Performance Improvement

“Premature optimization is the root of all evil.”

-- Donald Knuth

“Rules of Optimization:
• Rule 1: Don't do it.
• Rule 2 (for experts only): Don't do it yet.”

-- Michael A. Jackson
“Programming in the Large”

Design & Implement
- Program & programming style (done)
- Common data structures and algorithms (done)
- Modularity (done)
- Building techniques & tools (done)

Debug
- Debugging techniques & tools (done)

Test
- Testing techniques (done)

Maintain
- Performance improvement techniques & tools ← we are here
Goals of this Lecture

Help you learn about:
- How to use profilers to identify code hot-spots
- How to make your programs run faster

Why?
- In a large program, typically a small fragment of the code consumes most of the CPU time
- A power programmer knows how to identify such code fragments
- A power programmer knows techniques for improving the performance of such code fragments
Agenda

Should you optimize?
What should you optimize?
Optimization techniques
Techniques described in this lecture can answer:

• How slow is my program?
• Where is my program slow?
• Why is my program slow?
• How can I make my program run faster?
• How can I make my program use less memory?
Techniques described in this lecture can yield code that:

• Is less clear/maintainable
• Might confuse debuggers
• Might contain bugs
  • Requires regression testing

So…
When to Improve Performance

“The first principle of optimization is don’t.

Is the program good enough already? Knowing how a program will be used and the environment it runs in, is there any benefit to making it faster?”

-- Kernighan & Pike
Timing a Program

Run a tool to time program execution
  • E.g., Unix `time` command

```
$ time sort < bigfile.txt > output.txt
real    0m12.977s
user    0m12.860s
sys     0m0.010s
```

Output:
  • **Real**: Wall-clock time between program invocation and termination
  • **User**: CPU time spent executing the program
  • **System**: CPU time spent within the OS on the program’s behalf
Enable compiler speed optimization

```
gcc217 -Ox mysort.c -o mysort
```

- Compiler looks for ways to transform your code so that result is the same but it runs faster
- `x` controls how many transformations the compiler tries – see “man gcc” for details
  - `-O1`: optimize (default if no number is specified)
  - `-O2`: optimize more (longer compile time)
  - `-O3`: optimize yet more (including inlining)

**Warning: Speed optimization can affect debugging**
- e.g., Optimization eliminates variable ⇒ GDB cannot print value of variable
Now What?

So you’ve determined that your program is taking too long, even with compiler optimization enabled (and NDEBUG defined, etc.)

Is it time to rewrite the program?
Should you optimize?
What should you optimize?
Optimization techniques
Identifying Hot Spots

Spend time optimizing only the parts of the program that will make a difference!

Gather statistics about your program’s execution

- **Coarse-grained**: how much time did execution of a particular function call take?
  - Time individual function calls or blocks of code

- **Fine-grained**: how many times was a particular function called? How much time was taken by all calls to that function?
  - Use an execution profiler such as gprof
Timing Parts of a Program

Call a function to compute **wall-clock time** consumed

- Unix `gettimeofday()` returns time in seconds + microseconds

```c
#include <sys/time.h>

struct timeval startTime;
struct timeval endTime;
double wallClockSecondsConsumed;

gettimeofday(&startTime, NULL);
<execute some code here>
gettimeofday(&endTime, NULL);
wallClockSecondsConsumed =
    endTime.tv_sec - startTime.tv_sec +
    1.0E-6 * (endTime.tv_usec - startTime.tv_usec);
```

- Not defined by C90 standard
Call a function to compute **CPU time** consumed

- `clock()` returns CPU times in `CLOCKS_PER_SEC` units

```c
#include <time.h>

clock_t startClock;
clock_t endClock;
double cpuSecondsConsumed;

startClock = clock();
<execute some code here>
endClock = clock();
cpuSecondsConsumed =
    ((double)(endClock - startClock)) / CLOCKS_PER_SEC;
```

- Defined by C90 standard
Identifying Hot Spots

Spend time optimizing only the parts of the program that will make a difference!

Gather statistics about your program’s execution

- **Coarse-grained**: how much time did execution of a particular function call take?
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GPROF Example Program

Example program for GPROF analysis

- Sort an array of 10 million random integers
- Artificial: consumes lots of CPU time, generates no output

```c
#include <string.h>
#include <stdio.h>
#include <stdlib.h>

enum {MAX_SIZE = 10000000};
int a[MAX_SIZE];

void fillArray(int a[], int size)
{
    int i;
    for (i = 0; i < size; i++)
        a[i] = rand();
}

void swap(int a[], int i, int j)
{
    int temp = a[i];
    a[i] = a[j];
    a[j] = temp;
}
...
```
int partition(int a[], int left, int right)
{
    int first = left-1;
    int last = right;
    for (;;) {
        while (a[++first] < a[right])
            ;
        while (a[right] < a[--last])
            if (last == left)
                break;
            if (first >= last)
                break;
        swap(a, first, last);
    }
    swap(a, first, right);
    return first;
}
Example program for GPROF analysis (cont.)

```c
...  
void quicksort(int a[], int left, int right)
{  if (right > left)
   {  int mid = partition(a, left, right);
      quicksort(a, left, mid - 1);
      quicksort(a, mid + 1, right);
   }
}

int main(void)
{  fillArray(a, MAX_SIZE);
    quicksort(a, 0, MAX_SIZE - 1);
    return 0;
}
```
Using GPROF

Step 1: Instrument the program

```bash
gcc217 -pgmysort.c -o mysort
```

• Adds profiling code to mysort, that is…
• “Instruments” mysort

Step 2: Run the program

```bash
./mysort
```

• Creates file `gmon.out` containing statistics

Step 3: Create a report

```bash
gprof mysort > myreport
```

• Uses `mysort` and `gmon.out` to create textual report

Step 4: Examine the report

```bash
cat myreport
```
gprof Design

What's going on behind the scenes?

- `-pg` generates code to interrupt program many times per second
- Each time, records *where* the code was interrupted
- `gprof` uses symbol table to map back to function name
The GPROF Report

- Each line describes one function
  - **name**: name of the function
  - **%time**: percentage of time spent executing this function
  - **cumulative seconds**: [skipping, as this isn’t all that useful]
  - **self seconds**: time spent executing this function
  - **calls**: number of times function was called (excluding recursive)
  - **self s/call**: average time per execution (excluding descendants)
  - **total s/call**: average time per execution (including descendants)

<table>
<thead>
<tr>
<th>%time</th>
<th>cumulative</th>
<th>self seconds</th>
<th>seconds</th>
<th>calls</th>
<th>s/call</th>
<th>s/call</th>
<th>name</th>
</tr>
</thead>
<tbody>
<tr>
<td>84.54</td>
<td>2.27</td>
<td>2.27</td>
<td>6665307</td>
<td>0.00</td>
<td>0.00</td>
<td>0.00</td>
<td>partition</td>
</tr>
<tr>
<td>9.33</td>
<td>2.53</td>
<td>0.25</td>
<td>54328749</td>
<td>0.00</td>
<td>0.00</td>
<td>0.00</td>
<td>swap</td>
</tr>
<tr>
<td>2.99</td>
<td>2.61</td>
<td>0.08</td>
<td>1</td>
<td>0.08</td>
<td>2.61</td>
<td></td>
<td>quicksort</td>
</tr>
<tr>
<td>2.61</td>
<td>2.68</td>
<td>0.07</td>
<td>1</td>
<td>0.07</td>
<td>0.07</td>
<td></td>
<td>fillArray</td>
</tr>
</tbody>
</table>
### Call graph profile

<table>
<thead>
<tr>
<th>index</th>
<th>% time</th>
<th>self</th>
<th>children</th>
<th>called</th>
<th>name</th>
</tr>
</thead>
<tbody>
<tr>
<td>[1]</td>
<td>100.0</td>
<td>0.00</td>
<td>2.68</td>
<td>main</td>
<td>[1]</td>
</tr>
<tr>
<td></td>
<td></td>
<td>0.08</td>
<td>2.53</td>
<td>1/1</td>
<td>quicksort [2]</td>
</tr>
<tr>
<td></td>
<td></td>
<td>0.07</td>
<td>0.00</td>
<td>1/1</td>
<td>fillArray [5]</td>
</tr>
<tr>
<td>[2]</td>
<td>97.4</td>
<td>0.08</td>
<td>2.53</td>
<td>main</td>
<td>[1]</td>
</tr>
<tr>
<td></td>
<td></td>
<td>1+13330614</td>
<td>2.27</td>
<td>quicksort [2]</td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td>6665307/6665307</td>
<td>0.25</td>
<td>partition [3]</td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td>13330614</td>
<td></td>
<td>quicksort [2]</td>
<td></td>
</tr>
<tr>
<td>[3]</td>
<td>94.4</td>
<td>2.27</td>
<td>0.25</td>
<td>6665307</td>
<td>partition [3]</td>
</tr>
<tr>
<td></td>
<td></td>
<td>54328749/54328749</td>
<td>0.00</td>
<td>swap [4]</td>
<td></td>
</tr>
<tr>
<td>[4]</td>
<td>9.4</td>
<td>0.25</td>
<td>0.00</td>
<td>54328749</td>
<td>partition [3]</td>
</tr>
<tr>
<td></td>
<td></td>
<td>54328749</td>
<td></td>
<td>swap [4]</td>
<td></td>
</tr>
<tr>
<td>[5]</td>
<td>2.6</td>
<td>0.07</td>
<td>0.00</td>
<td>1</td>
<td>fillArray [5]</td>
</tr>
</tbody>
</table>

---

The GPROF Report (cont.)

22
Call graph profile (cont.)

- Each section describes one function
  - Which functions called it, and how much time was consumed?
  - Which functions it calls, how many times, and for how long?
- Usually overkill; we won’t look at this output in any detail
GPROF Report Analysis

Observations

• `swap()` is called very many times; each call consumes little time; 
  `swap()` consumes only 9% of the time overall

• `partition()` is called many times; each call consumes little time; 
  but `partition()` consumes 85% of the time overall

Conclusions

• To improve performance, try to make `partition()` faster

• Don’t even think about trying to make `fillArray()` or `quicksort()` faster
Agenda

Should you optimize?
What should you optimize?
Optimization techniques
Using Better Algs and DSs

Use a better algorithm or data structure

Example:
• Would a different sorting algorithm work better?

See COS 226…
• But only where it would help! Not worth using asymptotically efficient (but complex, hard-to-understand, and hard-to-maintain) algorithms and data structures in parts of code that don't matter!
Avoiding Repeated Computation

Before:

```c
int g(int x)
{   return f(x) + f(x) + f(x) + f(x);
}
```

After:

```c
int g(int x)
{   return 4 * f(x);
}
```
Q: Could a good compiler do this optimization for you?

Before:
```c
int g(int x)
{  return f(x) + f(x) + f(x) + f(x);
}
```

After:
```c
int g(int x)
{  return 4 * f(x);
}
```

A. Yes
B. Only sometimes
C. No
Aside: Side Effects as Blockers

Q: Could a good compiler do that for you?
A: Only sometimes…

Suppose \( f() \) has side effects?

\[
\text{int } g(\text{int } x) \\
\{ \text{ return } f(x) + f(x) + f(x) + f(x); \} \\
\]

\[
\text{int } g(\text{int } x) \\
\{ \text{ return } 4 * f(x); \} \\
\]

\[
\text{int } \text{counter } = 0; \\
\ldots \\
\text{int } f(\text{int } x) \\
\{ \text{ return } \text{counter++}; \} \\
\]

And \( f() \) might be defined in another file known only at link time!
Avoiding Repeated Computation

Before:

\[
\text{for (i = 0; i < n; i++)}
\]
\[
\text{for (j = 0; j < n; j++)}
\]
\[
a[n*i + j] = b[j];
\]

After:

\[
\text{for (i = 0; i < n; i++)}
\]
\[
\{
\text{ni = n * i;}
\text{for (j = 0; j < n; j++)}
\]
\[
a[ni + j] = b[j];
\}
\]
Q: Could a good compiler do this optimization for you?

Before:

```c
for (i = 0; i < n; i++)
    for (j = 0; j < n; j++)
        a[n*i + j] = b[j];
```

After:

```c
for (i = 0; i < n; i++)
{
    ni = n * i;
    for (j = 0; j < n; j++)
        a[ni + j] = b[j];
}
```

A. Yes
B. Only sometimes
C. No
Avoiding Repeated Computation

Before:

```c
for (i = 0; i < strlen(s); i++)
{  /* Do something with s[i] */
}
```

After:

```c
length = strlen(s);
for (i = 0; i < length; i++)
{  /* Do something with s[i] */
}
```

Could a good compiler do that for you?
void twiddle(int *p1, int *p2)
{  *p1 += *p2;
    *p1 += *p2;
}

void twiddle(int *p1, int *p2)
{  *p1 += *p2 * 2;
}
Q: Could a good compiler do this optimization for you?

Before:
```c
void twiddle(int *p1, int *p2)
{  *p1 += *p2;
   *p1 += *p2;
}
```

After:
```c
void twiddle(int *p1, int *p2)
{  *p1 += *p2 * 2;
}
```

A. Yes
B. Only sometimes
C. No
Aside: Aliases as Blockers

Q: Could a good compiler do that for you?
A: Not necessarily

What if \( p1 \) and \( p2 \) are aliases?

- What if \( p1 \) and \( p2 \) point to the same integer?
- First version: result is 4 times \( *p1 \)
- Second version: result is 3 times \( *p1 \)

Some compilers support \texttt{restrict} keyword
Inlining Function Calls

Before:

```c
void g(void)
{  /* Some code */
}
void f(void)
{  ...
    g();
    ...
}
```

After:

```c
void f(void)
{  ...
    /* Some code */
    ...
}
```

Beware: Can introduce redundant/cloned code
Some compilers support `inline` keyword

Could a good compiler do that for you?
Unrolling Loops

Original:

```c
for (i = 0; i < 6; i++)
a[i] = b[i] + c[i];
```

Maybe faster:

```c
for (i = 0; i < 6; i += 2)
{  a[i+0] = b[i+0] + c[i+0];
    a[i+1] = b[i+1] + c[i+1];
}
```

Maybe even faster:

```c
a[i+0] = b[i+0] + c[i+0];
a[i+1] = b[i+1] + c[i+1];
a[i+2] = b[i+2] + c[i+2];
a[i+3] = b[i+3] + c[i+3];
a[i+4] = b[i+4] + c[i+4];
a[i+5] = b[i+5] + c[i+5];
```

Could a good compiler do that for you?

Some compilers provide option, e.g. `-funroll-loops`
Using a Lower-Level Language

Rewrite code in a lower-level language

• As described in second half of course…
• Compose key functions in **assembly language** instead of C
  • Use registers instead of memory
  • Use instructions (e.g. `adc`) that compiler doesn’t know

Beware: Modern optimizing compilers generate fast code

• Hand-written assembly language code could be slower!
Summary

Steps to improve **execution (time)** efficiency:
- Don't do it.
- Don't do it yet.
- Time the code to make sure it's necessary
- Enable compiler optimizations
- Identify hot spots using profiling
- Use a better algorithm or data structure
- Tune the code