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2.2 MERGESORT

- *mergesort*
- *bottom-up mergesort*
- *sorting complexity*
- *divide-and-conquer*

Two classic sorting algorithms: mergesort and quicksort

Critical components in the world's computational infrastructure.

- Full scientific understanding of their properties has enabled us to develop them into practical system sorts.
- Quicksort honored as one of top 10 algorithms of 20th century in science and engineering.

Mergesort. [this lecture]

…

Quicksort. [next lecture]

…

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2.2 Mergesort

- *mergesort*
- *bottom-up mergesort*
- *sorting complexity*
- *divide-and-conquer*

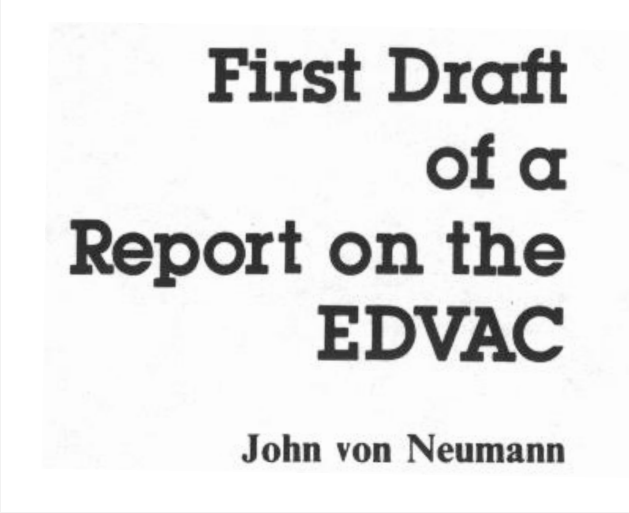
Mergesort

Basic plan.

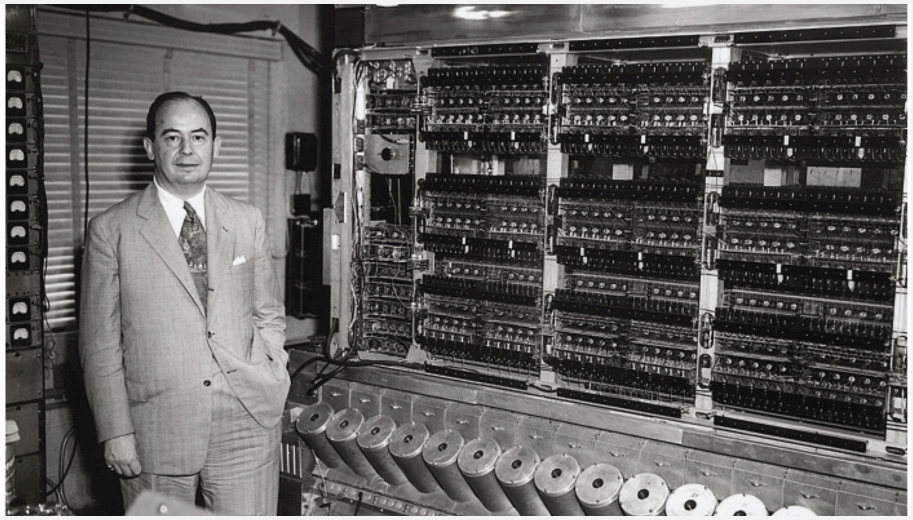
- Divide array into two halves.
- Recursively sort each half.
- Merge two halves.

input	M	E	R	G	E	S	O	R	T	E	X	A	M	P	L	E
sort left half	E	E	G	M	O	R	R	S	T	E	X	A	M	P	L	E
sort right half	E	E	G	M	O	R	R	S	A	E	E	L	M	P	T	X
merge results	A	E	E	E	E	G	L	M	M	O	P	R	R	S	T	X

Mergesort overview

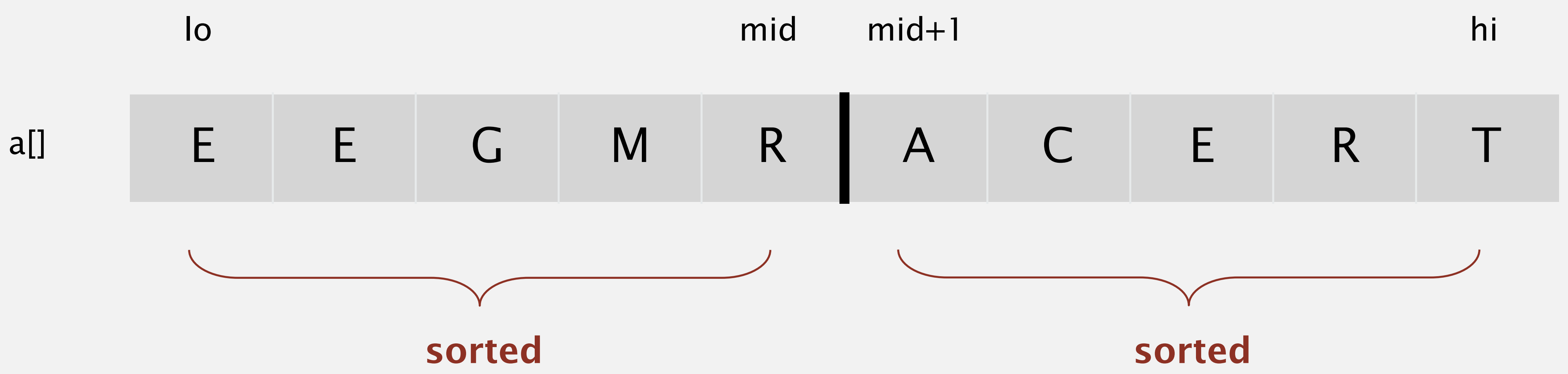

First Draft
of a
Report on the
EDVAC

John von Neumann



Abstract in-place merge demo

Goal. Given two sorted subarrays `a[lo]` to `a[mid]` and `a[mid+1]` to `a[hi]`, replace with sorted subarray `a[lo]` to `a[hi]`.



Mergesort: Transylvanian–Saxon folk dance

http://www.youtube.com/watch?v=XaqR3G_NVoo

Merging: Java implementation

```
private static void merge(Comparable[] a, Comparable[] aux, int lo, int mid, int hi)
{
   for (int k = lo; k <= hi; k++)
      aux[k] = a[k];

   int i = lo, j = mid+1;
   for (int k = lo; k <= hi; k++)
   {
      if      (i > mid)              a[k] = aux[j++];
      else if (j > hi)               a[k] = aux[i++];
      else if (less(aux[j], aux[i])) a[k] = aux[j++];
      else                           a[k] = aux[i++];
   }

}
```

copy

merge

	lo			i	mid			j		hi
aux[]	A	G	L	O	R	H	I	M	S	T

						k				
a[]	A	G	H	I	L	M				

Mergesort quiz 1

How many calls does `merge()` make to `less()` in order to merge two sorted subarrays, each of length $n / 2$, into a sorted array of length n?

A. $\sim \frac{1}{4} n$ to $\sim \frac{1}{2} n$

B. $\sim \frac{1}{2} n$

C. $\sim \frac{1}{2} n$ to $\sim n$

D. $\sim n$

merging two sorted arrays, each of length n/2

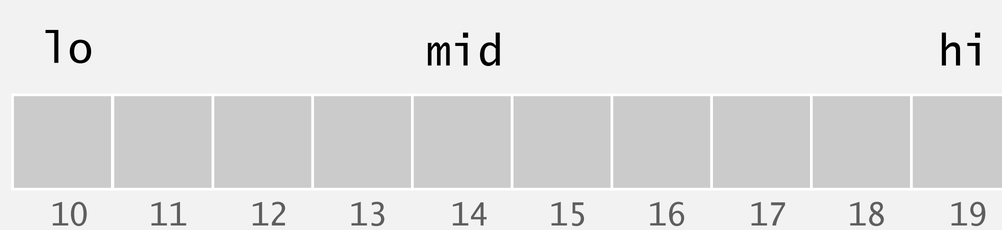
a_0	a_1	a_2	a_3	b_0	b_1	b_2	b_3

Mergesort: Java implementation

```java
public class Merge
{
   private static void merge(...)
   {  /* as before */  }

   private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi)
   {
      if (hi <= lo) return;
      int mid = lo + (hi - lo) / 2;
      sort(a, aux, lo, mid);
      sort(a, aux, mid+1, hi);
      merge(a, aux, lo, mid, hi);
   }

   public static void sort(Comparable[] a)
   {
      Comparable[] aux = new Comparable[a.length];
      sort(a, aux, 0, a.length - 1);
   }
}
```



Mergesort: trace

lo hi

a[]

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
	M	E	R	G	E	S	O	R	T	E	X	A	M	P	L	E
`merge(a, aux, 0, 0, 1)`	E	M	R	G	E	S	O	R	T	E	X	A	M	P	L	E
`merge(a, aux, 2, 2, 3)`	E	M	G	R	E	S	O	R	T	E	X	A	M	P	L	E
`merge(a, aux, 0, 1, 3)`	E	G	M	R	E	S	O	R	T	E	X	A	M	P	L	E
`merge(a, aux, 4, 4, 5)`	E	G	M	R	E	S	O	R	T	E	X	A	M	P	L	E
`merge(a, aux, 6, 6, 7)`	E	G	M	R	E	S	O	R	T	E	X	A	M	P	L	E
`merge(a, aux, 4, 5, 7)`	E	G	M	R	E	O	R	S	T	E	X	A	M	P	L	E
`merge(a, aux, 0, 3, 7)`	E	E	G	M	O	R	R	S	T	E	X	A	M	P	L	E
`merge(a, aux, 8, 8, 9)`	E	E	G	M	O	R	R	S	E	T	X	A	M	P	L	E
`merge(a, aux, 10, 10, 11)`	E	E	G	M	O	R	R	S	E	T	A	X	M	P	L	E
`merge(a, aux, 8, 9, 11)`	E	E	G	M	O	R	R	S	A	E	T	X	M	P	L	E
`merge(a, aux, 12, 12, 13)`	E	E	G	M	O	R	R	S	A	E	T	X	M	P	L	E
`merge(a, aux, 14, 14, 15)`	E	E	G	M	O	R	R	S	A	E	T	X	M	P	E	L
`merge(a, aux, 12, 13, 15)`	E	E	G	M	O	R	R	S	A	E	T	X	E	L	M	P
`merge(a, aux, 8, 11, 15)`	E	E	G	M	O	R	R	S	A	E	E	L	M	P	T	X
`merge(a, aux, 0, 7, 15)`	A	E	E	E	E	G	L	M	M	O	P	R	R	S	T	X

result after recursive call

Mergesort quiz 2

Which of the following subarray lengths will occur when running mergesort on an array of length 12?

A. $\{ 1, 2, 3, 4, 6, 8, 12 \}$

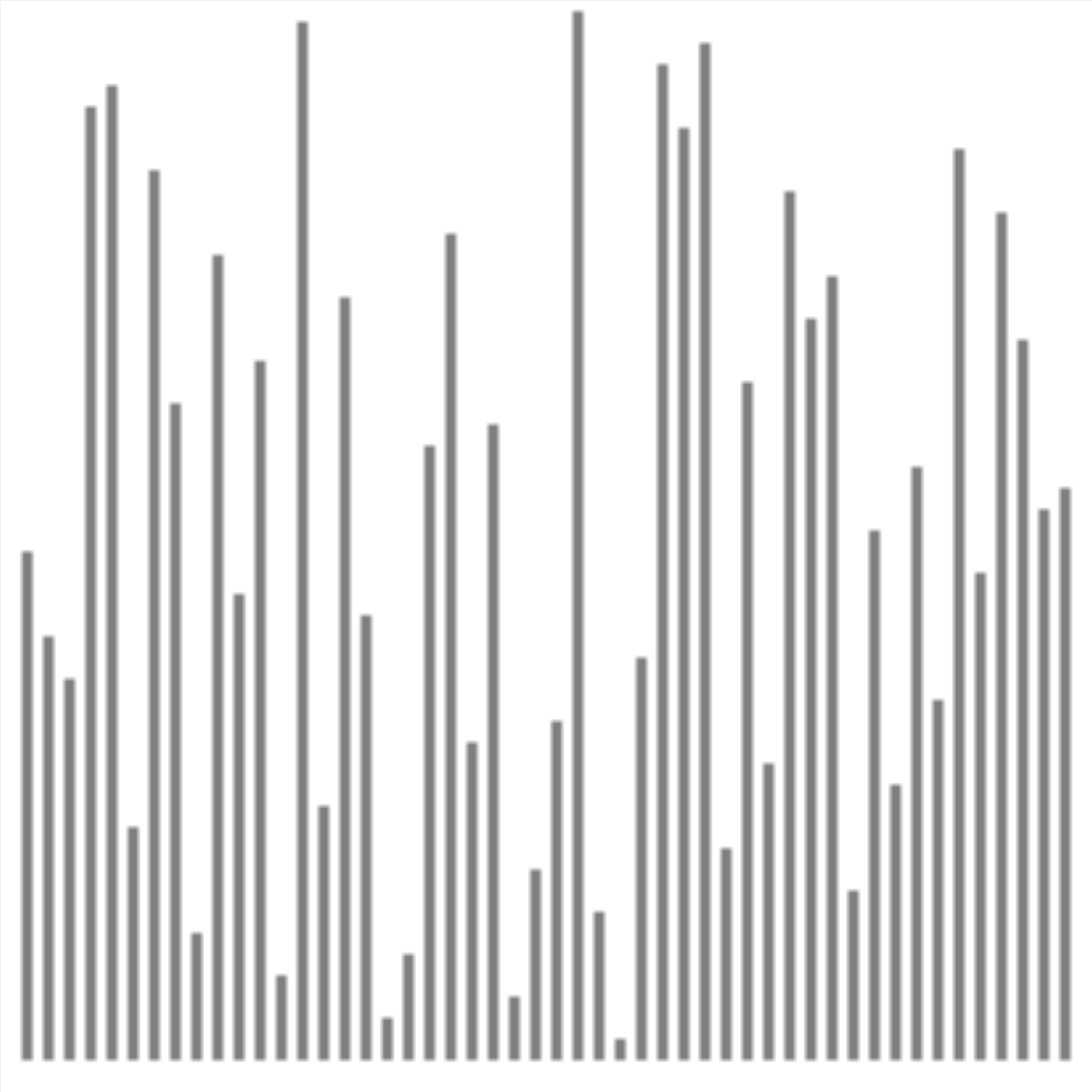
B. $\{ 1, 2, 3, 6, 12 \}$

C. $\{ 1, 2, 4, 8, 12 \}$

D. $\{ 1, 3, 6, 9, 12 \}$

Mergesort: animation

50 random items



http://www.sorting-algorithms.com/merge-sort

▲ algorithm position

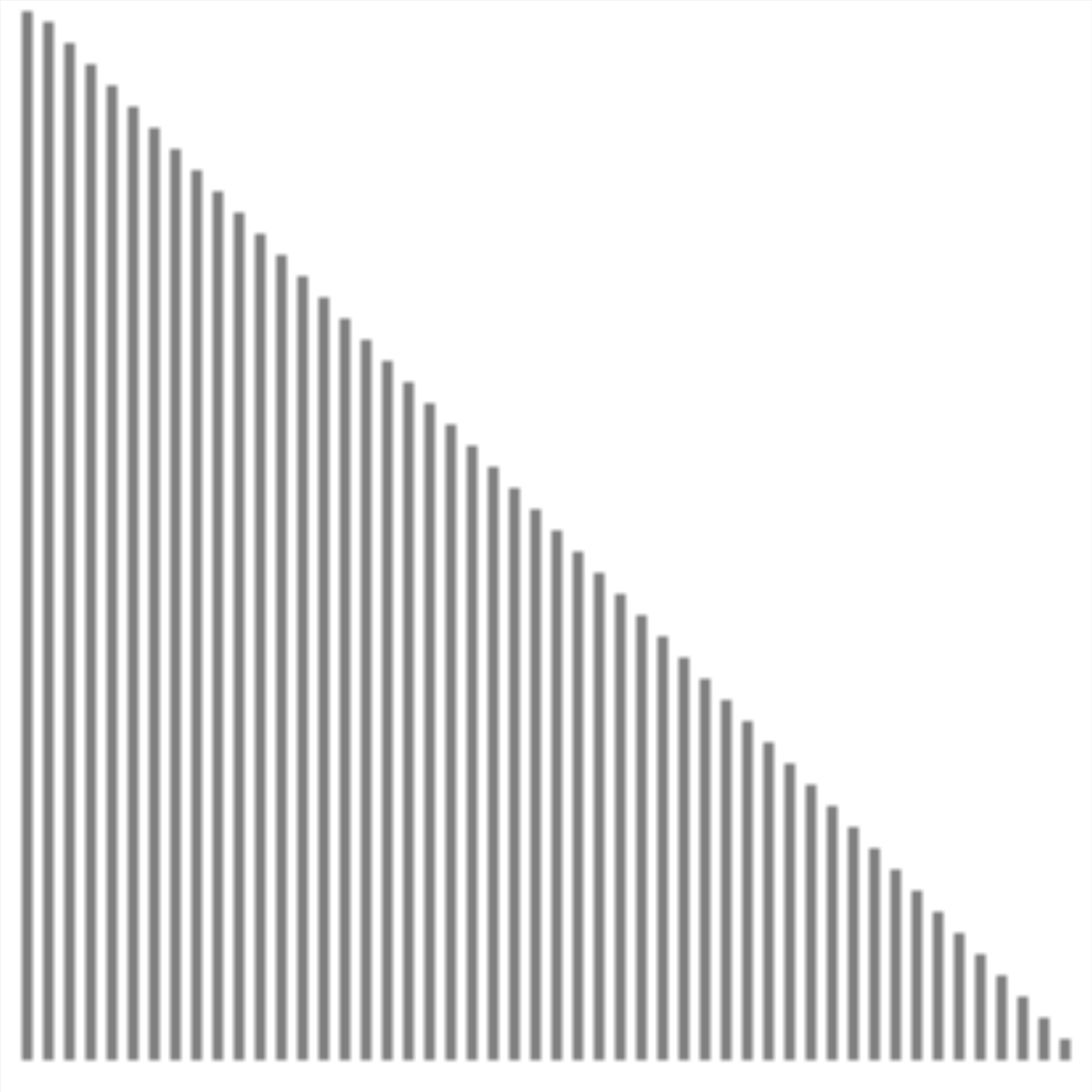
in order

current subarray

not in order

Mergesort: animation

50 reverse-sorted items



http://www.sorting-algorithms.com/merge-sort

▲ algorithm position

in order

current subarray

not in order

Mergesort: empirical analysis

Running time estimates:

- Laptop executes 10^8 compares/second.
- Supercomputer executes 10^{12} compares/second.

	insertion sort (n^2)			mergesort ($n \log n$)		
computer	**thousand**	**million**	**billion**	**thousand**	**million**	**billion**
home	instant	2.8 hours	317 years	instant	1 second	18 min
super	instant	1 second	1 week	instant	instant	instant

Bottom line. Good algorithms are better than supercomputers.

Mergesort analysis: number of compares

Proposition. Mergesort uses $\le n \lg n$ compares to sort any array of length n.

Pf sketch. The number of compares $C(n)$ to mergesort an array of length n satisfies the recurrence:

$$C(n) \le \underbrace{C(\lceil n/2 \rceil)}_{\text{left half}} + \underbrace{C(\lfloor n/2 \rfloor)}_{\text{right half}} + \underbrace{n - 1}_{\text{merge}} \quad \text{for } n > 1, \text{ with } C(1) = 0.$$

We assume n is a power of 2 and solve this simpler recurrence:

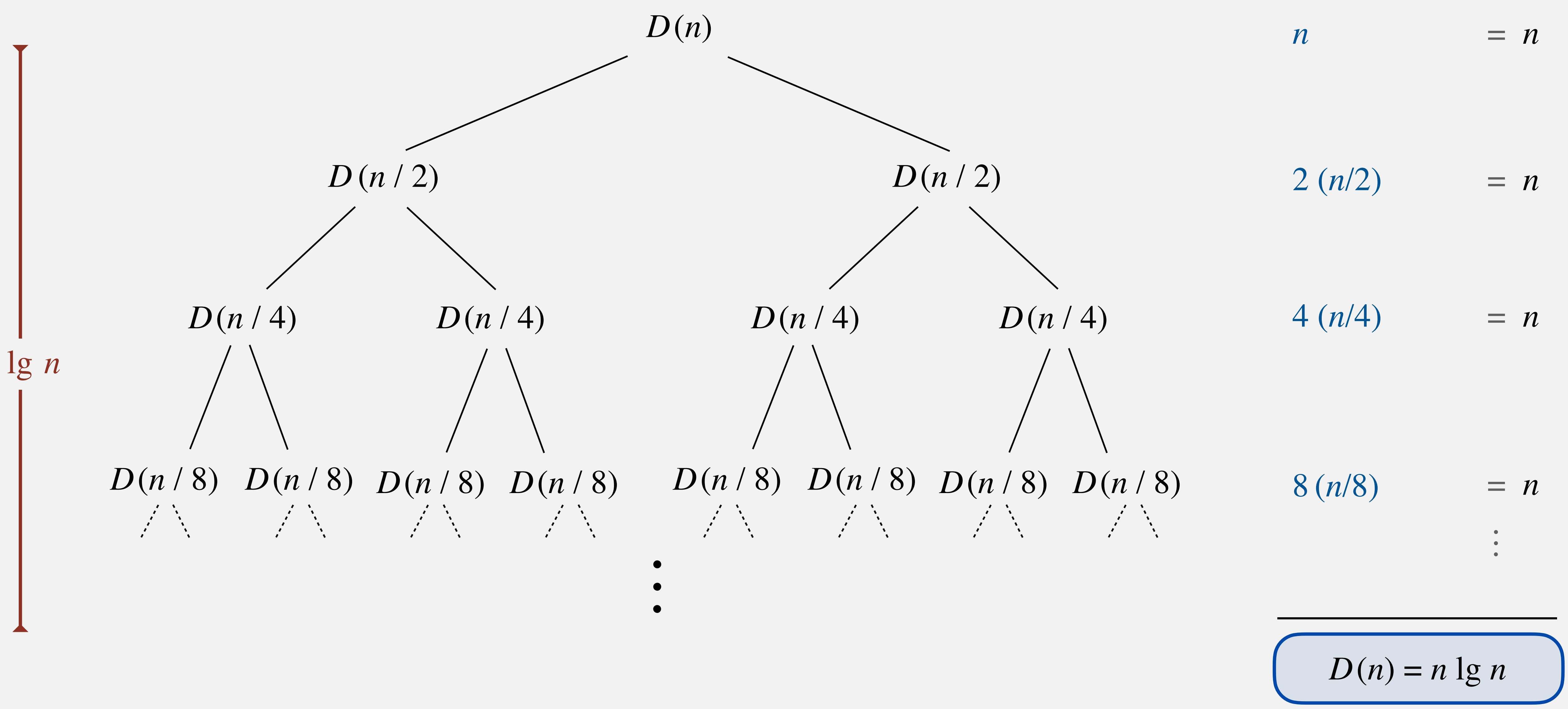
$$D(n) = 2\,D(n/2) + n, \text{ for } n > 1, \text{ with } D(1) = 0.$$

proposition holds for all n
(analysis cleaner in this case)

Divide-and-conquer recurrence

Proposition. If $D(n)$ satisfies $D(n) = 2\,D(n/2) + n$ for $n > 1$, with $D(1) = 0$, then $D(n) = n \lg n$.

Pf by picture. [assuming n is a power of 2]



Mergesort analysis: number of array accesses

Proposition. Mergesort uses $\le 6\, n \lg n$ array accesses to sort any array of length n.

Pf sketch. The number of array accesses $A(n)$ satisfies the recurrence:

$$A(n) \le A(\lceil n/2 \rceil) + A(\lfloor n/2 \rfloor) + 6n \quad \text{for } n > 1, \text{ with } A(1) = 0.$$

Key point. Any algorithm with the following structure takes $n \log n$ time:

```
public static void f(int n)
{
    if (n == 0) return;
    f(n/2);          ← solve two problems
    f(n/2);          ← of half the size
    linear(n);       ← do a linear amount of work
}
```

Notable examples. FFT, hidden-line removal, Kendall-tau distance, …

Mergesort analysis: memory

Proposition. Mergesort uses extra space proportional to n.

Pf. The array `aux[]` needs to be of length n for the last merge.

two sorted subarrays

A	C	D	G	H	I	M	N	U	V

B	E	F	J	O	P	Q	R	S	T

merged result

A	B	C	D	E	F	G	H	I	J	M	N	O	P	Q	R	S	T	U	V

essentially negligible

Def. A sorting algorithm is in-place if it uses $\leq c \log n$ extra memory.

Ex. Insertion sort and selection sort.

Challenge 1 (not hard). Use `aux[]` array of length $\sim \frac{1}{2}\, n$ instead of n.

Challenge 2 (very hard). In-place merge. [Kronrod 1969]

Mergesort quiz 3

Is our implementation of mergesort stable?

A. Yes.

B. No, but it can be easily modified to be stable.

C. No, mergesort is inherently unstable.

D. *I don't remember what stability means.*

a sorting algorithm is stable if it preserves the relative order of equal keys

input	C	A_1	B	A_2	A_3
sorted	A_3	A_1	A_2	B	C

not stable

Stability: mergesort

Proposition. Mergesort is stable.

```
public class Merge
{
   private static void merge(...)
   {  /* as before */  }

   private static void sort(Comparable[] a, Comparable[] aux, int lo, int hi)
   {
      if (hi <= lo) return;
      int mid = lo + (hi - lo) / 2;
      sort(a, aux, lo, mid);
      sort(a, aux, mid+1, hi);
      merge(a, aux, lo, mid, hi);
   }

   public static void sort(Comparable[] a)
   {  /* as before */  }
}
```

Pf. Suffices to verify that merge operation is stable.

Stability: mergesort

Proposition. Merge operation is stable.

```
private static void merge(...)
{
   for (int k = lo; k <= hi; k++)
      aux[k] = a[k];

   int i = lo, j = mid+1;
   for (int k = lo; k <= hi; k++)
   {
      if      (i > mid)              a[k] = aux[j++];
      else if (j > hi)               a[k] = aux[i++];
      else if (less(aux[j], aux[i])) a[k] = aux[j++];
      else                           a[k] = aux[i++];
   }
}
```

0	1	2	3	4
A_1	A_2	A_3	B	D

5	6	7	8	9	10
A_4	A_5	C	E	F	G

Pf. Takes from left subarray if equal keys.

Mergesort: practical improvement

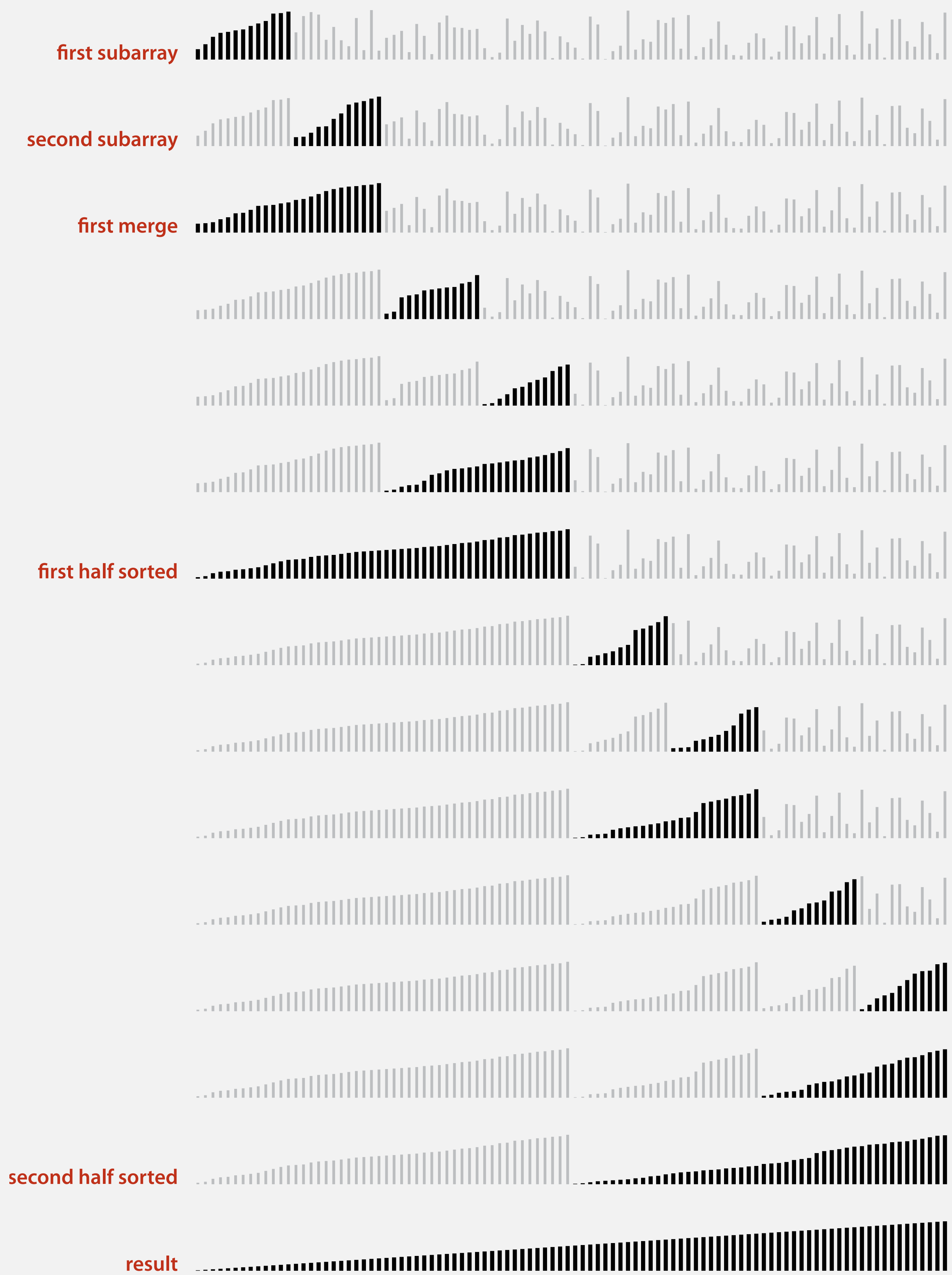
Use insertion sort for small subarrays.

- Mergesort has too much overhead for tiny subarrays.
- Cutoff to insertion sort for ≈ 10 items.

```
private static void sort(...)
{
   if (hi <= lo + CUTOFF - 1)
   {
      Insertion.sort(a, lo, hi);
      return;
   }

   int mid = lo + (hi - lo) / 2;
   sort (a, aux, lo, mid);
   sort (a, aux, mid+1, hi);
   merge(a, aux, lo, mid, hi);
}
```

Mergesort with cutoff to insertion sort: visualization



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2.2 Mergesort

- mergesort
- bottom-up mergesort
- sorting complexity
- divide-and-conquer

Bottom-up mergesort

Basic plan.

- Pass through array, merging subarrays of size 1.
- Repeat for subarrays of size 2, 4, 8, ...

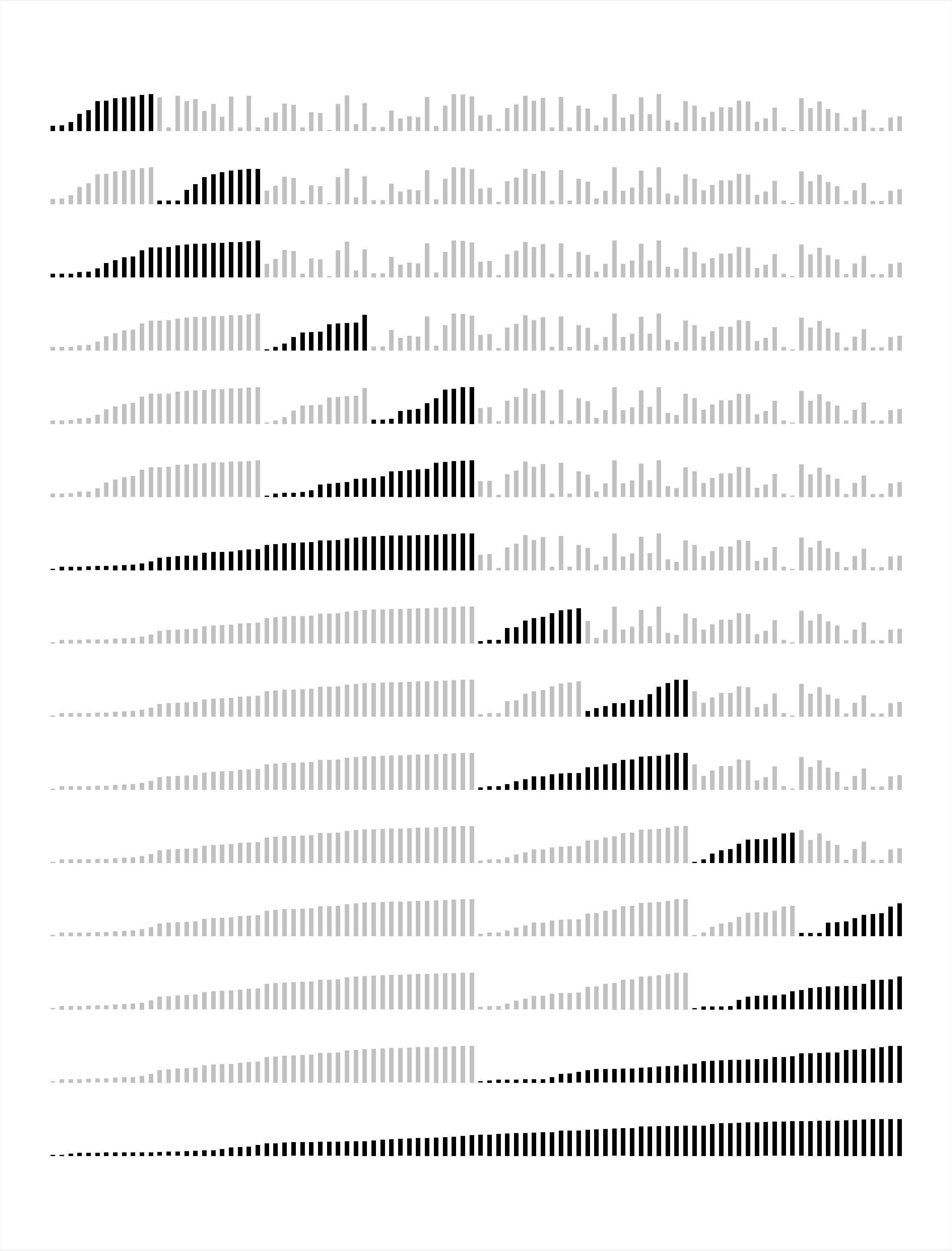
	a[i] 0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
sz = 1	M	E	R	G	E	S	O	R	T	E	X	A	M	P	L	E
`merge(a, aux, 0, 0, 1)`	E	M	R	G	E	S	O	R	T	E	X	A	M	P	L	E
`merge(a, aux, 2, 2, 3)`	E	M	G	R	E	S	O	R	T	E	X	A	M	P	L	E
`merge(a, aux, 4, 4, 5)`	E	M	G	R	E	S	O	R	T	E	X	A	M	P	L	E
`merge(a, aux, 6, 6, 7)`	E	M	G	R	E	S	O	R	T	E	X	A	M	P	L	E
`merge(a, aux, 8, 8, 9)`	E	M	G	R	E	S	O	R	E	T	X	A	M	P	L	E
`merge(a, aux, 10, 10, 11)`	E	M	G	R	E	S	O	R	E	T	A	X	M	P	L	E
`merge(a, aux, 12, 12, 13)`	E	M	G	R	E	S	O	R	E	T	A	X	M	P	L	E
`merge(a, aux, 14, 14, 15)`	E	M	G	R	E	S	O	R	E	T	A	X	M	P	E	L
sz = 2																
`merge(a, aux, 0, 1, 3)`	E	G	M	R	E	S	O	R	E	T	A	X	M	P	E	L
`merge(a, aux, 4, 5, 7)`	E	G	M	R	E	O	R	S	E	T	A	X	M	P	E	L
`merge(a, aux, 8, 9, 11)`	E	G	M	R	E	O	R	S	A	E	T	X	M	P	E	L
`merge(a, aux, 12, 13, 15)`	E	G	M	R	E	O	R	S	A	E	T	X	E	L	M	P
sz = 4																
`merge(a, aux, 0, 3, 7)`	E	E	G	M	O	R	R	S	A	E	T	X	E	L	M	P
`merge(a, aux, 8, 11, 15)`	E	E	G	M	O	R	R	S	A	E	E	L	M	P	T	X
sz = 8																
`merge(a, aux, 0, 7, 15)`	A	E	E	E	E	G	L	M	M	O	P	R	R	S	T	X

Bottom-up mergesort: Java implementation

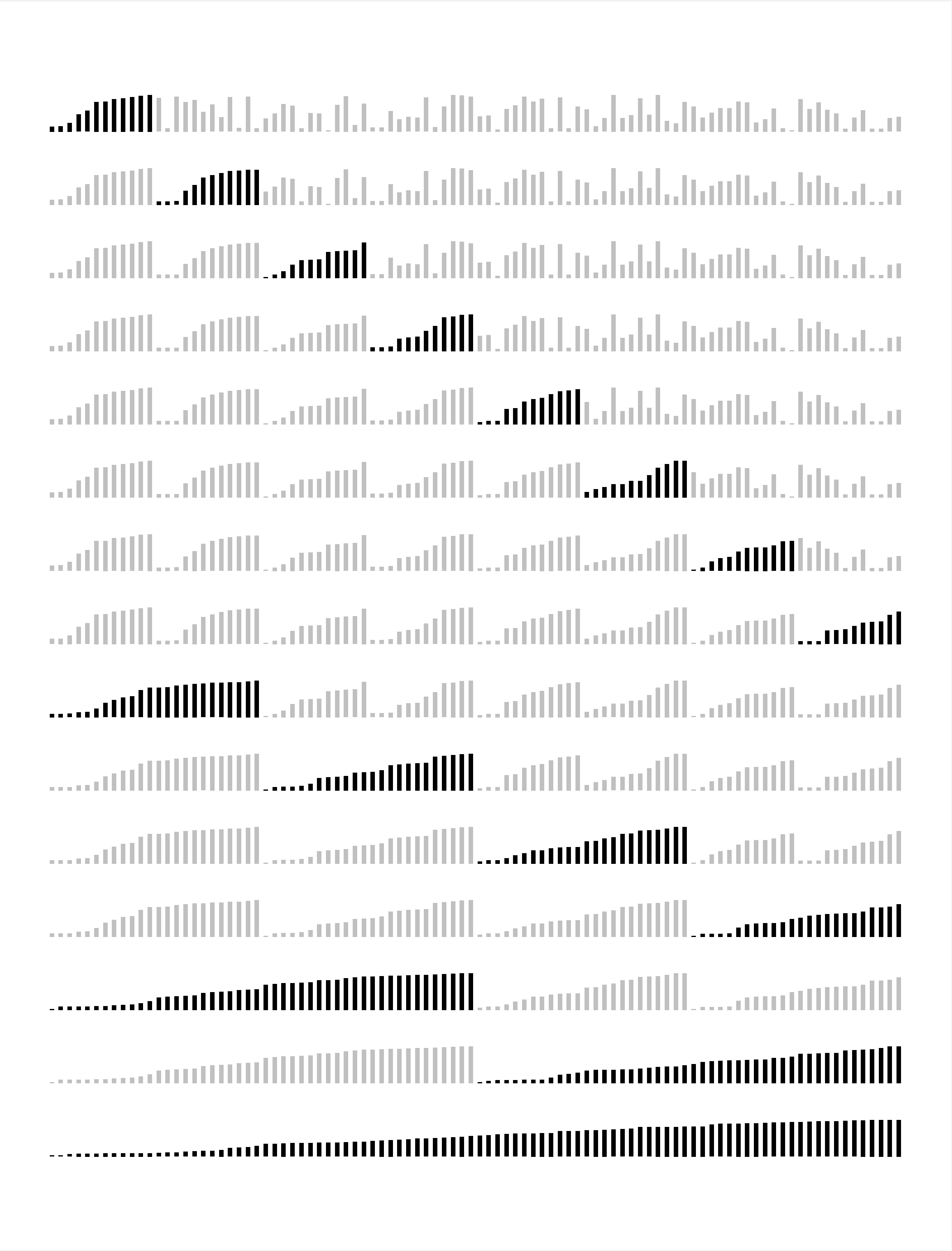
```
public class MergeBU
{
   private static void merge(...)
   {  /* as before */  }

   public static void sort(Comparable[] a)
   {
      int n = a.length;
      Comparable[] aux = new Comparable[n];
      for (int sz = 1; sz < n; sz = sz+sz)
         for (int lo = 0; lo < n-sz; lo += sz+sz)
            merge(a, aux, lo, lo+sz-1, Math.min(lo+sz+sz-1, n-1));
   }
}
```

Bottom line. Simple and non-recursive version of mergesort.

Mergesort: visualizations



top-down mergesort (cutoff = 12)



bottom-up mergesort (cutoff = 12)

Mergesort quiz 4

Which is faster in practice for $n = 2^{20}$, top–down mergesort or bottom–up mergesort?

A. Top-down (recursive) mergesort.

B. Bottom-up (non-recursive) mergesort.

C. No difference.

D. *I don't know.*

Natural mergesort

Idea. Exploit pre-existing order by identifying naturally occurring runs.

input

1	5	10	16	3	4	23	9	13	2	7	8	12	14

first run

1	5	10	16	3	4	23	9	13	2	7	8	12	14

second run

1	5	10	16	3	4	23	9	13	2	7	8	12	14

merge two runs

1	3	4	5	10	16	23	9	13	2	7	8	12	14

Tradeoff. Fewer passes vs. extra compares per pass to identify runs.

Timsort

- Natural mergesort.
- Use binary insertion sort to make initial runs (if needed).
- A few more clever optimizations.

```
Intro
-----
This describes an adaptive, stable, natural mergesort, modestly called
timsort (hey, I earned it <wink>).  It has supernatural performance on many
kinds of partially ordered arrays (less than lg(n!) comparisons needed, and
as few as n-1), yet as fast as Python's previous highly tuned samplesort
hybrid on random arrays.

In a nutshell, the main routine marches over the array once, left to right,
alternately identifying the next run, then merging it into the previous
runs "intelligently".  Everything else is complication for speed, and some
hard-won measure of memory efficiency.
...
```

Tim Peters

Consequence. Linear time on many arrays with pre-existing order.

Now widely used. Python, Java 7–11, GNU Octave, Android, ….

http://hg.openjdk.java.net/jdk7/jdk7/jdk/file/tip/src/share/classes/java/util/Arrays.java

Timsort bug (February 2015)

Envisage: Engineering Virtualized Services

Envisage About Envisage Follow Envisage Dissemination Log in

Proving that Android's, Java's and Python's sorting algorithm is broken (and showing how to fix it)

February 24, 2015 Envisage Written by Stijn de Gouw. $s

Tim Peters developed the Timsort hybrid sorting algorithm in 2002. It is a clever combination of ideas from merge sort and insertion sort, and designed to perform well on real world data. TimSort was first developed for Python, but later ported to Java (where it appears as java.util.Collections.sort and java.util.Arrays.sort) by Joshua Bloch (the designer of Java Collections who also pointed out that most binary search algorithms were broken). TimSort is today used as the default sorting algorithm for Android SDK, Sun's JDK and OpenJDK. Given the popularity of these platforms this means that the number of computers, cloud services and mobile phones that use TimSort for sorting is well into the billions.

http://envisage-project.eu/proving-android-java-and-python-sorting-algorithm-is-broken-and-how-to-fix-it

Timsort bug (May 2018)

JDK / JDK-8203864

Execution error in Java's Timsort

Details

Type:	Bug
Status:	RESOLVED
Priority:	3 P3
Resolution:	Fixed
Affects Version/s:	None
Fix Version/s:	11
Component/s:	core-libs
Labels:	None
Subcomponent:	java.util:collections
Introduced In Version:	6
Resolved In Build:	b20

Description

Carine Pivoteau wrote:
While working on a proper complexity analysis of the algorithm, we realised that there was an error in the last paper reporting such a bug (http://envisage-project.eu/wp-content/uploads/2015/02/sorting.pdf). This implies that the correction implemented in the Java source code (changing Timsort stack size) is wrong and that it is still possible to make it break. This is explained in full details in our analysis: https://arxiv.org/pdf/1805.08612.pdf.
We understand that coming upon data that actually causes this error is very unlikely, but we thought you'd still like to know and do something about it.
As the authors of the previous article advocated for, we strongly believe that you should consider modifying the algorithm as explained in their article (and as was done in Python) rather than trying to fix the stack size.

https://bugs.openjdk.java.net/browse/JDK-8203864

Sorting summary

	inplace?	stable?	best	average	worst	remarks
selection	✔		½ n^2	½ n^2	½ n^2	n exchanges
insertion	✔	✔	n	¼ n^2	½ n^2	use for small n or partially ordered
shell	✔		$n \log_3 n$	?	$c\, n^{3/2}$	tight code; subquadratic
merge		✔	½ $n \lg n$	$n \lg n$	$n \lg n$	$n \log n$ guarantee; stable
timsort		✔	n	$n \lg n$	$n \lg n$	improves mergesort when pre-existing order
?	✔	✔	n	$n \lg n$	$n \lg n$	holy sorting grail

n lg ϱ, where ϱ = # runs
(proved in August 2018)

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2.2 Mergesort

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Complexity of sorting

Computational complexity. Framework to study efficiency of algorithms for solving a particular problem X.

Model of computation. Allowable operations.
Cost model. Operation counts.
Upper bound. Cost guarantee provided by some algorithm for X.
Lower bound. Proven limit on cost guarantee of all algorithms for X.
Optimal algorithm. Algorithm with best possible cost guarantee for X.

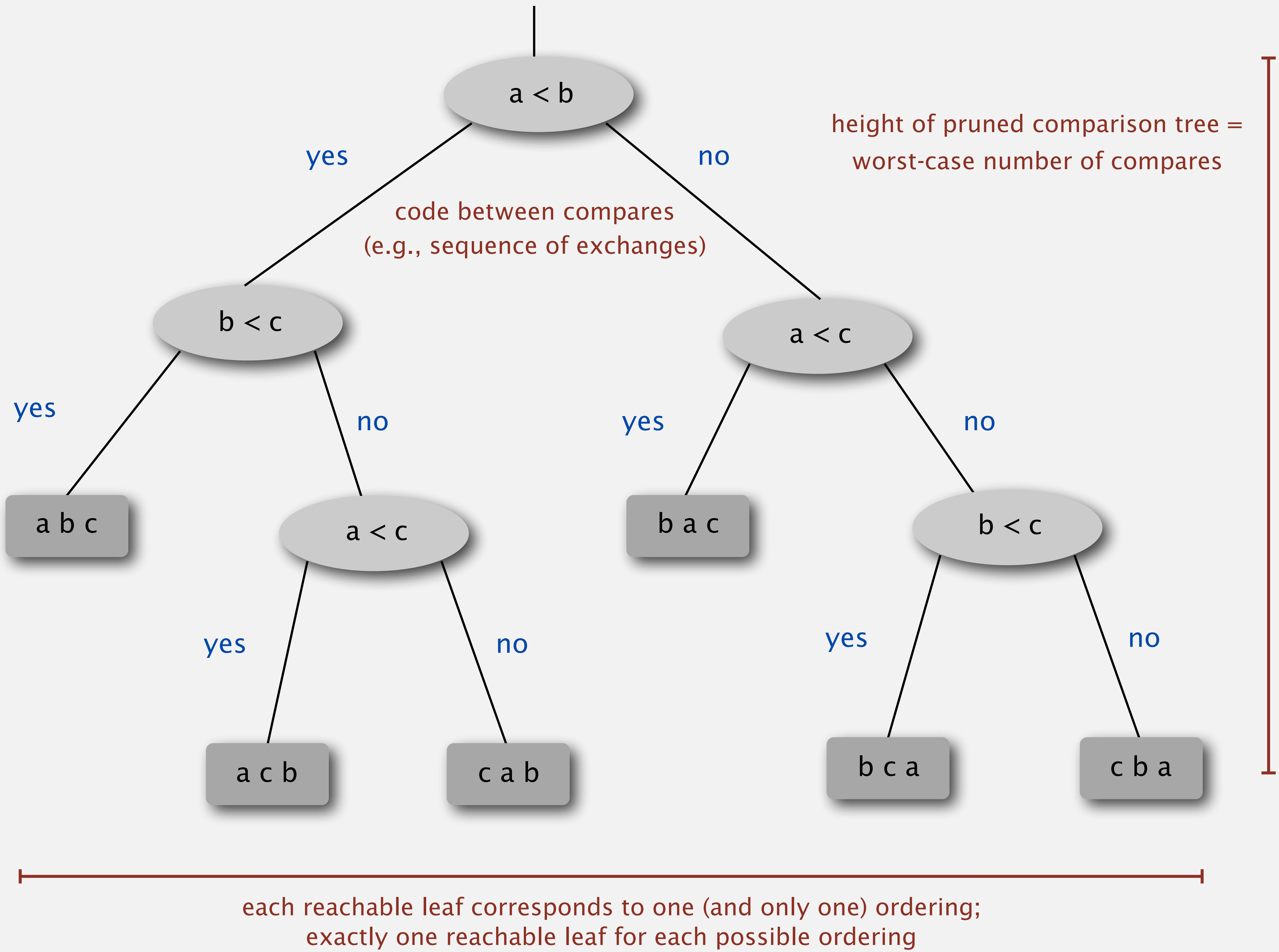
lower bound ~ upper bound

model of computation	*comparison tree*
cost model	*# compares*
upper bound	$\sim n \lg n$ *from mergesort*
lower bound	?
optimal algorithm	?

can access information only through compares (e.g., Java `Comparable` framework)

complexity of sorting

Comparison tree (for 3 distinct keys a, b, and c)

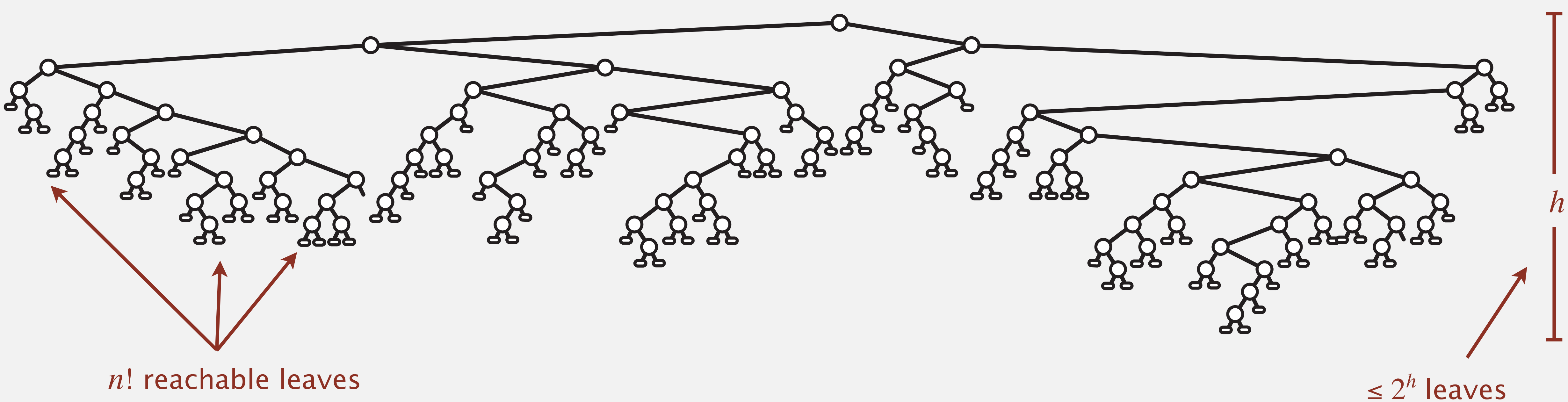


Compare-based lower bound for sorting

Proposition. Any compare-based sorting algorithm must make at least $\lg(n!) \sim n \lg n$ compares in the worst case.

Pf.

- Assume array consists of n distinct values a_1 through a_n.
- Worst-case number of compares = height h of pruned comparison tree.
- Binary tree of height h has $\leq 2^h$ leaves.
- $n!$ different orderings $\Rightarrow$ $n!$ reachable leaves.



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- Worst-case number of compares = height h of pruned comparison tree.
- Binary tree of height h has $\leq 2^h$ leaves.
- $n!$ different orderings $\Rightarrow$ $n!$ reachable leaves.

$$
\begin{aligned}
2^h &\geq \#\text{ reachable leaves } = n! \\
\Rightarrow \quad h &\geq \lg(n!) \\
&\sim n \lg n
\end{aligned}
$$

Stirling's formula

Complexity of sorting

Model of computation. Allowable operations.

Cost model. Operation count(s).

Upper bound. Cost guarantee provided by some algorithm for X.

Lower bound. Proven limit on cost guarantee of all algorithms for X.

Optimal algorithm. Algorithm with best possible cost guarantee for X.

model of computation	*comparison tree*
cost model	*# compares*
upper bound	$\sim n \lg n$
lower bound	$\sim n \lg n$
optimal algorithm	*mergesort*

complexity of sorting

First goal of algorithm design: optimal algorithms.

Complexity results in context

Compares? Mergesort is optimal with respect to number compares.

Space? Mergesort is not optimal with respect to space usage.

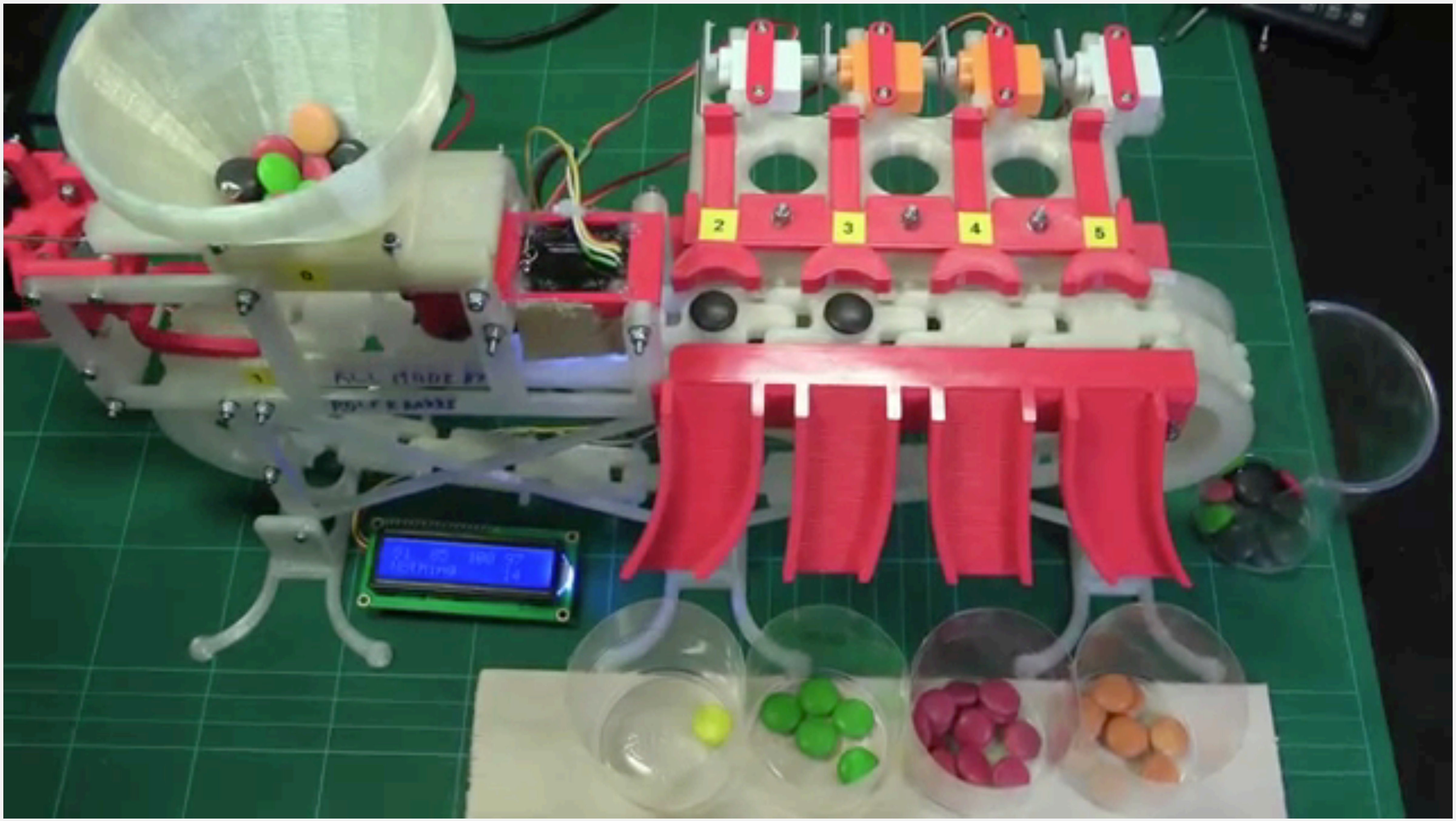
Lessons. Use theory as a guide.

Ex. Design sorting algorithm that guarantees $\sim \frac{1}{2}\, n \lg n$ compares?

Ex. Design sorting algorithm that is both time- and space-optimal?

Commercial break

Q. Why doesn't this Skittles sorter violate the sorting lower bound?



https://www.youtube.com/watch?v=tSEHDBSynVo

Complexity results in context (continued)

Lower bound may not hold if the algorithm can take advantage of:

- The initial order of the input array.
  Ex: insertion sort requires only a linear number of compares on partially sorted arrays.

- The distribution of key values.
  Ex: 3-way quicksort requires only a linear number of compares on arrays with a constant number of distinct keys. [stay tuned]

- The representation of the keys.
  Ex: radix sorts do not make any key compares;
  they access the data via character/digit compares. [stay tuned]

Big O notation (and cousins)

notation	provides	example	shorthand for
Tilde	leading term	$\sim \frac{1}{2} n^2$	$\frac{1}{2} n^2$ $\frac{1}{2} n^2 + 22 n \log n + 3 n$
Big Theta	order of growth	$\Theta(n^2)$	$\frac{1}{2} n^2$ $10 n^2$ $5 n^2 + 22 n \log n + 3 n$
Big O	upper bound	$O(n^2)$	$10 n^2$ $100 n$ $22 n \log n + 3 n$
Big Omega	lower bound	$\Omega(n^2)$	$\frac{1}{2} n^2$ n^5 $n^3 + 22 n \log n + 3 n$

Sorting lower bound

Interviewer. Give a formal description of the sorting lower bound for sorting an array of n elements.

https://algs4.cs.princeton.edu

2.2 Mergesort

- *mergesort*
- *bottom-up mergesort*
- *sorting complexity*
- *divide-and-conquer*

Sorting a linked list

Problem. Given a singly linked list, rearrange its nodes in sorter order.

Version 1. Linearithmic time, linear extra space.

Version 2. Linearithmic time, logarithmic (or constant) extra space.