“Too Many Cookies” Problem

- Roommates Lance and James want a bag of cookies in the room at all times, but don’t want to buy too many cookies.
- They buy cookies independently, using the following sequence:
  - Look in cabinet: Out of cookies
  - Leave for Wawa to buy cookies
  - Arrive at Wawa
  - Buy a bag of cookies
  - Arrive home and put cookies in cabinet

Using A Note?

James and Lance’s Cookie Optimization Algorithm:

```java
if (noCookies) { // check if roommate left a note
  if (noNote) {
    leave note;   // let them know you went to Wawa
    buy cookies;
    remove note;
  }
}
```

- Any issue with this approach?

<table>
<thead>
<tr>
<th>Time</th>
<th>James</th>
<th>Lance</th>
</tr>
</thead>
<tbody>
<tr>
<td>15:00</td>
<td>Look in cabinet: out of cookies</td>
<td></td>
</tr>
<tr>
<td>15:05</td>
<td>Leave for Wawa</td>
<td></td>
</tr>
<tr>
<td>15:10</td>
<td>Arrive at Wawa</td>
<td>Look in cabinet: out of cookies</td>
</tr>
<tr>
<td>15:15</td>
<td>Buy a bag of cookies</td>
<td>Leave for Wawa</td>
</tr>
<tr>
<td>15:20</td>
<td>Arrive home; put cookies away</td>
<td>Arrive at Wawa</td>
</tr>
<tr>
<td>15:25</td>
<td>Buy a bag of cookies</td>
<td></td>
</tr>
<tr>
<td></td>
<td>Arrive home; put cookies away</td>
<td></td>
</tr>
</tbody>
</table>

- Oh No! Too many cookies.
Using A Note?

Any issue with this approach?

James
if (noCookies) {
    if (noNote) {
        leave note;
        buy cookies;
        remove note;
    } else {
        buy cookies;
        remove note;
    }
}

Lance
if (noCookies) {
    if (noNote) {
        leave note;
        buy cookies;
        remove note;
    } else {
        buy cookies;
        remove note;
    }
}

Why Solution #1 Does Not Work

Threads can get context-switched at any time
Too many cookies!

Possible Solution #2: Leave Note First

James
leave noteA
if (noNoteB) {
    if (noCookies) {
        buy cookies
    }
}
remove noteA

Lance
leave noteB
if (noNoteA) {
    if (noCookies) {
        buy cookies
    }
}
remove noteB

Does this method work?

Didn’t buy cookies

Didn’t buy cookies

Possible Solution #3: One Spin-waits

Problem was that threads checked once and moved on
So have one of them spin-wait on the note

James
leave noteA
while (noteB) do nothing;
if (noCookies) {
    buy cookies;
    remove noteA
}

Lance
leave noteB
if (noNoteA) {
    if (noCookies) {
        buy cookies;
        remove noteB
    }
}

Would this fix the problem?
Yes, but complicated, different code for different threads, busy waiting wasteful, and not fair
Threads Example: Shared Counter

- Google gets millions of hits a day. Uses multiple threads (on multiple processors) to speed things up.
- Simple shared state error: each thread increments a shared counter to track the number of hits today:

  ```
  ... 
  hits = hits + 1; 
  ... 
  ```

- What happens when two threads execute this code concurrently?

Problem with Shared Counters

- One possible result: lost update!
  ```
  hits = 0
  ```
  
  ![Diagram showing time, thread T1, thread T2, reading, writing, and update operations]

- Another possible result: everything works!

  ```
  hits = 0
  ```
  
  ![Diagram showing another set of time, thread T1, thread T2, reading, writing, and update operations]

- Another possible result: everything works
- This is called a "race condition"

Race Conditions

- Race condition: accesses to shared state that can lead to a timing dependent error
  - Whether it happens depends on how threads are scheduled
- Difficult to avoid because:
  - Must make sure all possible schedules are safe.
  - Number of possible schedule permutations is huge.
  - One or more of them may be "bad"
  - They are intermittent
  - Timing dependent => small changes can hide or reveal bug
    - Adding a print statement
    - Running on a different machine
It’s Actually Even Worse

- Compilers reorder instruction issue within a thread
  - To optimize register usage and hence run code faster
- Hardware reorders instruction execution/completion
  - E.g. write buffers, etc
- All done to optimize execution speed
- But they don’t know about multiple threads and issues across them

Preventing Race Conditions: Atomicity

- Atomic unit = instruction sequence guaranteed to execute indivisibly (also called a “critical section”).
  - If two threads execute the same atomic unit at the same time, one thread will execute the whole sequence before the other begins

```
hits = 0
```
```
T1

time

T2
```
```
hits = hits + 1
```
```
hits = 2
```

- How to make multiple instrs seem like an atomic one?

Providing Atomicity

- Have hardware provide better primitives than atomic load and store.
- Build higher-level programming abstractions on this new hardware support.
- Example: locks

  ```
  Acquire(lock);
  hits = hits + 1;
  Release(lock);
  ```

  Critical section

These must be atomic operations — if two threads are waiting for the lock, and both see it is free, only one grabs it

Preventing Race Conditions: Atomicity

- Counter problem

  ```
  Acquire(lock);
  hits = hits + 1;
  Release(lock);
  ```

  Critical section
Preventing Race Conditions: Atomicity

- Cookies problem

```java
Acquire(lock);
if (noCookies)
    buy cookies;
Release(lock);
```

Desirable Properties:
1. At most one holder, or thread in critical section, at a time (safety)
2. If no one is holding the lock, an acquire gets the lock (progress)
3. If all lock holders finish and there are no higher priority waiters, waiter eventually gets the lock (progress)

Rules for Using Locks

- Lock is initially free
- Always acquire before accessing shared data structure
- Always release after finishing with shared data
  - Only the lock holder can release
- Don’t access shared data without lock

Some Definitions

- Synchronization:
  - Ensuring proper cooperation among threads
  - Mutual exclusion, event synchronization

- Mutual exclusion:
  - Ensuring that only one thread does a particular thing at a time. One thread doing it excludes another from doing it at the same time.

- Critical section:
  - Piece of code that only one thread can “be in” at a given time. Only one thread at a time will be allowed to get into the section of code.

- Lock:
  - Prevents someone from doing something
  - Lock before entering critical section, before accessing shared data
  - Unlock when leaving, after done accessing shared data
  - Wait if locked

- Event synchronization:
  - Making sure an event in one thread does not happen before/after an event in another thread

Implementing Mutual Exclusion (Locks)

What makes a good solution?

- Only one process/thread inside a critical section at a time
- No assumptions need to be made about CPU speeds
- A process/thread inside a critical section should not be blocked by any process outside the critical section
- No one waits forever

- Should work for multiprocessors
- Should allow same code for all processes/threads
Simple Lock Variables

```c
Acquire(lock) {
    while (lock.value == 1) {
    }
    lock.value = 1;
}

Thread 1: 
Acquire(lock) {
    while (lock.value == 1) {
{context switch)
    lock.value = 1;
} 
}

Thread 2: 
Acquire(lock) {
    while (lock.value == 1) {
{context switch)
    lock.value = 1;
{context switch)
}
```

Prevent Context Switches in Critical Section

- On a uniprocessor, operations are atomic as long as a context switch doesn’t occur
- Context switches are caused either by actions the thread takes (e.g. traps etc) or by external interrupts
- The former can be controlled
- Disable interrupts during certain portions of code?
  - Delay the handling of external events

Why Enable or Disable Interrupts

- Interrupts are important
  - Process I/O requests (e.g. keyboard)
  - Implement preemptive CPU scheduling
- Disabling interrupts can be helpful
  - Introduce uninterruptible code regions
  - Think sequentially most of the time
  - Delay handling of external events

Disabling Interrupts for Critical Section?

- **Acquire()** : disable interrupts
- **Release()** : enable interrupts

Issues:
- Critical sections can be arbitrarily long
"Disable Interrupts" to Implement Mutex

Acquire(lock) {
    disable interrupts;
    while (lock.value != 0) {
        lock.value = 1;
        enable interrupts;
    }
    lock.value = 1;
    enable interrupts;
}

- Don’t let acquire be interrupted before sets lock.value to 1
  - This was the problem when interrupts weren’t disabled
- Don’t disable interrupts for entire critical section
- Issues:
  - May disable interrupts forever

Fix "Disable Forever" problem?

Acquire(lock) {
    disable interrupts;
    while (lock.value != 0) {
        lock.value = 0;
        enable interrupts;
    }
    lock.value = 1;
    enable interrupts;
}
Release(lock) {
    disable interrupts;
    lock.value = 0;
    enable interrupts;
}

- Enable interrupts during spin loop
- Disable interrupts only when accessing lock.value
- Issues:
  - Cannot be interrupted after loop and before setting value

Another Implementation

Acquire(lock) {
    disable interrupts;
    if (lock.value != 0) {
        Enqueue me for lock;
        Yield();
    }
    lock.value = 1;
    enable interrupts;
}
Release(lock) {
    disable interrupts;
    if (anyone in queue) {
        Dequeue a thread;
        make it ready;
    }
    lock.value = 0;
    enable interrupts;
}

Avoid busy-waiting

Issues
- Interrupt based approaches don’t work for multiprocessors
- Cannot allow user code to disable interrupts

Atomic Operations

- A thread executing an atomic instruction can’t be preempted or interrupted while it’s doing it
- Atomic operations on same memory value are serialized
  - Even on multiprocessors!
  - Result is consistent with some sequential ordering of operations
  - Without atomic ops, simultaneous writes by different threads may produce a garbage value, or read that happens simultaneously with a write may read garbage value
- Don’t usually require special privileges, can be user level
Peterson's Algorithm

- See textbook

```c
int turn;
int interested[N];

void enter_region(int process)
{
    int other;
    other = 1 - process;
    interested[process] = TRUE; /* express interest */
    turn = other; /* give turn to other process */
    while(turn == process && interested[other] == TRUE);
    /* wait till other loses interest or gives me turn */
}
```

- 5 writes and 2 reads

Atomic Read-Modify-Write Instructions

- LOCK prefix in x86
  - Make a specific of set instructions atomic
  - Can be used to implement Test&Set
- Exchange (xchg, x86 architecture)
  - Swap register and memory
  - Atomic (even without LOCK)
- Fetch&Add or Fetch&Op
  - Atomic instructions for large shared memory multiprocessors
- Load linked and store conditional (LL-SC)
  - Two separate instructions (LL, SC) that are used together
  - Read value in one instruction (load linked)
    - Do some operations;
    - When time to store, check if value has been modified. If not, ok; otherwise, jump back to start

A Simple Solution with Test&Set

- Define TAS(lock)
  - If successfully set (wasn’t already set when tested but this operation set it), return 1;
  - Otherwise, return 0;
- Any issues with the following solution?

```c
Acquire(lock) {
    while (!TAS(lock.value))
        ;
}
```

```c
Release(lock) {
    while (!TAS(lock.guard))
        ;
    if (lock.value) {
        enqueue the thread;
        block and lock.guard = 0;
    } else {
        lock.value = 1;
        lock.guard = 0;
    }
}
```

Mutex with Less Waiting?

- Separate access to lock variable from value of it
Example: Protect a Shared Variable

```c
Acquire(lock); /* system call */
    count++;
Release(lock) /* system call */
```

- **Acquire(mutex) system call**
  - Pushing parameter, sys call # onto stack
  - Generating trap/interrupt to enter kernel
  - Jump to appropriate function in kernel
  - Verify process passed in valid pointer to mutex
  - Minimal spinning
  - Block and unblock process if needed
  - Get the lock
- **Execute “count++;”**
- **Release(mutex) system call**

Available Primitives and Operations

- **Test-and-set**
  - Works at either user or kernel level
- **System calls for block/unblock**
  - **Block** takes some token and goes to sleep
  - **Unblock** “wakes up” a waiter on token

Block and Unblock System Calls

- **Block( lock )**
  - Spin on lock.guard
  - Save the context to TCB
  - Enqueue TCB to lock.q
  - Clear lock.guard
  - Call scheduler

- **Unblock( lock )**
  - Spin on lock.guard
  - Dequeue a TCB from lock.q
  - Put TCB in ready queue
  - Clear lock.guard

Always Block

```c
Acquire(lock) {
    while (!TAS(lock.value))
        Block( lock );
}
```

```c
Release(lock) {
    lock.value = 0;
    Unblock( lock );
}
```

- **Good**
  - Acquire won’t make a system call if TAS succeeds
- **Bad**
  - TAS instruction locks the memory bus
  - Block/Unblock still has substantial overhead
Always Spin

```
Acquire(lock) {
    while (!TAS(lock.value))
        while (lock.value)
            ;
}
```

- Two spinning loops in Acquire()?  
- CPU
  - L1
  - L1
  - L2

Competitive Algorithms

- An algorithm is c-competitive if for every input sequence $\sigma$
  
  \[ C_A(\sigma) \leq c \times C_{opt}(\sigma) + k \]

  - $c$ is a constant
  - $C_A(\sigma)$ is the cost incurred by algorithm A in processing $\sigma$
  - $C_{opt}(\sigma)$ is the cost incurred by the optimal algorithm in processing $\sigma$

- What we want is to have $c$ as small as possible
  - Deterministic
  - Randomized

Optimal Algorithms

- What is the optimal solution to spin vs. block?
  - Know the future
  - Exactly when to spin and when to block
- But, we don’t know the future
  - There is no online optimal algorithm

- Offline optimal algorithm
  - Afterwards, derive exactly when to block or spin (“what if”)
  - Useful to compare against online algorithms

Constant Competitive Algorithms

```
Acquire(lock, N) {
    int i;
    while (!TAS(lock.value)) {
        i = N;
        while (!lock.value && i)
            i--;
        if (!i)
            Block(lock);
    }
```
Approximate Optimal Online Algorithms

- **Main idea**
  - Use past to predict future

- **Approach**
  - Random walk
    - Decrement N by a unit if the last Acquire() blocked
    - Increment N by a unit if the last Acquire() didn’t block
  - Recompute N each time for each Acquire() based on some lock-waiting distribution for each lock

- **Theoretical results**
  \[ E_{CA}(\sigma(P)) \leq \frac{e}{(e-1)} \times E_{C_{opt}}(\sigma(P)) \]

  The competitive factor is about 1.58.

---

The Big Picture

<table>
<thead>
<tr>
<th>Concurrent applications/software</th>
</tr>
</thead>
<tbody>
<tr>
<td>Shared Objects</td>
</tr>
<tr>
<td>----------------</td>
</tr>
<tr>
<td>High-Level Atomic API (portable)</td>
</tr>
<tr>
<td>Low-Level Atomic Ops (specific)</td>
</tr>
</tbody>
</table>

- **Summary**
  - **Disabling interrupts for mutex**
    - There are many issues
    - When making it work, it works for only uniprocessors
  - **Atomic instruction support for mutex**
    - Atomic load and stores are not good enough
    - Test&set and other instructions are the way to go
  - **Competitive spinning**
    - Spin at the user level most of the time
    - Make no system calls in the absence of contention
    - Have more threads than processors