[Docs] [txt pdf] [draft-ietf-http...] [Tracker] [Diff1] [Diff2] [Errata] Obsoleted by: 7230, 7231, 7232, 7233, 7234, 7235 DRAFT STANDARD Updated by: 2817, 5785, 6266, 6585 Errata Exist Network Working Group R. Fielding Request for Comments: 2616 UC Irvine Obsoletes: 2068 J. Gettys Category: Standards Track Compaq/W3C J. Mogul Compaq H. Frystyk W3C/MIT L. Masinter Xerox P. Leach Microsoft T. Berners-Lee

Hypertext Transfer Protocol -- HTTP/1.1

## **HTTP and the Web**

Kyle Jamieson Lecture 15 COS 461: Computer Networks

W3C/MIT June 1999

## Today

1. HTTP basics: headers, requests, responses

2. Web proxies; web caches

3. Web performance optimization

## **Two Forms of Header Formats**

- Fixed: Every field (type, length) defined
  - Fast parsing (good for hardware implementations)
  - Not human readable
  - Fairly static (IPv6 ~20 years to deploy)
  - E.g., Ethernet, IP, TCP headers
- Today: Variable length headers
  - Slower parsing (hard to implement in hardware)
  - Human readable
  - Extensible
  - E.g., HTTP (Web), SMTP (Email), XML

## **HTTP Basics (Overview)**

• HTTP over bidirectional byte stream (e.g. TCP)

#### Interaction

- Client looks up host (DNS)
- Client sends *request* message to server
- Server *response* message contains data or error
- Requests & responses are encoded in text
- HTTP protocol itself is *Stateless* 
  - HTTP maintains no info about past client requests
  - "Cookies" allow server to identify client and associate requests into a client session

## **HTTP Request**

#### • Request line

- Method
  - GET return URI
  - HEAD return headers only of GET response
  - POST send data to the server (forms, etc.)
- URL (relative)
  - E.g., /index.html
- HTTP version

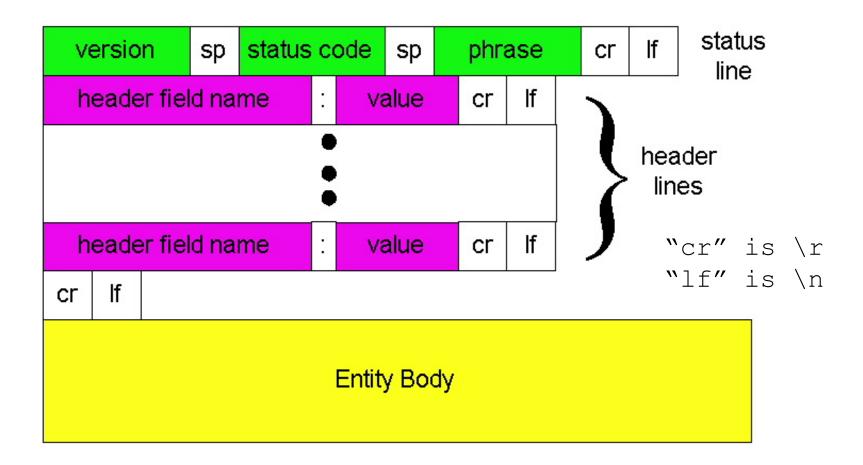
## HTTP Request (cont.)

- Request headers
  - Variable length, human-readable
  - Uses:
    - Authorization authentication info
    - Acceptable document types/encodings
    - From user email
    - If-Modified-Since
    - Referrer what caused this page to be requested
    - User-Agent client software
- Blank-line
- Body

## **HTTP Request Example**

- **GET /index.html HTTP/1.1**
- Host: www.example.com
- Accept-Language: en-us
- Accept-Encoding: gzip, deflate
- User-Agent: Mozilla/4.0 (compatible; MSIE 5.5; Windows NT 5.0)
- **Connection: Keep-Alive**

### **HTTP Response**



## **HTTP Response**

#### • Status-line

- HTTP version (now "1.1")
- 3 digit response code
  - 1XX informational
  - 2XX success
    - 200 OK
  - 3XX redirection
    - 301 Moved Permanently
    - 303 Moved Temporarily
    - 304 Not Modified
  - 4XX client error
    - 404 Not Found
  - 5XX server error
    - 505 HTTP Version Not Supported
- Reason phrase

## HTTP Response (cont.)

- Headers
  - Variable length, human-readable
  - Uses:
    - Location for redirection
    - Server server software
    - WWW-Authenticate request for authentication
    - Allow list of methods supported (get, head, etc)
    - Content-Encoding E.g x-gzip
    - Content-Length
    - Content-Type
    - Expires (caching)
    - Last-Modified (caching)
- Blank-line
- Body

## **HTTP Response Example**

HTTP/1.1 200 OK

Date: Tue, 27 Mar 2001 03:49:38 GMT

- Server: Apache/1.3.14 (Unix) (Red-Hat/Linux) mod\_ssl/2.7.1 OpenSSL/0.9.5a DAV/1.0.2 PHP/4.0.1pl2 mod\_perl/1.24
- Last-Modified: Mon, 29 Jan 2001 17:54:18 GMT
- **Accept-Ranges: bytes**
- **Content-Length: 4333**
- Keep-Alive: timeout=15, max=100
- **Connection: Keep-Alive**
- **Content-Type: text/html**

## How to Mark End of Message?

- Close connection
  - Only server can do this
  - One request per TCP connection. Hurts performance.
- Content-Length
  - Must know size of transfer in advance
- No body content. Double CRLF marks end
  - E.g., 304 never have body content
- Transfer-Encoding: chunked (HTTP/1.1)
  - After headers, each chunk is content length in hex, CRLF, then body. Final chunk is length 0.

#### **Example: Chunked Encoding**

- HTTP/1.1 200 OK <CRLF>
- Transfer-Encoding: chunked <CRLF>
  <CRLF>
- 25 <CRLF>
- This is the data in the first chunk <CRLF> 1A <CRLF> and this is the second one <CRLF>
- 0 <CRLF>
- Especially useful for dynamically-generated content, as length is not a priori known
  - Server would otherwise need to cache data until done generating, and then go back and fill-in length header before transmitting

## Web Proxies

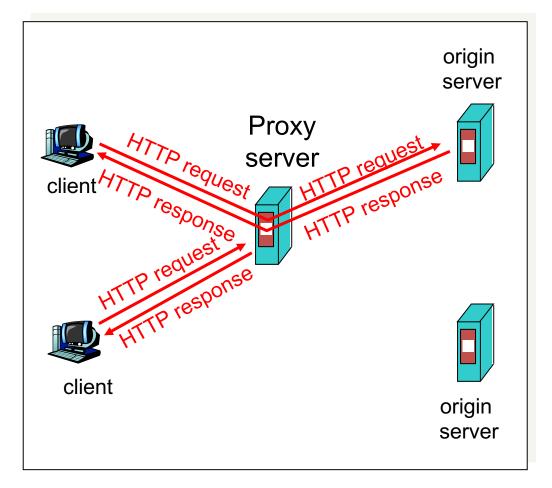
# **HTTP Caching**

### **Proxies**

- End host that acts a broker between client and server
  - Speaks to server on client's behalf
- Why?
  - Privacy
  - Content filtering
  - Caching!!!

# **Proxies (Cont.)**

- Accept requests from multiple clients
- Takes request and reissues it to server
- Takes response and forwards to client



## **HTTP Caching**

- Why cache?
  - Lot of objects don't change (images, js, css)
  - Reduce # of client connections
  - Reduce server load
  - Reduce overall network traffic; save \$\$\$

## Caching is Hard

- Significant fraction (>50%?) of distinct HTTP objects may be uncacheable
  - Dynamic data: Stock prices, scores, web cams
  - CGI scripts: results based on passed parameters
  - Cookies: results may be based on passed data
  - SSL: encrypted data is not cacheable
  - Advertising / analytics: owner wants to measure # hits
    - Random strings in content to ensure unique counting
- Yet significant fraction of HTTP bytes are cacheable
  - Images, video, CSS pages, etc.
- Want to limit staleness of cached objects

### How long should the client cache for?

- Clients (and proxies) cache documents
  - When should origin be checked for changes?
  - Every time? Every session? Date?
- HTTP includes caching information in headers
  - HTTP 0.9/1.0 used: "Expires: <date>"; "Pragma: no-cache"
  - HTTP/1.1 has "Cache-Control"
    - "No-Cache", "Max-age: <seconds>"
    - "ETag: <opaque value>

### Why the changes between 1.0 and 1.1?

#### • Timestamps

- Server hints when an object "Expires" (Expires: xxx)
- Server provides last modified date, client can check if that's still valid

- Problems
  - Client and server might not have synchronized clocks
  - Server replicas might not have synchronized clocks
  - Max-age solves this: relative seconds, not abs time

## What if cache expires?

- Store past expiry time (if room in cache)
- Upon client request, cache revalidates with server

#### GET / HTTP/1.1

Accept-Language: en-us

If-Modified-Since: Mon, 29 Jan 2001 17:54:18 GMT

Host: www.example.com

Connection: Keep-Alive

HTTP/1.1 304 Not Modified

Date: Tue, 27 Mar 2001 03:50:51 GMT

**Connection: Keep-Alive** 

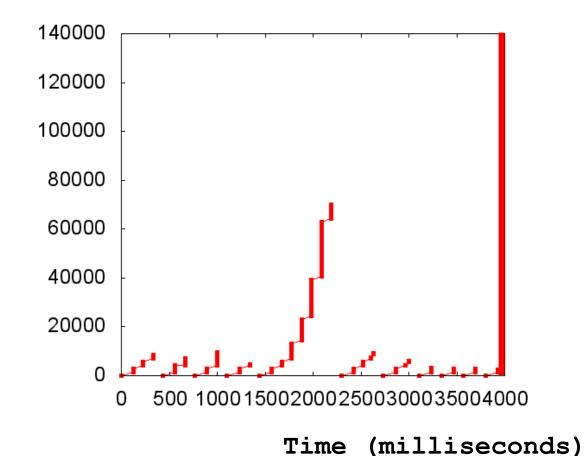
HTTP xfer = single object Web pages = many objects

#### nytimes.com

Filter 🔲 Hide data URLs 🔊 XHR JS CSS Img Media Font Doc WS Manifest Other 🗌 Only show requests with SameSite issues					
200 ms 400 ms 600 ms 800 ms 1000 ms 1200 ms	1400 ms 1600 m	s 1800 ms 2000	0 ms 2200 ms 2400 ms 2600 m	ms 2800 ms 3000 ms 3200 ms	s 3400 ms
	<u> </u>				
Name	Status	Туре	Initiator	Size Time	Waterfall
www.nytimes.com	200	document	Other	160 KB 7	47 ms
web-fonts.5810def60210a2fa7d0848f37e3fa048bb6147b1.css	200	stylesheet	(index)	9.8 KB	42 ms 🔰
global-f2dfe2d3172b0c4bd44703c796af9242.css	200	stylesheet	www.nytimes.com/:14	2.7 KB	37 ms 🕴
adslot-62ac0f8ce48e20d31a57.js	200	script	(index)	4.5 KB	28 ms 🕴
coronavirus-map-promo-master1050-v212.png	200	png	(index)	233 KB	27 ms
react_devtools_backend.js	200	script	injectGlobalHook.js:32	158 KB 2	52 ms
track	200	json	<u>VM6204:60</u>	0 B	44 ms 🕴
gpt.js	(blocked:other)	script	<u>(index):97</u>	0 B 1	28 ms
als?uri=https%3A%2F%2Fwww.nytimes.com%2Fpages%2Findex.html&typ=≺	200	xhr	<u>(index):115</u>	1.9 KB	55 ms
bidexchange.js?cid=8CU2553YN&dn=www.nytimes.com&https=1	(blocked:other)	script	<u>(index):117</u>	0 B 1	21 ms
apstag.js	(blocked:other)	script	(index):117	0 B 1	21 ms
adsbygoogle.js	(blocked:other)	script	<u>(index):117</u>	0 B 1	21 ms
build.js	200	script	(index)	115 KB	38 ms
vhs.min.js	200	script	(index)	148 KB 1	24 ms
coronavirus-us-cases-map-promo-1583277425489-master1050-v165.png	200	png	(index)	516 KB 1	20 ms
api.asp?sym=%24SP&duration=-1&fromDate=43831&toDatroundColor=FFFFF	200	png	www.nytimes.com/:589	25.9 KB 1	09 ms 📑
31hpvirus-tab5-videoLarge-v4.jpg	200	jpeg	www.nytimes.com/:724	113 KB	47 ms
31hpvirus-tab3-videoLarge-v2.jpg	200	jpeg	www.nytimes.com/:745	127 KB	51 ms
31virus-hp-queens-videoLarge.jpg	200	jpeg	www.nytimes.com/:766	141 KB	51 ms
31hpvirus-tabs11-videoLarge-v2.jpg	200	jpeg	www.nytimes.com/:787	80.4 KB	54 ms
e merlin_171163362_56095b9f-1896-4096-a591-5919fcacba8d-videoLarge.jpg	200	jpeg	www.nytimes.com/:808	101 KB	55 ms
31virus-hp-thailand-videoLarge.jpg	200	jpeg	www.nytimes.com/:829	87.6 KB	59 ms
🖬 31virus-hp-brazil-videoLarge.jpg	200	jpeg	www.nytimes.com/:850	109 KB	62 ms
e merlin_171041649_24c53f3c-ec8c-4f1b-b840-46b3ecf20diumAt2X.jpg?quality	200	webp	www.nytimes.com/:1071	83.1 KB	67 ms
S1VIRUS-DOCTORDISSENT1-threeByTwoMediumAt2X.jpg?quality=75&auto=we	200	webp	www.nytimes.com/:1071	44.5 KB	66 ms 🔄
033120evening-briefing-promo-square640.jpg?quality=75&auto=webp&disable=u	200	webp	www.nytimes.com/:1259	2.4 KB	72 ms
v2	200	xhr	<u>VM6204:75</u>	801 B	76 ms
the-daily-album-art-square320-v4.png	200	png	www.nytimes.com/:1259	40.2 KB	70 ms
look-review-album-art-v2-square320.jpg	200	jpeg	www.nytimes.com/:1259	23.1 KB	69 ms
30SCI-UNDERWATERFOREST-dive-threeByTwoMediumAt2X.jpg	200	jpeg	www.nytimes.com/:1259	198 KB	72 ms
franklin-normal-700.b44c88f09ca7ce914b836d4ae72891b8.woff2	200	font	www.nytimes.com/:135	20.2 KB	28 ms 🕴
franklin-normal-500.d6c06a3d84a57100edad5bf9b84ff739.woff2	200	font	www.nytimes.com/:135	19.9 KB	27 ms
cheltenham-normal-700.530cfb72378419eedb60da7e266ad5f1.woff2	200	font	www.nytimes.com/:135	28.1 KB	26 ms
imperial-normal-400.2531995fefd3b997f9c4d564ebe89268.woff2	200	font	www.nytimes.com/:135	28.7 KB	27 ms
tpc-check.html	(blocked:other)	document	VM6209:113	0 B	10 ms
120 requests 5.1 MB transferred 8.6 MB resources Finish: 3.53 s DOMCon	tentLoaded: 1.13 s Load	: 2.66 s			

120 requests 5.1 MB transferred 8.6 MB resources Finish: 3.53 s DOMContentLoaded: 1.13 s Load: 2.66 s

#### HTTP/1.0 fetching items: Received sequence number plot



Fetch an 8.5 Kbyte page with 10 embedded objects, most < 10 Kbyte All TCP connections stay in slow start, except for the large object

### How to handle many requests?

- Maximize goodput by reusing connections
  - Avoid connection (TCP) setup
  - Avoid TCP slow-start
- Client-server will maintain existing TCP connection for up to K idle seconds

GET / HTTP/1.1

Host: www.example.com

**Connection: Keep-Alive** 

HTTP/1.1 200 OK

Date: Tue, 27 Mar 2001 03:50:51 GMT

**Connection: Keep-Alive** 

#### Three approaches to multiple requests

Parallel Connections

Conn 1:

- Request 1
- Response 1

Conn 2:

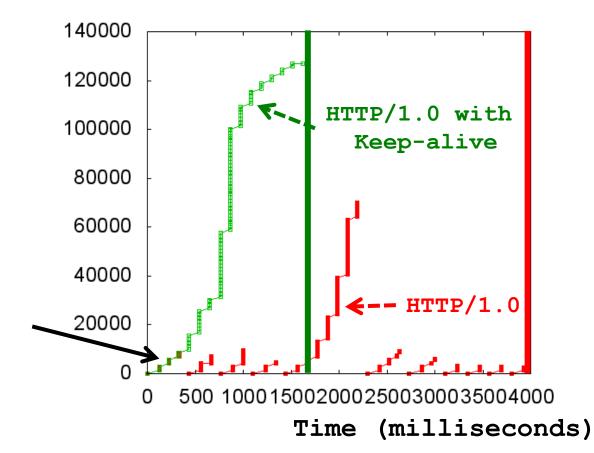
- Request 2
- Response 2

Persistent Connections

Conn 1:

- Request 1
- Response 1
- Request 2
- Response 2
- Request 3
- Response 3

#### Persistent connections avoid unnecessary slow starts



Fetch an 8.5 Kbyte page with 10 embedded objects, most < 10 Kbyte Leave TCP connection open after server response, next HTTP request reuses it **Only incur one slow start, but takes an RTT to issue next request** 

### Three approaches to multiple requests

Parallel Connections

Conn 1:

- Request 1
- Response 1

Conn 2:

- Request 2
- Response 2

Persistent Connections

Conn 1:

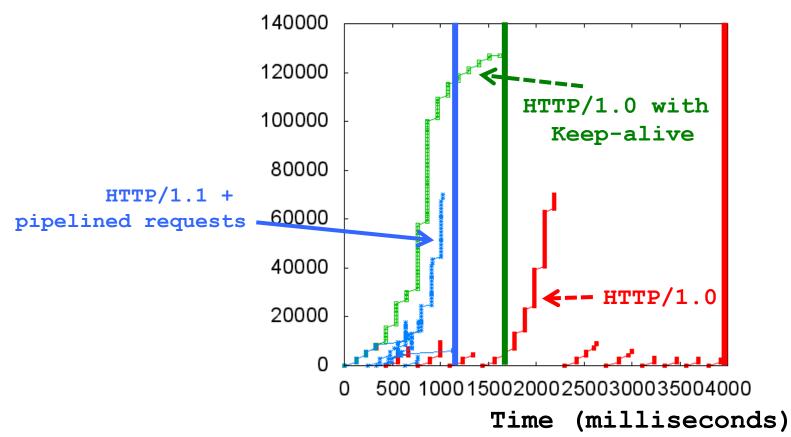
- Request 1
- Response 1
- Request 2
- Response 2
- Request 3
- Response 3

Pipelined Connections

#### Conn 1:

- Request 1
- Request 2
- Request 3
- Response 1
- Response 2
- Response 3

## Pipelined + Parallel Connections overlap RTTs



Fetch an 8.5 Kbyte page with 10 embedded objects, most < 10 Kbyte Send multiple HTTP requests simultaneously

**Overlaps RTTs of all requests** 

### What are challenges with pipelining?

Head-of-line blocking

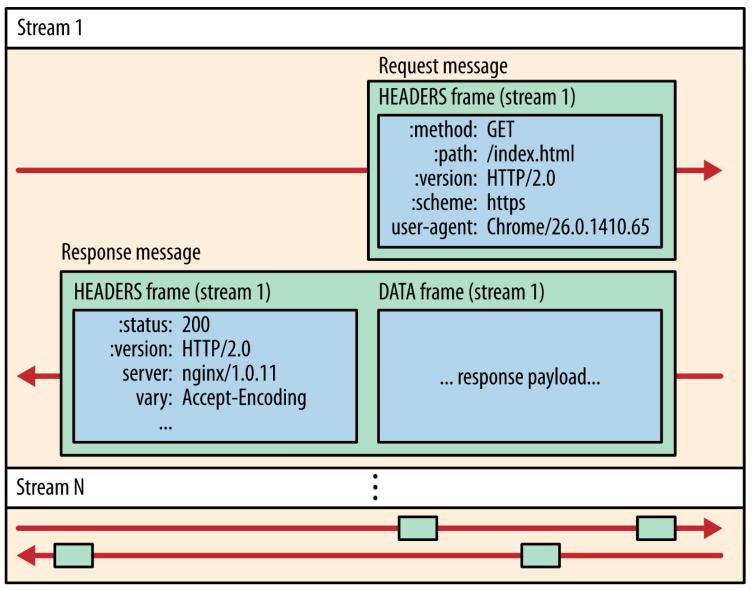
Small xfers can "block" behind large xfer

- No reordering
  - HTTP response does not "identify" which request it's in response to; obvious in simple request/response
- Can behave *worse* than parallel + persistent
  - Can send expensive query 1 on conn 1, while sending many cheap queries on conn 2

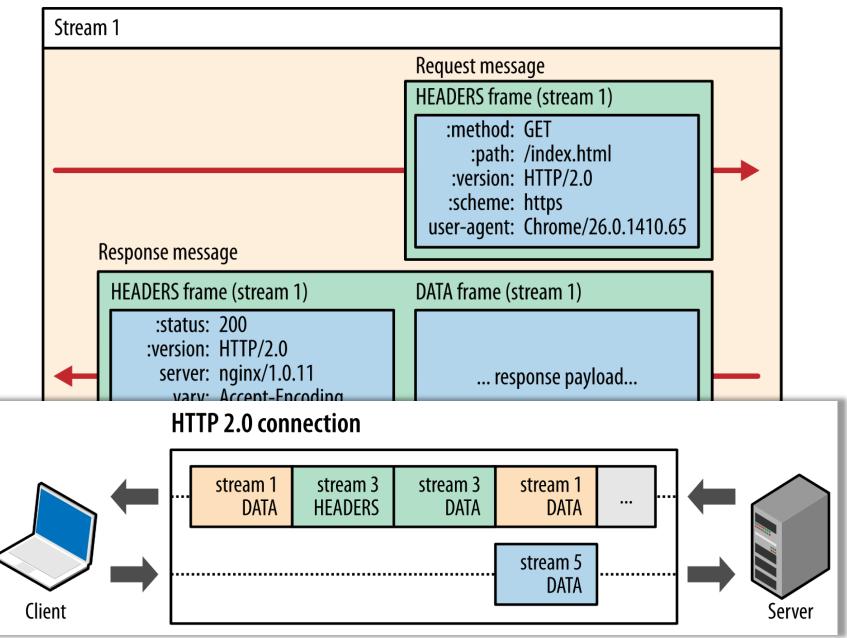
## Google's SPDY → HTTP/2 Standard

- Server "push" for content
  - One client request, multiple responses
  - After all, server knows that after parsing HTML, client will immediately request embedded URLs
- Better pipelining and xfer
  - Multiplexing multiple xfers w/o HOL blocking
  - Request prioritization
  - Header compression

#### Connection



#### Connection



### Summary

• HTTP: dominant application layer protocol for the web

HTTP caching had a limited impact (CDNs next)

 Recent optimization and evolution of HTTP for performance and efficiency