# Algorithms

 $\checkmark$ 

#### ROBERT SEDGEWICK | KEVIN WAYNE

## 2.1 ELEMENTARY SORTS

rules of the game

selection sort

insertion sort

shuffling

stability

comparators

https://algs4.cs.princeton.edu

Robert Sedgewick | Kevin Wayne

Algorithms

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Ex. Student records in a university.



#### Sort. Rearrange array of *n* items in ascending order by key.

Andrews	3	А	(664) 480–0023	097 Little	
Battle	4	С	(874) 088–1212	121 Whitman	
Chen	3	А	(991) 878–4944	308 Blair	
Furia	1	А	(766) 093–9873	101 Brown	
Gazsi	4	В	(800) 867–5309	101 Brown	
Kanaga	3	В	(898) 122–9643	22 Brown	
Rohde	2	А	(232) 343–5555	343 Forbes	

Sorting is a well-defined problem if and only if there is a total order.

A total order is a binary relation  $\leq$  that satisfies:

- Totality: either  $v \le w$  or  $w \le v$  or both.
- Transitivity: if both  $v \le w$  and  $w \le x$ , then  $v \le x$ .
- Antisymmetry: if both  $v \le w$  and  $w \le v$ , then v = w.

#### Examples.

Video name 🗢	Views* +	
"Despacito" <sup>[6]</sup>	2,993,700,000	
"See You Again" <sup>[11]</sup>	2,894,000,000	
"Gangnam Style" <sup>[17]</sup>	803,700,000	
"Baby" <sup>[41]</sup>	245,400,000	
"Bad Romance" <sup>[146]</sup>	178,400,000	
"Charlie Bit My Finger" <sup>[136]</sup>	128,900,000	
"Evolution of Dance" <sup>[131]</sup>	118,900,000	

numerical order (descending)

Intern	ational Departures			
light No	Destination	Time	Gate	Remarks
CX7183	Berlin	7:50	A-11	Gate closing
QF3474	London	7:50	A-12	Gate closing
3A372	Paris	7:55	B-10	Boarding
AY6554	New York	8:00	C-33	Boarding
(L3160	San Francisco	8:00	F-15	Boarding
BA8903	Manchester	8:05	B-12	Gate lounge open
3A710	Los Angeles	8:10	C-12	Check-in open
QF3371	Hong Kong	8:15	F-10	Check-in open
MA4866	Barcelona	8:15	F-12	Check-in at kiosks
CX7221	Copenhagen	8:20	G-32	Check-in at kiosks

#### chronological order

	❹ 46% ■
Groups All Contacts	+
Q Search	A
All Howers	B
Ally Kazmucha	L E F
Amanda	G H
Amanda Jozaitis	I J K
Amanda VanVoorhis	L
Amy Bruemmer	N O P
Amy M	Q
Amy Riehle	S T U
Andrew Wray	v w
Andy Hynek	X Y
Anil Kumar	z #

#### lexicographic order

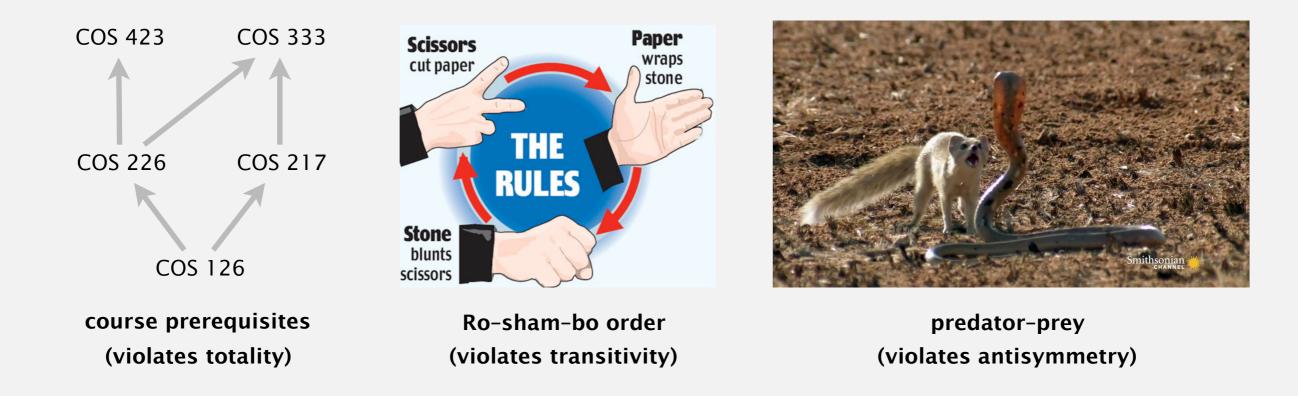
### Total order

Sorting is a well-defined problem if and only if there is a total order.

A total order is a binary relation  $\leq$  that satisfies:

- Totality: either  $v \le w$  or  $w \le v$  or both.
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- Antisymmetry: if both  $v \le w$  and  $w \le v$ , then v = w.

#### Non-examples.



### Sample sort clients

Goal. Single function that sorts any type of data (that has a total order). Ex 1. Sort strings in alphabetical order.

```
public class StringSorter
{
   public static void main(String[] args)
   {
      String[] a = StdIn.readAllStrings();
      Insertion.sort(a);
      for (int i = 0; i < a.length; i++)
         StdOut.println(a[i]);
   }
}
      % more words3.txt
      bed bug dad yet zoo ... all bad yes
      % java StringSorter < words3.txt
      all bad bed bug dad ... yes yet zoo
      [suppressing newlines]
```

Goal. Single function that sorts any type of data (that has a total order).Ex 2. Sort random real numbers in ascending order.

seems artificial (stay tuned for an application)

```
public class Experiment
{
    public static void main(String[] args)
    {
        int n = Integer.parseInt(args[0]);
        Double[] a = new Double[n];
        for (int i = 0; i < n; i++)
            a[i] = StdRandom.uniform();
        Insertion.sort(a);
        for (int i = 0; i < n; i++)
            StdOut.println(a[i]);
    }
}</pre>
```

% java Experiment 10 0.08614716385210452 0.09054270895414829 0.10708746304898642 0.21166190071646818 0.363292849257276 0.460954145685913 0.5340026311350087 0.7216129793703496 0.9003500354411443 0.9293994908845686 Goal. Single function that sorts any type of data (that has a total order).Ex 3. Sort the files in a given directory by filename.

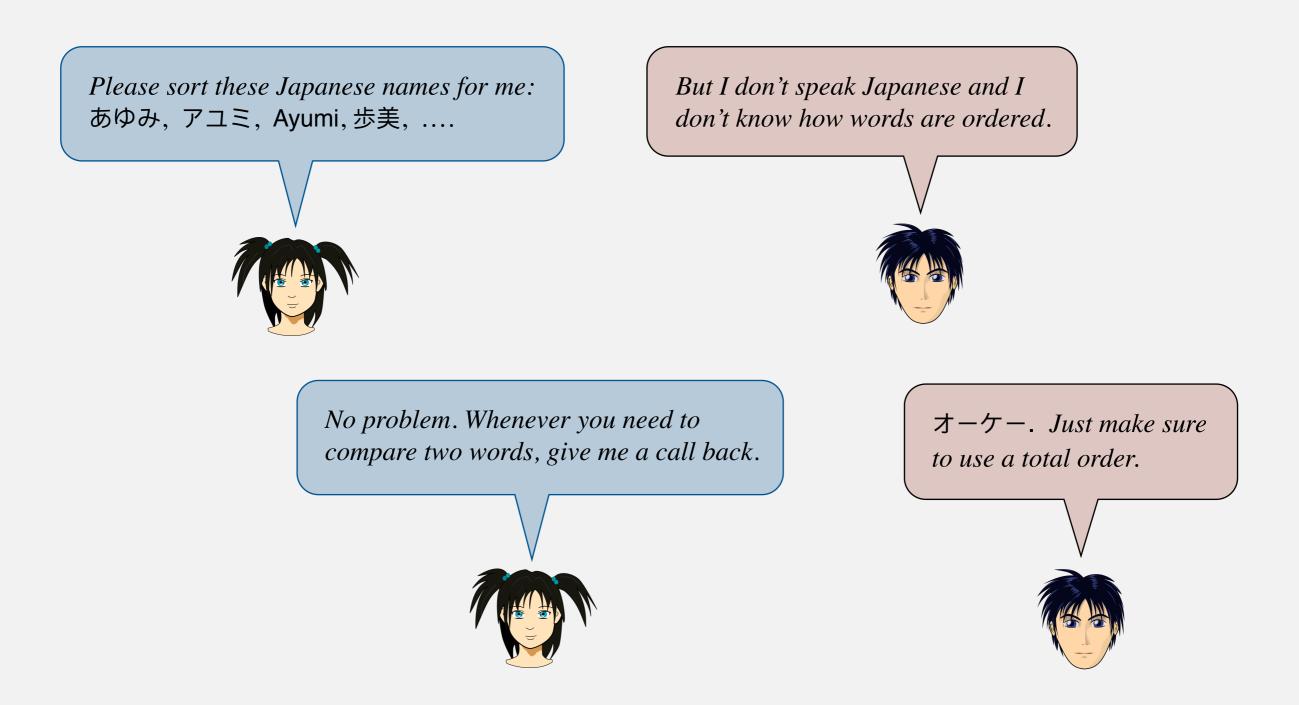
```
import java.io.File;
public class FileSorter
{
    public static void main(String[] args)
    {
        File directory = new File(args[0]);
        File[] files = directory.listFiles();
        Insertion.sort(files);
        for (int i = 0; i < files.length; i++)
            StdOut.println(files[i].getName());
     }
}
```

% java FileSorter .
Insertion.class
Insertion.java
InsertionX.class
InsertionX.java
Selection.class
Selection.java
Shell.class
Shell.java
ShellX.class
ShellX.java

## How can a single function sort any type of data?

Goal. Single function that sorts any type of data (that has a total order).

**Solution.** Callback = reference to executable code.



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**Solution.** Callback = reference to executable code.

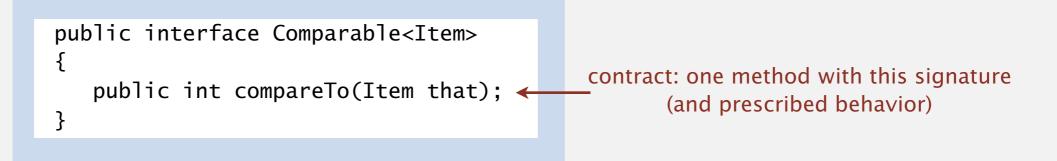
- Client passes array of objects to sort() function.
- The sort() function calls object's compareTo() function as needed.

#### Implementing callbacks.

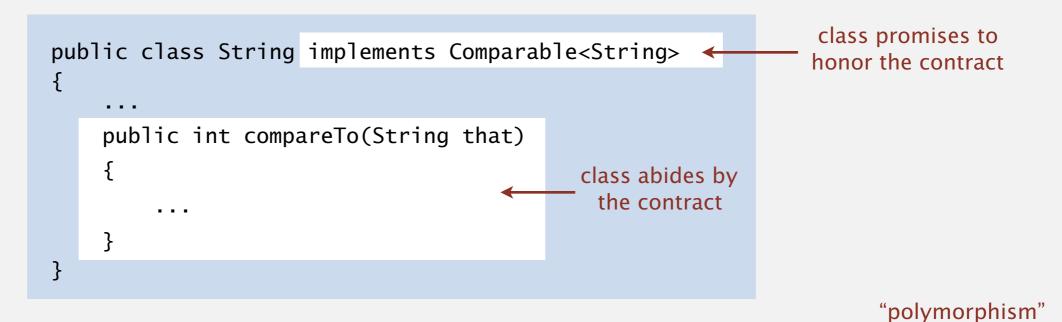
- Java: interfaces.
- C: function pointers.
- C++: class-type functors.
- C#: delegates.
- Python, Perl, ML, Javascript: first-class functions.

#### Java interfaces

Interface. A type that defines a set of methods that a class can provide.



Class that implements interface. Must implement all interface methods.

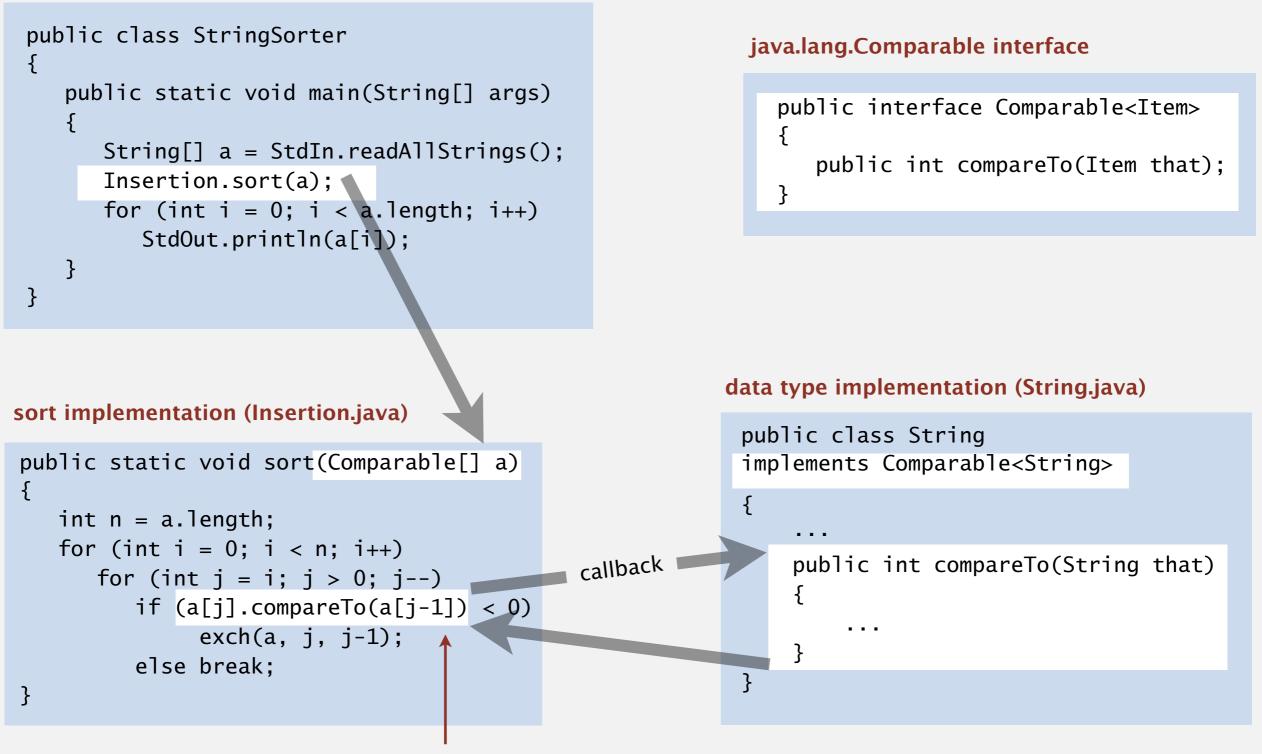


Impact.

- You can treat any String object as an object of type Comparable.
- On a Comparable object, you can invoke (only) the compareTo() method.
- Enables callbacks.

### Callbacks in Java: roadmap

#### client (StringSorter.java)



key point: client code does not depend upon type of data to be sorted



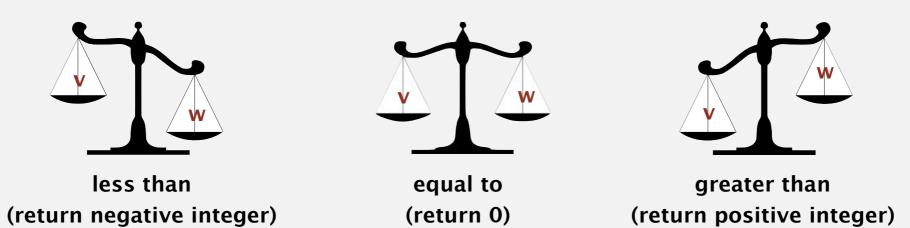
Suppose that the Java architects left out implements Comparable<String> in the class declaration for String. Which would be the effect?

- A. String.java won't compile.
- B. StringSorter.java won't compile.
- **C.** Insertion.java won't compile.
- **D.** Insertion.java will throw an exception.

#### Implement compareTo() so that v.compareTo(w)

- Returns a
  - negative integer if v is less than w
  - positive integer if v is greater than w
  - zero if v is equal to w
- Defines a total order.

- v.compareTo(w) <= 0 means v is less than or equal to w
- Throws an exception if incompatible types (or either is null).



Built-in comparable types. Integer, Double, String, Date, File, ... User-defined comparable types. Implement the Comparable interface.

### Implementing the Comparable interface

Date data type. Simplified version of java.util.Date.

```
public class Date implements Comparable<Date>
{
   private final int month, day, year;
   public Date(int m, int d, int y)
   {
      month = m;
                                                           can compare Date objects
                                                           only to other Date objects
      day = d;
      year = y;
   }
   public int compareTo(Date that)
   {
      if (this.year < that.year ) return -1;
      if (this.year > that.year ) return +1;
      if (this.month < that.month) return -1;
      if (this.month > that.month) return +1;
      if (this.day < that.day ) return -1;
      if (this.day > that.day ) return +1;
      return 0;
   }
}
```

## 2.1 ELEMENTARY SORTS

# rules of the game selection sort

insertion sort

comparators

shuffling

stability

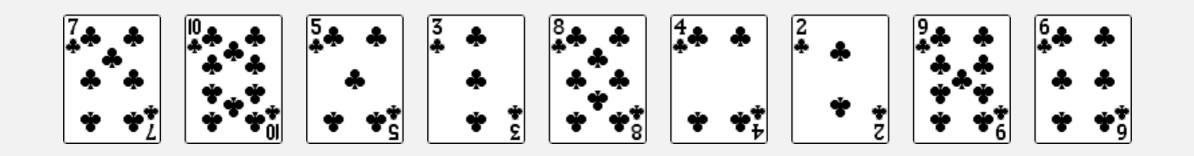
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### Selection sort demo

- In iteration i, find index min of smallest remaining entry.
- Swap a[i] and a[min].







Algorithm. ↑ scans from left to right.

Invariants.

- No entry to right of ↑ is smaller than any entry to the left of ↑.



### Selection sort inner loop

- To maintain algorithm invariants:
  - Move the pointer to the right.



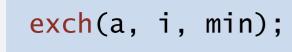
• Identify index of minimum entry on right.

```
int min = i;
for (int j = i+1; j < n; j++)
    if (less(a[j], a[min]))
        min = j;
```



in final order

• Exchange into position.





Helper functions. Refer to data only through compares and exchanges.

Less. Is item v less than w?

```
private static boolean less(Comparable v, Comparable w)
{ return v.compareTo(w) < 0; }</pre>
```

Exchange. Swap item in array a[] at index i with the one at index j.

```
private static void exch(Object[] a, int i, int j)
{
    Object swap = a[i];
    a[i] = a[j];
    a[j] = swap;
}
```

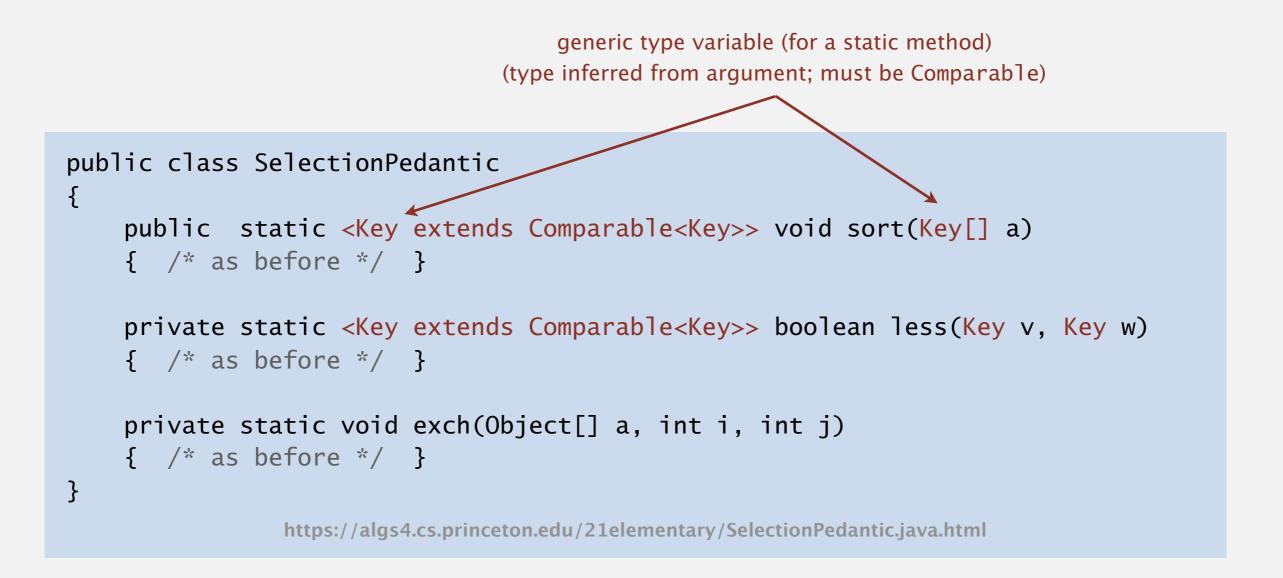
```
public class Selection
Ł
   public static void sort(Comparable[] a)
   {
      int n = a.length;
      for (int i = 0; i < n; i++)
      {
         int min = i;
         for (int j = i+1; j < n; j++)
            if (less(a[j], a[min]))
               min = j;
         exch(a, i, min);
      }
   }
   private static boolean less(Comparable v, Comparable w)
   { /* see previous slide */ }
   private static void exch(Object[] a, int i, int j)
   { /* see previous slide */ }
}
```

https://algs4.cs.princeton.edu/21elementary/Selection.java.html

#### **Oops.** The compiler complains.

**Q.** How to appease the compiler?

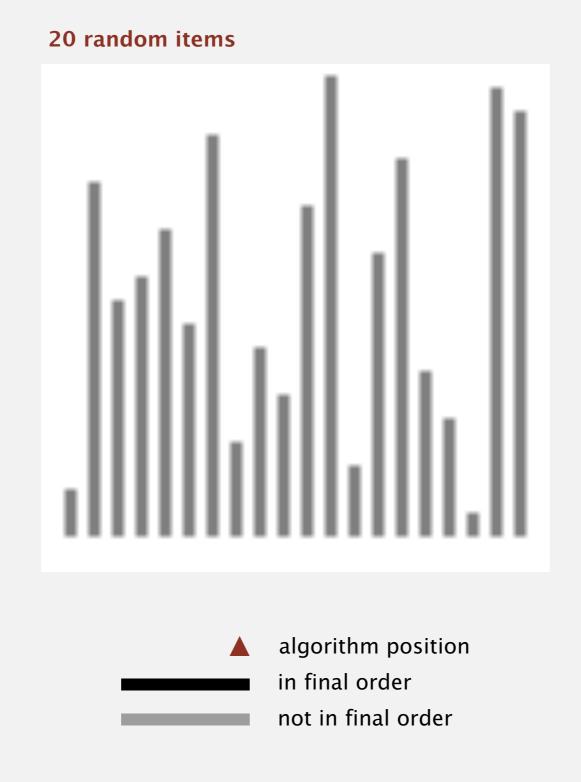
#### Pedantic (type-safe) version. Compiles without any warnings.



and Assignment 3

Remark. Use type-safe version in system code (but not in lecture).

#### Selection sort: animations



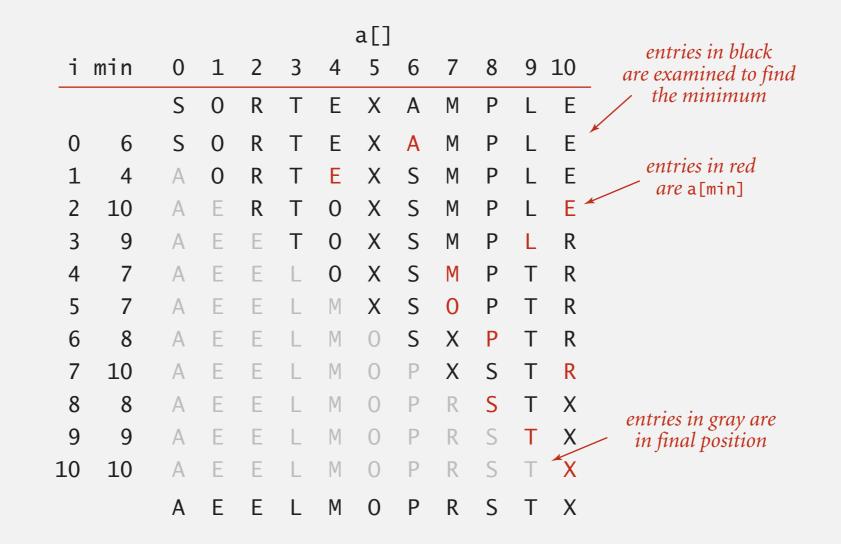
http://www.sorting-algorithms.com/selection-sort



How many compares does selection sort make to sort an array of *n* distinct items in reverse order?

- **A.** ~ n**B.** ~  $1/4 n^2$
- **C.** ~  $1/2 n^2$
- **D.**  $\sim n^2$

**Proposition.** Selection sort makes  $(n-1) + (n-2) + ... + 1 + 0 \sim n^2/2$  compares and *n* exchanges to sort any array of *n* items.



Running time insensitive to input. Quadratic time, even if input is sorted. Data movement is minimal. Linear number of exchanges—exactly *n*.

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# rules of the game selection sort \_\_\_\_\_

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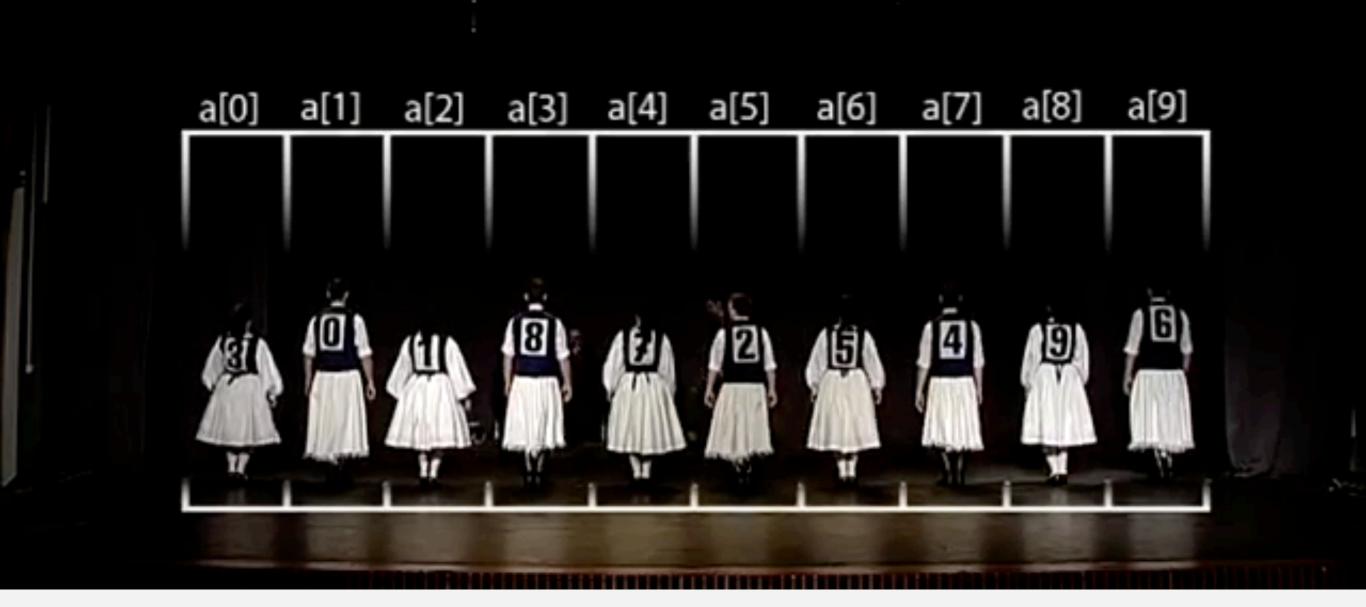
stability

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• In iteration i, swap a[i] with each larger entry to its left.



https://www.youtube.com/watch?v=ROalU379I3U

Algorithm. ↑ scans from left to right.

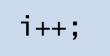
Invariants.



#### Insertion sort: inner loop

To maintain algorithm invariants:

• Move the pointer to the right.



in order for a for

 Moving from right to left, exchange a[i] with each larger entry to its left.



```
public class Insertion
{
   public static void sort(Comparable[] a)
   {
      int n = a.length;
      for (int i = 0; i < n; i++)
         for (int j = i; j > 0; j--)
            if (less(a[j], a[j-1]))
               exch(a, j, j-1);
            else break;
   }
   private static boolean less(Comparable v, Comparable w)
   { /* as before */ }
   private static void exch(Object[] a, int i, int j)
   { /* as before */ }
}
```

https://algs4.cs.princeton.edu/21elementary/Insertion.java.html



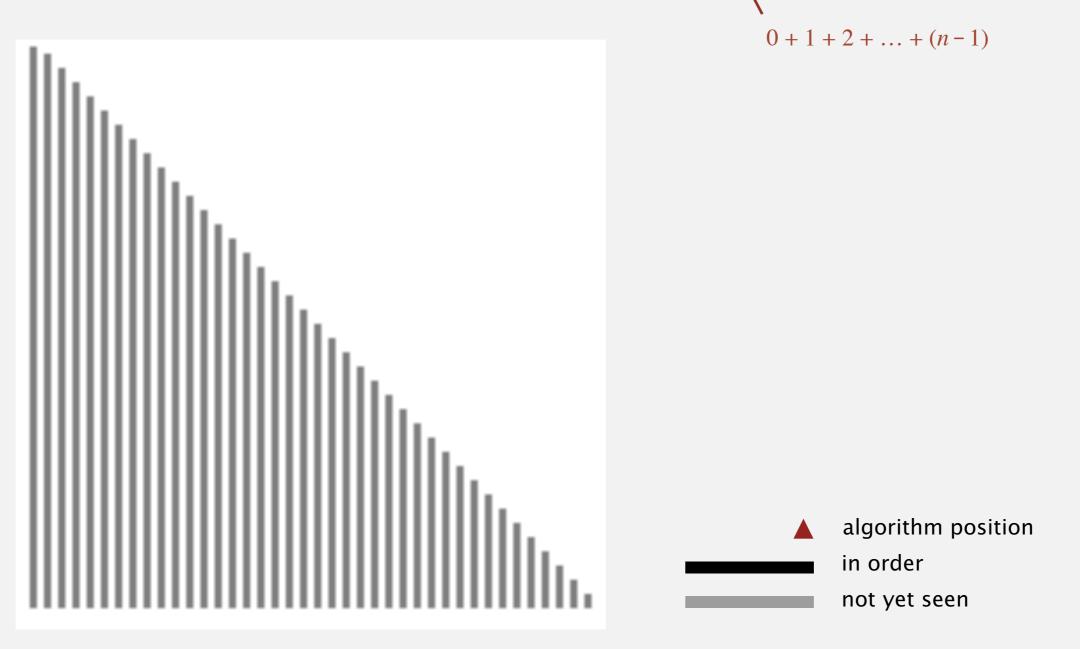
How many compares does insertion sort make to sort an array of *n* distinct keys in reverse order?

- A. ~ nB. ~  $1/4 n^2$ C. ~  $1/2 n^2$
- **D.** ~  $n^2$

#### Insertion sort: analysis

Worst case. Insertion sort makes ~  $\frac{1}{2} n^2$  compares and ~  $\frac{1}{2} n^2$  exchanges to sort an array of *n* distinct keys in reverse order.

Pf. Exactly *i* compares and exchanges in iteration *i*.

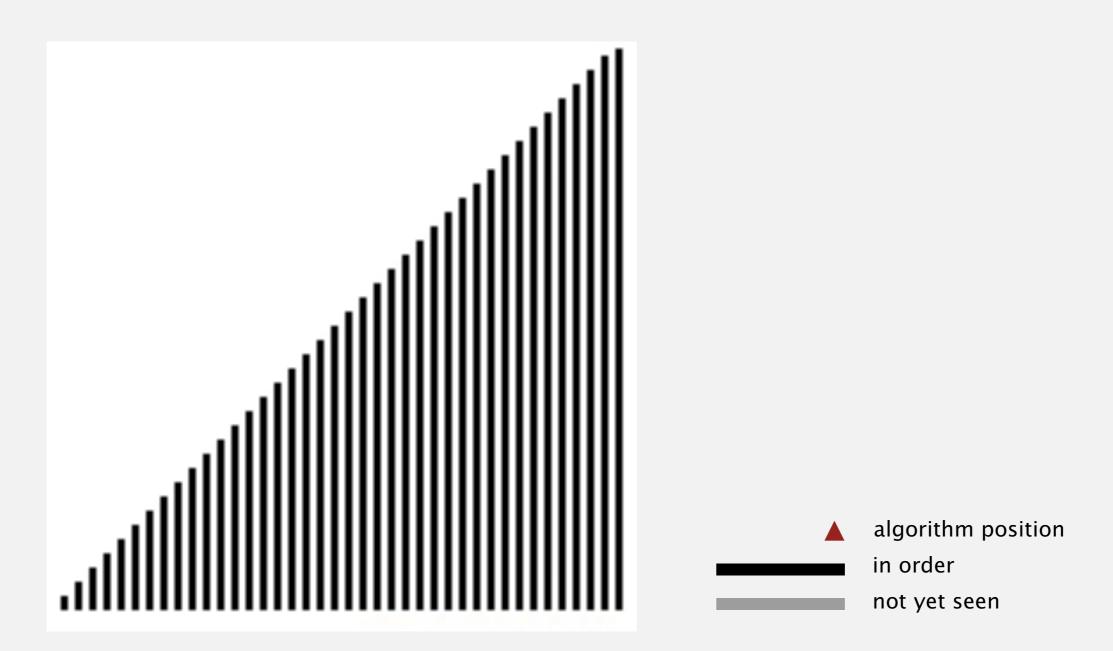




# Which is faster in practice to sort an array of n items, selection sort or insertion sort?

- A. Selection sort.
- **B.** Insertion sort.
- C. No significant difference.
- D. It depends.

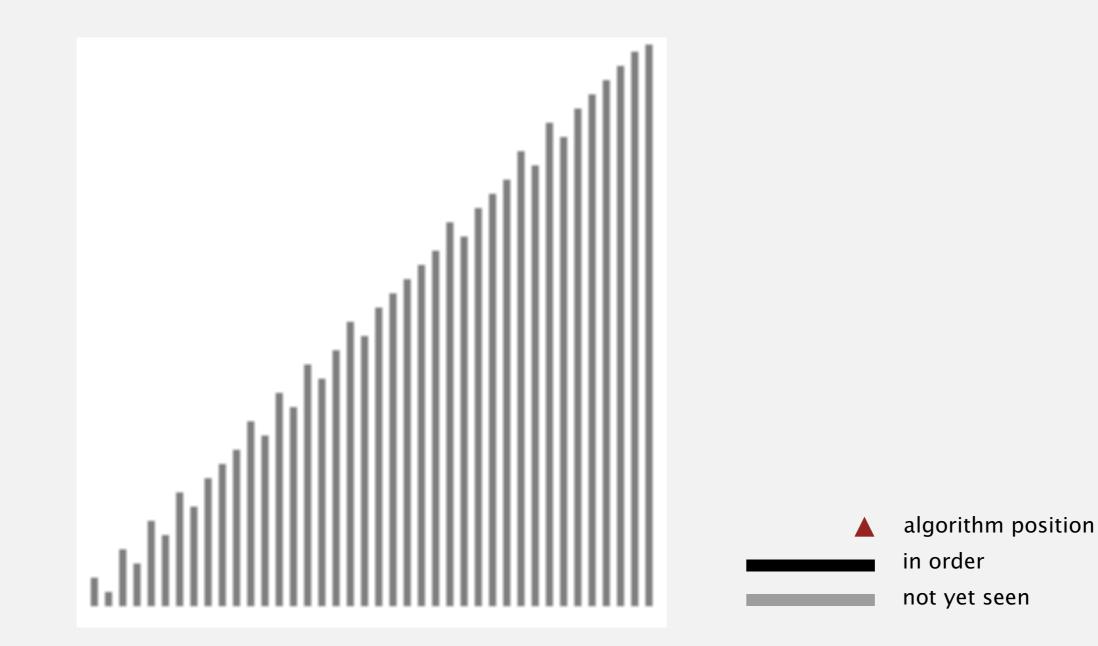
Best case. Insertion sort makes n-1 compares and 0 exchanges to sort an array of n distinct keys in ascending order.



#### Insertion sort: analysis

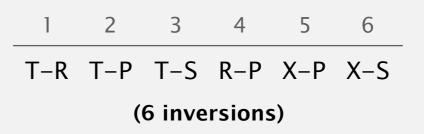
Good case. Insertion sort is linear time for "partially sorted" arrays.

Q. What do we mean by partially sorted?



**Def.** An inversion is a pair of keys that are out of order.

AEELMOTRXPS



**Def.** A family of arrays is partially sorted if the number of inversions is  $\leq c n$ .

- Ex 1. A sorted array.
- Ex 2. A subarray of length 10 appended to a sorted subarray of length *n*.

 $\leq 10n + 45$  inversions

Proposition. Insertion sort runs in linear time on partially sorted arrays. Pf. exchange decreases number of inversions by 1

- Number of exchanges in insertion sort = number of inversions.
- Number of compares  $\leq$  number of exchanges + (n 1).

each compare in iteration *i* triggers one exchange (except possibly last one in iteration)

### Insertion sort: practical improvements

Half exchanges. Shift items over (instead of exchanging).

- Eliminates unnecessary data movement.
- No longer uses only less() and exch() to access data.

A C H H I M N N P Q X Y K B I N A R Y

Binary insertion sort. Use binary search to find insertion point.

- Number of compares  $\sim n \log_2 n$ .
- But still a quadratic number of array accesses.

binary search for first key > K

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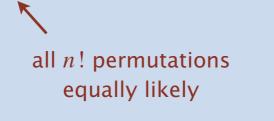
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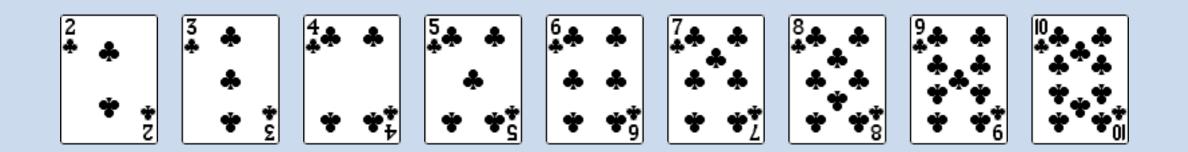
https://algs4.cs.princeton.edu

# **INTERVIEW QUESTION: SHUFFLE AN ARRAY**



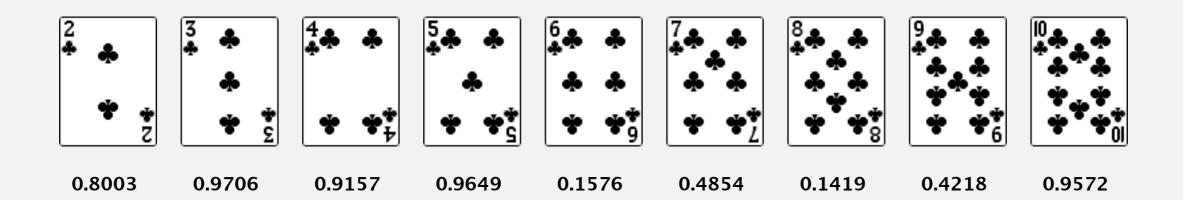
Goal. Rearrange array so that result is a uniformly random permutation.





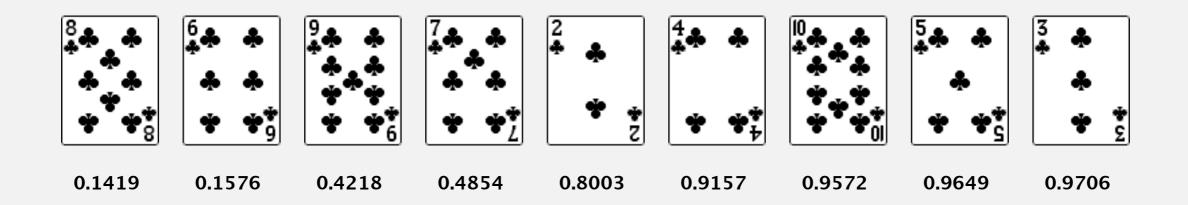
## Shuffling by sorting

- Generate a random real number for each array entry.
- Sort the array.



## Shuffling by sorting

- Generate a random real number for each array entry.
- Sort the array.



Proposition. Shuffle sort produces a uniformly random permutation.

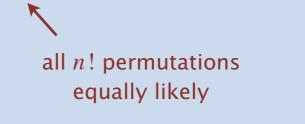
Application. Shuffle columns in a spreadsheet.

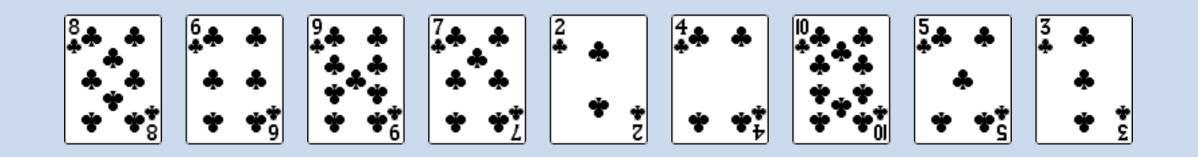
 assuming real numbers are uniformly random (and no ties)

## **INTERVIEW QUESTION: SHUFFLE AN ARRAY**



Goal. Rearrange array so that result is a uniformly random permutation.





Shuffling by sorting.

- Quadratic time (with insertion sort or selection sort).
- Linearithmic time (with mergesort).

Challenge. Design a linear-time algorithm (without sorting).

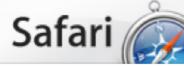


Microsoft antitrust probe by EU. Microsoft agreed to provide a randomized ballot screen for users to select browser.

### Select your web browser(s)



A fast new browser from Google. Try it now!



Safari for Windows from Apple, the world's most innovative browser.



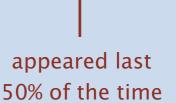
Your online security is Firefox's top priority. Firefox is free, and made to help you get the most out of the



The fastest browser on Earth. Secure, powerful and easy to use, with excellent privacy protection.



Designed to help you take control of your privacy and browse with confidence. Free from Microsoft.



# SHUFFLING HALL OF SHAME (PLANETPOKER.COM)



#### Texas hold'em poker. Software must shuffle electronic cards.



#### How We Learned to Cheat at Online Poker: A Study in Software Security



Ivy league school room draw. Students assigned random room draw times.

#### U. claims error in room draw process, provides compensation to affected students

By <u>Rebecca Han</u> | Apr 26, 2019

🔒 Print



Photo Credit: Jon Ort / The Daily Princetonian

# 2.1 ELEMENTARY SORTS

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rules of the game

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## Different orderings

**Q.** When might we need to define different sort orderings?

		0.00 K/s 💎 🏰 78% 🖥 16	54
S htt	Eighth Floor	0	:
Pł	Eleventh Floor	0	
0	Fifth Floor	0	
Er	First Floor	0	
Co	Fourteenth Floor	0	
t "	Fourth Floor	0	
W	Ninth Floor	0	
	Second Floor	۲	
	Seventh Floor	0	
	Sixth Floor	0	
	Tenth Floor	0	
<	Third Floor	0	

## Sort music library by artist



13	Take My Breath Away	BERLIN	4:13	Top Gun – Soundtrack	1
14	Circle Of Friends	Better Than Ezra	3:27	Empire Records	
15	Dancing With Myself	Billy Idol	4:43	Don't Stop	
16	Rebel Yell	Billy Idol	4:49	Rebel Yell	
17	✓ Piano Man	Billy Joel	5:36	Greatest Hits Vol. 1	I
18	Pressure	Billy Joel	3:16	Greatest Hits, Vol. II (1978 - 1985) (Disc 2)	I
19	The Longest Time	Billy Joel	3:36	Greatest Hits, Vol. II (1978 - 1985) (Disc 2)	I
20	✓ Atomic	Blondie	3:50	Atomic: The Very Best Of Blondie	I
21	✓ Sunday Girl	Blondie	3:15	Atomic: The Very Best Of Blondie	I
22	✓ Call Me	Blondie	3:33	Atomic: The Very Best Of Blondie	I
23	☑ Dreaming	Blondie	3:06	Atomic: The Very Best Of Blondie	I
24	✓ Hurricane	Bob Dylan	8:32	Desire	I
25	The Times They Are A-Changin'	Bob Dylan	3:17	Greatest Hits	I
26	Livin' On A Prayer	Bon Jovi	4:11	Cross Road	I
27	Beds Of Roses	Bon Jovi	6:35	Cross Road	I
28	Runaway	Bon Jovi	3:53	Cross Road	I
29	Rasputin (Extended Mix)	Boney M	5:50	Greatest Hits	I
30	Have You Ever Seen The Rain	Bonnie Tyler	4:10	Faster Than The Speed Of Night	I
31	Total Eclipse Of The Heart	Bonnie Tyler	7:02	Faster Than The Speed Of Night	I
32	Straight From The Heart	Bonnie Tyler	3:41	Faster Than The Speed Of Night	I
33	Holding Out For A Hero	Bonny Tyler	5:49	Meat Loaf And Friends	
34	🛛 Dancing In The Dark 📀	Bruce Springsteen 🕤	4:05	Born In The U.S.A.	l
35	Thunder Road	Bruce Springsteen	4:51	Born To Run	
36	Born To Run	Bruce Springsteen	4:30	Born To Run	
37	✓ Jungleland	Bruce Springsteen	9:34	Born To Run	H
	I Turn Turn Turn (To Funnthin	The Burde	2.57	Forract Cump The Soundtrack (Dice 3)	-

## Sort music library by song name

		Cross Road Bon Jovi		
	Name 🔺	Artist		Album
1	☑ Alive	Pearl Jam	5:41	
2	All Over The World	Pixies	5:27	
3	All Through The Night	Cyndi Lauper		She's So Unusual
4	✓ Allison Road	Gin Blossoms		New Miserable Experience
5	🗹 Ama, Ama, Ama Y Ensancha El	Extremoduro		Deltoya (1992)
6	And We Danced	Hooters		Nervous Night
7	As I Lay Me Down	Sophie B. Hawkins		Whaler
8	Atomic	Blondie		Atomic: The Very Best Of Blondie
9	Automatic Lover	Jay-Jay Johanson		Antenna
10	Baba O'Riley	The Who		Who's Better, Who's Best
11	☑ Beautiful Life	Ace Of Base		The Bridge
12	Beds Of Roses	Bon Jovi 📀		Cross Road
13	Black	Pearl Jam	5:44	
14	Bleed American	Jimmy Eat World		Bleed American
15	Borderline	Madonna		The Immaculate Collection
16	Born To Run	Bruce Springsteen		Born To Run
17	Both Sides Of The Story	Phil Collins		Both Sides
18	Bouncing Around The Room	Phish The Cure		A Live One (Disc 1)
19 20	✓ Boys Don't Cry ✓ Brat	The Cure Green Day		Staring At The Sea: The Singles 1979–1985 Insomniac
20	Breakdown	Deerheart		Deerheart
22	Bring Me To Life (Kevin Roen Mix)	Evanescence Vs. Pa	9:48	beemeant
23	Californication	Red Hot Chili Pepp	1:40	
24	Call Me	Blondie		Atomic: The Very Best Of Blondie
25	Can't Get You Out Of My Head	Kylie Minogue		Fever
26	Celebration	Kool & The Gang		Time Life Music Sounds Of The Seventies – C
20	Chainers Chainers	Sukhwinder Singh		Rombay Droams
				) 4 ) 4

### Comparable interface: review

Comparable interface: sort using a type's natural order.

```
public class Date implements Comparable<Date>
{
  private final int month, day, year;
  public Date(int m, int d, int y)
   {
     month = m;
     day = d;
     year = y;
   }
                                                            natural order
  public int compareTo(Date that)
   {
      if (this.year < that.year ) return -1;
      if (this.year > that.year ) return +1;
      if (this.month < that.month) return -1;
      if (this.month > that.month) return +1;
      if (this.day < that.day ) return -1;
      if (this.day > that.day ) return +1;
      return 0;
   }
}
```

### Comparator interface

Comparator interface: sort using an alternate order.

```
public interface Comparator<Item>
{
    public int compare(Item v, Item w);
}
```

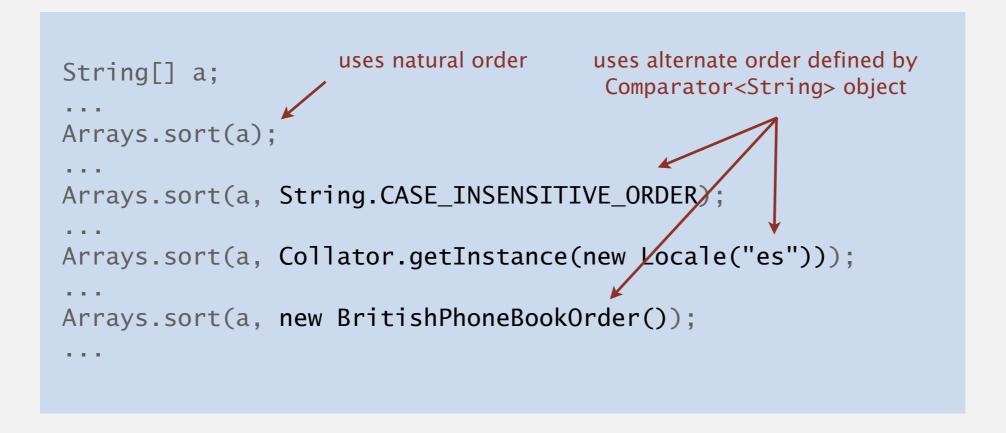
#### Required property. Must be a total order.

string order	example		
natural order	Now is the time pre-1994 order for		
case insensitive	is Now the time		
Spanish language	café cafetero cuarto churro nube ñoño		
British phone book	M <mark>cK</mark> inley M <mark>ac</mark> kintosh		

### Comparator interface: system sort

#### To use with Java system sort:

- Create Comparator object.
- Pass as second argument to Arrays.sort().



Bottom line. Decouples the definition of the data type from the definition of what it means to compare two objects of that type.

#### To implement a comparator:

- Define a (nested) class that implements the Comparator interface.
- Implement the compare() method.
- Provide client access to Comparator.

#### To implement a comparator:

}

- Define a (nested) class that implements the Comparator interface.
- Implement the compare() method.
- Provide client access to Comparator.

```
import java.util.Comparator;
public class Student
ł
   private final String name;
   private final int section;
   private static class SectionOrder implements Comparator<Student>
   {
      public int compare(Student v, Student w)
         return(Integer.compare(v.section, w.section);)
                                                          }
   }
   public static Comparator<Student> bySectionOrder()
                                                           useful library
      return new SectionOrder(); }
                                                             method
```

#### To implement a comparator:

- Define a (nested) class that implements the Comparator interface.
- Implement the compare() method.
- Provide client access to Comparator.

#### Insertion.sort(a, Student.byNameOrder());

Andrews	3	А	(664) 480-0023	097 Little
Battle	4	С	(874) 088–1212	121 Whitman
Chen	3	А	(991) 878–4944	308 Blair
Fox	3	А	(884) 232–5341	11 Dickinson
Furia	1	А	(766) 093–9873	101 Brown
Gazsi	4	В	(800) 867–5309	101 Brown
Kanaga	3	В	(898) 122–9643	22 Brown
Rohde	2	А	(232) 343–5555	343 Forbes

#### Insertion.sort(a, Student.bySectionOrder());

Furia	1	А	(766) 093–9873	101 Brown
Rohde	2	А	(232) 343–5555	343 Forbes
Andrews	3	А	(664) 480-0023	097 Little
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# 2.1 ELEMENTARY SORTS

skipped in lecture Isee precepti

rules of the game

selection sort

insertion sort

comparators

# Algorithms

Robert Sedgewick | Kevin Wayne

https://algs4.cs.princeton.edu

stability

shuffling

#### A typical application. First, sort by name; then sort by section.

Andrews	3	А	(664) 480–0023	097 Little
Battle	4	С	(874) 088–1212	121 Whitman
Chen	3	А	(991) 878–4944	308 Blair
Fox	3	А	(884) 232–5341	11 Dickinson
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Rohde	2	А	(232) 343–5555	343 Forbes

#### Selection.sort(a, Student.byNameOrder());

Selection.sort(a, Student.bySectionOrder());

Furia	1	А	(766) 093–9873	101 Brown
Rohde	2	А	(232) 343–5555	343 Forbes
Chen	3	А	(991) 878–4944	308 Blair
Fox	3	А	(884) 232–5341	11 Dickinson
Andrews	3	А	(664) 480-0023	097 Little
Kanaga	3	В	(898) 122–9643	22 Brown
Gazsi	4	В	(800) 867–5309	101 Brown
Battle	4	С	(874) 088–1212	121 Whitman

@#%&@! Students in section 3 no longer sorted by name.

A stable sort preserves the relative order of items with equal keys.



#### Which sorting algorithm(s) are stable?

- A. Selection sort.
- **B.** Insertion sort.
- C. Both A and B.
- **D.** Neither A nor B.

### Stability: insertion sort

Proposition. Insertion sort is stable.

```
public class Insertion
{
     public static void sort(Comparable[] a)
     {
          int n = a.length;
          for (int i = 0; i < n; i++)
                for (int j = i; j > 0 && less(a[j], a[j-1]); j--)
                     exch(a, j, j-1);
     }
}
                                             j
                                                      0 1 2 3 4
                                       i
                                       0
                                          \mathbf{0} \quad \mathbf{B}_1 \quad \mathbf{A}_1 \quad \mathbf{A}_2 \quad \mathbf{A}_3 \quad \mathbf{B}_2
                                              0 \quad A_1 \quad B_1 \quad A_2 \quad A_3 \quad B_2
                                       1
                                       2 1 A<sub>1</sub> A<sub>2</sub> B<sub>1</sub> A<sub>3</sub> B<sub>2</sub>
                                       3 \qquad 2 \qquad A_1 \quad A_2 \quad A_3 \quad B_1 \quad B_2
                                       4 \quad 4 \quad A_1 \quad A_2 \quad A_3 \quad B_1 \quad B_2
                                                     A_1 \quad A_2 \quad A_3 \quad B_1 \quad B_2
```

Pf. Equal items never move past each other.

### Stability: selection sort

#### Proposition. Selection sort is not stable.

```
public class Selection
{
    public static void sort(Comparable[] a)
    {
        int n = a.length;
        for (int i = 0; i < n; i++)
        {
            int min = i;
            for (int j = i+1; j < n; j++)
                if (less(a[j], a[min]))
                        min = j;
            exch(a, i, min);
        }
    }
}</pre>
```

i	min	0	1	2	
0	2	Bı	B <sub>2</sub>	А	
1	1	А	<b>B</b> <sub>2</sub>	Bı	
2	2	А	B <sub>2</sub>	Bı	
		А	B <sub>2</sub>	Bı	

Pf by counterexample. Long-distance exchange can move an equal item past another one.