Concurrency in Go

9/15/22
Go Resources

https://tour.golang.org/list
https://play.golang.org
https://gobyexample.com
Today’s Precept...

1. Two synchronization mechanisms
   a. Locks
   b. Channels
2. Mapreduce
Two synchronization mechanisms

**Locks** - limit access to a critical section

**Channels** - pass information across processes using a queue
Example: Bank account

Bob

Read $b = 100$

$b = b + 10$

Write $b = 110$

Bank Account

100

110

120

Alice

Read $b = 110$

$b = b + 10$

Write $b = 120$
Example: Bank account

Bob

Read $b = 100$

$b = b + 10$

Write $b = 110$

Alice

Read $b = 100$

$b = b + 10$

Write $b = 110$
What went wrong?

Changes to balance are not *atomic*

```java
func Deposit(amount) {
    read balance
    balance = balance + amount
    write balance
}
```
What went wrong?

Changes to balance are not **atomic**

```java
func Deposit(amount) {
    lock balanceLock
    read balance
    balance = balance + amount
    write balance
    unlock balanceLock
}
```
package account

import "sync"

type Account struct {
    balance int
}

func NewAccount(init int) Account {
    return Account{balance: init}
}

func (a *Account) Deposit(v int) {
    a.mu.Lock()
    defer a.mu.Unlock()
    a.balance += v
}

func (a *Account) CheckBalance() int {
    a.mu.Lock()
    defer a.mu.Unlock()
    return a.balance
}

func (a *Account) Withdraw(v int) {
    a.mu.Lock()
    defer a.mu.Unlock()
    a.balance -= v
}

func (a *Account) Deposit(v int) {
    a.mu.Lock()
    defer a.mu.Unlock()
    a.balance += v
}
package account

import "sync"

type Account struct {
    balance int
    mu sync.Mutex
}

func NewAccount(init int) Account {
    return Account{balance: init}
}

func (a *Account) Deposit(v int) {
    a.mu.Lock()
    defer a.mu.Unlock()
    a.balance += v
}

func (a *Account) CheckBalance() int {
    a.mu.Lock()
    defer a.mu.Unlock()
    return a.balance
}

func (a *Account) Withdraw(v int) {
    a.mu.Lock()
    defer a.mu.Unlock()
    a.balance -= v
}

func (a *Account) Deposit(v int) {
    a.mu.Lock()
    defer a.mu.Unlock()
    a.balance += v
}
Read Write Locks in Go

package account

import "sync"

type Account struct {
    balance int
}

func NewAccount(init int) Account {
    return Account{balance: init}
}

func (a *Account) Deposit(v int) {
    a.rwLock.Lock()
    defer a.rwLock.Unlock()
    a.balance += v
}

func (a *Account) CheckBalance() int {
    a.rwLock.RLock()
    defer a.rwLock.RUnlock()
    return a.balance
}

func (a *Account) Withdraw(v int) {
    a.rwLock.Lock()
    defer a.rwLock.Unlock()
    a.balance -= v
}

func (a *Account) Deposit(v int) {
    a.rwLock.Lock()
    defer a.rwLock.Unlock()
    a.balance += v
}
package account

import "sync"

type Account struct {
    balance int
    rwLock sync.RWMutex
}

func NewAccount(init int) Account {
    return Account{balance: init}
}

func (a *Account) CheckBalance() int {
    a.rwLock.RLock()
    defer a.rwLock.RUnlock()
    return a.balance
}

func (a *Account) Withdraw(v int) {
    a.rwLock.Lock()
    defer a.rwLock.Unlock()
    a.balance -= v
}

func (a *Account) Deposit(v int) {
    a.rwLock.Lock()
    defer a.rwLock.Unlock()
    a.balance += v
}
Two Solutions to the Same Problem

**Locks:**

Multiple threads can reference same memory location

Use lock to ensure only one thread is updating it at any given time

<table>
<thead>
<tr>
<th>T1</th>
<th>T2</th>
<th>T3</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
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</tbody>
</table>

| 0x1000: | 100 |

**Channels:**

Data item initially stored in channel

Threads must request item from channel, make updates, and return item to channel

<table>
<thead>
<tr>
<th>T1</th>
<th>T2</th>
<th>T3</th>
</tr>
</thead>
<tbody>
<tr>
<td>100</td>
<td>110</td>
<td></td>
</tr>
</tbody>
</table>

C
package account

type Account struct {
    // Fill in Here
}

func NewAccount(init int) Account {
    // Fill in Here
}

func (a *Account) CheckBalance() int {
    // What goes Here?
}

func (a *Account) Withdraw(v int) {
    // ???
}

func (a *Account) Deposit(v int) {
    // ???
}
package account

type Account struct {
    balance chan int
}

func NewAccount(init int) Account {
    a := Account{
        balance: make(chan int, 1)
    }
    a.balance <- init
    return a
}

func (a *Account) CheckBalance() int {
    // What goes Here?
}

func (a *Account) Withdraw(v int) {
    // ???
}

func (a *Account) Deposit(v int) {
    // ???
}
package account

type Account struct {
    balance chan int
}

func NewAccount(init int) Account {
    a := Account{
        balance: make(chan int, 1)
    }
    a.balance <- init
    return a
}

func (a *Account) CheckBalance() int {
    bal := <-a.balance
    a.balance <- bal
    return bal
}

func (a *Account) Withdraw(v int) {
    // ???
}

func (a *Account) Deposit(v int) {
    // ???
}
Bank Account Code (using channels)

```go
cpackage account
type Account struct {
    balance chan int
}

func NewAccount(init int) Account {
    a := Account{
        balance: make(chan int, 1)
    }
    a.balance <- init
    return a
}

func (a *Account) CheckBalance() int {
    bal := <-a.balance
    a.balance <- bal
    return bal
}

func (a *Account) Withdraw(v int) {
    bal := <-a.balance
    a.balance <- (bal - v)
}

func (a *Account) Deposit(v int) {
    //???
}
```
package account

type Account struct {
    balance chan int
}

func NewAccount(init int) Account {
    a := Account{
        balance: make(chan int, 1)
    }
    a.balance <- init
    return a
}

func (a *Account) CheckBalance() int {
    bal := <-a.balance
    a.balance <- bal
    return bal
}

func (a *Account) Withdraw(v int) {
    bal := <-a.balance
    a.balance <- (bal - v)
}

func (a *Account) Deposit(v int) {
    bal := <-a.balance
    a.balance <- (bal + v)
}
Go channels

**Channels** also allow us to safely communicate between **goroutines**

```go
result := make(chan int, numWorkers)

// Launch workers
for i := 0; i < numWorkers; i++ {
    go func() {
        doWork()
        result <- i
    }()
}

// Wait until all worker threads have finished
for i := 0; i < numWorkers; i++ {
    handleResult(<-result)
}
fmt.Println("Done!")
```
Go channels

Easy to express asynchronous RPC

Awkward to express this using locks

```go
result := make(chan int, numServers)

// Send query to all servers
for i := 0; i < numServers; i++ {
    go func() {
        resp := // ... send RPC to server
        result <- resp
    }()
}

// Return as soon as the first server responds
handleResponse(<-result)
```
Select statement

`select` allows a goroutine to wait on multiple channels at once

```go
for {
    select {
        case money := <-dad:
            buySnacks(money)
        case money := <-mom:
            buySnacks(money)
    }
}
```
Select statement

`select` allows a goroutine to wait on multiple channels at once

```go
for {
    select {
        case money := <-dad:
            buySnacks(money)
        case money := <-mom:
            buySnacks(money)
        case default:
            starve()
            time.Sleep(5 * time.Second)
    }
}
```
Handle timeouts using select

// Asynchronously request an answer from server, timing out after X seconds
result := make(chan int)
timeout := make(chan bool)

// Ask server
go func() {
    response := // ... send RPC
    result <- response
}()

// Start timer
go func() {
    time.Sleep(5 * time.Second)
    timeout <- true
}()

// Wait on both channels
select {
    case res := <-result:
        handleResult(res)
    case <-timeout:
        fmt.Println("Timeout!")
}
Exercise: Implementing a mutex using channels

type Lock struct {
    // ???
}

func NewLock() Lock {
    // ???
}

func (l *Lock) Lock() {
    // ???
}

func (l *Lock) Unlock() {
    // ???
}
Exercise: Implementing a mutex using channels

type Lock struct {
    ch chan bool
}

func NewLock() Lock {
    // ???
}

func (l *Lock) Lock() {
    // ???
}

func (l *Lock) Unlock() {
    // ???
}
Exercise: Implementing a mutex using channels

type Lock struct {
    ch chan bool
}

func NewLock() Lock {
    lock := Lock{make(chan bool, 1)}
    lock.ch <- true
    return lock
}

func (l *Lock) Lock() {
    // ???
}

func (l *Lock) Unlock() {
    // ???
}
Exercise: Implementing a mutex using channels

type Lock struct {
    ch chan bool
}

func NewLock() Lock {
    lock := Lock{make(chan bool, 1)}
    lock.ch <- true
    return lock
}

func (l *Lock) Lock() {
    <-lock.ch
}

func (l *Lock) Unlock() {
    // ???
}
Exercise: Implementing a mutex using channels

```go
type Lock struct {
    ch chan bool
}

func NewLock() Lock {
    lock := Lock{make(chan bool, 1)}
    lock.ch <- true
    return lock
}

func (l *Lock) Lock() {
    <-lock.ch
}

func (l *Lock) Unlock() {
    lock.ch <- true
}
```
Mutexes vs. Semaphores

**Mutexes** allow 1 process to enter critical section at a time. Allows at most \( n \) concurrent accesses.

**Semaphores** allow up to \( N \) processes to enter critical section simultaneously.
Outline

Two synchronization mechanisms

- Locks
- Channels

Mapreduce
How much wood would a woodchuck chuck if a woodchuck could chuck wood?

how: 1, much: 1, wood: 2, would: 1, a: 2, woodchuck: 2, chuck: 2, if: 1, could: 1
Application: Word count

Locally: tokenize and put words in a hash map

How do you parallelize this?

Partition the document into \( n \) partitions.

Build \( n \) hash maps, one for each partition

Merge the \( n \) hash maps (by key)
How do you do this in a distributed environment?
When in the Course of human events, it becomes necessary for one people to dissolve the political bands which have connected them with another, and to assume, among the Powers of the earth, the separate and equal station to which the Laws of Nature and of Nature's God entitle them, a decent respect to the opinions of mankind requires that they should declare the causes which impel them to the separation.
When in the Course of human events, it becomes necessary for one people to dissolve the political bands which have connected them with another, and to assume, among the Powers of the earth, the separate and equal station to which the Laws of Nature and of Nature's God entitle them, a decent respect to the opinions of mankind requires that they should declare the causes which impel them to the separation.

**Partition**
When in the Course of human events, it becomes necessary for one people to dissolve the political bands which have connected them with another, and to assume, among the Powers of the earth, the separate and equal station to which the Laws of Nature and of Nature's God entitle them, a decent respect to the opinions of mankind requires that they should declare the causes which impel them to the separation.

Partition
When in the Course of human events, it becomes necessary for one people to dissolve the political bands which have connected them with another, and to assume, requires that they should declare the causes which impel them to the separation.

Nature and of Nature’s God entitle them, a decent respect to the opinions of mankind among the Powers of the earth, the separate and equal station to which the Laws of
when: 1, in: 1,
the: 1, course: 1,
of: 1, human: 1,
events: 1, it: 1

dissolve: 1, the: 2,
political: 1, bands: 1,
which: 1, have: 1,
connected: 1, them: 1 ...

requires: 1, that: 1,
they: 1, should: 1,
declare: 1, the: 1,
causes: 1, which: 1 ...

nature: 2, and: 1, of: 2,
god: 1, entitle: 1, them: 1,
decent: 1, respect: 1,
mankind: 1, opinion: 1 ...

among: 1, the: 2,
powers: 1, of: 2,
earth: 1, separate: 1,
equal: 1, and: 1 ...

Compute word counts locally
when: 1, in: 1,
the: 1, course: 1,
of: 1, human: 1,
events: 1, it: 1
dissolve: 1, the: 2,
political: 1, bands: 1,
which: 1, have: 1,
connected: 1, them: 1 ...

requires: 1, that: 1,
they: 1, should: 1,
declare: 1, the: 1,
causes: 1, which: 1 ...

among: 1, the: 2,
powers: 1, of: 2,
earth: 1, separate: 1,
equal: 1, and: 1 ...

dissolve: 1, the: 2,

nature: 2, and: 1, of: 2,
god: 1, entitle: 1, them: 1,
decent: 1, respect: 1,
mankind: 1, opinion: 1 ...

Compute word counts locally

Now what...
How to merge results?
Merging results computed locally

Several options

Don’t merge — requires additional computation for correct results

Send everything to one node — what if data is too big? Too slow…

Partition key space among nodes in cluster (e.g. [a-e], [f-j], [k-p] …)

1. Assign a key space to each node
2. Split local results by the key spaces
3. Fetch and merge results that correspond to the node’s key space
when: 1, in: 1,
the: 1, course: 1,
of: 1, human: 1,
events: 1, it: 1

dissolve: 1, the: 2,
political: 1, bands: 1,
which: 1, have: 1,
connected: 1, them: 1 ...

requires: 1, that: 1,
they: 1, should: 1,
declare: 1, the: 1,
causes: 1, which: 1 ...

among: 1, the: 2,
powers: 1, of: 2,
earth: 1, separate: 1,
equal: 1, and: 1 ...

nature: 2, and: 1, of: 2,
god: 1, entitle: 1, them: 1,
decent: 1, respect: 1,
mankind: 1, opinion: 1 ...
when: 1, the: 1,
in: 1, it: 1, human: 1,
course: 1, events: 1,
of: 1

bands: 1, dissolve: 1,
connected: 1, have: 1,
political: 1, the: 1,
them: 1, which: 1

causes: 1, declare: 1,
requires: 1, should: 1,
that: 1, they: 1, the: 1,
which: 1

nature: 2, of: 2,
mankind: 1, opinion: 1,
etitle: 1, and: 1,
decent: 1, god: 1,
them: 1, respect: 1,
among: 1, and: 1,
equal: 1, earth: 1,
separate: 1, the: 2,
powers: 1, of: 2

Split local results by key space
when: 1, the: 1, that: 1,
they: 1, the: 1, which: 1,
them: 1, the: 2, the: 1,
them: 1, which: 1

bands: 1, dissolve: 1,
connected: 1, course: 1,
events: 1, among: 1, and: 1,
equal: 1, earth: 1, entitle: 1,
and: 1, decent: 1, causes: 1,
declare: 1

requires: 1, should: 1,
respect: 1, separate: 1

god: 1, have: 1,
in: 1, it: 1,
human: 1,
powers: 1, of: 2,
nature: 2, of: 2,
mankind: 1, of: 1,
opinion: 1, political: 1

Note the duplicates...
Mapreduce

Partition dataset into many chunks

**Map stage:** Each node processes one or more chunks locally

**Reduce stage:** Each node fetches and merges partial results from all other nodes
Mapreduce Interface

\textbf{map}(key, value) \rightarrow \text{list}(<k', v'>)

- Apply function to (key, value) pair
- Outputs list of intermediate pairs

\textbf{reduce}(key, \text{list}<value>) \rightarrow <k', v'>

- Applies aggregation function to values
- Outputs result
Mapreduce: Word count

map(key, value):
   // key = document name
   // value = document contents
   for each word w in value:
      emit (w, 1)

reduce(key, values):
   // key = the word
   // values = number of occurrences of that word
   count = sum(values)
   emit (key, count)
Mapreduce: Word count

How much wood would a woodchuck chuck if a woodchuck could chuck wood?

(a, 1), (woodchuck, 1), (would, 1), (a, 1) (wood, 1), (if, 1) (woodchuck, 1), (could, 1), (chuck, 1), (wood, 1)

map

combine

shuffle

reduce

A woodchuck would chuck a lot of wood if a woodchuck could chuck wood.

(a, 1), (woodchuck, 1), (would, 1), (chuck, 1), (a, 1) (lot, 1), (of, 1), (wood, 1), (if, 1), (a, 1), (woodchuck, 1), (could, 1), (chuck, 1), (wood, 1)

<table>
<thead>
<tr>
<th>how</th>
<th>1</th>
</tr>
</thead>
<tbody>
<tr>
<td>much</td>
<td>1</td>
</tr>
<tr>
<td>wood</td>
<td>2</td>
</tr>
<tr>
<td>would</td>
<td>1</td>
</tr>
<tr>
<td>a</td>
<td>2</td>
</tr>
<tr>
<td>woodchuck</td>
<td>2</td>
</tr>
<tr>
<td>chuck</td>
<td>2</td>
</tr>
<tr>
<td>if</td>
<td>1</td>
</tr>
<tr>
<td>could</td>
<td>1</td>
</tr>
</tbody>
</table>

how | 1
much | 1
wood | 4
would | 2
a | 5
woodchuck | 4

chuck | 4
if | 2
could | 2
lot | 1
of | 1
Why is implementing MapReduce hard?

- Failure is common
  - Even if each machine is available $p = 99.999\%$ of the time, a datacenter with $n = 100,000$ machines still encounters failures $(1-p^n) = 63\%$ of the time

- Data skew causes unbalanced performance across cluster
  - Problems occur at scale.
  - Hard to debug!