

# COS 217: Introduction to Programming Systems

## Performance Improvement

“Premature optimization is the root of all evil.”

– Donald Knuth

“Rules of Optimization:

Rule 1: Don't do it.

Rule 2 (for experts only): Don't do it yet.”

– Michael A. Jackson



**PRINCETON UNIVERSITY**



# “Programming in the Large”

## Design & Implement

- Program & programming style (done)
- Common data structures and algorithms (done)
- Modularity (done)
- Building techniques & tools (done)

## Debug

- Debugging techniques & tools (done)

## Test

- Testing techniques (done)

## Maintain

- Performance improvement techniques & tools ← we are here



# Goals of this Lecture

## Help you learn about:

- How to use profilers to identify code hot-spots
- How to make your programs run faster



## Why?

- In a large program, typically a small fragment of the code consumes most of the CPU time
  - Identifying that fragment is likely to identify the source of inadequate performance
- Part of “programming maturity” is being able to recognize common approaches for improving the performance of such code fragments
- Part of “programming maturity” is also being able to recognize what is worth your time to improve and what is already “good enough”

# Agenda



**Should you optimize?**

What should you optimize?

Optimization techniques



# Performance Improvement Pros

Techniques described in this lecture can answer:



How slow is my code?



Where is it slow?



Why is it slow?

Similar techniques (not discussed) can address:

- How can I make my program use less memory?

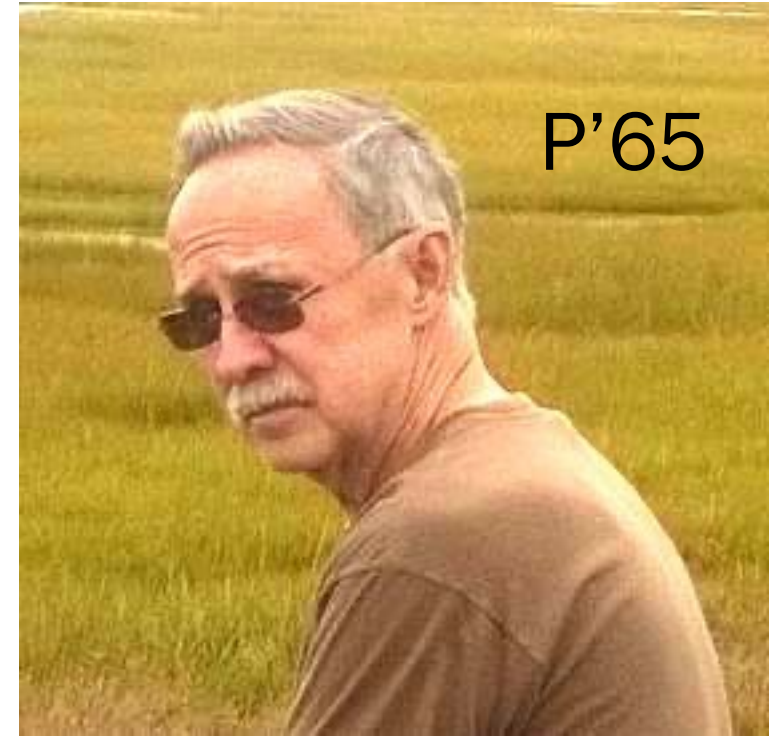


# Performance Improvement Cons

Techniques described in this lecture can yield code that:

- Is less clear/maintainable
- Might confuse debuggers
- Might contain bugs
  - Requires regression testing

So...





# When to Improve Performance

“The first principle of optimization is

*don't.*

Is the program good enough already?  
Knowing how a program will be used  
and the environment it runs in,  
is there any benefit to making it faster?”

– Kernighan & Pike



# Timing a Program

Run a tool to time program execution

- E.g., Unix `time` command

```
$ time sort < bigfile.txt > output.txt  
real    0m12.977s  
user    0m12.860s  
sys     0m0.010s
```

Output:

- **Real:** Wall-clock time between program invocation and termination
- **User:** CPU time spent executing the program
- **System:** CPU time spent within the OS on the program's behalf





# Enabling Compiler Optimization

Enable compiler speed optimization

```
gcc217 -Ox mysort.c -o mysort
```

- Compiler looks for ways to transform your code so that result is the same but it runs faster
- **X** controls how many transformations the compiler tries – see “man gcc” for details
  - **-O0**: do not optimize (default if **-O** not specified)
  - **-O1**: optimize (default if **-O** but no number is specified)
  - **-O2**: optimize more (longer compile time)
  - **-O3**: optimize yet more (including inlining)

**Warning: Speed optimization can affect debugging**

- e.g., Optimization eliminates variable  $\Rightarrow$  GDB cannot print value of variable



# Now What?

So you've determined that your program is taking too long, even with compiler optimization enabled (and NDEBUG defined, etc.)

Is it time to completely rewrite the program?



# Agenda



Should you optimize?

**What should you optimize?**

Optimization techniques



# Identifying Hot Spots

Spend time optimizing only the parts of the program that will make a difference!

Gather statistics about your program's execution

- **Coarse-grained:** how much time did execution of a particular function call take?
  - Time individual function calls or blocks of code
- **Fine-grained:** how many times was a particular function called?  
How much time was taken by all calls to that function?
  - Use an execution profiler such as gprof



# Timing Parts of a Program

Call a function to compute **wall-clock time** consumed

- Unix `gettimeofday()` returns time in seconds + microseconds

```
#include <sys/time.h>

struct timeval startTime;
struct timeval endTime;
double wallClockSecondsConsumed;

gettimeofday(&startTime, NULL);
<execute some code here>
gettimeofday(&endTime, NULL);
wallClockSecondsConsumed =
    endTime.tv_sec - startTime.tv_sec +
    1.0E-6 * (endTime.tv_usec - startTime.tv_usec);
```

- Not defined by C90 standard



# Timing Parts of a Program (cont.)

Call a function to compute **CPU time** consumed

- `clock()` returns CPU times in `CLOCKS_PER_SEC` units

```
#include <time.h>

clock_t startClock;
clock_t endClock;
double cpuSecondsConsumed;

startClock = clock();
<execute some code here>
endClock = clock();
cpuSecondsConsumed =
    ((double)(endClock - startClock)) / CLOCKS_PER_SEC;
```

- Defined by C90 standard



# Identifying Hot Spots

Spend time optimizing only the parts of the program that will make a difference!

Gather statistics about your program's execution

- *Coarse-grained:* how much time did execution of a particular function call take?
  - Time individual function calls or blocks of code
- *Fine-grained:* how many times was a particular function called?  
How much time was taken by all calls to that function?
  - Use an **execution profiler** such as gprof



# GPROF Example Program

## Example program for GPROF analysis

- Sort an array of 10 million random integers
- Artificial: consumes lots of CPU time, generates no output

```
#include <string.h>
#include <stdio.h>
#include <stdlib.h>

enum { MAX_SIZE = 10000000 };
int a[MAX_SIZE];

void fillArray(int a[], int size)
{
    int i;
    for (i = 0; i < size; i++)
        a[i] = rand();
}

void swap(int a[], int i, int j)
{
    int temp = a[i];
    a[i] = a[j];
    a[j] = temp;
}
```

```
int part(int a[], int left, int right)
{
    int first = left-1;
    int last = right;
    for (;;) {
        while (a[++first] < a[right])
            ;
        while (a[right] < a[--last])
            if (last == left)
                break;
        if (first >= last)
            break;
        swap(a, first, last);
    }
    swap(a, first, right);
    return first;
}
```





# GPROF Example Program (cont.)

## Example program for GPROF analysis

- Sort an array of 10 million random integers
- Artificial: consumes lots of CPU time, generates no output

```
void quicksort(int a[], int left, int right)
{
    if (right > left) {
        int mid = part(a, left, right);
        quicksort(a, left, mid - 1);
        quicksort(a, mid + 1, right);
    }
}

int main(void)
{
    fillArray(a, MAX_SIZE);
    quicksort(a, 0, MAX_SIZE - 1);
    return 0;
}
```



# Using GPROF

## Step 1: Instrument the program

```
gcc217 -pg mysort.c -o mysort
```

- Adds profiling code to mysort, that is...
- “Instruments” mysort

## Step 2: Run the program

```
./mysort
```

- Creates file gmon.out containing statistics

## Step 3: Create a report

```
gprof mysort > myreport
```

- Uses mysort and gmon.out to create textual report

## Step 4: Examine the report

```
more myreport
```

# gprof Design



## What's going on behind the scenes?

- `-pg` generates code to interrupt program many times per second
- Each time, records *where* the code was interrupted
- `gprof` uses symbol table to map back to function name



# The GPROF Report

% time	cumulative seconds	self seconds	calls	self s/call	total s/call	name
84.54	2.27	2.27	6665307	0.00	0.00	part
9.33	2.53	0.25	54328749	0.00	0.00	swap
2.99	2.61	0.08	1	0.08	2.61	quicksort
2.61	2.68	0.07	1	0.07	0.07	fillArray

- Each line describes one function
  - **name**: name of the function
  - **%time**: percentage of time spent executing this function
  - **cumulative seconds**: [skipping, as this isn't all that useful]
  - **self seconds**: time spent executing this function
  - **calls**: number of times function was called (excluding recursive)
  - **self s/call**: average time per execution (excluding descendants)
  - **total s/call**: average time per execution (including descendants)



# The GPROF Report (cont.)

## Call graph profile

index	% time	self	children	called	name
[1]	100.0	0.00	2.68		<spontaneous> main [1]
		0.08	2.53	1/1	quicksort [2]
		0.07	0.00	1/1	fillArray [5]
-----					
				13330614	quicksort [2]
[2]	97.4	0.08	2.53	1/1	main [1]
		0.08	2.53	1+13330614	quicksort [2]
		2.27	0.25	6665307/6665307	part [3]
				13330614	quicksort [2]
-----					
		2.27	0.25	6665307/6665307	quicksort [2]
[3]	94.4	2.27	0.25	6665307	part [3]
		0.25	0.00	54328749/54328749	swap [4]
-----					
		0.25	0.00	54328749/54328749	part [3]
[4]	9.4	0.25	0.00	54328749	swap [4]
-----					
		0.07	0.00	1/1	main [1]
[5]	2.6	0.07	0.00	1	fillArray [5]
-----					



# GPROF Report Analysis

% time	cumulative seconds	self seconds	calls	self s/call	total s/call	name
84.54	2.27	2.27	6665307	0.00	0.00	part
9.33	2.53	0.25	54328749	0.00	0.00	swap
2.99	2.61	0.08	1	0.08	2.61	quicksort
2.61	2.68	0.07	1	0.07	0.07	fillArray

## Observations:

- `swap()` is called many times; each call consumes little time; in all, `swap()` consumes only 9% of the time overall
- `part()` is called fewer times; each call consumes little time, but clearly more than `swap()`, since `part()` consumes 85% of the time overall

## Conclusions:

- To improve performance, try to make `partition()` faster
- Don't even think about trying to make `fillArray()` or `quicksort()` faster

# Agenda



Should you optimize?

What should you optimize?

**Optimization techniques**



# Using Better Algs and DSs

Use a better algorithm or data structure

- e.g., would a different sorting algorithm work better?

#include COS 226 (or should we say, import 226?)

- But only where it would really help!

Not worth using asymptotically-efficient algorithms and data structures that are complex, hard to understand, hard to debug, or hard to maintain if they will not make any difference anyway!



# Optimization Strategy: Avoid Repeated Computation



Before:

```
int g(int x)
{
    return f(x) + f(x) + f(x) + f(x);
}
```

After:

```
int g(int x)
{
    return 4 * f(x);
}
```



# four fs' sake



Q: Could a good compiler do this optimization for you?

Before:

```
int g(int x)
{
    return f(x) + f(x) + f(x) + f(x);
}
```

After:

```
int g(int x)
{
    return 4 * f(x);
}
```

- A. Yes
- B. Only sometimes
- C. No

Answer: only sometimes!



# Side Effects as Blockers

```
int g(int x)
{
    return f(x) + f(x) + f(x) + f(x);
}
```

```
int g(int x)
{
    return 4 * f(x);
}
```

Suppose `f ( )` has **side effects**?

```
int counter = 0;
...
int f(int x)
{
    return counter++;
}
```

And `f ( )` might be defined in another file known only at link time!



Q: Could a good compiler do this optimization for you?

Before:

```
for (i = 0; i < n; i++)  
  for (j = 0; j < n; j++)  
    a[n*i + j] = b[j];
```

After:

```
for (i = 0; i < n; i++) {  
  int ni = n * i;  
  for (j = 0; j < n; j++)  
    a[ni + j] = b[j];  
}
```

- A. Yes
- B. Only sometimes
- C. No

Probably A.



# Wasn't this in the A2 grading rubric?

Before:

```
for (i = 0; i < strlen(s); i++) {  
    /* Do something with s[i] */  
}
```

After:

```
length = strlen(s);  
for (i = 0; i < length; i++) {  
    /* Do something with s[i] */  
}
```

Could a good compiler do that for you?



# Sydney Bristow asks ...



Q: Could a good compiler do this optimization for you?

Before:

```
void twiddle(int *p1, int *p2)
{
    *p1 += *p2;
    *p1 += *p2;
}
```

- A. Yes
- B. Only sometimes
- C. No

After:

```
void twiddle(int *p1, int *p2)
{
    *p1 += *p2 * 2;
}
```

C.

# ALIASes as Blockers



```
void twiddle(int *p1, int *p2)
{
    *p1 += *p2;
    *p1 += *p2;
}
```

```
void twiddle(int *p1, int *p2)
{
    *p1 += *p2 * 2;
}
```

What if **p1** and **p2** are **aliases**?

- What if **p1** and **p2** point to *the same* integer?
- First version: result is 4 times **\*p1**
- Second version: result is 3 times **\*p1**

C99 supports the **restrict** keyword

- e.g., `int * restrict p1`



# Inlining Function Calls

Before:

```
void g(void)
{
    /* Some code */
}
void f(void)
{
    ...
    g();
    ...
}
```

After:

```
void f(void)
{
    ...
    /* Some code */
    ...
}
```



Beware: Can introduce redundant/cloned code, making maintenance more difficult.

Some compilers support `inline` keyword





# Unrolling Loops

Original:

```
for (i = 0; i < 6; i++)  
    a[i] = b[i] + c[i];
```

Maybe  
faster:

```
for (i = 0; i < 6; i += 2) {  
    a[i] = b[i] + c[i];  
    a[i+1] = b[i+1] + c[i+1];  
}
```

Maybe  
even  
faster:

```
a[i] = b[i] + c[i];  
a[i+1] = b[i+1] + c[i+1];  
a[i+2] = b[i+2] + c[i+2];  
a[i+3] = b[i+3] + c[i+3];  
a[i+4] = b[i+4] + c[i+4];  
a[i+5] = b[i+5] + c[i+5];
```



Some compilers provide option, e.g. `-funroll-loops`



# Using a Lower-Level Language

## Rewrite code in a lower-level language

- Use registers instead of memory
- Use instructions (e.g. `adc`) that compiler doesn't know
- Gee, where have I seen this before...?

## Beware: Modern optimizing compilers generate fast code!

- Hand-written assembly language code could be slower!

# Summary



## Steps to improve **execution (time)** efficiency:

- Don't do it.
- Don't do it yet.
- Time the code to make sure it's necessary
- Enable compiler optimizations
- Identify hot spots using profiling
- Use a better algorithm or data structure
- Identify common inefficiencies and bad idioms
- Fine-tune the code