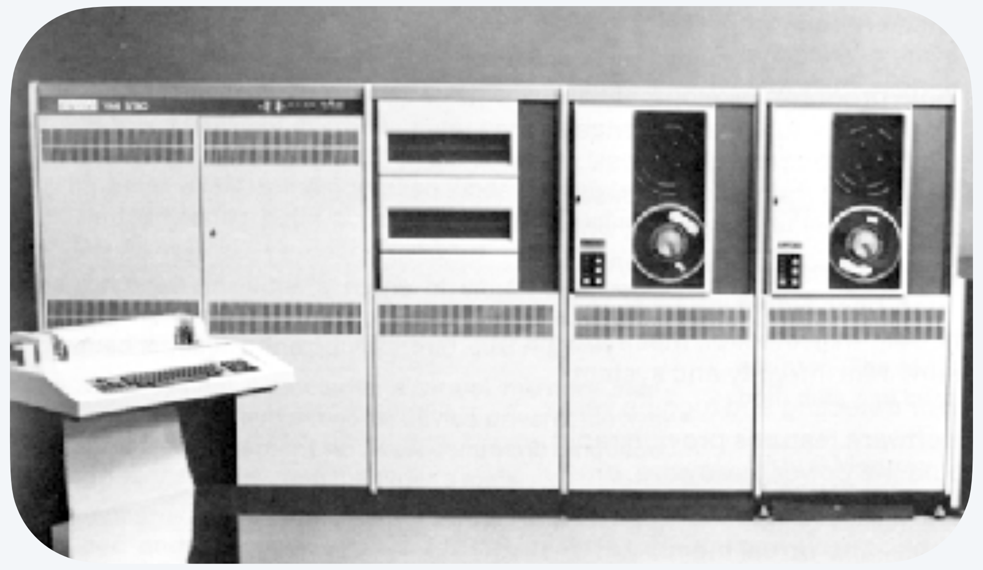
System dependent effects.

- Hardware: CPU, memory, cache, …
- Software: compiler, interpreter, garbage collector, …
- System: operating system, network, other apps, …

determines leading coefficient a in power law $T(n) = a \times n^b$

inferring running time?

1970s



10,000× *faster*

2020s
(Macbook Pro M2)

Bad news. Sometimes difficult to get accurate measurements.

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Mathematical models for running time

Total running time: sum of frequency × cost for all operations.

- Frequency depends on algorithm and input data.
- Cost depends on CPU, compiler, operating system, …

Warning. No general-purpose method (e.g., halting problem).

Example: 1-SUM

Q. How many operations as a function of input size n ?

```
int count = 0;
for (int i = 0; i < n; i++)
   if (a[i] == 0)
      count++;
```

operation	cost (ns) †	frequency
variable declaration	2/5	2
assignment statement	1/5	2
less than compare	1/5	$n + 1$
equal to compare	1/10	n
array access	1/10	n
increment	1/10	n to $2\,n$

in practice, depends on caching, bounds checking, … (*see* COS 217)

tedious to count exactly

† *representative estimates (with some poetic license)*

Simplification 1: cost model

Cost model. Use some elementary operation as a proxy for running time.

array accesses, compares, API calls, floating-point operations, …

```
int count = 0;
for (int i = 0; i < n; i++)
   if (a[i] == 0)
      count++;
```

"inner loop"

operation	cost (ns) †	frequency
variable declaration	2/5	2
assignment statement	1/5	2
less than compare	1/5	$n + 1$
equal to compare	1/10	n
array access	1/10	n
increment	1/10	n to $2\,n$

cost model = array accesses

Simplification 2: asymptotic notations

Tilde notation. Discard lower-order terms.

Big Theta notation. Discard lower-order terms and leading coefficient.

← *rigorous definitions involve limits*

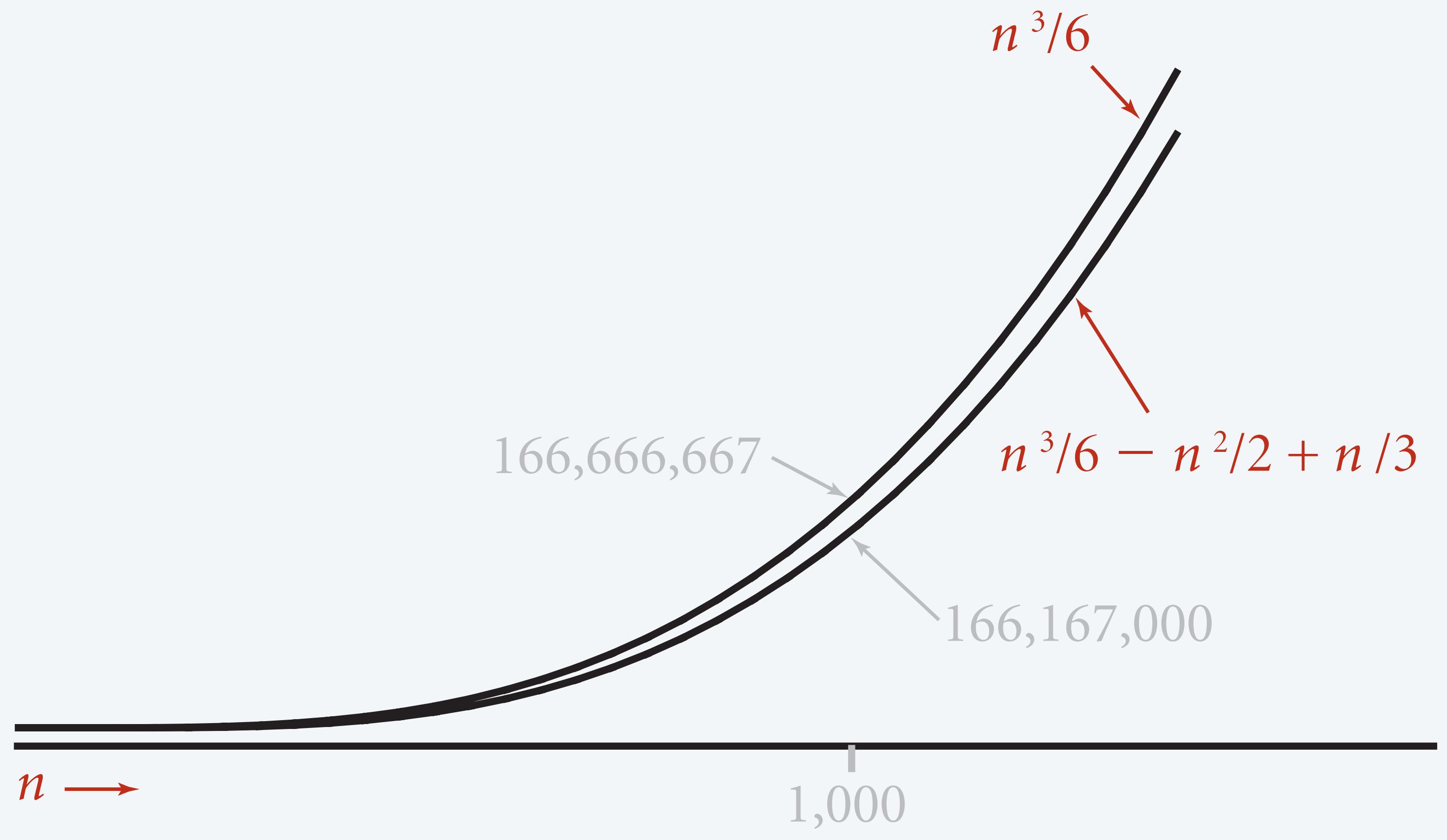
function	tilde notation	big Theta
$4\,n^5 + 20\,n^3 + 1600$	$\sim 4\,n^5$	$\Theta(n^5)$
$0.01\,n^2 + 100\,n^{4/3} - 56$	$\sim 0.01\,n^2$	$\Theta(n^2)$
$8\log^2 n + 7\,n$	$\sim 7\,n$	$\Theta(n)$
$10\,n + 3\,n\log n$	$\sim 3\,n\log n$	$\Theta(n\log n)$
$2^n + n^5$	$\sim 2^n$	$\Theta(2^n)$

← *"order of growth"*

$$\tfrac{1}{6}\,n^3 \underbrace{-\ \tfrac{1}{2}\,n^2 + \tfrac{1}{3}\,n}_{}$$

discard lower-order terms

(*e.g.*, n = 1,000: 166.667 *million* vs. 166.167 *million*)



Leading-term approximation

Rationale.

- When n is large, lower-order terms are negligible.
- When n is small, we don't care.

Which of the following correctly describes the function $f(n) = n \log n + 0.6\, n^2 + 10\, n$?

A. $\sim 10\, n$

B. $\sim n \log n$

C. $\sim n^2$

D. $\Theta(n \log n)$

E. $\Theta(n^2)$

How many array accesses as a function of n?

```
int count = 0;
for (int i = 0; i < n; i++)
   for (int j = i+1; j < n; j++)
      if (a[i] + a[j] == 0)
         count++;
```

"inner loop"

A. $\frac{1}{2}\, n\, (n - 1)$

B. $n\, (n - 1)$

C. $2\, n^2$

D. $2\, n\, (n - 1)$

Example: two-sum

Q. How many operations as a function of input size n ?

```
int count = 0;
for (int i = 0; i < n; i++)
   for (int j = i+1; j < n; j++)
      if (a[i] + a[j] == 0)
         count++;
```

$$\underbrace{\overset{\substack{i=0\\ j=1,\,\ldots,\,n-1\\ \downarrow}}{(n-1)} + \overset{\substack{i=1\\ j=2,\,\ldots,\,n-1\\ \downarrow}}{(n-2)} + \overset{\substack{i=2\\ j=3,\,\ldots,\,n-1\\ \downarrow}}{(n-3)} + \ldots + \overset{\substack{i=n-3\\ j=n-2,\,n-1\\ \downarrow}}{2} + \overset{\substack{i=n-2\\ j=n-1\\ \downarrow}}{1} + \overset{\substack{i=n-1\\ \text{no } j\\ \downarrow}}{0}}_{\frac{1}{2}\,n\,(n-1)}$$

Proof:

$$(n-1) + (n-2) + (n-3) + \ldots + 2 + 1 + 0 \;=\; (n-1)\,n/2$$

Loops analysis:

If loops are independent, analyze separately and multiply.

Else, write a sum and use a formula to simplify.

Example: two-sum

Q. How many operations as a function of input size n ?

```
int count = 0;
for (int i = 0; i < n; i++)
   for (int j = i+1; j < n; j++)
      if (a[i] + a[j] == 0)
         count++;
```

$$\underbrace{(n-1) + (n-2) + (n-3) \ldots + 2 + 1 + 0}_{\frac{1}{2}\, n\, (n-1)}$$

operation	cost (ns) †	frequency
variable declaration	2/5	$n + 2$
assignment statement	1/5	$n + 2$
less than compare	1/5	$\frac{1}{2}\,(n+1)\,(n+2)$
equal to compare	1/10	$\frac{1}{2}\,n\,(n-1)$
array access	1/10	$n\,(n-1)$
increment	1/10	$\frac{1}{2}\,n\,(n+1)$ to n^2

tedious to count exactly

$1/4\, n^2 + 13/20\, n + 13/10$ ns

to

$3/10\, n^2 + 3/5\, n + 13/10$ ns

Example: 2-SUM

Q. Approximately how many operations as a function of input size n ?

```
int count = 0;
for (int i = 0; i < n; i++)
   for (int j = i+1; j < n; j++)
      if (a[i] + a[j] == 0)
         count++;
```

$$\underbrace{(n-1) + (n-2) + (n-3) \ldots + 2 + 1 + 0}_{\frac{1}{2}\, n\,(n-1)}$$

operation	cost (ns) †	frequency
variable declaration	2/5	$\Theta(n)$
assignment statement	1/5	$\Theta(n)$
less than compare	1/5	$\Theta(n^2)$
equal to compare	1/10	$\Theta(n^2)$
array access	1/10	$\Theta(n^2)$
increment	1/10	$\Theta(n^2)$

Example: 3-SUM

Q. Approximately how many array accesses as a function of input size n ?

```
int count = 0;
for (int i = 0; i < n; i++)
   for (int j = i+1; j < n; j++)
      for (int k = j+1; k < n; k++)
         if (a[i] + a[j] + a[k] == 0)
            count++;
```

see COS 240

$$\binom{n}{3} = \frac{n(n-1)(n-2)}{3!} \sim \frac{1}{6}n^3$$

in general, for r nested loops of this type,
the innermost loop happens $\Theta(n^r)$ times

A1. $\sim \frac{1}{2}n^3$ array accesses.

A2. $\Theta(n^3)$ array accesses.

Bottom line. Use cost model and asymptotic notation to simplify analysis.

Common order-of-growth classifications

order of growth	emoji	name	typical code framework	description	example
$\Theta(1)$	😍	**constant**	`a = b + c;`	statement	*add two numbers*
$\Theta(\log n)$	😎	**logarithmic**	`for (int i = n; i > 0; i /= 2)` `{  ...  }`	divide in half	*binary search*
$\Theta(n)$	😁	**linear**	`for (int i = 0; i < n; i++)` `{  ...  }`	single loop	*find the maximum*
$\Theta(n \log n)$	😀	**linearithmic**	*mergesort*	divide and conquer	*mergesort*

$\Theta(n^2)$	😕	**quadratic**	`for (int i = 0; i < n; i++)`
			`  for (int j = 0; j < n; j++)`
			`    {  ...  }`
				double loop	*check all pairs*

$\Theta(n^3)$	☹️	**cubic**	`for (int i = 0; i < n; i++)`
			`  for (int j = 0; j < n; j++)`
			`    for (int k = 0; k < n; k++)`
			`  {  ...  }`
				triple loop	*check all triples*

$\Theta(2^n)$	👿	**exponential**	*towers of Hanoi*	exhaustive search	*check all subsets*

Some useful discrete sums and approximations

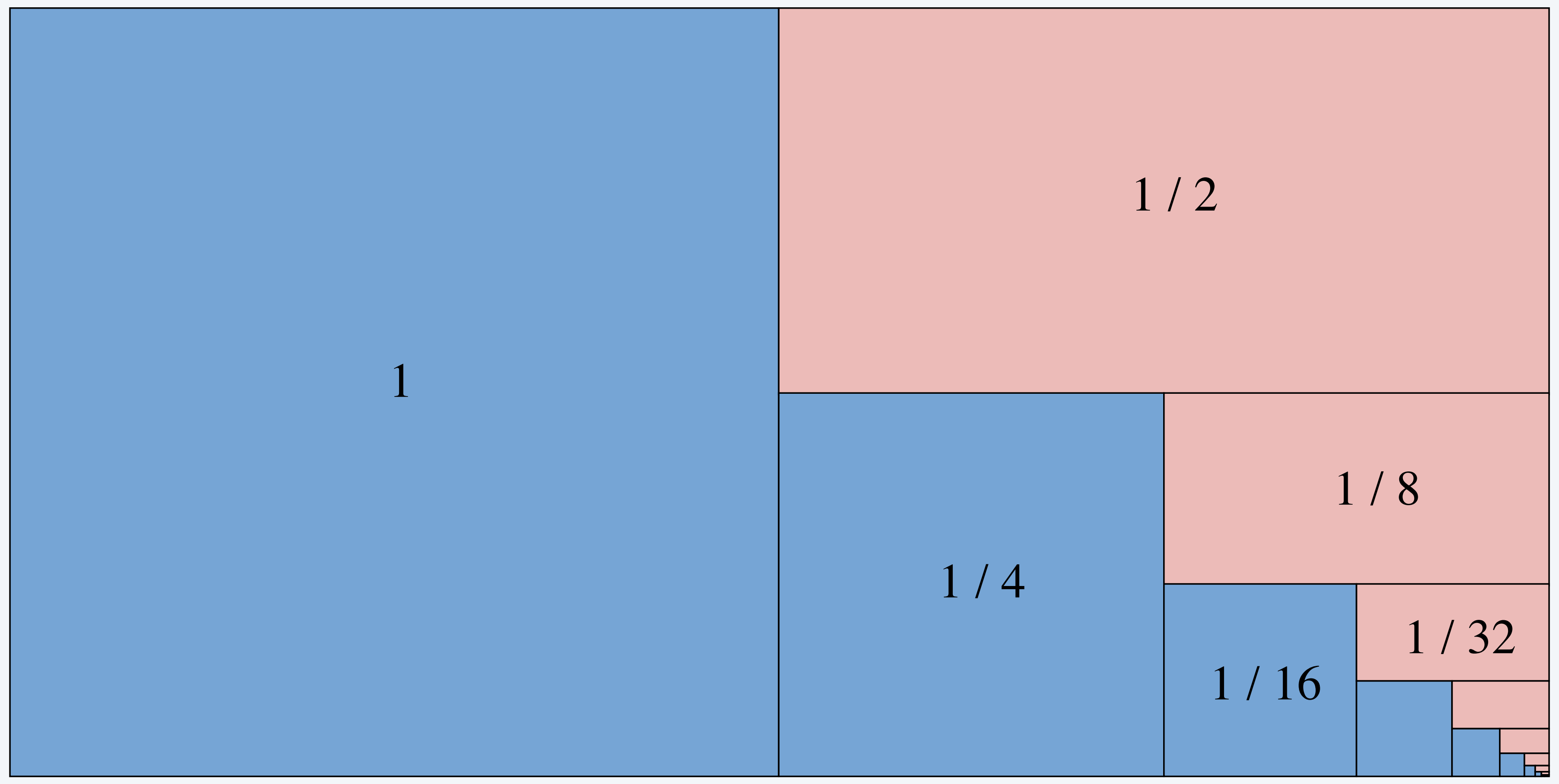
Triangular sum. $1 + 2 + 3 + \ldots + n \sim \frac{1}{2} n^2$

Geometric sum. $1 + 2 + 4 + 8 + \ldots + n = 2n - 1$

n a power of 2

in general, for $r > 1$,
$1 + r + r^2 + r^3 + \ldots + n = \Theta(n)$

Geometric sum′. $n + \frac{n}{2} + \frac{n}{4} + \ldots + 1 = 2n - 1$



$$1 + \frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \ldots$$

Approximately how many array accesses as a function of n?

```
int count = 0;
for (int i = 0; i < n; i++)
   for (int j = i+1; j < n; j++)
      for (int k = 1; k <= n; k = k*2)
         if (a[i] + a[j] >= a[k])
            count++;
```

A. $\sim n^2 \log_2 n$

B. $\sim 3/2\ n^2 \log_2 n$

C. $\sim 1/2\, n^3$

D. $\sim 3/2\, n^3$

What is the order of growth of the running time as a function of n ?

```
int count = 0;
for (int i = n; i >= 1; i = i/2)
   for (int j = 1; j <= i; j++)
      count++;
```

A. $\Theta(n)$

B. $\Theta(n \log n)$

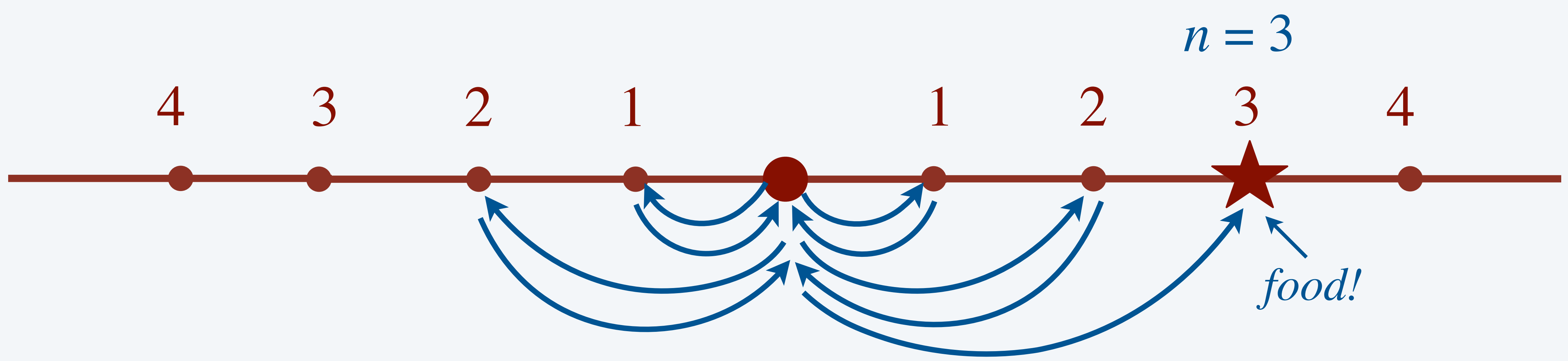
C. $\Theta(n^2)$

D. $\Theta(2^n)$

Example: Rat (f22 midterm exam)

A rat in a sewer pipe is searching for food. If the nearest food source is n steps to the right of its starting location, how many steps will it take to reach it using the given strategy?

Strategy 1: Take 1 step right, return to start, take 1 step left, return to start.
Repeat with 2, 3, 4, 5… steps until food found.



A rat in a sewer pipe is searching for food. If the nearest food source is n steps to the right of its starting location, how many steps will it take to reach it using the given strategy?

assume n is a power of 2

Strategy 2: Take 1 step right, return to start, take 1 step left, return to start.
Repeat with 2, 4, 8, 16… steps until food found.

A. $\Theta(\log n)$

B. $\Theta(n)$

C. $\Theta(n \log n)$

D. $\Theta(n^2)$

E. $\Theta(2^n)$

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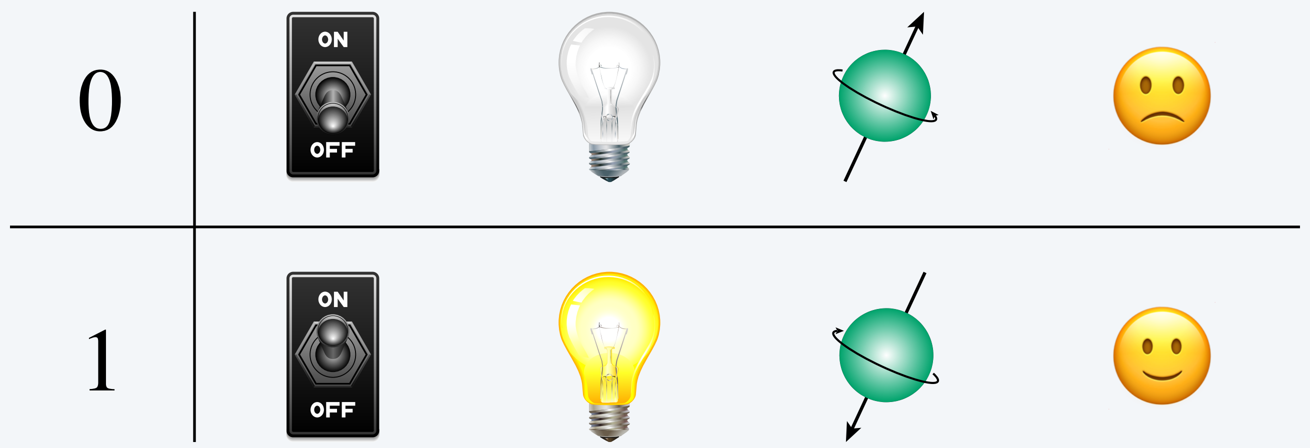
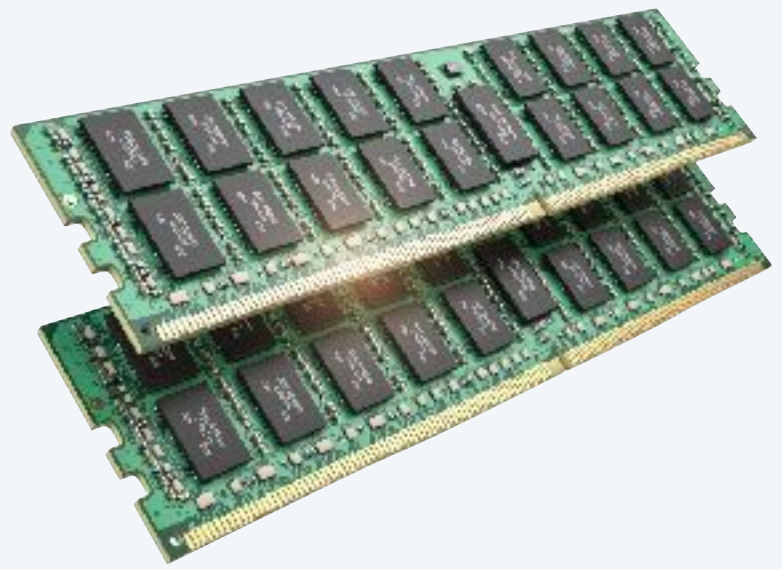
Memory basics

Bit. 0 or 1.



term	symbol	quantity
byte	B	8 bits
kilobyte	KB	1000 bytes
megabyte	MB	1000^2 bytes
gigabyte	GB	1000^3 bytes
terabyte	TB	1000^4 bytes

some define using powers of 2
(MB = 2^{10} *bytes*)



64-bit machine. We assume a 64-bit machine with 8-byte pointers.

some JVMs "compress" pointers
to 4 bytes to avoid this cost

Typical memory usage for primitive types and arrays

type	bytes
`boolean`	1
`byte`	1
`char`	2
`int`	4
`float`	4
`long`	8
`double`	8

primitive types

type	bytes
`boolean[]`	$\sim n$
`int[]`	$\sim 4n$
`double[]`	$\sim 8n$

array overhead = 24 bytes

one-dimensional arrays (length n)

type	bytes
`boolean[][]`	$\sim 1\ n^2$
`int[][]`	$\sim 4\ n^2$
`double[][]`	$\sim 8\ n^2$

two-dimensional arrays (n-by-n)

type	bytes
`reference (pointer)`	8

64-bit machine

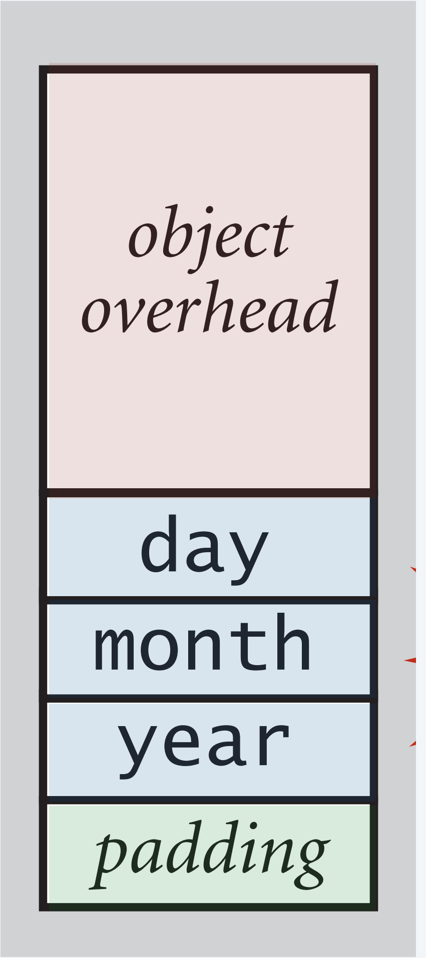
Typical memory usage for objects in Java

Objects memory = sum of memory for instance variables + overheads

Ex. Each *Date* object uses 32 bytes of memory.

```
public class Date {
   private int day;
   private int month;
   private int year;

   ...
}
```



16 bytes (object overhead)

4 bytes (`int`)

4 bytes (`int`)

4 bytes (`int`)

4 bytes (padding, round to a multiply of 8 bytes)

32 bytes

Array declaration is 8 bytes.

```
Date[] dates;
```
← *reference*

When *dates* contains n elements, it uses $\Theta(n)$ bytes.

Credits

image	source	license
Charles Babbage	The Illustrated London News	public domain
Babbage Enginine in Operation	xRez Studio	
Algorithm for the Analytic Engine	Ada Lovelace	public domain
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Alan Turing	Science Museum, London	

A final thought

“ *It is convenient to have a measure of the amount of work involved in a computing process, even though it be a very crude one. We may count up the number of times that various elementary operations are applied in the whole process and then give them various weights.*” — Alan Turing (1947)