

### 2.1 Elementary Sorts

- rules of the game
- selection sort
- insertion sort
- binary search

Robert Sedgewick I Kevin Wayne
https://algs4.cs.princeton.edu

### 2.1 Elementary Sorts

- rules of the game
- selection sort
- insertion sort
- binary search

Robert Sedgewick I Kevin Wayne
https://algs4.cs.princeton.edu

## Sorting problem

Problem. Given an array of $n$ elements, rearrange in ascending order by key.

|  | Last | First | House | Year |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | Longbottom | Neville | Gryffindor | 1998 |  |
|  | Weasley | Ron | Gryffindor | 1998 |  |
|  | Abbott | Hannah | Hufflepuff | 1998 |  |
| element $\longrightarrow$ | Potter | Harry | Gryffindor | 1998 |  |
|  | Chang | Cho | Ravenclaw | 1997 |  |
|  | Granger | Hermione | Gryffindor | 1998 |  |
| key $\longrightarrow$ | Malfoy | Draco | Slytherin | 1998 | sorting hat |
|  | Diggory | Cedric | Hufflepuff | 1996 |  |
|  | Weasley | Ginny | Gryffindor | 1999 |  |
|  | Parkinson | Pansy | Slytherin | 1998 |  |

## Sorting problem

Problem. Given an array of $n$ elements, rearrange in ascending order by key.


## Sorting problem

Sorting is a well-defined problem if there is a binary relation $\leq$ that satisfies:

- Totality: either $v \leq w$ or $w \leq v$ or both.
- Transitivity: if both $v \leq w$ and $w \leq x$, then $v \leq x$.

Examples.

| International Departures |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| Flight No | Destination | Time | Gate | Remarks |
| CX7183 | Berlin | 77:50 | A-11 | Gate closed |
| [FF3474 | London | 77:50 | A-12 | Gote closed |
| BA372 | Paris | 77:55 | B-10 | Boarding |
| AY6554 | New York | 8:00 | c-33 | Boarding |
| KL3150 | San Froncisco | 8:00 | F-15 | Boording |
| EA8903 | Manchester | 8:05 | B-12 | See ticket desk |
| Ba710 | Los Angeles | 8:10 | C-12 | Check-in open |
| QF3371 | Hong Kong | 8:15 | F-10 | Check-in open |
| MA4865 | Barcelona | $8: 15$ | F-12 | Check-in ot kiosks |
| cx7221 | Copenhogen | 8:20 | 6-32 | Check-in ot kiosks |


alphabetical order

| No. $\boldsymbol{\wedge}$ | Video name | Views (billions) |
| :---: | :--- | :---: |
| 1. | "Baby Shark Dance"[3] | 10.15 |
| 2. | "Despacito"[6] | 7.73 |
| 3. | "Johny Johny Yes Papa"[12] | 6.15 |
| 4. | "Shape of You"[13] | 5.61 |
| 5. | "See You Again"[15] | 5.41 |

numerical order (descending)

## Sorting problem

Sorting is a well-defined problem if there is a binary relation $\leq$ that satisfies:

- Totality: either $v \leq w$ or $w \leq v$ or both.
- Transitivity: if both $v \leq w$ and $w \leq x$, then $v \leq x$.

Non-examples.

course prerequisites (violates totality)


Ro-sham-bo order (violates transitivity)

```
~/cos226/sort> jshel1
Math.sqrt(-1.0) <= Math.sqrt(-1.0);
false
```

the <= operator for double (irreflexive, which violates totality)

## Sample sort clients

Goal. General-purpose sorting function.
Ex 1. Sort strings in lexicographic order. alphabetical order, using Unicode character ordering

```
public class StringSorter {
    public static void main(String[] args) {
        String[] a = StdIn.readAl1Strings();
        Insertion.sort(a);
        for (int i = 0; i < a.length; i++)
            StdOut.println(a[i]);
    }
}
```

```
~/cos226/sort> more words3.txt
bed bug dad yet zoo ... all bad yes
~/cos226/sort> java StringSorter < words3.txt
a11 bad bed bug dad
yes yet zoo
[suppressing newlines]
```


## Sample sort clients

## Goal. General-purpose sorting function.

Ex 2. Sort real numbers in numerical order (ascending).

```
public class Experiment {
    public static void main(String[] args) {
        int n = Integer.parseInt(args[0]);
        Doub7e[] a = new Doub7e[n];
        for (int i = 0; i < n; i++)
            a[i] = StdRandom.uniformDoub7e();
        Insertion.sort(a);
        for (int i = 0; i < n; i++)
            StdOut.println(a[i]);
    }
}
```

```
~/cos226/sort> java Experiment 10
0.08614716385210452
0.09054270895414829
0.10708746304898642
0.21166190071646818
0.363292849257276
0.460954145685913
0.5340026311350087
0.7216129793703496
0.9003500354411443
0.9293994908845686
```


## Sample sort clients

Goal. General-purpose sorting function.
Ex 3. Sort playing cards by suit and rank.

```
public class Deck {
    public static void main(String[] args) {
        int n = Integer.parseInt(args[0]);
        PlayingCard[] cards = deal(n);
        Insertion.sort(cards);
        draw(cards);
    }
}
```


## How can a single function sort any type of data？

Goal．General－purpose sorting function．

Please sort these Japanese names for me： あゆみ，アユミ，Ayumi，歩美，


But I don＇t speak Japanese and I don＇t know how words are ordered．


## Callbacks

Goal. General-purpose sorting function.

Solution. Callback $=$ reference to executable code passed to other code and later executed.

- Client passes array of objects to sort() function.
- The sort() function calls object's compareTo() method as needed.

> in effect, client passes compareTo() method to sort() function; the callback occurs when sort() invokes compareTo()

Implementing callbacks.

- Java: interfaces.
- Python, ML, Javascript: first-class functions.
- C\#: delegates.
- C: function pointers.
- C++: class-type functors.


## Review: Java interfaces

Interface. A set of related methods that define some behavior (partial API) for a class.

## interface (java.lang.Comparable)

```
public interface Comparable<Item> {
```

    public int compareTo(Item that);
    \}
$\qquad$ contract: method with this signature (and prescribed behavior)

Class that implements interface. Must implement all interface methods.


## Callbacks in Java: roadmap

## client (StringSorter.java)

```
public class StringSorter {
    public static void main(String[] args) {
        String[] a = StdIn.readAl1Strings();
        Insertion.sort(a);
```


## interface (Comparable.java)

```
public interface Comparable<Item>
    int compareTo(Item that);
}
```

String[] is a subtype
of Comparab7e[]
sort implementation (Insertion.java)
key point: sorting code does not
depend upon type of data to be sorted

## Elementary sorts: quiz 1

Suppose that the Java architects left out implements Comparable<String> in the class declaration for String. What would be the effect?
A. Compile-time error in String.java.
B. Compile-time error in StringSorter.java.
C. Compile-time error in Insertion.java.
D. Run-time exception in Insertion.java.

## Comparable API

Implement compareTo() so that v.compareTo(w)

- Returns a negative integer if $v$ is less than $w$.

API requirement:

- Returns a positive integer if $v$ is greater than $w$.
- Returns zero if $v$ is equal to $w$.
- Throws an exception if incompatible types (or either is nul1).

$v$ is less than $w$ (return negative integer)

$v$ is equal to $w$ (return 0)

$v$ is greater than $w$ (return positive integer)

Built-in comparable types. Integer, Double, String, java.util.Date, ...
User-defined comparable types. Implement the Comparable interface.

## Implementing the Comparable interface

Date data type. Simplified version of java.util. Date.

```
public class Date implements Comparable<Date>
    private final int month, day, year;
    public Date(int m, int d, int y) {
        month = m;
        day = d;
        year = y;
    }
    public int compareTo(Date that) {
        if (this.year < that.year ) return -1;
        if (this.year > that.year ) return +1;
        if (this.month < that.month) return -1;
        if (this.month > that.month) return +1;
        if (this.day < that.day ) return -1;
        if (this.day > that.day ) return +1;
        return 0;
    }
}
```


### 2.1 Elementary Sorts

- rules of the game
- selection sort

Algorithms

Robert Sedgewick | Kevin Wayn
https://algs4.cs.princeton.edu

## Selection sort demo

Algorithm. For each index $i$ from 0 to $n-1$ :

- Find index min of smallest remaining element.
- Swap elements at indices $i$ and min.

initial array


## Selection sort: visualization

Visualization. Sort vertical bars by length.


A algorithm position
in order
not yet seen

## Selection sort invariants

Algorithm. For each index $i$ from 0 to $n-1$ :

- Find index min of smallest remaining element.
- Swap elements at indices $i$ and min.

Invariants.
before iteration $\mathbf{i}$

| $\longmapsto$ sorted | min |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| smallest $i$ elements |  | $\ldots$ |  |  |

after iteration i


## Two useful sorting primitives (and a cost model)

Helper functions. Refer to data only through compares and exchanges. e.g., no calls to equa1s ()


Compare. Is item v less than item w?


Exchange. Swap array entries a[i] and $\mathrm{a}[\mathrm{j}]$.

```
private static void exch(Object[] a, int i, int j) {
    Object swap = a[i];
    a[i] = a[j];
    a[j] = swap; Java arrays are "covariant"
}
    (e.g., String[] is a subtype of Object[])
```


## Selection sort: Java implementation

```
public class Selection
    public static void sort(Comparable[] a) {
        int n = a.length;
        for (int i = 0; i < n; i++)
            int min = i;
            for (int j = i+1; j < n; j++)
                if (less(a[j], a[min]))
                min = j;
        exch(a, i, min);
    }
    private static boolean less(Comparable v, Comparable w) {
        see previous stide */
    }
    private static void exch(Object[] a, int i, int j) {
        /* see previous slide */
    }
}
```


## Elementary sorts: quiz 2

How many compares to selection sort an array of $\boldsymbol{n}$ distinct items in reverse order?
A. $\sim n$
B. $\sim 1 / 4 n^{2}$
C. $\sim 1 / 2 n^{2}$
D. $\sim n^{2}$

## Selection sort: mathematical analysis

Proposition. Selection sort makes $(n-1)+(n-2)+\ldots+1+0 \sim 1 / 2 n^{2}$ compares and $n$ exchanges to sort any array of $n$ items.


Running time insensitive to input. $\Theta\left(n^{2}\right)$ compares. $\qquad$ even if input array is sorted

Data movement is minimal. $\Theta(n)$ exchanges.
In place. $\Theta(1)$ extra space.

### 2.1 Elementary Sorts

- rules of the game
- selectión sort
- insertion sort
- binary search

Robert Sedgewick I Kevin Wayne
https://algs4.cs.princeton.edu

Algorithm. For each index $i=0$ to $n-1$ :

- Let $x$ be the element at index $i$.
- Repeatedly exchange $x$ with each larger element to its immediate left.

initial array


## Insertion sort invariants

Algorithm. For each index $i=0$ to $n-1$ :

- Let $x$ be the element at index $i$.
- Repeatedly exchange $x$ with each larger element to its immediate left.

Invariants.

## before iteration $\mathbf{i}$

|  |  |  |  |
| :---: | :---: | :---: | :---: |
|  |  |  | sorted |
| $\leq x$ | $>x$ | $x$ | $\ldots$ |

after iteration $\mathbf{i}$

| sorted |  |  | untouched |  |
| :---: | :---: | :---: | :---: | :---: |
|  | $x$ | $>x$ | $\ldots$ |  |

## Insertion sort: Java implementation

```
public class Insertion
    public static void sort(Comparable[] a) {
        int n = a.length;
        for (int i = 0; i < n; i++)
            for (int j = i; j > 0; j--)
                if (less(a[j], a[j-1]))
                exch(a, j, j-1);
            else break;
    }
    private static boolean less(Comparable v, Comparable w) {
        /* as before */
    }
    private static void exch(Object[] a, int i, int j) {
        /* as before */
    }
}
```

https://algs4.cs.princeton.edu/21elementary/Insertion.java.htmI

## Elementary sorts: quiz 4

How many compares to insertion sort an array of $\boldsymbol{n}$ distinct keys in reverse order?
A. $\sim n$
B. $\sim 1 / 4 n^{2}$
C. $\sim 1 / 2 n^{2}$
D. $\sim n^{2}$

Worst case. Insertion sort makes $\sim 1 / 2 n^{2}$ compares and $\sim 1 / 2 n^{2}$ exchanges to sort an array of $n$ distinct keys in reverse order.
Pf. Exactly $i$ compares and exchanges in iteration $i$.

$$
0+1+2+\ldots+(n-1) \sim 1 / 2 n^{2}
$$



A algorithm position
in order
not yet seen

Best case. Insertion sort makes $n-1$ compares and 0 exchanges to sort an array of $n$ distinct keys in ascending order.


- algorithm position
in order
not yet seen

Good case. Insertion sort takes $\Theta(n)$ time on "partially sorted" arrays.
Q. Can we formalize what we mean by partially sorted?
A. Yes, in terms of "inversions" (see textbook).


A algorithm position
in order
not yet seen

## Insertion sort: practical improvements

Half exchanges. Shift items over (instead of exchanging).

- Same compares; fewer array accesses.
- No longer uses only less() and exch() to access data.

A C H H I M N P Q X Y K B I N A R Y

Binary insertion sort. Use binary search to find insertion point.

- Now, worst-case number of compares $\sim n \log _{2} n$.
- But still makes $\Theta\left(n^{2}\right)$ array accesses in worst case.

```
A C H H I M N P Q X Y K B I N A R Y
```


### 1.4 Analysis of Algorithms

- rules of the game
$\rightarrow$ selectión sort
- insertion sort
- binary search

Robert Sedgewick \| Kevin Wayne
https://algs4.cs.princeton.edu

## Binary search

Goal. Given a sorted array and a search key, find index of the search key in the array?

Binary search. Compare search key with middle entry.

- Too small, go left.
- Too big, go right.
- Equal, found.
sorted array

| 6 | 13 | 14 | 25 | 33 | 43 | 51 | 53 | 64 | 72 | 84 | 93 | 95 | 96 | 97 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 |

## Binary search: implementation

Trivial to implement?

- First binary search published in 1946.
- First bug-free one in 1962.
- Bentley experiment: 90\% of programmers implement it incorrectly.
- Bug in Java's Arrays.binarySearch() discovered in 2006.
and in $C, C++, \ldots$

Extra, Extra - Read All About It: Nearly All Binary Searches and Mergesorts are Broken
Friday, June 02, 2006
Posted by Joshua Bloch, Software Engineer
I remember vividly Jon Bentley's first Algorithms lecture at CMU, where he asked all of us incoming Ph.D. students to write a binary search, and then dissected one of our implementations in front of the class. Of course it was broken, as were most of our implementations. This made a real impression on me, as did the treatment of this material in his wonderful Programming Pearls (Addison-Wesley, 1986; Second Edition, 2000). The key lesson was to carefully consider the invariants in your programs.


[^0]Binary search: implementation

Invariant. If key appears in array a[] , then $\mathrm{a}[1 \mathrm{o}] \leq$ key $\leq \mathrm{a}[\mathrm{hi}]$.

```
public static int binarySearch(String[] a, String key) {
    int lo = 0, hi = a.length - 1;
    while (lo <= hi) { whynot mid = (10 + hi) / 2?
        int mid = (lo + hi) >>> 1;
        int compare = key.compareTo(a[mid]);
        if (compare < 0) hi = mid - 1;
        else if (compare > 0) 10 = mid + 1;
        else return mid;
    }
    return -1;
}
```


## Binary search: analysis

Proposition. Binary search makes at most $1+\log _{2} n$ compares to search in any sorted array of length $n$.

Pf.

- Each iteration of while loop:
- calls compareTo() once
- decreases the length of remaining subarray by at least a factor of 2
$\longleftarrow$ can happen at most $1+\log _{2} n$ times. Why?
$n \rightarrow n / 2 \rightarrow n / 4 \rightarrow n / 8 \rightarrow \cdots \rightarrow 2 \rightarrow 1$
slightly better than $2 \times$,
due to elimination of a [mid] from subarray
(or early termination of while loop)


3-Sum. Given an array of $n$ distinct integers, count number of triples that sum to 0 .

Version 0. $\Theta\left(n^{3}\right)$ time in worst case.
Version 1. $\Theta\left(n^{2} \log n\right)$ time in worst case.
Version 2. $\Theta\left(n^{2}\right)$ time in worst case.

Note. For full credit, use only $\Theta(1)$ extra space.

## Summary

Comparable interface. Java framework for comparing items.

Selection sort. $\Theta\left(n^{2}\right)$ compares; $\Theta(n)$ exchanges.
Insertion sort. $\Theta\left(n^{2}\right)$ compares and exchanges in the worst case.

Binary search. Search a sorted array using $\Theta(\log n)$ compares in worst case.

## Credits

| image/video | source | license |
| :---: | :---: | :---: |
| Sorting Hat | Hannah Hill | CC BY-NC 4.0 |
| Airport Departures | $\underline{\text { Adobe Stock }}$ | $\underline{\text { education license }}$ |
| iPhone Contacts | $\underline{\text { StackOverflow }}$ |  |
| Playing Cards | $\underline{\text { Google Code }}$ | public domain |
| Rock, Paper, Scissors | $\underline{\text { Daily Mail }}$ |  |
| Anime Boy | $\underline{\text { freesvg.org }}$ | public domain |
| Anime Girl | $\underline{\text { freesvg.org }}$ | public domain |
| Balance | $\underline{\text { Adobe Stock }}$ | $\underline{\text { education license }}$ |
| Jon Bentley | $\underline{\text { Amazon }}$ |  |
| Binary vs. Sequential Search | $\underline{\text { Silicon Valley S6E4 }}$ |  |
| Insertion Sort Dance | $\underline{\text { AlgoRythmics }}$ |  |
| Lecture Slides © Copyright 2023 Robert Sedgewick and Kevin Wayne |  |  |




[^0]:    https://ai.googleblog.com/2006/06/extra-extra-read-all-about-it-nearly.html

