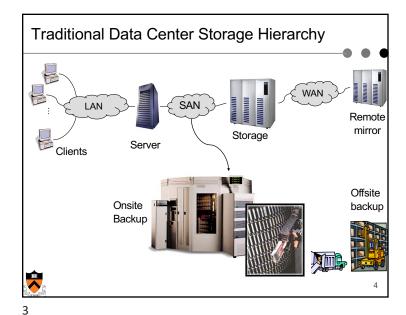
COS 318: Operating Systems

File Systems: Networked,
Abstractions and Protection

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(http://www.cs.princeton.edu/courses/cos318/)

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Topics

- What's behind the file system: Networked Storage hierarchy
- More on the file system abstraction
- File system protection

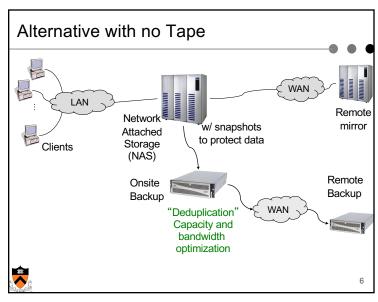
Evolved Data Center Storage Hierarchy

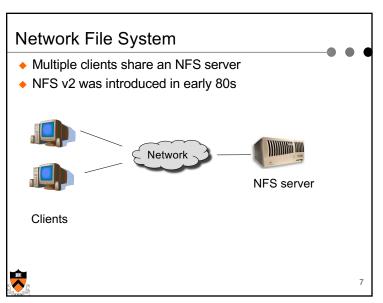
Network
Attached
Storage
(NAS)

Onsite
Backup

Offsite
backup

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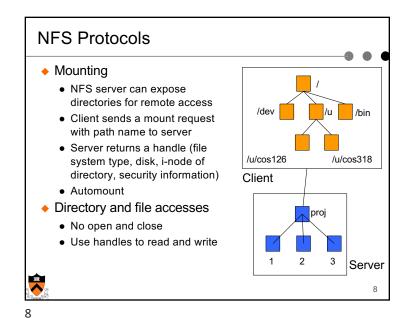


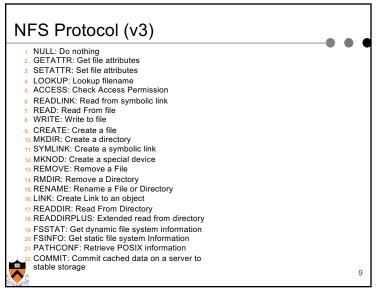


"Public Cloud" Storage Hierarchy

WAN
Interfaces

Examples: Google GFS, Spanner,
Apple Cloud, Amazon S3, Box,
Dropbox, Mozy, etc

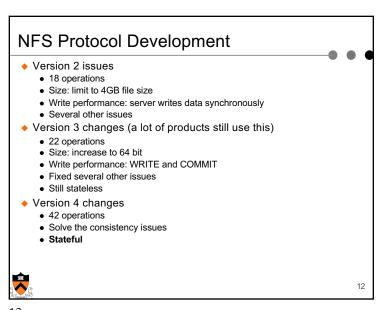


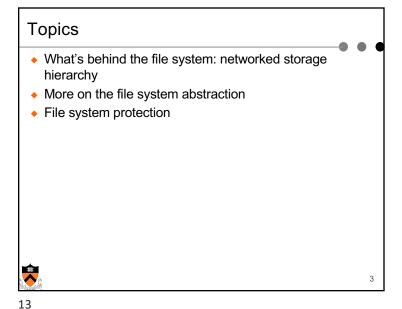


NFS Client Caching Issues Consistency among multiple client caches Client cache contents may not be up-to-date Multiple writes can happen simultaneously Solutions Expiration Read-only file and directory data (expire in 60 seconds) Data written by the client machine (write back in 30 seconds) No shared caching A file can be cached at only one client cache Network lock manager Sequential consistency (one writer or N readers)

NFS Architecture

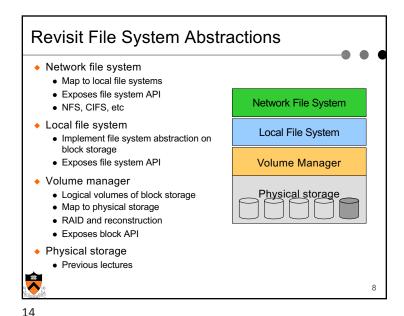
Client kernel
Virtual file system
Virtual file system
NFS
FS
FS
Buffer cache
Buffer cache
Network





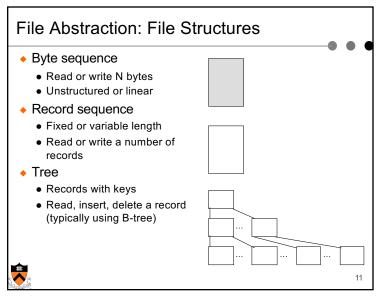
Volume Manager · Group multiple storage partitions into a logical volume Virtualization of capacity and performance No need to deal with physical disk or sector numbers ◆ Read(vol#, block#, buf, n) · Reliable block storage • Include RAID, tolerating device failures Provide error detection at block level Remote abstraction · Block storage in the cloud · Remote volumes for disaster recovery · Remote mirrors can be split or merged for backups How to implement? • OS kernel: Windows, OSX, Linux, etc. Storage subsystem: EMC, Hitachi, HP, IBM, NetApp

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File versus Block Abstractions Disk/Volume abstraction File abstraction Block oriented Byte oriented Block numbers Named files No protection among users of Users protected from each the system other Data might be corrupted if Robust to machine failures machine crashes Support file systems, database • Emulate block storage systems, etc. interface

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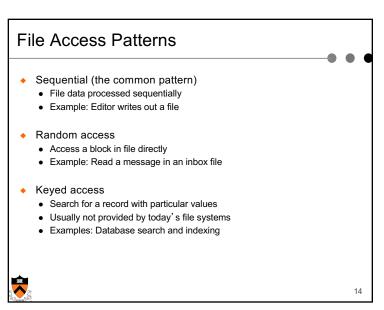


File Abstraction: File Operations Operations for "sequence of bytes" files Create: create a file (mapping from a name to a file) Delete: delete a file Open: including authentication Close: done with accessing a file Seek: jump to a particular location in a file Read: read some bytes from a file Write: write some bytes to a file A few more operations on directories: later Implementation challenges Keep disk accesses low Keep space overhead low

File Abstraction: File Types

ASCII
Binary data
Record
Tree
An Unix executable file
header: magic number, sizes, entry point, flags
text
data
relocation bits
symbol table

Devices
Character special files (to model terminals, printers)
Block special files (to model disks)
Everything else in the system



File system abstraction

- Directory
 - · Group of named files or subdirectories
 - Mapping from file name to file metadata location
- Path
 - String that uniquely identifies file or directory
 - Ex: /cse/www/education/courses/cse451/12au
- Links
 - Hard link: link from name to metadata location
 - Soft link: link from name to alternate name
- Mount
 - Mapping from name in one file system to root of another



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VM Page Table vs. File System Metadata

Page table

- Manage the mappings of an address space
- Map virtual to physical page #
- Check access permission and illegal addressing
- TLB does it all in one cycle

File metadata

- Manage the mappings of files
- Map byte offset to disk block address
- Check access permission and illegal addressing
- Implemented in software, may cause I/Os



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File System vs. Virtual Memory

- Similarity
 - Location transparency
 - Size "obliviousness"
 - Protection
- File system is easier than VM in some ways
 - File system mappings can be slow
 - Files are dense and mostly sequential, while page tables deal with sparse address spaces and random accesses
- File system is more difficult than VM in some ways
 - Each layer of translation causes potential I/Os
 - Memory space for caching is never enough
 - File size range vary: many < 10k, some > GB
 - Implementation must be reliable



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Topics

- What's behind the file system: Storage hierarchy
- More on file system abstraction
- File system protection



Protection: Policy vs. Mechanism

- Policy is about what
- Mechanism is about how
- A security policy defines acceptable and unacceptable behaviors. Examples:
 - A given user can only allocate 4GB of disk storage
 - · No one but root can write to the password file
 - · A user is not allowed to read others' mail files
- A protection system is the mechanism to enforce a security policy
 - Same set of choices, no matter what policies
- ◆ Principle of least privilege



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Authentication

- Usually done with passwords
 - Relatively weak, because you must remember them
- Passwords are stored in an encrypted form
 - Use a "secure hash" (one way only)
- Issues
 - Passwords should be obscure, to prevent "dictionary attacks"
 - Each user has many passwords
- Alternatives?



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Protection Mechanisms



- Identity check
 - · Unix: password
 - Credit card: last 4 digits of credit card # + SSN + zipcode
 - · Airport: driver's license or passport
- Authorization
 - Determine if x is allowed to do y
 - Need a simple database
- Access enforcement
 - Enforce authorization decision
 - Must make sure there are no loopholes



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Protection Domain

- Once identity known, provides rules
 - E.g. what is Bob allowed to do?
 - E.g. who can do what to file A?
- Protection matrix: domains and resources

	File A	Printer B	File C
Domain 1	R	W	RW
Domain 2	RW	W	
Domain 3	R		RW



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By Columns: Access Control Lists (ACLs)

- Each object has a list of <user, privilege> pairs
- ◆ ACL is simple, implemented in most systems
 - Owner, group, world
- Implementation considerations
 - Stores ACLs in each file
 - · Use login authentication to identify
 - Kernel implements ACLs
- Any issues?



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Access Enforcement

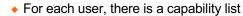
- Use a trusted party to
 - Enforce access controls
 - Protect authorization information
- Kernel is the trusted party
 - This part of the system can do anything it wants
 - If there is a bug, the entire system could be destroyed
 - Want it to be as small & simple as possible
- Security is only as strong as the weakest link in the protection system



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By Rows: Capabilities



- A lists of <object, privilege> pairs
- Capabilities provide both naming and protection
 - Can only "see" an object if you have a capability
- Implementation considerations
 - Architecture support
 - · Capabilities stored in the kernel
 - Capabilities stored in the user space in encrypted format
- Issues?



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Some Easy Attacks

- Abuse of valid privilege
 - On Unix, super-user can do anything
 - · Read your mail, send mail in your name, etc.
 - If you delete the code for COS318 project 5, your partner is not happy
- Spoiler/Denial of service (DoS)
 - Use up all resources and make system crash
 - Run shell script to: "while(1) { mkdir foo; cd foo; }"
- Listener
 - Passively watch network traffic



No Perfect Protection System

- Cannot prevent bad things, can only make it difficult to do them
- ◆ There are always ways to defeat protection
 - burglary, bribery, blackmail, bludgeoning, etc.
- Every system has holes



Summary

- Storage hierarchy can be complex
 - Reliability, security, performance and cost
 - Many things are hidden
- Key storage layers above hardware
 - Volume or block storage
 - Local file system
 - Network file system
- Protection
 - ACL is the default in file systems
 - More protection is needed in the cloud

