

# COS 217: Introduction to Programming Systems

## Processes



**PRINCETON UNIVERSITY**



# Agenda



## Processes

Illusion: Private address space  
Illusion: Private control flow



## Process management in C

Creating new processes  
Waiting for termination  
Executing new programs



## Unix Process Control

Exceptions  
Signals



# Processes

## Program

- Executable code
- A static entity

## Process

- An instance of a program in execution
- A dynamic entity: has a time dimension
- Each process runs one program
  - E.g. the process with Process ID 12345 might be running emacs
- One program can run in multiple processes
  - E.g. PID 12345 might be running emacs, and PID 23456 might also be running emacs – for the same user or for a different user



# Processes Significance

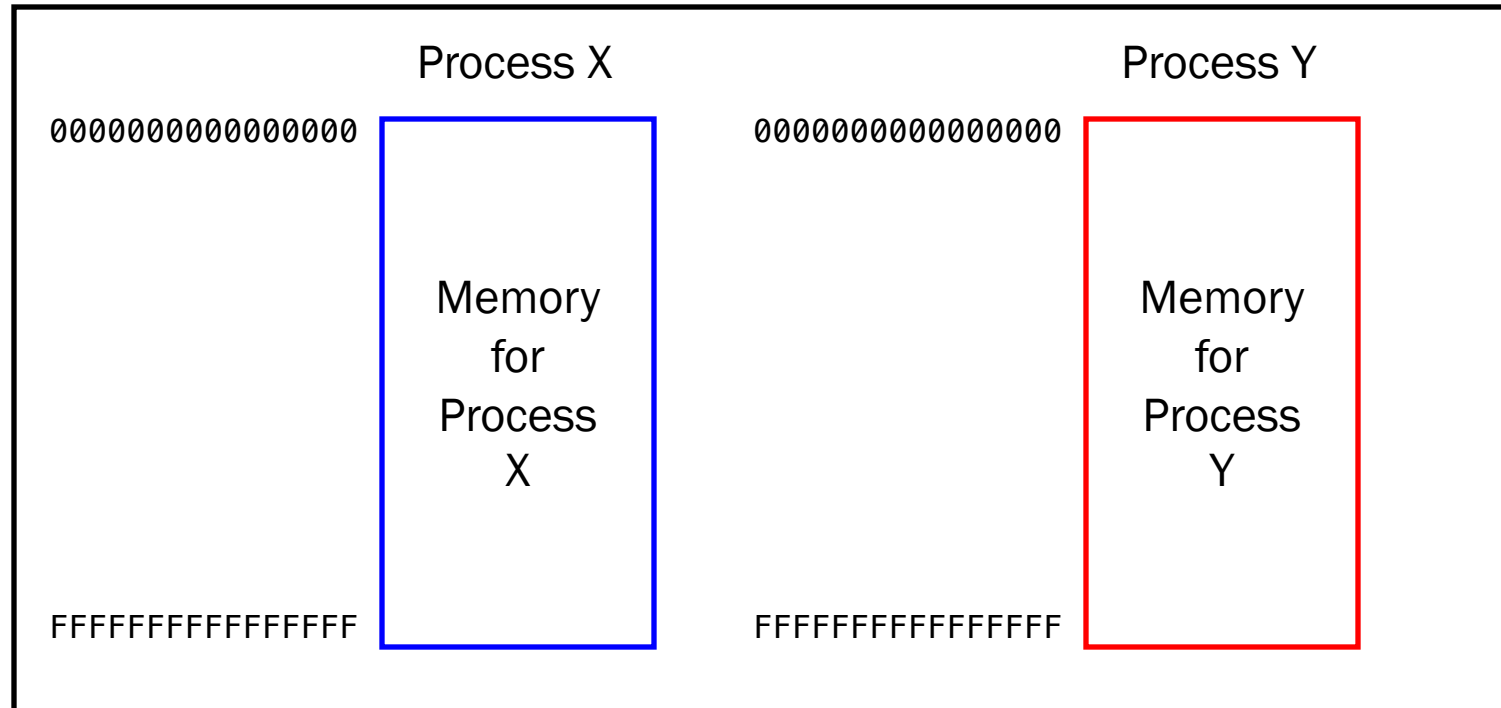
Process abstraction provides two key illusions:

- Processes believe they have a *private address space*
- Processes believe they have *private control flow*

**Process is a profound abstraction in computer science**



# Private Address Space: Illusion

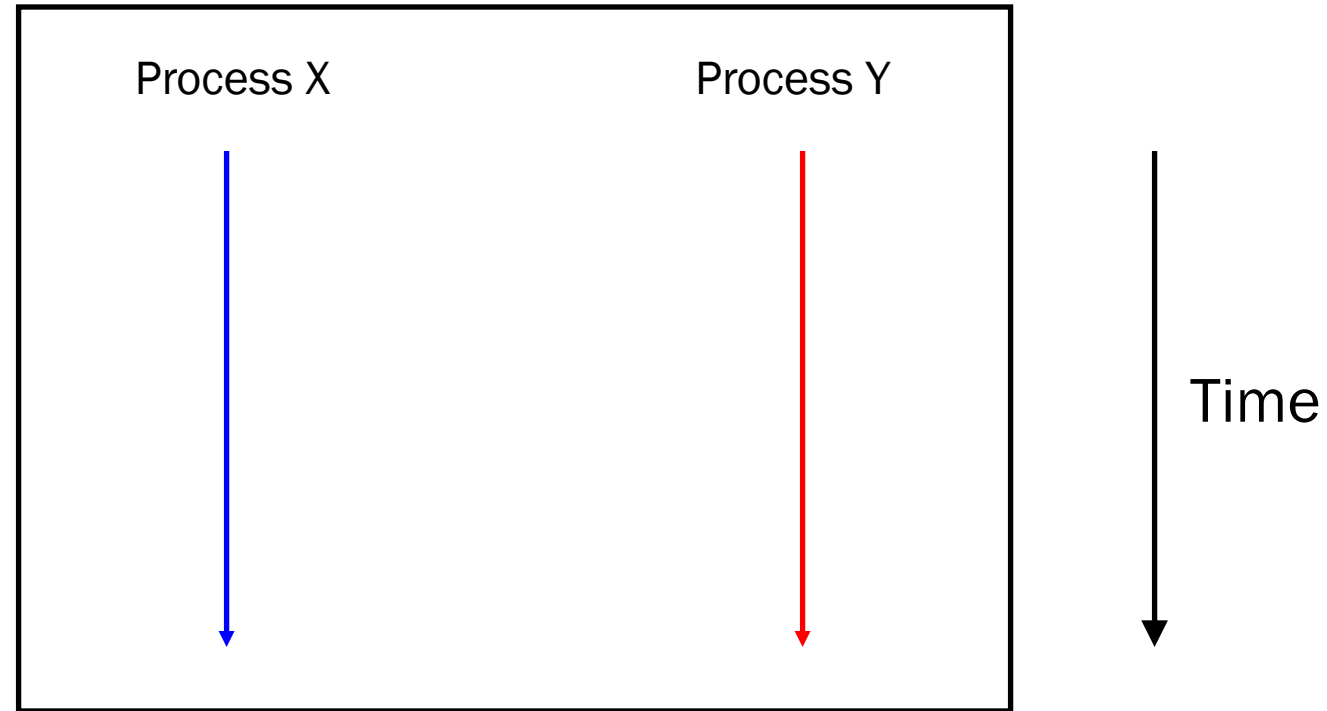


Hardware and OS give each application process the illusion that it is the *only* process using memory

- Enables multiple simultaneous instances of one program!



# Private Control Flow: Illusion

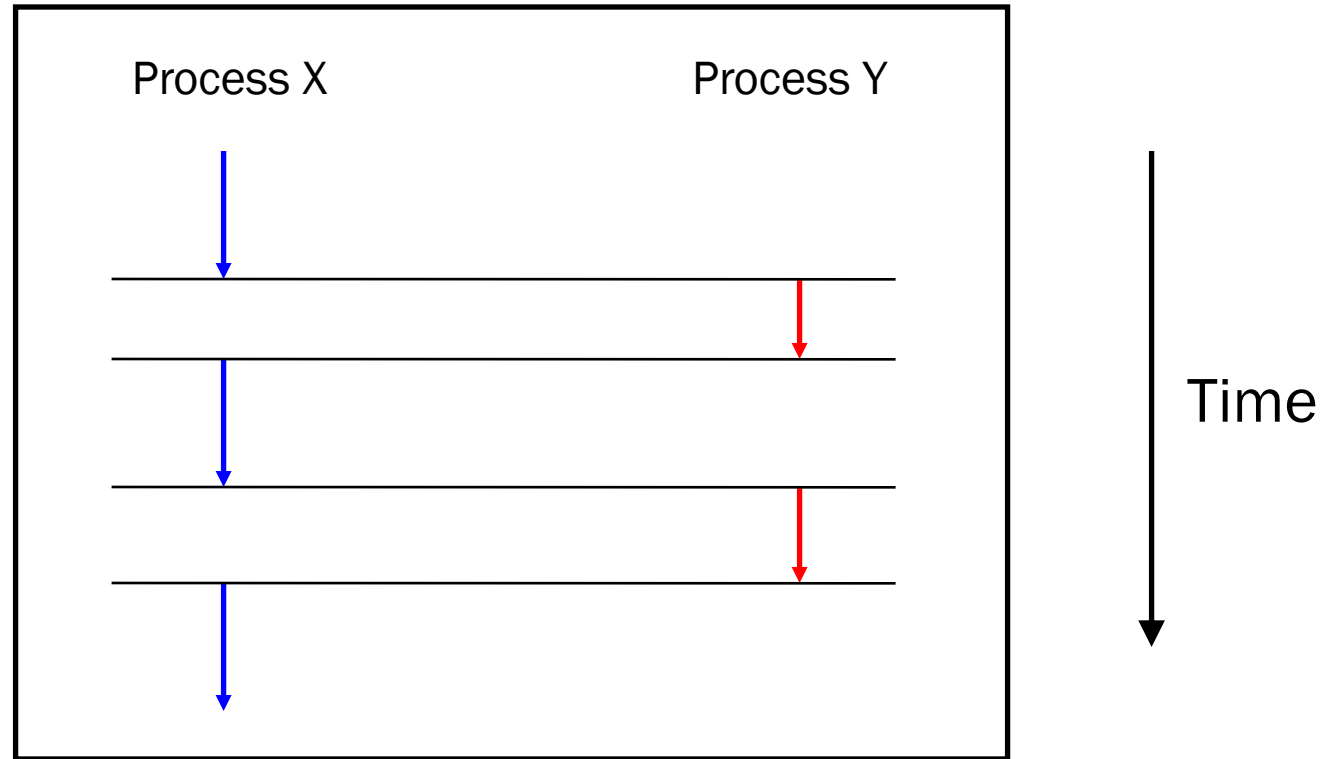


Simplifying assumption: only one CPU / core

Hardware and OS give each application process the illusion that it is the *only* process running on the CPU



# Private Control Flow: Reality



Multiple processes are **time-sliced** (possibly across multiple CPUs / cores) to run concurrently

OS occasionally **preempts** running process to give other processes their fair share of CPU time



# Process Status

More specifically...

At any time, a process has a **status**:

- **Running**: a CPU is executing instructions for the process
- **Ready**: Process is ready for OS to assign it to a CPU
- **Blocked**: Process is waiting for some requested service (typically I/O) to finish

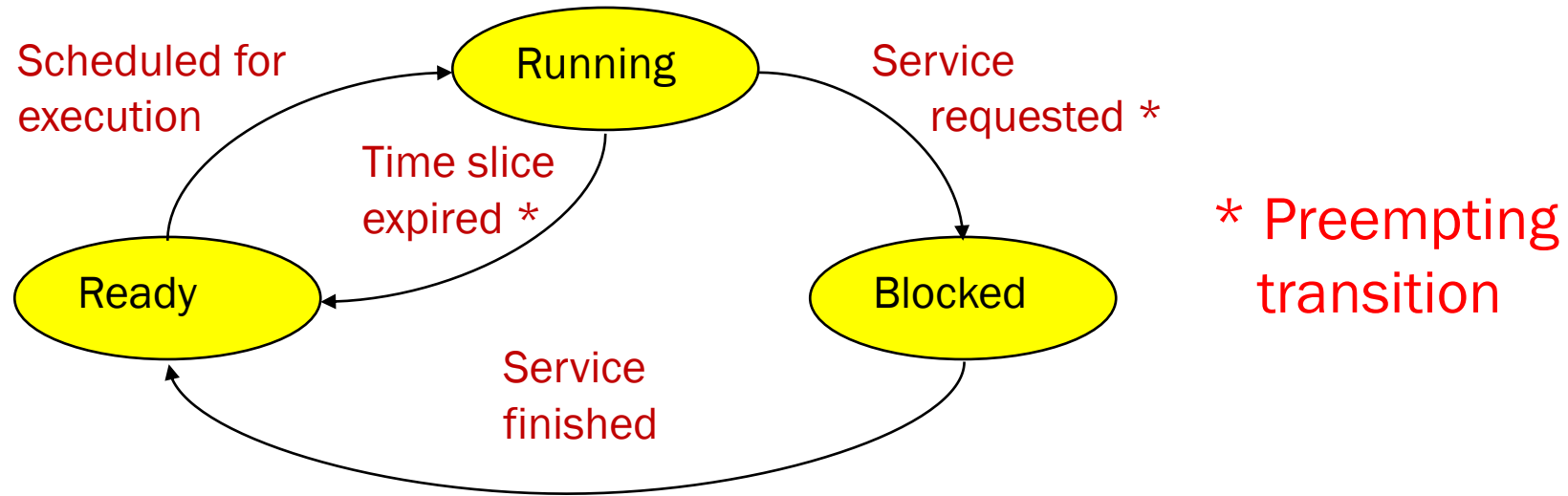
Modern machines may have multiple CPUs or “cores”, but the same principles apply if  $\#processes > \#cores$

- For simplicity, we will speak of “the” CPU





# Process Status Transitions



**Scheduled for execution:** OS selects some process from ready set and assigns CPU to it

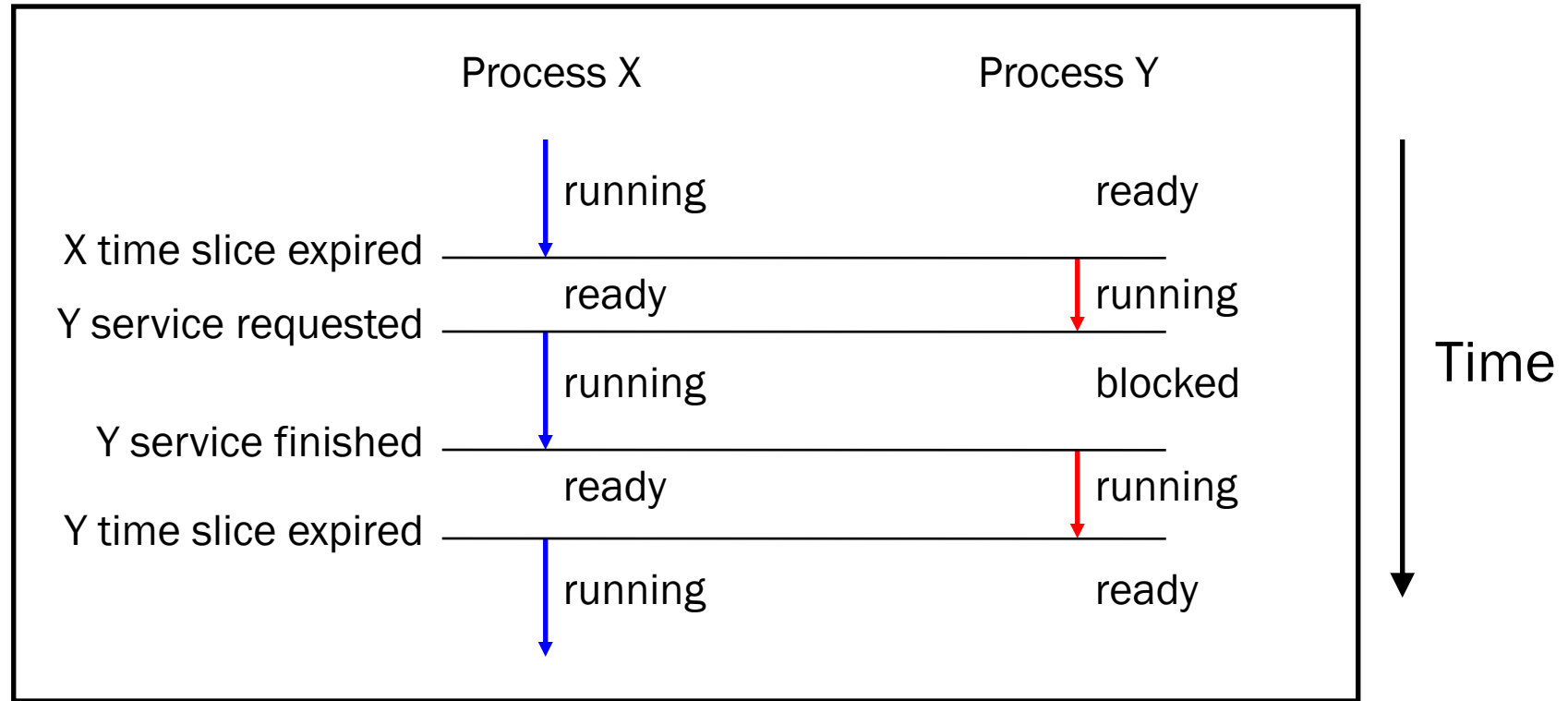
**Time slice expired:** OS moves running process to ready set because process consumed its fair share of CPU time

**Service requested:** OS moves running process to blocked set because it requested a (time consuming) system service (often I/O)

**Service finished:** OS moves blocked process to ready set because the requested service finished



# Process Status Transitions Over Time



Throughout its lifetime, a process's status switches between running, ready, and blocked



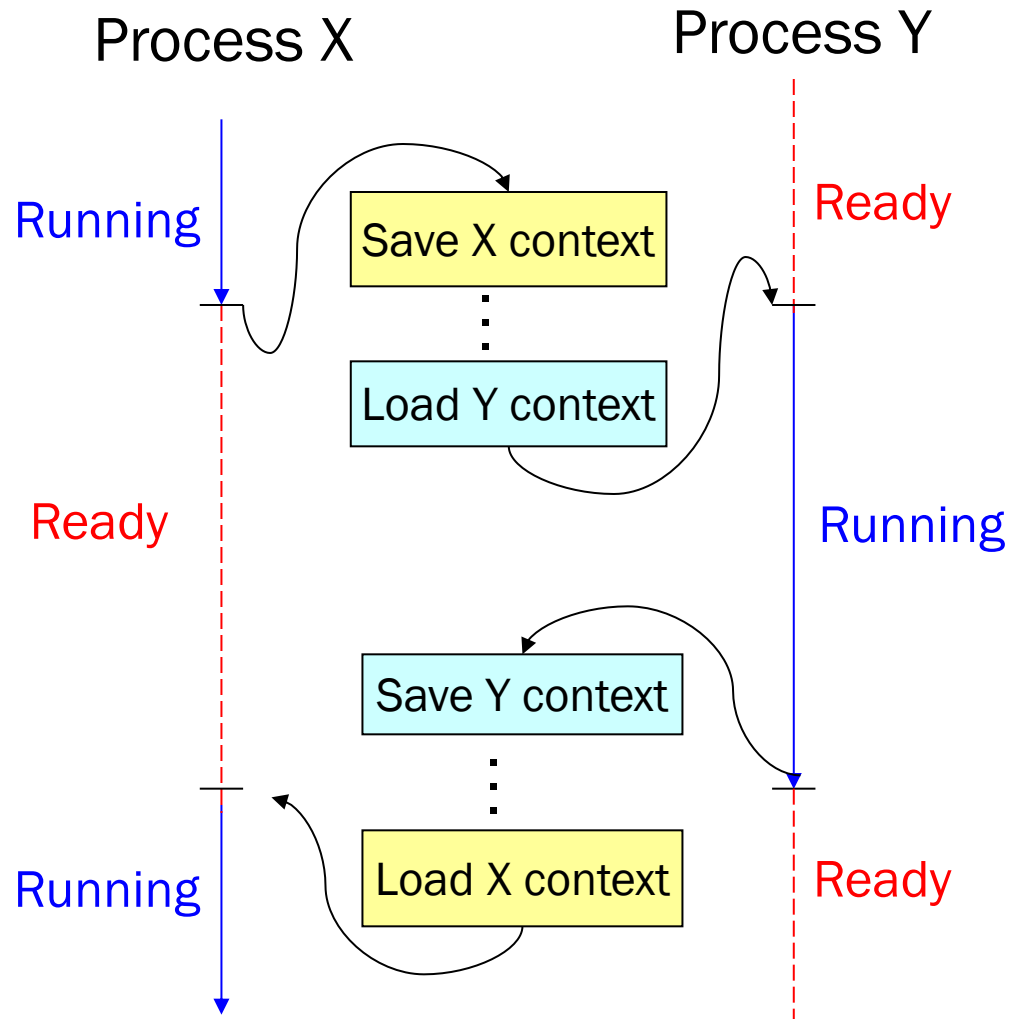
# Process Contexts

Each process has a **context**

- The process's state, that is...
- Contents of registers
- Memory contents
  - TEXT, RODATA, DATA, BSS, HEAP, and STACK
  - More specifically, **page tables** mapping this process's VM
  - Even more specifically, **pointer** to page tables used for this process



# Context Switch



## Context switch:

- OS saves context of running process
- OS loads context of some ready process
- OS passes control to newly restored process
  
- Manageable cost: just registers, plus an indication of which page tables to use

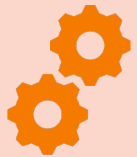


# Agenda



## Processes

Illusion: Private address space  
Illusion: Private control flow



## Process management in C

Creating new processes  
Waiting for termination  
Executing new programs



## Unix Process Control

Exceptions  
Signals

# System-Level Process Management Functions



Function	Description
<code>exit()</code>	Terminate the process
<code>fork()</code>	Create a child process
<code>wait()</code>	Wait for child process termination
<code>execvp()</code>	Execute a program in current process



# Why Create New Processes?

## Why create a new process?

- Scenario 1: Program wants to run an additional instance of itself
  - E.g., **web server** receives request; creates additional instance of itself to handle the request; original instance continues listening for requests
- Scenario 2: Program wants to run a different program
  - E.g., **shell** receives a command; creates an additional instance of itself; additional instance overwrites itself with requested program to handle command; original instance continues listening for commands

## How to create a new process?

- A “parent” process **forks** a “child” process
- (Optionally) child process overwrites itself with a new program, after performing appropriate setup



# fork System-Level Function

```
pid_t fork(void);
```

- Create a new process by duplicating the calling process
  - New (child) process is an exact duplicate\* of the calling (parent) process
- \* Almost – the call to `fork` has a different return value (wait 1 slide)

`fork()` is called once in parent process

`fork()` returns twice

- Once in parent process
- Once in child process





# fork and Return Values

Return value of fork has meaning

- In child, fork( ) returns 0
- In parent, fork( ) returns process id of child



**BEST WAY TO  
RECOGNIZE TWINS**

```
pid = fork();  
  
if (pid == 0) {  
    /* executed in child */  
    ...  
} else {  
    /* executed in parent */  
    ...  
}
```



# Programs With Processes

Parent process and child process run **concurrently**

- Two CPUs available ⇒
  - Parent process and child process run in **parallel**
- Fewer than two CPUs available ⇒
  - Parent process and child process time-sliced to run **serially**
  - OS provides the **illusion** of parallel execution

Reality: Each ArmLab computer has 96 CPUs

- But each student who is logged in might be concurrently running sshd, bash, emacs, make, gcc217, etc.

Simplifying assumption: there is only one CPU

- We'll speak of “which process gets **the CPU**”
- But which process gets the CPU first? Unknown!



# Simple fork Example

```
#include <stdio.h>
#include <unistd.h>
int main(void)
{  printf("one\n");
   fork();
   printf("two\n");
   return 0;
}
```

What is the output?



# Simple fork Example Trace 1 (1)

Parent prints “one”

```
#include <stdio.h>
#include <unistd.h>
int main(void)
{   printf("one\n");
    fork();
    printf("two\n");
    return 0;
}
```



# Simple fork Example Trace 1 (2)

Parent forks child

```
#include <stdio.h>
#include <unistd.h>
int main(void)
{ printf("one\n");
  fork();
  printf("two\n");
  return 0;
}
```

Executing concurrently

```
#include <stdio.h>
#include <unistd.h>
int main(void)
{ printf("one\n");
  fork();
  printf("two\n");
  return 0;
}
```



# Simple fork Example Trace 1 (3)

OS gives CPU to child; child prints “two”

```
#include <stdio.h>
#include <unistd.h>
int main(void)
{ printf("one\n");
  fork();
  printf("two\n");
  return 0;
}
```

Executing concurrently

```
#include <stdio.h>
#include <unistd.h>
int main(void)
{ printf("one\n");
  fork();
  printf("two\n");
  return 0;
}
```



# Simple fork Example Trace 1 (4)

Child exits

```
#include <stdio.h>
#include <unistd.h>
int main(void)
{ printf("one\n");
  fork();
  printf("two\n");
  return 0;
}
```

Executing concurrently

```
#include <stdio.h>
#include <unistd.h>
int main(void)
{ printf("one\n");
  fork();
  printf("two\n");
  return 0;
}
```



# Simple fork Example Trace 1 (5)

OS gives CPU to parent; parent prints “two”

```
#include <stdio.h>
#include <unistd.h>
int main(void)
{  printf("one\n");
   fork();
   printf("two\n");
   return 0;
}
```





# Simple fork Example Trace 1 (6)

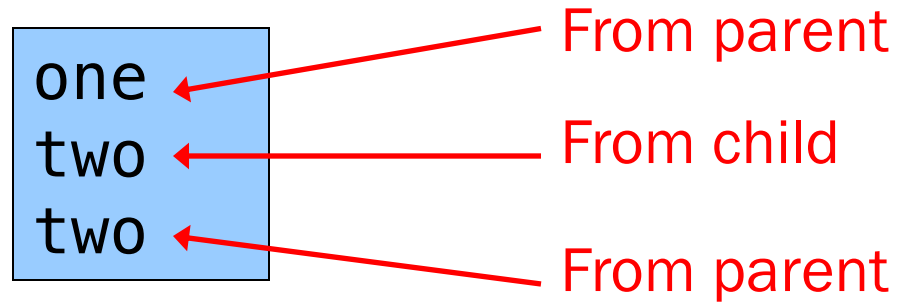
OS gives CPU to parent; parent prints “two”

```
#include <stdio.h>
#include <unistd.h>
int main(void)
{ printf("one\n");
  fork();
  printf("two\n");
  return 0;
}
```



# Simple fork Example Trace 1 Output

Output:





# Simple fork Example Trace 2 (1)

Parent prints "one"

```
#include <stdio.h>
#include <unistd.h>
int main(void)
{   printf("one\n");
    fork();
    printf("two\n");
    return 0;
}
```



# Simple fork Example Trace 2 (2)

Parent forks child

```
#include <stdio.h>
#include <unistd.h>
int main(void)
{ printf("one\n");
  fork();
  printf("two\n");
  return 0;
}
```

Executing concurrently

```
#include <stdio.h>
#include <unistd.h>
int main(void)
{ printf("one\n");
  fork();
  printf("two\n");
  return 0;
}
```



# Simple fork Example Trace 2 (3)

OS gives CPU to parent; parent prints "two"

```
#include <stdio.h>
#include <unistd.h>
int main(void)
{ printf("one\n");
  fork();
  printf("two\n");
  return 0;
}
```

Executing concurrently

```
#include <stdio.h>
#include <unistd.h>
int main(void)
{ printf("one\n");
  fork();
  printf("two\n");
  return 0;
}
```



# Simple fork Example Trace 2 (4)

Parent exits

```
#include <stdio.h>
#include <unistd.h>
int main(void)
{ printf("one\n");
  fork();
  printf("two\n");
  return 0;
}
```

Executing concurrently

```
#include <stdio.h>
#include <unistd.h>
int main(void)
{ printf("one\n");
  fork();
  printf("two\n");
  return 0;
}
```



## Simple fork Example Trace 2 (5)

OS gives CPU to child; child prints “two”

```
#include <stdio.h>
#include <unistd.h>
int main(void)
{   printf("one\n");
    fork();
    printf("two\n");
    return 0;
}
```



# Simple fork Example Trace 2 (6)

Child exits

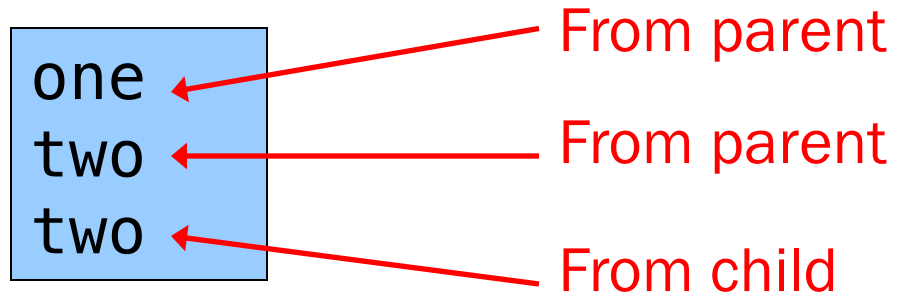
```
#include <stdio.h>
#include <unistd.h>
int main(void)
{ printf("one\n");
  fork();
  printf("two\n");
  return 0;
}
```





# Simple fork Example Trace 2 Output

Output:





# iClicker Question



Q: What is the output of this program?

- A. child: 0  
parent: 2
- B. parent: 2  
child: 0
- C. child: 0  
parent: 1
- D. parent: 2  
child: 1
- E. A or B

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

The answer is E.

See following slides.



# fork Example Trace 1 (1)

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

x = 1



# fork Example Trace 1 (2)

## Parent forks child

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

x = 1

Executing concurrently

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

x = 1



# fork Example Trace 1 (3)

Assume OS gives CPU to child

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

x = 1

0  
Executing concurrently

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

x = 1



# fork Example Trace 1 (4)

Child decrements its x, and prints "child: 0"

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

x = 1

Executing concurrently

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

x = 0



# fork Example Trace 1 (5)

Child exits; OS gives CPU to parent

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

x = 1

Executing concurrently

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

x = 0



# fork Example Trace 1 (6)

In parent, fork() returns process id of child

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

x = 1

Process id of child





# fork Example Trace 1 (7)

Parent increments its x, and prints “parent: 2”

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

x = 2



# fork Example Trace 1 (8)

## Parent exits

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

x = 2



# fork Example Trace 1 Output

Example trace 1 output:

```
Child: 0  
Parent: 2
```



# fork Example Trace 2 (1)

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

x = 1



# fork Example Trace 2 (2)

## Parent forks child

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

x = 1

Executing concurrently

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

x = 1



# fork Example Trace 2 (3)

Assume OS gives CPU to parent

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;
  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

Process ID  
of child

x = 1

Executing concurrently

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;
  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

x = 1



# fork Example Trace 2 (4)

Parent increments its x and prints "parent: 2"

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

x = 2

Executing concurrently

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

x = 1



# fork Example Trace 2 (5)

Parent exits; OS gives CPU to child

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

**x = 2**

Executing concurrently

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

**x = 1**





# fork Example Trace 2 (6)

In child, fork() returns 0

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;
  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

0

x = 1



# fork Example Trace 2 (7)

Child decrements its x and prints “child: 0”

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

x = 0



# fork Example Trace 2 (8)

## Child exits

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
int main(void)
{ pid_t pid;
  int x = 1;

  pid = fork();
  if (pid == 0)
  { x--;
    printf("child: %d\n", x);
    exit(0);
  }
  else
  { x++;
    printf("parent: %d\n", x);
    exit(0);
  }
}
```

x = 0



# fork Example Trace 2 Output

Example trace 2 output:

```
Parent: 2  
Child: 0
```

```
armlab01:~$ for i in `seq 1 10000`; do ./fpe | head -n 1; done | sort | uniq -c  
56 child: 0  
9944 parent: 2
```



# wait System-Level Function

## Problem:

- How to control execution order?

## Solution:

- Parent calls `wait()`

```
pid_t wait(int *status);
```

- Suspends execution of the calling process until one of its children terminates
- If `status` is not `NULL`, stores status information in the `int` to which it points; this integer can be inspected with macros [see man page for details].
- On success, returns the process ID of the terminated child
- On error, returns `-1`



# iClicker Question



Q: What is the output of this program?

- A. child  
parent
- B. parent  
child
- C. something other than A or B
- D. A or B
- E. A or C

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>
int main(void)
{ pid_t pid;
  pid = fork();
  if (pid == 0)
  { printf("child\n");
    exit(0);
  }
  wait(NULL);
  printf("parent\n");
  return 0;
}
```

The answer is A.

See following slides.



# wait Example Trace 1 (1)

Parent forks child

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>
int main(void)
{   pid_t pid;
    pid = fork();
    if (pid == 0)
    {   printf("child\n");
        exit(0);
    }
    wait(NULL);
    printf("parent\n");
    return 0;
}
```

Executing concurrently

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>
int main(void)
{   pid_t pid;
    pid = fork();
    if (pid == 0)
    {   printf("child\n");
        exit(0);
    }
    wait(NULL);
    printf("parent\n");
    return 0;
}
```



# wait Example Trace 1 (2)

OS gives CPU to parent

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>
int main(void)
{ pid_t pid;
  pid = fork();
  if (pid == 0)
  { printf("child\n");
    exit(0);
  }
  wait(NULL);
  printf("parent\n");
  return 0;
}
```

Executing concurrently

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>
int main(void)
{ pid_t pid;
  pid = fork();
  if (pid == 0)
  { printf("child\n");
    exit(0);
  }
  wait(NULL);
  printf("parent\n");
  return 0;
}
```





# wait Example Trace 1 (3)

Parent calls wait ( )

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>
int main(void)
{ pid_t pid;
  pid = fork();
  if (pid == 0)
  { printf("child\n");
    exit(0);
  }
  wait(NULL);
  printf("parent\n");
  return 0;
}
```

Executing concurrently

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>
int main(void)
{ pid_t pid;
  pid = fork();
  if (pid == 0)
  { printf("child\n");
    exit(0);
  }
  wait(NULL);
  printf("parent\n");
  return 0;
}
```



# wait Example Trace 1 (4)

OS gives CPU to child

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>
int main(void)
{ pid_t pid;
  pid = fork();
  if (pid == 0)
  { printf("child\n");
    exit(0);
  }
  wait(NULL);
  printf("parent\n");
  return 0;
}
```

Executing concurrently

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>
int main(void)
{ pid_t pid;
  pid = fork();
  if (pid == 0)
  { printf("child\n");
    exit(0);
  }
  wait(NULL);
  printf("parent\n");
  return 0;
}
```



# wait Example Trace 1 (5)

Child prints “child” and exits

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>
int main(void)
{ pid_t pid;
  pid = fork();
  if (pid == 0)
  { printf("child\n");
    exit(0);
  }
  wait(NULL);
  printf("parent\n");
  return 0;
}
```

Executing concurrently

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>
int main(void)
{ pid_t pid;
  pid = fork();
  if (pid == 0)
  { printf("child\n");
    exit(0);
  }
  wait(NULL);
  printf("parent\n");
  return 0;
}
```



# wait Example Trace 1 (6)

Parent returns from call of wait(), prints “parent”, exits

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>
int main(void)
{ pid_t pid;
  pid = fork();
  if (pid == 0)
  { printf("child\n");
    exit(0);
  }
  wait(NULL);
  printf("parent\n");
  return 0;
}
```



# wait Example Trace 1 Output

Example trace 1 output

```
child  
parent
```



# wait Example Trace 2 (1)

Parent forks child

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>
int main(void)
{   pid_t pid;
    pid = fork();
    if (pid == 0)
    {   printf("child\n");
        exit(0);
    }
    wait(NULL);
    printf("parent\n");
    return 0;
}
```

Executing concurrently

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>
int main(void)
{   pid_t pid;
    pid = fork();
    if (pid == 0)
    {   printf("child\n");
        exit(0);
    }
    wait(NULL);
    printf("parent\n");
    return 0;
}
```



# wait Example Trace 2 (2)

OS gives CPU to child

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>
int main(void)
{ pid_t pid;
  pid = fork();
  if (pid == 0)
  { printf("child\n");
    exit(0);
  }
  wait(NULL);
  printf("parent\n");
  return 0;
}
```

Executing concurrently

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>
int main(void)
{ pid_t pid;
  pid = fork();
  if (pid == 0)
  { printf("child\n");
    exit(0);
  }
  wait(NULL);
  printf("parent\n");
  return 0;
}
```



# wait Example Trace 2 (3)

Child prints “child” and exits

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>
int main(void)
{ pid_t pid;
  pid = fork();
  if (pid == 0)
  { printf("child\n");
    exit(0);
  }
  wait(NULL);
  printf("parent\n");
  return 0;
}
```

Executing concurrently

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>
int main(void)
{ pid_t pid;
  pid = fork();
  if (pid == 0)
  { printf("child\n");
    exit(0);
  }
  wait(NULL);
  printf("parent\n");
  return 0;
}
```





# wait Example Trace 2 (4)

OS gives CPU to parent

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>
int main(void)
{ pid_t pid;
  pid = fork();
  if (pid == 0)
  { printf("child\n");
    exit(0);
  }
  wait(NULL);
  printf("parent\n");
  return 0;
}
```



# wait Example Trace 2 (5)

Parent calls `wait ( )`; returns immediately

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>
int main(void)
{  pid_t pid;
   pid = fork();
   if (pid == 0)
   {  printf("child\n");
      exit(0);
   }
   wait(NULL);
   printf("parent\n");
   return 0;
}
```



# wait Example Trace 2 (6)

Parent prints “parent” and exits

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>
int main(void)
{ pid_t pid;
  pid = fork();
  if (pid == 0)
  { printf("child\n");
    exit(0);
  }
  wait(NULL);
  printf("parent\n");
  return 0;
}
```



# wait Example Trace 2 Output

Example trace 2 output

```
child  
parent
```

Same as trace 1 output!



# execvp System-Level Function

Problem: How to execute a new program?

- Usually, in the newly-created child process

Solution: `execvp( )`

```
int execvp(const char *file, char *const argv[]);
```

- Replaces the current process image with a new process image
- Provides an array of pointers to null-terminated strings that represent the argument list available to the new program
  - The first argument, by convention, should point to the program's filename
  - The array of pointers must be terminated by a NULL pointer



# execvp System-Level Function

Example: Execute “cat readme”

```
char *newCmd;  
char *newArgv[3];  
newCmd = "cat";  
newArgv[0] = "cat";  
newArgv[1] = "readme";  
newArgv[2] = NULL;  
execvp(newCmd, newArgv);
```

- First argument: name of program to be executed
- Second argument: argv to be passed to main() of new program
  - Must begin with program name, end with NULL



# execvp Failure

## fork()

- If successful, returns **two** times
  - Once in parent, once in child

## execvp()

- If successful, returns **zero** times
  - Calling program is overwritten with new program
- Corollary:
  - If `execvp()` returns, then it must have failed

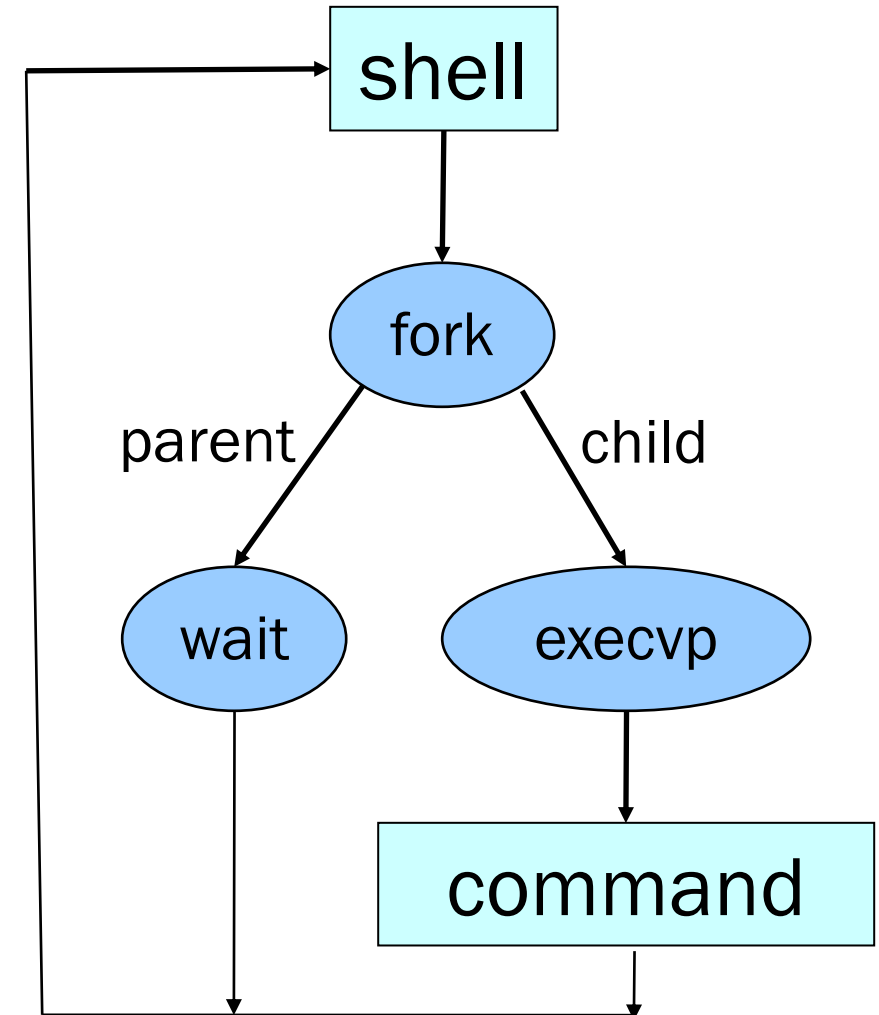
```
char *newCmd;  
char *newArgv[3];  
newCmd = "cat";  
newArgv[0] = "cat";  
newArgv[1] = "readme";  
newArgv[2] = NULL;  
execvp(newCmd, newArgv);  
fprintf(stderr, "exec failed\n");  
exit(EXIT_FAILURE);
```



# Shell Structure

- Parent (shell) reads & parses the command line
- **Parent forks child and waits**
- **Child calls execvp to execute command**
- Child exits, parent returns from wait and repeats

```
while (1) {  
    Parse command line  
    Assign values to somepgm, someargv  
    pid = fork();  
    if (pid == 0) {  
        /* in child */  
        execvp(somepgm, someargv);  
        fprintf(stderr, "exec failed\n");  
        exit(EXIT_FAILURE);  
    }  
    /* in parent */  
    wait(NULL);  
}
```







# Aside: system Function

## Common combination of operations

- `fork()` to create a new child process
- `execvp()` to execute new program in child process
- `wait()` in the parent process for the child to complete

## Single call that combines all three

- `int system(const char *cmd);`

## Example:

```
#include <stdlib.h>
int main(void)
{  system("cat readme");
  return 0;
}
```



# Agenda



## Processes

Illusion: Private address space  
Illusion: Private control flow



## Process management in C

Creating new processes  
Waiting for termination  
Executing new programs



## Unix Process Control

Exceptions  
Signals



# Exceptions

## Exception

- An abrupt change in control flow of a running program corresponding to a change in process state

\*Note: Exceptions in OS  $\neq$  exceptions in Java

Implemented using  
**try/catch** and **throw** statements



# Synchronous Exceptions

Some exceptions are **synchronous**

- Occur as result of actions of executing program
- Examples:
  - **System call:** Application requests I/O
  - **System call:** Application requests more heap memory
  - Application pgm attempts integer division by 0
  - Application pgm attempts to access privileged memory
  - Application pgm accesses variable that is not in physical memory



# Asynchronous Exceptions

Some exceptions are **asynchronous**

- Do not occur (directly) as result of actions of executing program

• Examples:

- User presses key on keyboard



- Packet received over network



- Disk controller finishes reading data

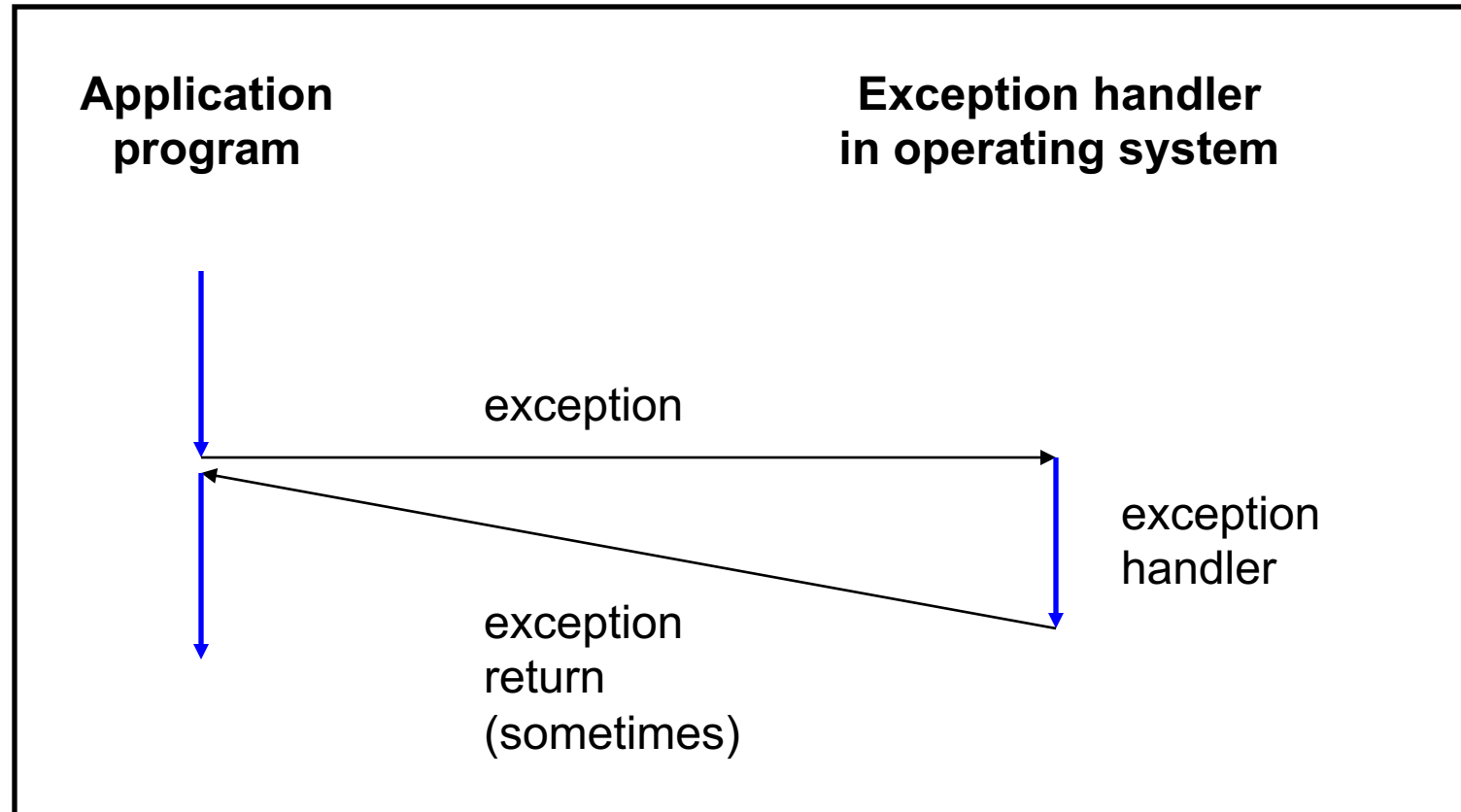


- Hardware timer expires





# Exceptional Control Flow





# Exceptions vs. Function Calls

Handling an exception is **similar to** calling a function

- Control transfers from original code to other code
- Other code executes
- Control returns to some instruction in original code

Handling an exception is **different from** calling a function

- CPU saves **additional data**
  - E.g. values of all registers
- CPU pushes data onto **OS's stack**, not application pgm's stack
- Handler runs in **kernel/privileged mode**, not in **user mode**
  - Handler can execute all instructions and access all memory
- Control **might return** to some instruction in original code
  - Sometimes control returns to **next** instruction
  - Sometimes control returns to **current** instruction
  - Sometimes control does not return at all!

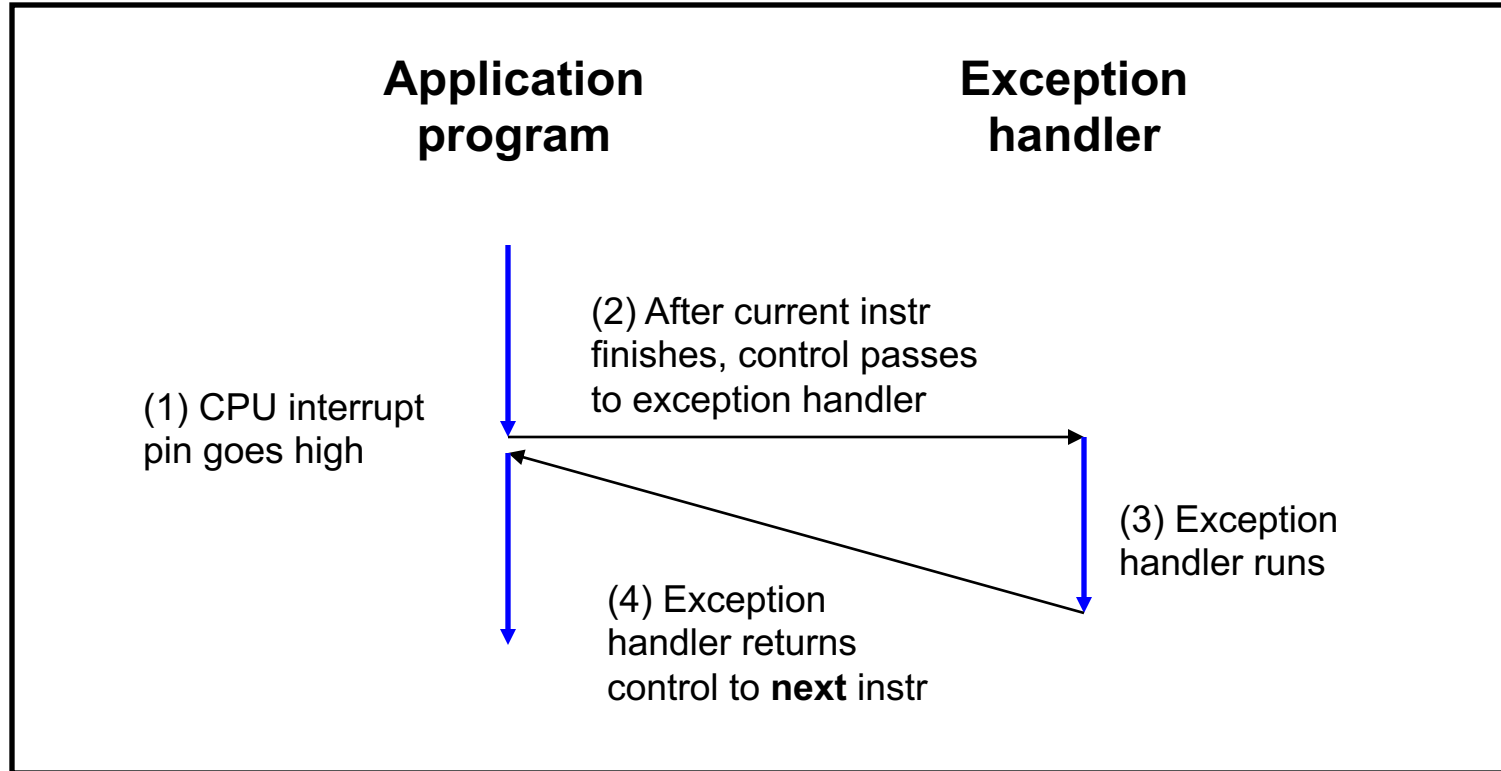
# Classes of Exceptions



There are 4 classes of exceptions...



# (1) Interrupts



**Occurs when:** External (off-CPU) device requests attention

**Examples:**

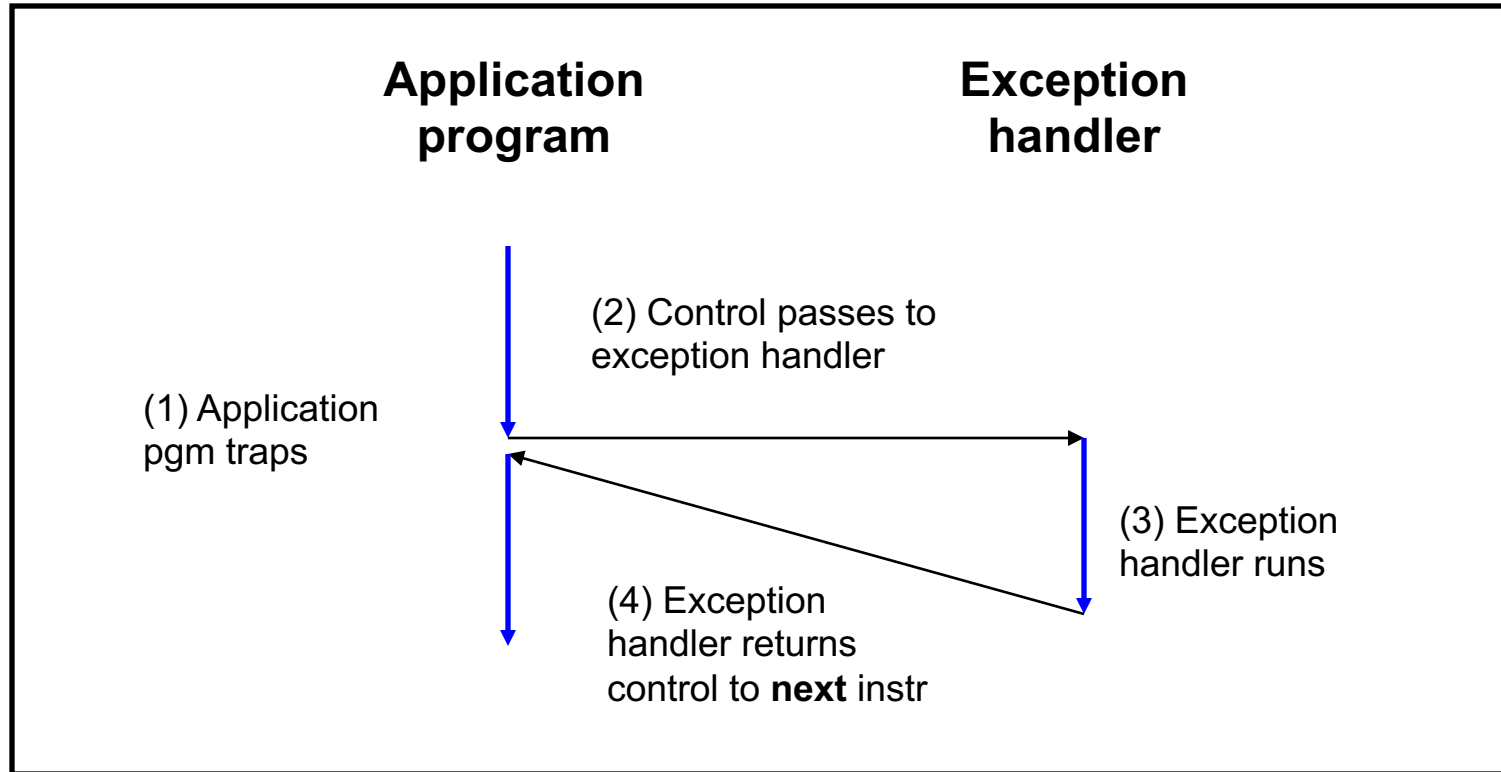
User presses key

Disk controller finishes reading/writing data

Network packet arrives



## (2) Traps



**Occurs when:** Application pgm requests OS service

**Examples:**

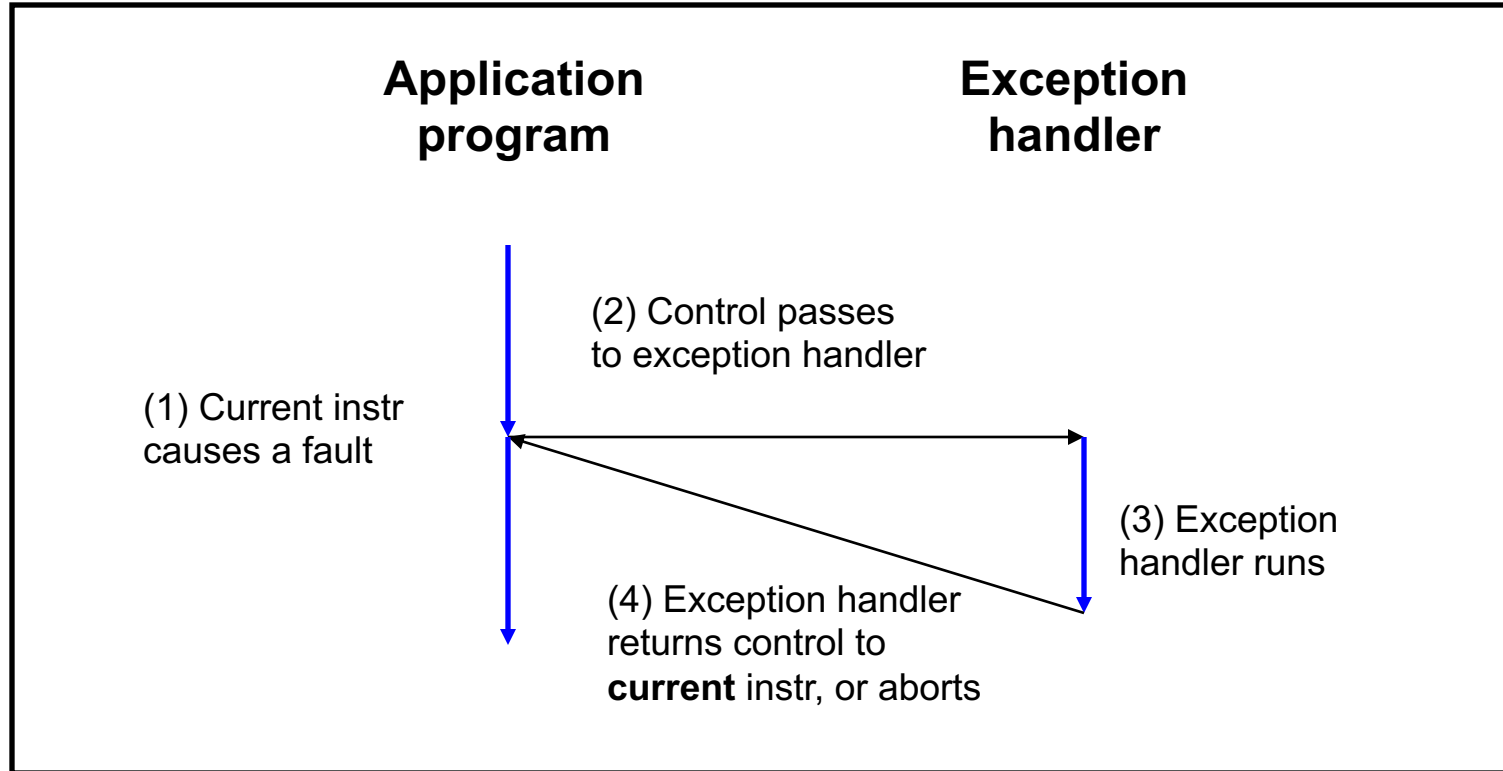
Application pgm requests I/O

Application pgm requests more heap memory

Traps provide a function-call-like interface between application pgm and OS



# (3) Faults



**Occurs when:** Application pgm causes a (possibly recoverable) error

**Examples:**

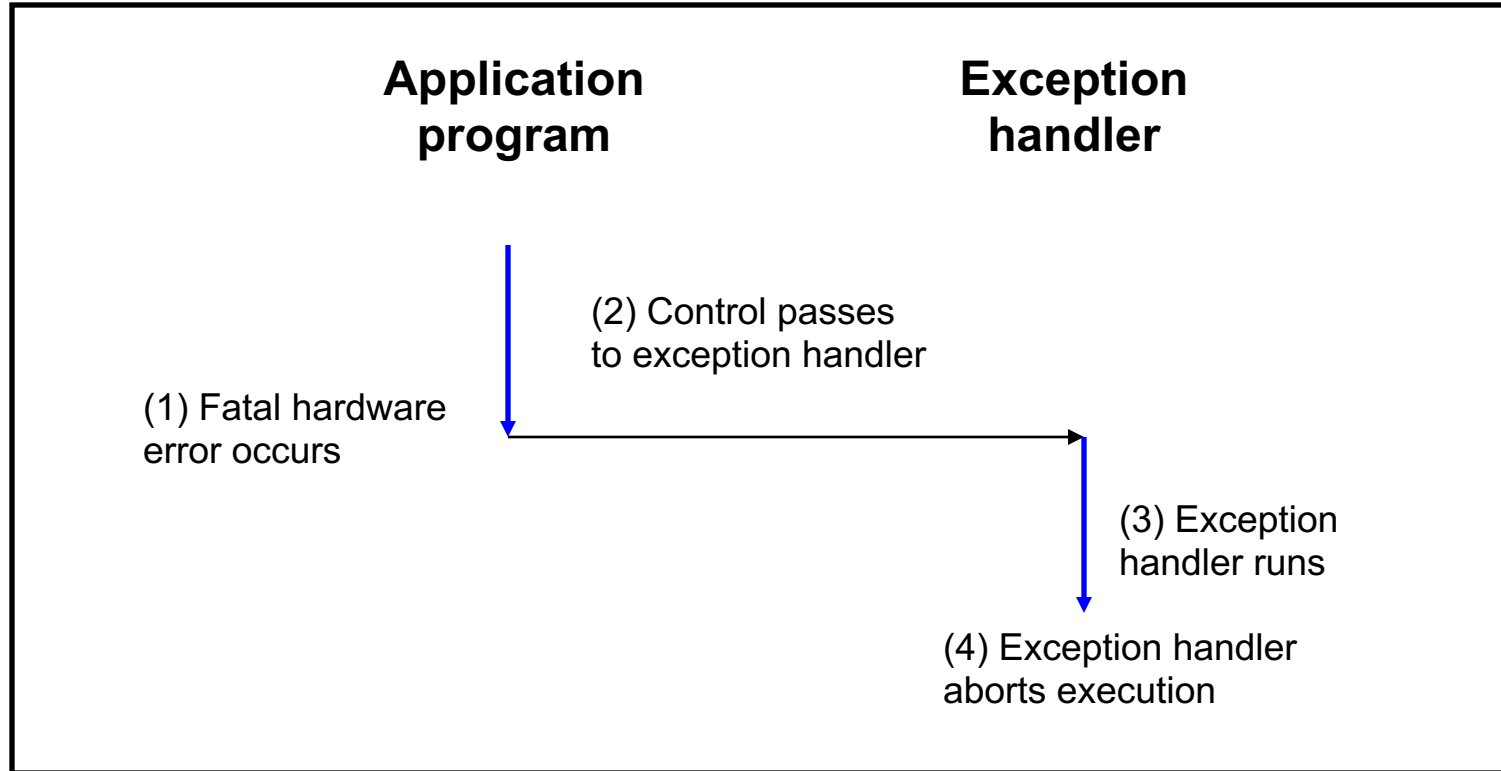
Application pgm divides by 0

Application pgm accesses privileged memory (seg fault)

Application pgm accesses data that is not in physical memory (page fault)



# (4) Aborts



**Occurs when:** HW detects a non-recoverable error

**Example:**

Parity check indicates corruption of memory bit (overheating, cosmic ray!, etc.)



# Summary of Exception Classes

Class	Occurs when	Asynch /Synch	Return Behavior
<b>Interrupt</b>	External device requests attention	Asynch	Return to next instr
<b>Trap</b>	Application pgm requests OS service	Synch	Return to next instr
<b>Fault</b>	Application pgm causes (maybe recoverable) error	Synch	Return to current instr (maybe)
<b>Abort</b>	HW detects non-recoverable error	Synch-ish	Do not return



# Process Control Examples

Exactly what happens when you:

## Type Ctrl-c?

- Keystroke generates **interrupt**
- OS handles interrupt
- OS sends process a 2/SIGINT **signal**

## Type Ctrl-z?

- Keystroke generates **interrupt**
- OS handles interrupt
- OS sends process a 20/SIGTSTP **signal**



# Signals Overview

**Signal:** A notification of an exception

Typical signal sequence:

- Process P is executing
- Exception occurs (interrupt, trap, fault, or abort)
- OS gains control of CPU
- OS wishes to inform process P that something happened
- OS **sends** a signal to process P
  - OS sets a bit in **pending bit vector** of process P
  - Indicates that OS is sending a signal of type X to process P
  - A signal of type X is **pending** for process P



# Signals Overview (cont.)

## Typical signal sequence (cont.):

- Sometime later...
- OS is ready to give CPU back to process P
- OS checks `pending` for process P, sees that signal of type X is pending
- OS forces process P to **receive** signal of type X
  - OS clears bit in process P's `pending`
- Process P executes action for signal of type X
  - Normally process P executes **default action** for that signal
  - If **signal handler** was installed for signal of type X, then process P executes signal handler
  - Action might terminate process P; otherwise...
- Process P resumes where it left off





# Sending Signals via Keystrokes

User can send three signals from keyboard:

- **Ctrl-c** ⇒ **2/SIGINT** signal
  - Default action is “terminate”
- **Ctrl-z** ⇒ **20/SIGTSTP** signal
  - Default action is “stop until next 18/SIGCONT”
- **Ctrl-\** ⇒ **3/SIGQUIT** signal
  - Default action is “terminate”



# Examples of Non-keyboard Signals

## Process makes illegal memory reference

- Segmentation fault occurs
- OS gains control of CPU
- OS sends 11/SIGSEGV signal to process
- Process receives 11/SIGSEGV signal
- Default action for 11/SIGSEGV signal is “terminate”



<https://xkcd.com/371/>



# Signals, signals everywhere

List of the predefined signals, learn many details with these commands:

```
$ kill -l
 1) SIGHUP          2) SIGINT          3) SIGQUIT        4) SIGILL
 5) SIGTRAP        6) SIGABRT        7) SIGBUS         8) SIGFPE
 9) SIGKILL        10) SIGUSR1       11) SIGSEGV       12) SIGUSR2
13) SIGPIPE       14) SIGALRM       15) SIGTERM       17) SIGCHLD
18) SIGCONT       19) SIGSTOP       20) SIGTSTP       21) SIGTTIN
22) SIGTTOU       23) SIGURG        24) SIGXCPU       25) SIGXFSZ
26) SIGVTALRM     27) SIGPROF       28) SIGWINCH      29) SIGIO
30) SIGPWR        31) SIGSYS        34) SIGRTMIN      35) SIGRTMIN+1
36) SIGRTMIN+2   37) SIGRTMIN+3   38) SIGRTMIN+4   39) SIGRTMIN+5
40) SIGRTMIN+6   41) SIGRTMIN+7   42) SIGRTMIN+8   43) SIGRTMIN+9
44) SIGRTMIN+10  45) SIGRTMIN+11  46) SIGRTMIN+12  47) SIGRTMIN+13
48) SIGRTMIN+14  49) SIGRTMIN+15  50) SIGRTMAX-14  51) SIGRTMAX-13
52) SIGRTMAX-12  53) SIGRTMAX-11  54) SIGRTMAX-10  55) SIGRTMAX-9
56) SIGRTMAX-8   57) SIGRTMAX-7   58) SIGRTMAX-6   59) SIGRTMAX-5
60) SIGRTMAX-4   61) SIGRTMAX-3   62) SIGRTMAX-2   63) SIGRTMAX-1
64) SIGRTMAX

$ man 7 signal
```



# Handling Signals

Each signal type has a default action

- For most signal types, default action is “terminate”

A program can **install** a **signal handler**

- To change action of (almost) any signal type



# Installing a Signal Handler

## `signal()` function

- `sighandler_t signal(int iSig, sighandler_t pfHandler);`
- Install function `pfHandler` as the handler for signals of type `iSig`
- `pfHandler` is a function pointer:  
`typedef void (*sighandler_t)(int);`
- Return the old handler on success, `SIG_ERR` on error
- After call, `(*pfHandler)` is invoked whenever process receives a signal of type `iSig`



# SIG\_DFL

Predefined value: **SIG\_DFL**

Use as argument to `signal()` to restore default action

```
int main(void)
{
    ...
    signal(SIGINT, somehandler);
    ...
    signal(SIGINT, SIG_DFL);
    ...
}
```

Subsequently, process will handle 2/SIGINT signals using default action for 2/SIGINT signals (“terminate”)



# SIG\_IGN

Predefined value: **SIG\_IGN**

Use as argument to `signal()` to ignore signals

```
int main(void)
{
    ...
    signal(SIGINT, SIG_IGN);
    ...
}
```

Subsequently, process will ignore 2/SIGINT signals



# Signal Handling Example 1

Program testsignal.c:

```
#define _GNU_SOURCE /* Use modern handling style */
#include <stdio.h>
#include <signal.h>

static void myHandler(int iSig)
{ printf("In myHandler with argument %d\n", iSig);
}

int main(void)
{ signal(SIGINT, myHandler);
  printf("Entering an infinite loop\n");
  for (;;)
    ;
  return 0; /* Never get here. */
}
```

```
armlab01:~/Test$ ./testsignal
Entering an infinite loop
^CIn myHandler with argument 2
^CIn myHandler with argument 2
^CIn myHandler with argument 2
^CIn myHandler with argument 2
^CIn myHandler with argument 2
^CIn myHandler with argument 2
^CIn myHandler with argument 2
^CIn myHandler with argument 2
^CIn myHandler with argument 2
^CIn myHandler with argument 2
^Z
[1]+  Stopped                  ./testsignal
armlab01:~/Test$ fg
./signal
^CIn myHandler with argument 2
^CIn myHandler with argument 2
^CIn myHandler with argument 2
^CIn myHandler with argument 2
^CIn myHandler with argument 2
^CQuit
```





# Signal Handling Example 2

Program generates lots of temporary data

- Stores the data in a temporary file
- Must delete the file before exiting

```
...
int main(void)
{  FILE *psFile;
   psFile = fopen("temp.txt", "w");
   ...
   fclose(psFile);
   remove("temp.txt");
   return 0;
}
```



# Example 2 Problem

What if user types Ctrl-c?

- OS sends a 2/SIGINT signal to the process
- Default action for 2/SIGINT is “terminate”

**Problem:** The temporary file is not deleted

- Process terminates before `remove("temp.txt")` is executed

**Challenge:** Ctrl-c could happen at any time

- Which line of code will be interrupted???

**Solution:** Install a signal handler

- Define a “clean up” function to delete the file
- Install the function as a signal handler for 2/SIGINT



# Example 2 Solution

```
...
static FILE *psFile; /* Must be global. */
static void cleanup(int iSig)
{  fclose(psFile);
   remove("temp.txt");
   exit(0);
}
int main(void)
{  ...
   psFile = fopen("temp.txt", "w");
   signal(SIGINT, cleanup);
   ...
   cleanup(0);
   return 0; /* Never get here. */
}
```

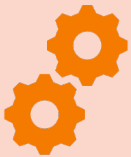


# Agenda



## Processes

Illusion: Private address space  
Illusion: Private control flow



## Process management in C

Creating new processes  
Waiting for termination  
Executing new programs



## Unix Process Control

Exceptions  
Signals