[Docs] [txt pdf] [draft-ietf-http...] [Tracker] [Diff1] [Diff2] [Errata] Obsoleted by: <u>7230</u>, <u>7231</u>, <u>7232</u>, <u>7233</u>, <u>7234</u>, <u>7235</u> DRAFT STANDARD Updated by: 2817, 5785, 6266, 6585 Errata Exist Network Working Group R. Fielding Request for Comments: 2616 UC Irvine Obsoletes: 2068 J. Gettys Category: Standards Track Compaq/W3C J. Mogul Compaq H. Frystyk W3C/MIT L. Masinter Xerox P. Leach Microsoft T. Berners-Lee W3C/MIT

Hypertext Transfer Protocol -- HTTP/1.1

HTTP and the Web

Kyle Jamieson
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Today

1. HTTP basics: headers, requests, responses

2. Web proxies; web caches

3. Web performance optimization

Two Forms of Header Formats

- Fixed: Every field (type, length) defined
 - Fast parsing (good for hardware implementations)
 - Not human readable
 - Fairly static (IPv6 ~20 years to deploy)
 - E.g., Ethernet, IP, TCP headers
- Today: Variable length headers
 - Slower parsing (hard to implement in hardware)
 - Human readable
 - Extensible
 - E.g., HTTP (Web), SMTP (Email), XML

HTTP Basics (Overview)

HTTP over bidirectional byte stream (e.g. TCP)

Interaction

- Client looks up host (DNS)
- Client sends request message to server
- Server response message contains data or error
- Requests & responses are encoded in text

HTTP protocol itself is Stateless

- HTTP maintains no info about past client requests
- "Cookies" allow server to identify client and associate requests into a client session

HTTP Request

Request line

- Method
 - GET return URI
 - HEAD return headers only of GET response
 - POST send data to the server (forms, etc.)
- URL (relative)
 - E.g., /index.html
- HTTP version

HTTP Request (cont.)

- Request headers
 - Variable length, human-readable
 - Uses:
 - Authorization authentication info
 - Acceptable document types/encodings
 - From user email
 - If-Modified-Since
 - Referrer what caused this page to be requested
 - User-Agent client software
- Blank-line
- Body

HTTP Request Example

GET /index.html HTTP/1.1

Host: www.example.com

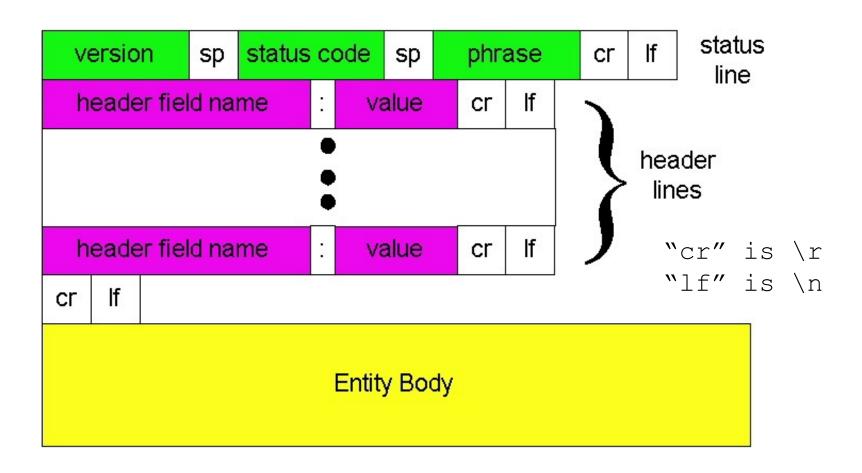
Accept-Language: en-us

Accept-Encoding: gzip, deflate

User-Agent: Mozilla/4.0 (compatible; MSIE 5.5; Windows NT 5.0)

Connection: Keep-Alive

HTTP Response



HTTP Response

- Status-line
 - HTTP version (now "1.1")
 - 3 digit response code
 - 1XX informational
 - 2XX success
 - 200 OK
 - 3XX redirection
 - 301 Moved Permanently
 - 303 Moved Temporarily
 - 304 Not Modified
 - 4XX client error
 - 404 Not Found
 - 5XX server error
 - 505 HTTP Version Not Supported
 - Reason phrase

HTTP Response (cont.)

Headers

- Variable length, human-readable
- Uses:
 - Location for redirection
 - Server server software
 - WWW-Authenticate request for authentication
 - Allow list of methods supported (get, head, etc)
 - Content-Encoding E.g x-gzip
 - Content-Length
 - Content-Type
 - Expires (caching)
 - Last-Modified (caching)
- Blank-line
- Body

HTTP Response Example

HTTP/1.1 200 OK

Date: Tue, 27 Mar 2001 03:49:38 GMT

Server: Apache/1.3.14 (Unix) (Red-Hat/Linux) mod_ssl/2.7.1

OpenSSL/0.9.5a DAV/1.0.2 PHP/4.0.1pl2 mod_perl/1.24

Last-Modified: Mon, 29 Jan 2001 17:54:18 GMT

Accept-Ranges: bytes

Content-Length: 4333

Keep-Alive: timeout=15, max=100

Connection: Keep-Alive

Content-Type: text/html

••••

How to Mark End of Message?

- Close connection
 - Only server can do this
 - One request per TCP connection. Hurts performance.
- Content-Length
 - Must know size of transfer in advance
- No body content. Double CRLF marks end
 - E.g., 304 never have body content
- Transfer-Encoding: chunked (HTTP/1.1)
 - After headers, each chunk is content length in hex, CRLF, then body. Final chunk is length 0.

Example: Chunked Encoding

```
HTTP/1.1 200 OK <CRLF>
Transfer-Encoding: chunked <CRLF>
<CRLF>
25 <CRLF>
This is the data in the first chunk <CRLF>
1A <CRLF>
and this is the second one <CRLF>
0 <CRLF>
```

- Especially useful for dynamically-generated content, as length is not a priori known
 - Server would otherwise need to cache data until done generating,
 and then go back and fill-in length header before transmitting

Web Proxies

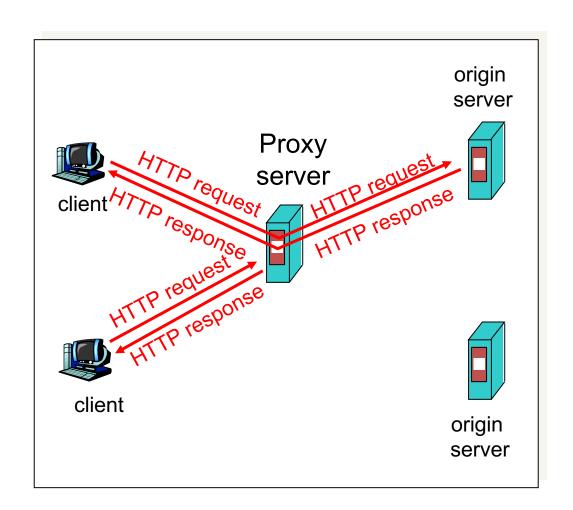
HTTP Caching

Proxies

- End host that acts a broker between client and server
 - Speaks to server on client's behalf
- Why?
 - Privacy
 - Content filtering
 - Caching!!!

Proxies (Cont.)

- Accept requests from multiple clients
- Takes request and reissues it to server
- Takes response and forwards to client



HTTP Caching

- Why cache?
 - Lot of objects don't change (images, js, css)
 - Reduce # of client connections
 - Reduce server load
 - Reduce overall network traffic; save \$\$\$

Caching is Hard

- Significant fraction (>50%?) of distinct HTTP objects may be uncacheable
 - Dynamic data: Stock prices, scores, web cams
 - CGI scripts: results based on passed parameters
 - Cookies: results may be based on passed data
 - SSL: encrypted data is not cacheable
 - Advertising / analytics: owner wants to measure # hits
 - Random strings in content to ensure unique counting
- Yet significant fraction of HTTP bytes are cacheable
 - Images, video, CSS pages, etc.
- Want to limit staleness of cached objects

How long should the client cache for?

- Clients (and proxies) cache documents
 - When should origin be checked for changes?
 - Every time? Every session? Date?
- HTTP includes caching information in headers
 - HTTP 0.9/1.0 used: "Expires: <date>"; "Pragma: no-cache"
 - HTTP/1.1 has "Cache-Control"
 - "No-Cache", "Max-age: <seconds>"
 - "ETag: <opaque value>

Why the changes between 1.0 and 1.1?

Timestamps

- Server hints when an object "Expires" (Expires: xxx)
- Server provides last modified date, client can check if that's still valid

Problems

- Client and server might not have synchronized clocks
- Server replicas might not have synchronized clocks
- Max-age solves this: relative seconds, not abs time

What if cache expires?

- Store past expiry time (if room in cache)
- Upon client request, cache revalidates with server

GET / HTTP/1.1

Accept-Language: en-us

If-Modified-Since: Mon, 29 Jan 2001 17:54:18 GMT

Host: www.example.com

Connection: Keep-Alive

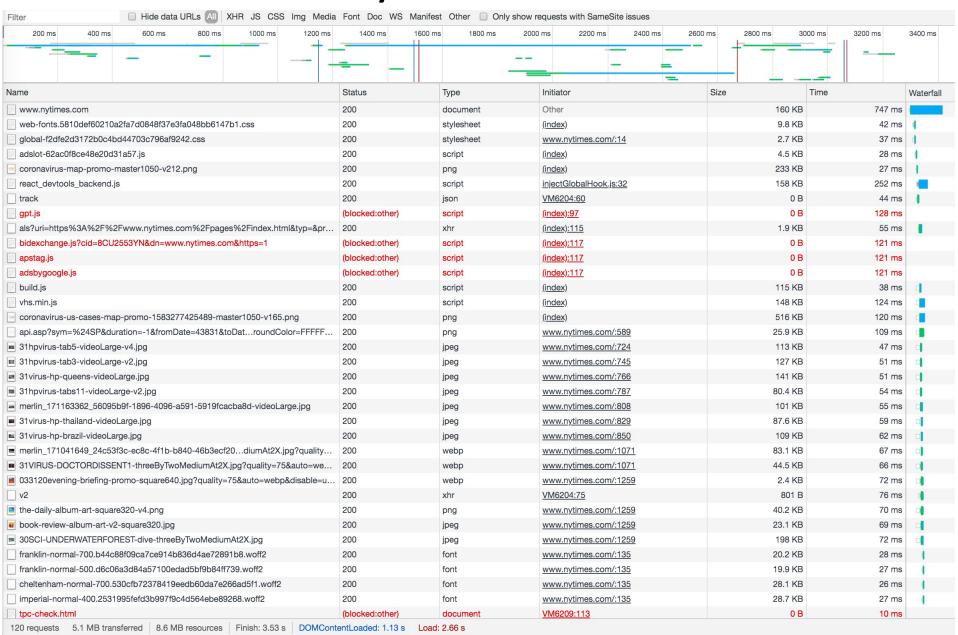
HTTP/1.1 304 Not Modified

Date: Tue, 27 Mar 2001 03:50:51 GMT

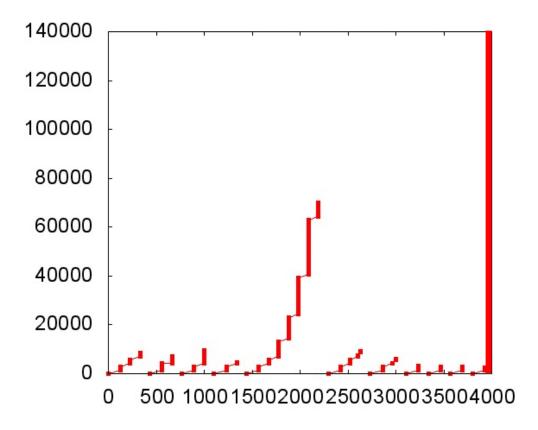
Connection: Keep-Alive

HTTP xfer = single object Web pages = many objects

nytimes.com



HTTP/1.0 fetching items: Received sequence number plot



Time (milliseconds)

Fetch an 8.5 Kbyte page with 10 embedded objects, most < 10 Kbyte All TCP connections stay in slow start, except for the large object

How to handle many requests?

- Maximize goodput by reusing connections
 - Avoid connection (TCP) setup
 - Avoid TCP slow-start
- Client-server will maintain existing TCP connection for up to K idle seconds

GET / HTTP/1.1

Host: www.example.com

Connection: Keep-Alive

HTTP/1.1 200 OK

Date: Tue, 27 Mar 2001 03:50:51 GMT

Connection: Keep-Alive

Three approaches to multiple requests

Parallel Connections

Persistent Connections

Conn 1:

- Request 1
- Response 1

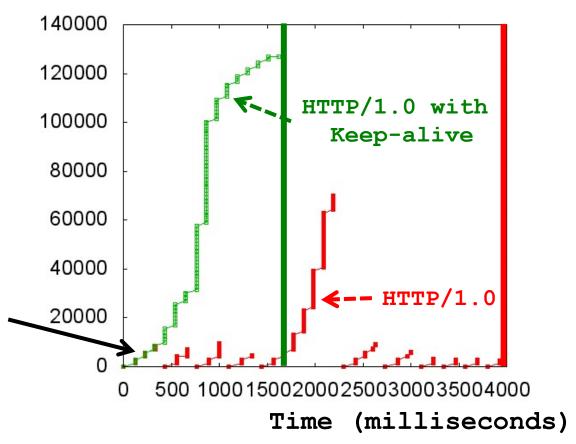
Conn 2:

- Request 2
- Response 2

Conn 1:

- Request 1
- Response 1
- Request 2
- Response 2
- Request 3
- Response 3

Persistent connections avoid unnecessary slow starts



Fetch an 8.5 Kbyte page with 10 embedded objects, most < 10 Kbyte Leave TCP connection open after server response, next HTTP request reuses it Only incur one slow start, but takes an RTT to issue next request

Three approaches to multiple requests

Parallel Connections

Persistent Connections

Pipelined Connections

Conn 1:

- Request 1
- Response 1

Conn 2:

- Request 2
- Response 2

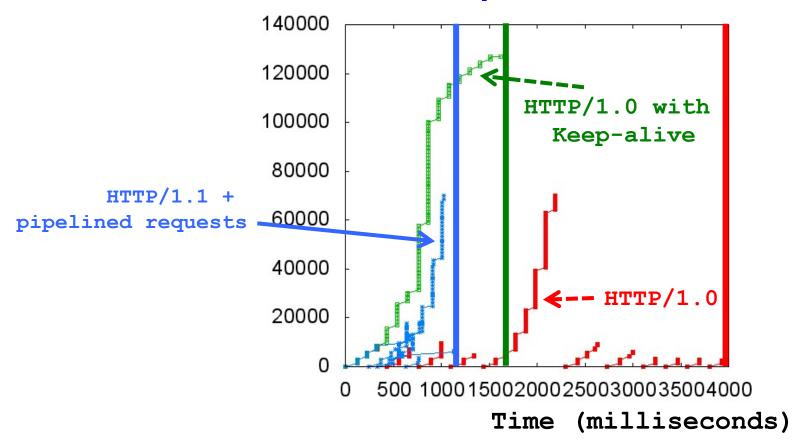
Conn 1:

- Request 1
- Response 1
- Request 2
- Response 2
- Request 3
- Response 3

Conn 1:

- Request 1
- Request 2
- Request 3
- Response 1
- Response 2
- Response 3

Pipelined + Parallel Connections overlap RTTs



Fetch an 8.5 Kbyte page with 10 embedded objects, most < 10 Kbyte Send multiple HTTP requests simultaneously

Overlaps RTTs of all requests

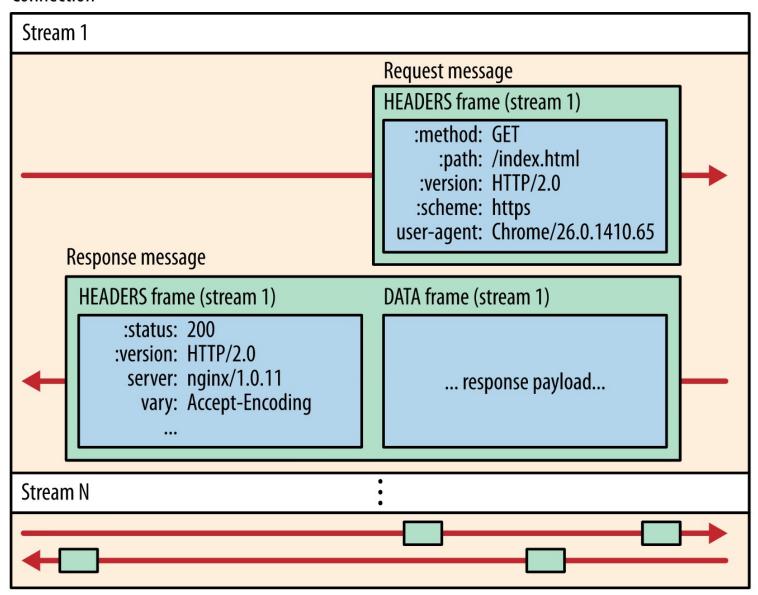
What are challenges with pipelining?

- Head-of-line blocking
 - Small xfers can "block" behind large xfer
- No reordering
 - HTTP response does not "identify" which request it's in response to; obvious in simple request/response
- Can behave worse than parallel + persistent
 - Can send expensive query 1 on conn 1, while sending many cheap queries on conn 2

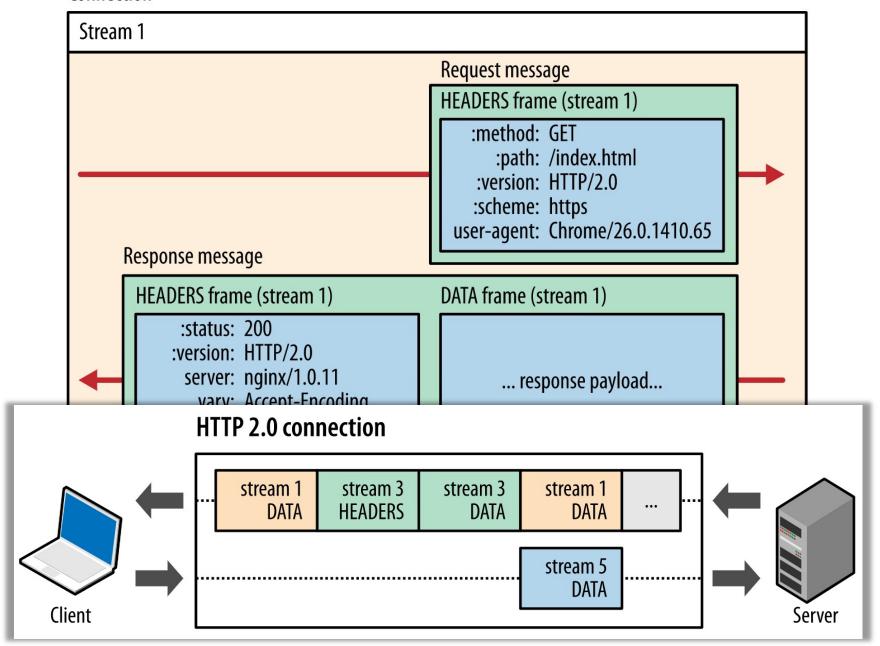
Google's SPDY -> HTTP/2 Standard

- Server "push" for content
 - One client request, multiple responses
 - After all, server knows that after parsing HTML, client will immediately request embedded URLs
- Better pipelining and xfer
 - Multiplexing multiple xfers w/o HOL blocking
 - Request prioritization
 - Header compression

Connection



Connection



Summary

 HTTP: dominant application layer protocol for the web

HTTP caching had a limited impact (CDNs next)

 Recent optimization and evolution of HTTP for performance and efficiency