Robert Sedgewick \| Kevin Wayne
https://algs4.cs.princeton.edu

### 2.1 Elementary Sorts

- rules of the game
- selection sort
- insertion sort
- shuffling
- comparators
- stability


### 2.1 Elementary Sorts

- rules of the game
- selection sort


## Algorithms

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## Sorting problem

## Ex. Student records in a university.

|  | Chen | 3 | A | $(991) 878-4944$ | 308 Blair |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | Rohde | 2 | A | $(232) 343-5555$ | 343 Forbes |
|  | Gazsi | 4 | B | $(800) 867-5309$ | 101 Brown |
|  | Furia | 1 | A | $(766) 093-9873$ | 101 Brown |
|  | Kanaga | 3 | B | $(898) 122-9643$ | 22 Brown |
|  | Andrews | 3 | A | $(664) 480-0023$ | 097 Little |

Sort. Rearrange array of $n$ items in ascending order by key.

| Andrews | 3 | A | $(664) 480-0023$ | 097 Little |
| :---: | :---: | :---: | :---: | :---: |
| Battle | 4 | C | $(874) 088-1212$ | 121 Whitman |
| Chen | 3 | A | $(991) 878-4944$ | 308 Blair |
| Furia | 1 | A | $(766) 093-9873$ | 101 Brown |
| Gazsi | 4 | B | $(800) 867-5309$ | 101 Brown |
| Kanaga | 3 | B | $(898) 122-9643$ | 22 Brown |
| Rohde | 2 | A | $(232) 343-5555$ | 343 Forbes |

## Total order

Sorting is a well-defined problem if and only if there is a total order.

## A total order is a binary relation $\leq$ that satisfies:

- Totality: either $v \leq w$ or $w \leq v$ or both.
- Transitivity: if both $v \leq w$ and $w \leq x$, then $v \leq x$.
- Antisymmetry: if both $v \leq w$ and $w \leq v$, then $v=w$.

Examples.

| Video name | Views* |
| :--- | ---: |
| "Despacito"[6] | $2,993,700,000$ |
| "See You Again"[11] | $2,894,000,000$ |
| "Gangnam Style"[17] | $803,700,000$ |
| "Baby"[41] | $245,400,000$ |
| "Bad Romance"[146] | $178,400,000$ |
| "Charlie Bit My Finger"[136] | $128,900,000$ |
| "Evolution of Dance"[131] | $118,900,000$ |

numerical order (descending)


Ally Kazmucha
Amanda
Amanda Jozaitis
Amanda VanVoorhis
Amy Bruemmer
Amy M
Amy Riehle
Andrew Wray
Andy Hynek
Anil Kumar

## Total order

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- Antisymmetry: if both $v \leq w$ and $w \leq v$, then $v=w$.

Non-examples.

course prerequisites (violates totality)


Ro-sham-bo order (violates transitivity)

predator-prey
(violates antisymmetry)

## Sample sort clients

Goal. Single function that sorts any type of data (that has a total order). Ex 1. Sort strings in alphabetical order.

```
public class StringSorter
{
    public static void main(String[] args)
    {
        String[] a = StdIn.readAl1Strings();
        Insertion.sort(a);
        for (int i = 0; i < a.length; i++)
            StdOut.println(a[i]);
    }
}
% more words3.txt
bed bug dad yet zoo ... al1 bad yes
% java StringSorter < words3.txt
a11 bad bed bug dad ... yes yet zoo
[suppressing new7ines]
```


## Sample sort clients

Goal. Single function that sorts any type of data (that has a total order).
Ex 2. Sort random real numbers in ascending order.
seems artificial (stay tuned for an application)

```
public class Experiment
{
    public static void main(String[] args)
    {
        int n = Integer.parseInt(args[0]);
            Doub7e[] a = new Double[n];
            for (int i = 0; i < n; i++)
            a[i] = StdRandom.uniform();
            Insertion.sort(a);
            for (int i = 0; i < n; i++)
            StdOut.println(a[i]);
    }
}
```

\% java Experiment 10
0.08614716385210452
0.09054270895414829
0.10708746304898642
0.21166190071646818
0.363292849257276
0.460954145685913
0.5340026311350087
0.7216129793703496
0.9003500354411443
0.9293994908845686

## Sample sort clients

Goal. Single function that sorts any type of data (that has a total order).
Ex 3. Sort the files in a given directory by filename.

```
import java.io.File;
public class FileSorter
{
    public static void main(String[] args)
    {
        File directory = new File(args[0]);
        File[] files = directory.listFiles();
        Insertion.sort(files);
        for (int i = 0; i < files.length; i++)
            StdOut.println(files[i].getName());
    }
}
```

```
% java FileSorter .
Insertion.class
Insertion.java
InsertionX.class
InsertionX.java
Selection.class
Selection.java
Shel1.class
She11.java
Shel1X.class
She11X.java
```


## How can a single function sort any type of data？

Goal．Single function that sorts any type of data（that has a total order）．

Solution．Callback＝reference to executable code．

```
Please sort these Japanese names for me:
あゆみ,アユミ,Ayumi, 歩美,
```

> But I don't speak Japanese and I don't know how words are ordered.


## Callbacks

Goal. Single function that sorts any type of data (that has a total order).

Solution. Callback = reference to executable code.

- Client passes array of objects to sort() function.
- The sort() function calls object's compareTo() function as needed.

Implementing callbacks.

- Java: interfaces.
- C: function pointers.
- C++: class-type functors.
- C\#: delegates.
- Python, Perl, ML, Javascript: first-class functions.


## Java interfaces

Interface. A type that defines a set of methods that a class can provide.

```
public interface Comparable<Item>
{
    public int compareTo(Item that);
\(\square\) contract: one method with this signature public int compareTo(Item that); \(\longleftarrow \quad\) (and prescribed behavior)
}
```

Class that implements interface. Must implement all interface methods.


Impact.

- You can treat any String object as an object of type Comparab7e.
- On a Comparable object, you can invoke (only) the compareTo() method.
- Enables callbacks.


## Callbacks in Java: roadmap

client (StringSorter.java)

```
        public class StringSorter
{
    public static void main(String[] args)
    {
        String[] a = StdIn.readA11Strings();
        Insertion.sort(a);
        for (int i = 0; i < a.length; i++)
            StdOut.println(a[i]);
    }
}
```

sort implementation (Insertion.java)
public static void sort(Comparable[] a)
\{
int $\mathrm{n}=\mathrm{a}$. ${ }^{\text {length; }}$
for (int $\mathbf{i}=0 ; \mathbf{i}<n ; i++$ )
for (int $\mathrm{j}=\mathrm{i} ; \mathrm{j}>0$; $\mathrm{j}-$-)
callback
if (a[j].compareTo $(a[j-1])<0$ )
exch(a, j, j-1);
else break;
\}
java.lang.Comparable interface
public interface Comparable<Item> \{ public int compareTo(Item that); \}
data type implementation (String.java)
public class String
implements Comparable<String>
\{
public int compareTo(String that)
\{
\}
\}
key point: client code does not

Elementary sorts: quiz 1

Suppose that the Java architects left out implements Comparable<String> in the class declaration for String. Which would be the effect?
A. String.java won't compile.
B. StringSorter.java won't compile.
C. Insertion.java won't compile.
D. Insertion.java will throw an exception.

## Comparable API

Implement compareTo() so that v.compareTo(w)

- Returns a
- negative integer if $v$ is less than $w$
- positive integer if $v$ is greater than $w$
- zero if $v$ is equal to $w$
v.compareTo (w) <= 0
means $v$ is less than or equal to $w$
- Defines a total order.
- Throws an exception if incompatible types (or either is null).

less than
(return negative integer)

equal to (return 0)

greater than
(return positive integer)

Built-in comparable types. Integer, Double, String, Date, File, ...
User-defined comparable types. Implement the Comparable interface.

## Implementing the Comparable interface

Date data type. Simplified version of java.util.Date.

```
public class Date implements Comparable<Date>
{
    private final int month, day, year;
    public Date(int m, int d, int y)
    {
        month = m;
        day = d;
        year = y;
    }
    public int compareTo(Date that)
    {
        if (this.year < that.year ) return -1;
        if (this.year > that.year ) return +1;
        if (this.month < that.month) return -1;
        if (this.month > that.month) return +1;
        if (this.day < that.day ) return -1;
        if (this.day > that.day ) return +1;
        return 0;
    }
}
```


### 2.1 Elementary Sorts

## - rules of the game

- selection sort


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## Selection sort demo

- In iteration i, find index min of smallest remaining entry.
- Swap a[i] and a[min].
initial


## Selection sort

Algorithm. $\uparrow$ scans from left to right.

Invariants.

- Entries the left of $\uparrow$ (including $\uparrow$ ) fixed and in ascending order.
- No entry to right of $\uparrow$ is smaller than any entry to the left of $\uparrow$.



## Selection sort inner loop

To maintain algorithm invariants:

- Move the pointer to the right.

$$
i++
$$



- Identify index of minimum entry on right.

```
int min = i;
for (int j = i+1; j < n; j++)
    if (less(a[j], a[min]))
        min = j;
```



- Exchange into position.

```
exch(a, i, min);
```



## Two useful sorting abstractions

Helper functions. Refer to data only through compares and exchanges.

Less. Is item v less than w?

```
private static boolean less(Comparable v, Comparable w)
{ return v.compareTo(w) < 0; }
```

Exchange. Swap item in array $a[]$ at index $i$ with the one at index $j$.

```
private static void exch(Object[] a, int i, int j)
{
    Object swap = a[i];
    a[i] = a[j];
    a[j] = swap;
}
```


## Selection sort: Java implementation

```
public class Selection
{
    public static void sort(Comparable[] a)
    {
        int n = a.length;
        for (int i = 0; i < n; i++)
        {
            int min = i;
            for (int j = i+1; j < n; j++)
                if (less(a[j], a[min]))
                        min = j;
        exch(a, i, min);
        }
    }
    private static boolean less(Comparable v, Comparable w)
    { /* see previous slide */ }
    private static void exch(Object[] a, int i, int j)
    { /* see previous slide */ }
}
```


## Generic methods

Oops. The compiler complains.

```
% javac-algs4 Selection.java
Selection.java:83: warning: [unchecked] unchecked cal1 to
compareTo(T) as a member of the raw type java.lang.Comparable
    return (v.compareTo(w) < 0);
    ^
```

1 warning
Q. How to appease the compiler?

## Generic methods

Pedantic (type-safe) version. Compiles without any warnings.

and Assignment 3

Remark. Use type-safe version in system code (but not in lecture).

## Selection sort: animations

20 random items


A algorithm position
in final order
not in final order
http:/ /www.sorting-algorithms.com/ selection-sort

Elementary sorts: quiz 2

How many compares does selection sort make to sort an array of $n$ distinct items in reverse order?
A. $\sim n$
B. $\sim 1 / 4 n^{2}$
C. $\sim 1 / 2 n^{2}$
D. $\sim n^{2}$

## Selection sort: mathematical analysis

Proposition. Selection sort makes $(n-1)+(n-2)+\ldots+1+0 \sim n^{2} / 2$ compares and $n$ exchanges to sort any array of $n$ items.


Running time insensitive to input. Quadratic time, even if input is sorted.
Data movement is minimal. Linear number of exchanges-exactly $n$.

### 2.1 Elementary Sorts

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Insertion sort demo

- In iteration i, swap a[i] with each larger entry to its left.



## Insertion sort

Algorithm. $\uparrow$ scans from left to right.

Invariants.

- Entries to the left of $\uparrow$ (including $\uparrow$ ) are in ascending order.
- Entries to the right of $\uparrow$ have not yet been seen.



## Insertion sort: inner loop

To maintain algorithm invariants:

- Move the pointer to the right.

```
i++;
```



- Moving from right to left, exchange a[i] with each larger entry to its left.

```
for (int j = i; j > 0; j--)
    if (less(a[j], a[j-1]))
        exch(a, j, j-1);
    else break;
```



```
public class Insertion
{
    public static void sort(Comparable[] a)
    {
        int n = a.length;
        for (int i = 0; i < n; i++)
            for (int j = i; j > 0; j--)
            if (less(a[j], a[j-1]))
                exch(a, j, j-1);
            else break;
    }
    private static boolean less(Comparable v, Comparable w)
    { /* as before */ }
    private static void exch(Object[] a, int i, int j)
    { /* as before */ }
}
```


## Elementary sorts: quiz 3

How many compares does insertion sort make to sort an array of $\boldsymbol{n}$ distinct keys in reverse order?
A. $\sim n$
B. $\sim 1 / 4 n^{2}$
C. $\sim 1 / 2 n^{2}$
D. $\sim n^{2}$

## Insertion sort: analysis

Worst case. Insertion sort makes $\sim 1 / 2 n^{2}$ compares and $\sim 1 / 2 n^{2}$ exchanges to sort an array of $n$ distinct keys in reverse order.

Pf. Exactly $i$ compares and exchanges in iteration $i$.

$$
0+1+2+\ldots+(n-1)
$$



A algorithm position
in order
not yet seen

Elementary sorts: quiz 4

## Which is faster in practice to sort an array of $\boldsymbol{n}$ items, selection sort or insertion sort?

A. Selection sort.
B. Insertion sort.
C. No significant difference.
D. It depends.

## Insertion sort: analysis

Best case. Insertion sort makes $n-1$ compares and 0 exchanges to sort an array of $n$ distinct keys in ascending order.


A algorithm position
in order
not yet seen

## Insertion sort: analysis

Good case. Insertion sort is linear time for "partially sorted" arrays.
Q. What do we mean by partially sorted?


A algorithm position in order not yet seen

## Insertion sort: partially sorted arrays

Def. An inversion is a pair of keys that are out of order.
A E E L M O T R X P S

| 1 | 2 | 3 | 4 | 5 | 6 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| T-R | T-P | T-S | R-P | X-P | X-S |
|  | ( $\mathbf{l}$ inversions) |  |  |  |  |

Def. A family of arrays is partially sorted if the number of inversions is $\leq c n$.

- Ex 1. A sorted array.
- Ex 2. A subarray of length 10 appended to a sorted subarray of length $n$.
$\leq 10 n+45$ inversions

Proposition. Insertion sort runs in linear time on partially sorted arrays.
Pf. exchange decreases number of inversions by 1

- Number of exchanges in insertion sort = number of inversions.
- Number of compares $\leq$ number of exchanges $+(n-1)$.


## Insertion sort: practical improvements

Half exchanges. Shift items over (instead of exchanging).

- Eliminates unnecessary data movement.
- No longer uses only less() and exch() to access data.
ACHHIMNNPQXYKBINARY

Binary insertion sort. Use binary search to find insertion point.

- Number of compares $\sim n \log _{2} n$.
- But still a quadratic number of array accesses.
ACHHIMNNPQXYKBINARY
binary search for first key > K


### 2.1 Elementary Sorts

- rules of the game
- selection sort


## Algorithms

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## INTERVIEW QUESTION: SHUFFLE AN ARRAY

Goal. Rearrange array so that result is a uniformly random permutation.
all $n$ ! permutations equally likely

## Shuffling by sorting

- Generate a random real number for each array entry.
- Sort the array.



## Shuffling by sorting

- Generate a random real number for each array entry.
- Sort the array.


Proposition. Shuffle sort produces a uniformly random permutation.

Application. Shuffle columns in a spreadsheet.

## INTERVIEW QUESTION: SHUFFLE AN ARRAY

Goal. Rearrange array so that result is a uniformly random permutation.
all $n$ ! permutations
equally likely


Shuffling by sorting.

- Quadratic time (with insertion sort or selection sort).
- Linearithmic time (with mergesort).

Challenge. Design a linear-time algorithm (without sorting).

## Shuffling hall of shame (Microsofi)

Microsoft antitrust probe by EU. Microsoft agreed to provide a randomized ballot screen for users to select browser.

## Select your web browser(s)

A fast new browser from Google. Try it now!

## Safari <br> 

Safari for Windows from Apple, the world's most innovative browser.
s

Your online security is Firefox's top priority. Firefox is free, and made to help you get the most out of the
mozilla firefox
ur online is

The fastest browser on Earth. Secure, powerful and easy to use, with excellent privacy protection.

Windows
Internet Explorer 8
Designed to help you take control of your privacy and browse with confidence. Free from Microsoft.

appeared last
$50 \%$ of the time

## SHuFFING HALL OF SHAME (PLANETPOKER.COM)

Texas hold'em poker. Software must shuffle electronic cards.


How We Learned to Cheat at Online Poker: A Study in Software Security

## Shuffling hall of shame (Princeton)

Ivy league school room draw. Students assigned random room draw times.
U. claims error in room draw process, provides compensation to affected students

By Rebecca Han | Apr 26, 2019


Photo Credit: Jon Ort / The Daily Princetonian

### 2.1 Elementary Sorts

- rules of the game
- selection sort


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## Different orderings

Q. When might we need to define different sort orderings?

| $0.00 \mathrm{~K} / \mathrm{s} \mathbf{\Sigma}: 4 \%$ 78\% $\quad 16.54$ |  |
| :---: | :---: |
| Eighth Floor | $\bigcirc$ |
| Eleventh Floor | $\bigcirc$ |
| Fifth Floor | $\bigcirc$ |
| First Floor | $\bigcirc$ |
| Fourteenth Floor | $\bigcirc$ |
| Fourth Floor | $\bigcirc$ |
| Ninth Floor | $\bigcirc$ |
| Second Floor | $\bigcirc$ |
| Seventh Floor | $\bigcirc$ |
| Sixth Floor | $\bigcirc$ |
| Tenth Floor | $\bigcirc$ |
| Third Eloor | $\bigcirc$ |

Sort music library by artist


Sort music library by song name


## Comparable interface: review

Comparab7e interface: sort using a type's natural order.

```
public class Date implements Comparable<Date>
{
    private final int month, day, year;
    public Date(int m, int d, int y)
    {
        month = m;
        day = d;
        year = y;
    }
    public int compareTo(Date that)
    {
        if (this.year < that.year ) return -1;
        if (this.year > that.year ) return +1;
        if (this.month < that.month) return -1;
        if (this.month > that.month) return +1;
        if (this.day < that.day ) return -1;
        if (this.day > that.day ) return +1;
        return 0;
    }
}
```


## Comparator interface

Comparator interface: sort using an alternate order.

```
pub1ic interface Comparator<Item>
{
    public int compare(Item v, Item w);
}
```

Required property. Must be a total order.

| string order | example |
| :---: | :---: |
| natural order | Now is the time |
| case insensitive | is Now the time digraphs ch and II and rr |
| Spanish language | café cafetero cuarto churro nube ñoño |
| British phone book | McKinley Mackintosh |

## Comparator interface: system sort

To use with Java system sort:

- Create Comparator object.
- Pass as second argument to Arrays.sort().


Bottom line. Decouples the definition of the data type from the definition of what it means to compare two objects of that type.

## Comparator interface: implementing

To implement a comparator:

- Define a (nested) class that implements the Comparator interface.
- Implement the compare() method.
- Provide client access to Comparator.

```
import java.util.Comparator;
public class Student
{
    private final String name;
    private final int section;
    -.. one Comparator for the class
    privatestatic clas& NameOrder implements Comparator<Student>
    {
        public int compare(Student v, Student w)
        { return v.mame.compareTo(w.name); }
    }
    pub1ic static Comparator<Student> byNameOrder()
    { return new NameOrder(); }
}

\section*{Comparator interface: implementing}

To implement a comparator:
- Define a (nested) class that implements the Comparator interface.
- Implement the compare() method.
- Provide client access to Comparator.
```

import java.util.Comparator;
public class Student
{
private final String name;
private final int section;
private static class SectionOrder implements Comparator<Student>
{
public int compare(Student v, Student w)
{ return Integer.compare(v.section, w.section); }
}
pub1ic static Comparator<Student> bySectionOrder()
{ return new SectionOrder(); } method
useful library
}

```

\section*{Comparator interface: implementing}

To implement a comparator:
- Define a (nested) class that implements the Comparator interface.
- Implement the compare() method.
- Provide client access to Comparator.

Insertion.sort(a, Student.byNameOrder());
\begin{tabular}{|c|c|c|c|c|}
\hline Andrews & 3 & A & \((664) 480-0023\) & 097 Little \\
\hline Battle & 4 & C & \((874) 088-1212\) & 121 Whitman \\
\hline Chen & 3 & A & \((991) 878-4944\) & 308 Blair \\
\hline Fox & 3 & A & \((884) 232-5341\) & 11 Dickinson \\
\hline Furia & 1 & A & \((766) 093-9873\) & 101 Brown \\
\hline Gazsi & 4 & B & \((800) 867-5309\) & 101 Brown \\
\hline Kanaga & 3 & B & \((898) 122-9643\) & 22 Brown \\
\hline Rohde & 2 & A & \((232) 343-5555\) & 343 Forbes \\
\hline
\end{tabular}

Insertion.sort(a, Student.bySectionOrder());
\begin{tabular}{|c|c|c|c|c|}
\hline Furia & 1 & A & \((766)\) 093-9873 & 101 Brown \\
\hline Rohde & 2 & A & \((232) 343-5555\) & 343 Forbes \\
\hline Andrews & 3 & A & \((664) 480-0023\) & 097 Little \\
\hline Chen & 3 & A & \((991) 878-4944\) & 308 Blair \\
\hline Fox & 3 & A & \((884) 232-5341\) & 11 Dickinson \\
\hline Kanaga & 3 & B & \((898) 122-9643\) & 22 Brown \\
\hline Battle & 4 & C & \((874) 088-1212\) & 121 Whitman \\
\hline Gazsi & 4 & B & \((800) 867-5309\) & 101 Brown \\
\hline
\end{tabular}

\subsection*{2.1 Elementary Sorts}

\section*{- rules of the game}

\section*{Algorithms}

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insertion sort
p shiufling
- comparałors
- stability

\section*{Stability}

A typical application. First, sort by name; then sort by section.

Selection.sort(a, Student.byNameOrder());
\begin{tabular}{|c|c|c|c|c|}
\hline Andrews & 3 & A & \((664) 480-0023\) & 097 Little \\
\hline Battle & 4 & C & \((874) 088-1212\) & 121 Whitman \\
\hline Chen & 3 & A & \((991) 878-4944\) & 308 Blair \\
\hline Fox & 3 & A & \((884) 232-5341\) & 11 Dickinson \\
\hline Furia & 1 & A & \((766) 093-9873\) & 101 Brown \\
\hline Gazsi & 4 & B & \((800) 867-5309\) & 101 Brown \\
\hline Kanaga & 3 & B & \((898) 122-9643\) & 22 Brown \\
\hline Rohde & 2 & A & \((232) 343-5555\) & 343 Forbes \\
\hline
\end{tabular}

Selection.sort(a, Student.bySectionOrder());
\begin{tabular}{|c|c|c|c|c|}
\hline Furia & 1 & A & \((766) 093-9873\) & 101 Brown \\
\hline Rohde & 2 & A & \((232) 343-5555\) & 343 Forbes \\
\hline Chen & 3 & A & \((991) 878-4944\) & 308 Blair \\
\hline Fox & 3 & A & \((884) 232-5341\) & 11 Dickinson \\
\hline Andrews & 3 & A & \((664) 480-0023\) & 097 Little \\
\hline Kanaga & 3 & B & \((898) 122-9643\) & 22 Brown \\
\hline Gazsi & 4 & B & \((800) 867-5309\) & 101 Brown \\
\hline Battle & 4 & C & \((874) 088-1212\) & 121 Whitman \\
\hline
\end{tabular}
@\#\%\&@! Students in section 3 no longer sorted by name.

A stable sort preserves the relative order of items with equal keys.

\section*{Elementary sorts: quiz 5}

Which sorting algorithm(s) are stable?
A. Selection sort.
B. Insertion sort.
C. Both A and B .
D. Neither A nor B.

\section*{Stability: insertion sort}

Proposition. Insertion sort is stable.
```

public class Insertion
{
public static void sort(Comparable[] a)
{
int n = a.length;
for (int i = 0; i < n; i++)
for (int j = i; j > 0 \&\& less(a[j], a[j-1]); j--)
exch(a, j, j-1);
}
}

| i | j | 0 | 1 | 2 | 3 | 4 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 0 | $\mathrm{~B}_{1}$ | $A_{1}$ | $A_{2}$ | $A_{3}$ | $\mathrm{~B}_{2}$ |
| 1 | 0 | $A_{1}$ | $\mathrm{~B}_{1}$ | $A_{2}$ | $A_{3}$ | $\mathrm{~B}_{2}$ |
| 2 | 1 | $A_{1}$ | $A_{2}$ | $\mathrm{~B}_{1}$ | $A_{3}$ | $\mathrm{~B}_{2}$ |
| 3 | 2 | $A_{1}$ | $A_{2}$ | $A_{3}$ | $\mathrm{~B}_{1}$ | $\mathrm{~B}_{2}$ |
| 4 | 4 | $A_{1}$ | $A_{2}$ | $A_{3}$ | $B_{1}$ | $B_{2}$ |
|  |  | $A_{1}$ | $A_{2}$ | $A_{3}$ | $B_{1}$ | $B_{2}$ |

```

Pf. Equal items never move past each other.

\section*{Stability: selection sort}

Proposition. Selection sort is not stable.
```

public class Selection
{
public static void sort(Comparable[] a)
{
int n = a.length;
for (int i = 0; i < n; i++)
{
int min = i;
for (int j = i+1; j < n; j++)
if (less(a[j], a[min]))
min = j;
exch(a, i, min);
}
}
}

```

Pf by counterexample. Long-distance exchange can move an equal item past another one.```

