Modules and Representation Invariants

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In previous classes:

Reasoning about individual OCaml expressions.

Now:

Reasoning about Modules (abstract types + collections of values)

A Signature for Sets

```
module type SET =
  sig
    type ‘a set
    val empty : ‘a set
    val mem : ‘a -> ‘a set -> bool
    val add : ‘a -> ‘a set -> ‘a set
    val rem : ‘a -> ‘a set -> ‘a set
    val size : ‘a set -> int
    val union : ‘a set -> ‘a set -> ‘a set
    val inter : 'a set -> ‘a set -> 'a set
  end
```

Sets as Lists

```
module Set1 : SET =
  struct
    type 'a set = 'a list
    let empty = []
    let mem = List.mem
    let add x l = x :: l
    let rem x l = List.filter ((<>) x) l
    let rec size l =
      match l with
      | [] -> 0
      | h::t -> size t + (if mem h t then 0 else 1)
    let union l1 l2 = l1 @ l2
    let inter l1 l2 = List.filter (fun h -> mem h l2) l1
  end
```

Very slow in many ways!

Sets as Lists without Duplicates

```
module Set2 : SET =
  struct
    type 'a set = 'a list
    let empty = []
    let mem = List.mem
    (* add:  check if already a member *)
    let add x l = if mem x l then l else x::l
    let rem x l = List.filter ((<>) x) l
    (* size:  list length is number of unique elements *)
    let size l = List.length l
    (* union: discard duplicates *)
    let union l1 l2 = List.fold_left
          (fun a x -> if mem x l2 then a else x::a) l2 l1
    let inter l1 l2 = List.filter (fun h -> mem h l2) l1
  end
```

Back to Sets

The interesting operation:

```
(* size:  list length is number of unique elements *)
let size (l:'a set) : int = List.length l
```

Why does this work? It depends on an invariant:

All lists supplied as an argument contain no duplicates.

A *representation invariant* is a property that holds of all values of a particular (abstract) type.

Implementing Representation Invariants

For lists with no duplicates:

```
(* checks that a list has no duplicates *)
let rec inv (s : 'a set) : bool =
    match s with
      [] -> true
    | hd::tail -> not (mem hd tail) && inv tail

let rec check (s : ‘a set) (m:string) : ‘a set =
  if inv s then
    s
  else
    failwith m
```

Debugging with Representation Invariants

As a precondition on input sets:

```
(* size:  list length is number of unique elements *)
let size (s:'a set) : int =
  ignore (check s "size:  bad set input");
  List.length s
```

Debugging with Representation Invariants

As a precondition on input sets:

```
(* size:  list length is number of unique elements *)
let size (s:'a set) : int =
  ignore (check s "size:  bad set input");
  List.length s
```

As a postcondition on output sets:

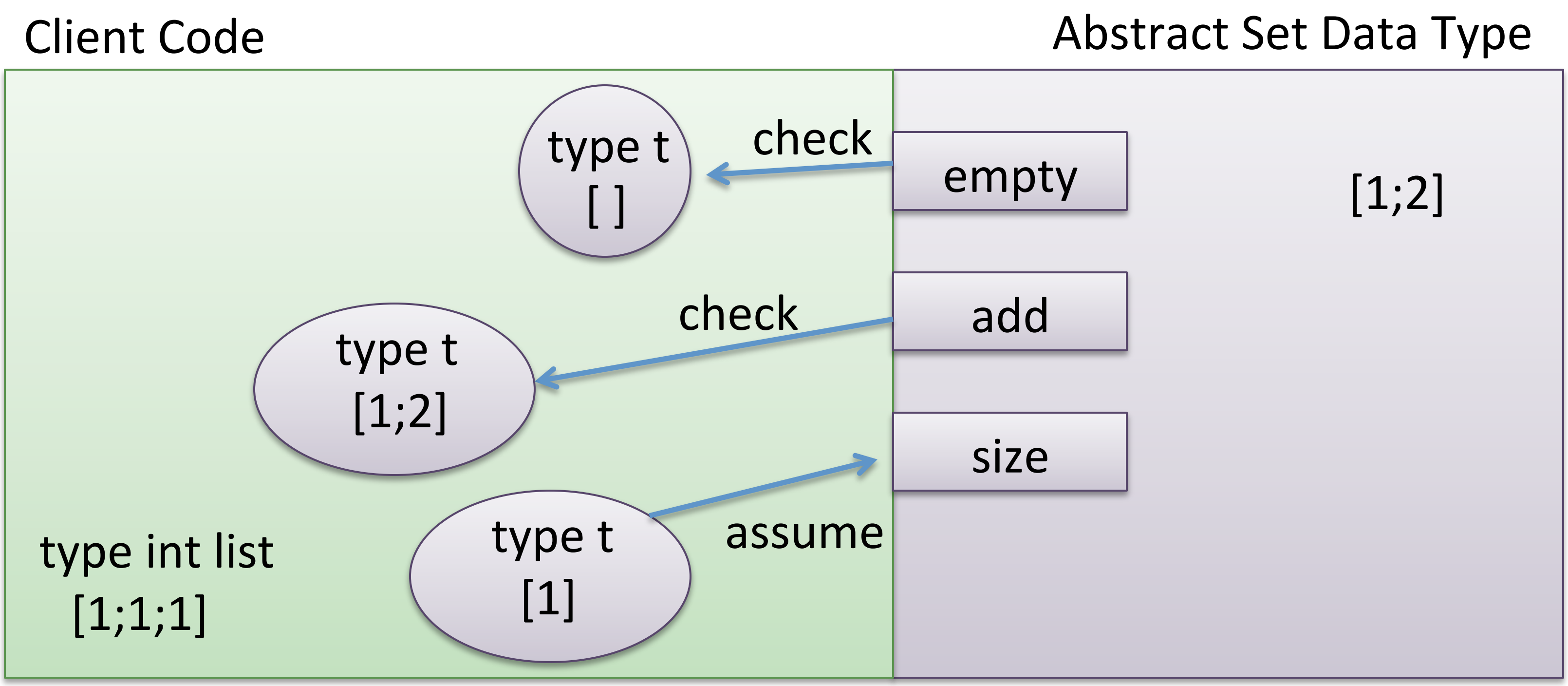
```
(* add x to set s *)
let add x s =
  let s = if mem x s then s else x::s in
  check s "add: bad set output"
```

A Signature for Sets

```
module type SET =
  sig
    type 'a set
    val empty : 'a set
    val mem : 'a -> 'a set -> bool
    val add : 'a -> 'a set -> 'a set
    val rem : 'a -> 'a set -> 'a set
    val size : 'a set -> int
    val union : 'a set -> 'a set -> 'a set
    val inter : 'a set -> 'a set -> 'a set
  end
```

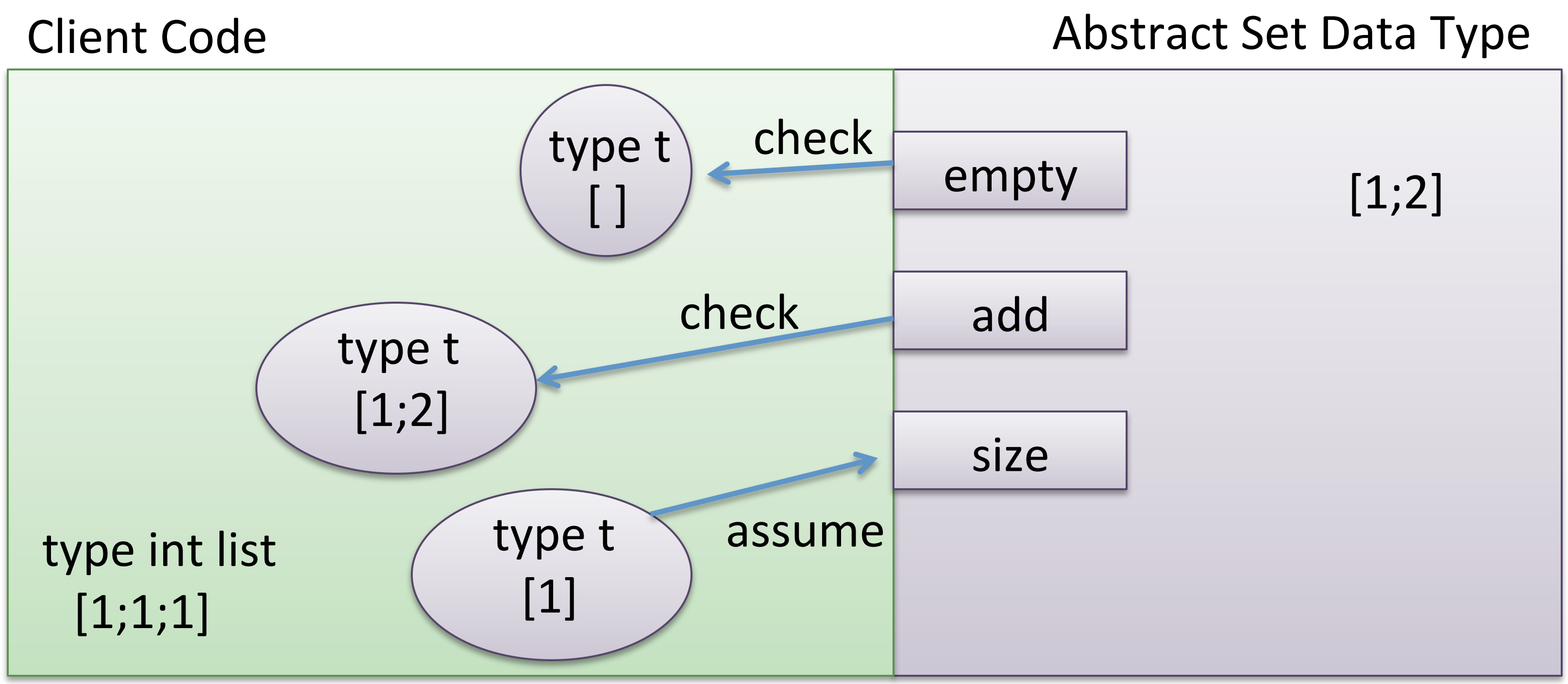
Suppose we check all the red values satisfy our invariant leaving the module, do we have to check the blue values entering the module satisfy our invariant?

Representation Invariants Pictorially



When debugging, we can check our invariant each time we construct a value of abstract type. We then get to assume the invariant on input to the module.

Representation Invariants Pictorially



When proving, we prove our invariant holds each time we construct a value of abstract type and release it to the client. We *get to assume* the invariant holds on input to the module.

Such a proof technique is *highly modular*: Independent of the client!

Repeating myself

You may

assume the invariant inv(i) for module inputs i with abstract type

provided you

prove the invariant inv(o) for all module outputs o with abstract type

Back to Sets

The interesting operation:

```
(* size:  list length is number of unique elements *)
let size (l:'a set) : int = List.length l
```

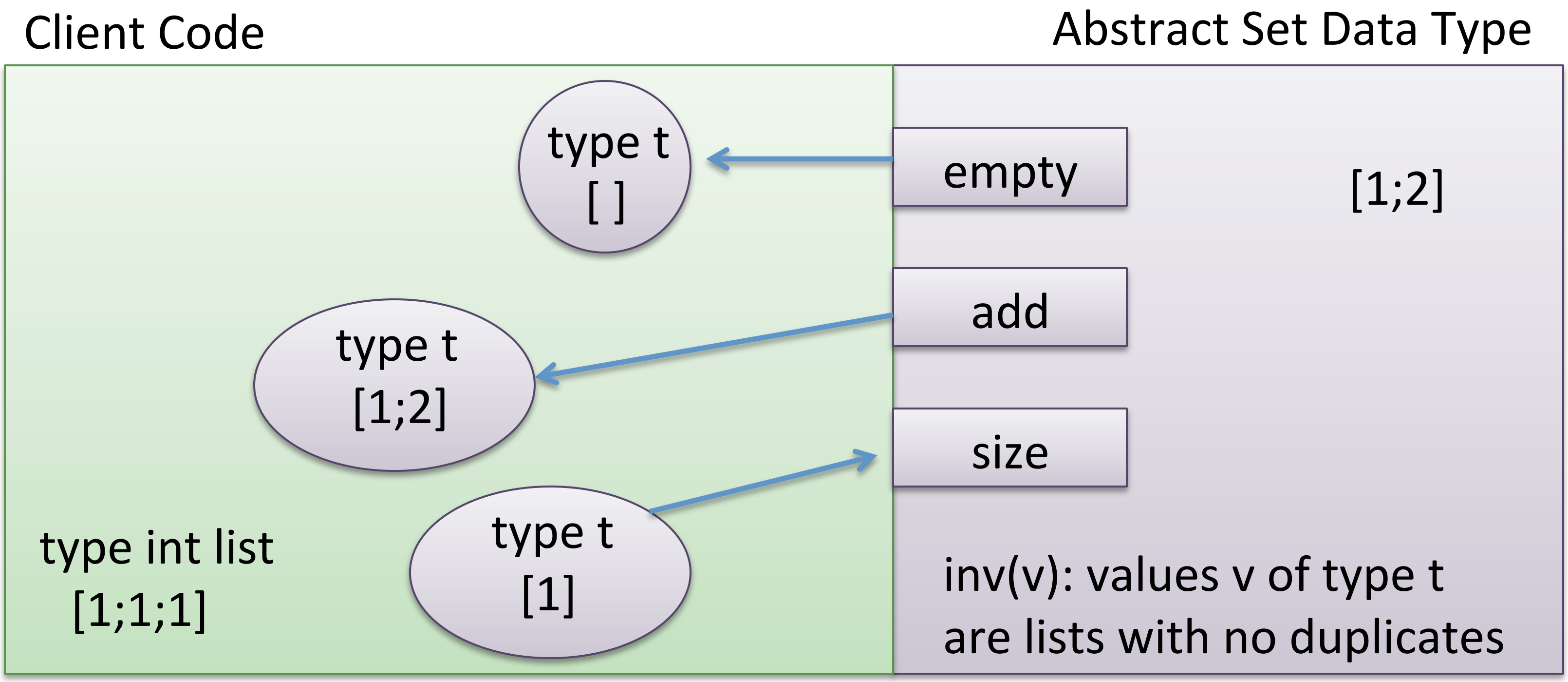
Why does this work? It depends on an invariant:

All lists supplied as an argument contain no duplicates.

How is this invariant enforced? By using abstract types. Every value of the abstract type 'a set satisfies the invariant. Internally, the module knows that the 'a set is 'a list and can establish the invariant, but externally clients don't know that and can't mess with established invariants.

Representation Invariants

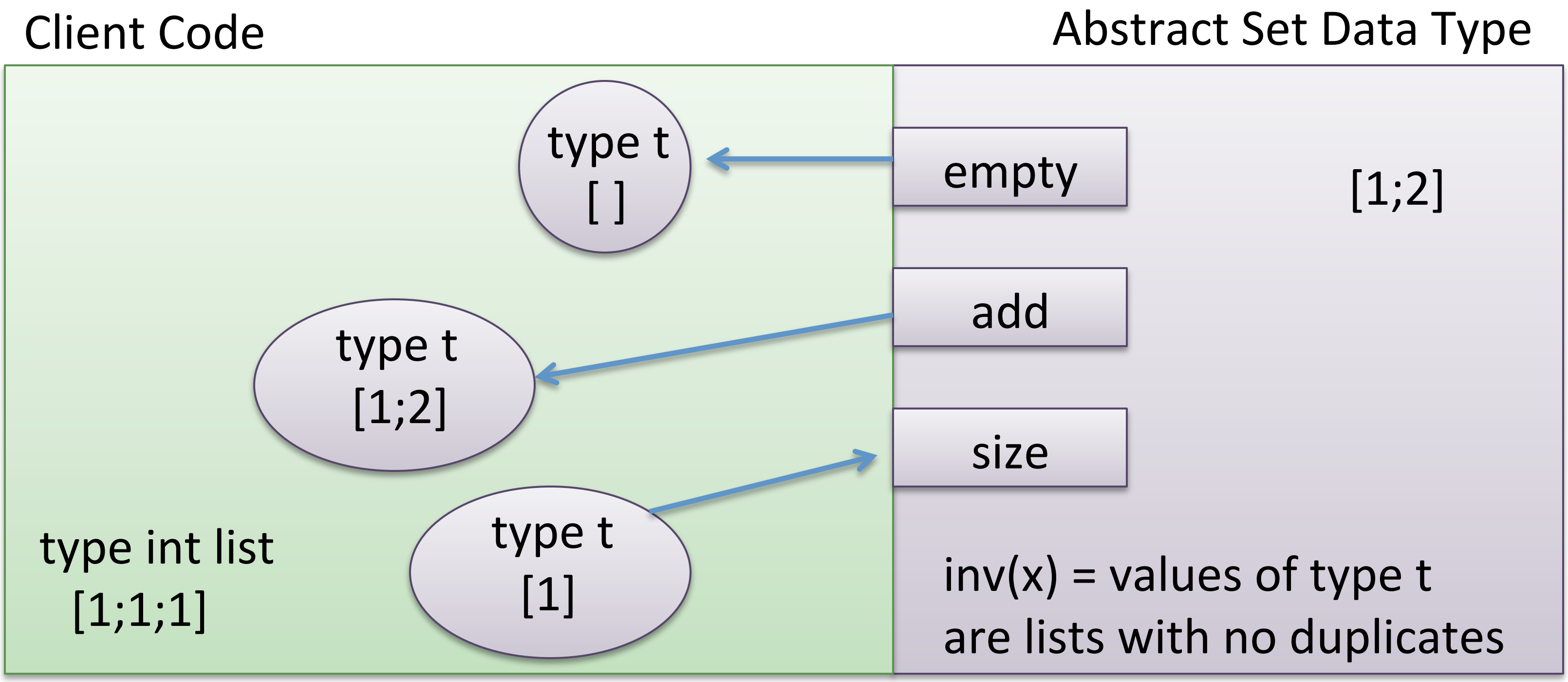
A *representation invariant inv(v) for abstract type t* is a property of all data values v with abstract data type t



- Invariants on abstract types are *local* to the ADT because they talk about the representation. Client code doesn't know or care what the invariant is.
- However, client code *preserves the invariant* because it can't mess with values of abstract type directly.

Representation Invariants

A *representation invariant inv(v) for abstract type t* is a property of all data values v with abstract data type t



- Because Clients can't mess with the invariants on abstract types, ADT code gets to *assume the invariant for inputs with abs. type* provided it *proves the invariant for outputs with abs. type*
- These proofs are *modular*: Done in isolation in the ADT module

Establishing Representation Invariants

E.g., when it comes to the size function:

```
(* signature *)
val size : 'a set -> int

(* implementation: length is # of distinct elements *)
let size l = List.length l
```

If we want to assume all arguments to size have no duplicates, then:

- we have to ensure that our client can only pass us a list with no dups
- clients get their values of type 'a set from our module, hence we have to ensure other functions in our module only produce lists with no duplicates
  - empty, add, rem, union, intersect
- typically the proof that a function produces elements that satisfy inv depend on assumptions that function inputs satisfy inv
  - add, rem, union, intersect

PROVING THE REP INVARIANT FOR THE SET ADT

Representation Invariants

Representation Invariant for sets without duplicates:

```
let rec inv (l : 'a set) : bool =
    match l with
      [] -> true
    | hd::tail -> not (mem hd tail) && inv tail
```

Definition of empty:

```
let empty : ‘a set = []
```

Proof Obligation:

```
inv (empty) == true
```

Proof:

```
   inv (empty)
== inv []
== match [] with [] -> true | hd::tail -> ...
== true
```

Representation Invariants

Representation Invariant for sets without duplicates:

```
let rec inv (l : 'a set) : bool =
    match l with
      [] -> true
    | hd::tail -> not (mem hd tail) && inv tail
```

Checking add:

```
let add (x:'a) (l:'a set) : 'a set =
  if mem x l then l else x::l
```

Proof obligation:

for all x:’a and for all l:’a set,

if inv(l) then inv (add x l)

prove invariant on output

assume invariant on input

Representation Invariants

```
let rec inv (l : 'a
    match l with
      [] -> true
    | hd::tail -> not (mem hd tail) && inv tail
```

```
let add (x:'a) (l:'a set) : 'a set =
  if mem x l then l else x::l
```

Theorem: for all x:'a and for all l:'a set, if inv(l) then inv (add x l)

Proof:

(1) pick an arbitrary x and l. (2) assume inv(l).

Break in to two cases:

-- one case when mem x l is true

-- one case where mem x l is false

Representation Invariants

```
let rec inv (l : 'a
    match l with
       [] -> true
     | hd::tail -> not (mem hd tail) && inv tail
```

```
let add (x:'a) (l:'a set) : 'a set =
   if mem x l then l else x::l
```

Theorem: for all x:'a and for all l:'a set, if inv(l) then inv (add x l)

Proof:

(1) pick an arbitrary x and l. (2) assume inv(l).

case 1: assume (3): mem x l == true:

inv (add x l)	
== inv (if mem x l then l else x::l)	(eval)
== inv (l)	(by (3), eval)
== true	(by (2))

Representation Invariants

```
let rec inv (l : 'a set) : bool =
  match l with
    [] -> true
  | hd::tail -> not (mem hd tail) && inv tail
```

```
let add (x:'a) (l:'a set) : 'a set =
  if mem x l then l else x::l
```

Theorem: for all x:'a and for all l:'a set, if inv(l) then inv (add x l)

Proof:

(1) pick an arbitrary x and l. (2) assume inv(l).

case 2: assume (3) not (mem x l) == true:

inv (add x l)	
== inv (if mem x l then l else x::l)	(eval)
== inv (x::l)	(by (3))
== not (mem x l) && inv (l)	(by eval)
== true && inv(l)	(by (3))
== true && true	(by (2))
== true	(eval)

Representation Invariants

Representation Invariant for sets without duplicates:

```
let rec inv (l : 'a set) : bool =
    match l with
      [] -> true
    | hd::tail -> not (mem hd tail) && inv tail
```

Checking rem:

```
let rem (x:'a) (l:'a set) : 'a set =
  List.filter ((<>) x) l
```

Proof obligation?

for all x:'a and for all l:'a set,

if inv(l) then inv (rem x l)

prove invariant on output

assume invariant on input

Representation Invariants

Representation Invariant for sets without duplicates:

```
let rec inv (l : 'a set) : bool =
    match l with
      [] -> true
    | hd::tail -> not (mem hd tail) && inv tail
```

Checking size:

```
let size (l:'a set) : int =
  List.length l
```

Proof obligation?

no obligation – does not produce value with type ‘a set

Representation Invariants

Representation Invariant for sets without duplicates:

```
let rec inv (l : 'a set) : bool =
    match l with
      [] -> true
    | hd::tail -> not (mem hd tail) && inv tail
```

Checking union:

```
let union (l1:'a set) (l2:'a set) : 'a set =
  ...
```

Proof obligation?

for all l1:'a set and for all l2:'a set,

if inv(l1) and inv(l2) then inv (union l1 l2)

assume invariant on input

prove invariant on output

Representation Invariants

Representation Invariant for sets without duplicates:

```
let rec inv (l : 'a set) : bool =
    match l with
      [] -> true
    | hd::tail -> not (mem hd tail) && inv tail
```

Checking inter:

```
let inter (l1:'a set) (l2:'a set) : 'a set =
  ...
```

Proof obligation?

for all l1:’a set and for all l2:’a set,

if inv(l1) and inv(l2) then inv (inter l1 l2)

assume invariant on input

prove invariant on output

Representation Invariants: a Few Types

- Given a module with abstract type t
- Define an invariant Inv(x)
- Assume arguments to functions satisfy Inv
- Prove results from functions satisfy Inv

```
sig
  type t

  val value : t

  val constructor : int -> t

  val transform : int -> t -> t

  val destructor : t -> int

end
```

- val value : t — prove: Inv (value)
- val constructor : int -> t — prove: for all x:int, Inv (constructor x)
- val transform : int -> t -> t — prove:
  for all x:int,
  for all v:t,
  if Inv(v)
  then Inv (transform x v)
- val destructor : t -> int — assume Inv(t))

REPRESENTATION INVARIANTS FOR HIGHER TYPES

Representation Invariants: More Types

What about more complex types?

eg: for abstract type t, consider: val op : t * t -> t option

Basic concept: Assume arguments are "valid"; Prove results "valid"

We know what it means to be a "valid" value v for abstract type t:

- Inv(v) must be true

What is a valid pair? v is valid for type s1 * s2 if

- (1) fst v is valid for type s1, and
- (2) snd v is valid for type s2

Equivalently: (v1, v2) is valid for type s1 * s2 if

- (1) v1 is valid for type s1, and
- (2) v2 is valid for type s2

Representation Invariants: More Types

What is a valid pair? v is valid for type s1 * s2 if

- (1) fst v is valid for s1, and
- (2) snd v is valid for s2

eg: for abstract type t, consider: val op : t * t -> t

must prove to establish rep invariant:
  for all x : t * t,
    if Inv(fst x) and Inv(snd x) then
    Inv (op x)

Equivalent Alternative:

must prove to establish rep invariant:
  for all x1:t, x2:t
    if Inv(x1) and Inv(x2) then
    Inv (op (x1, x2))

Representation Invariants: More Types

Another Example:

```
val v : t * (t -> t)
```

must prove both to satisfy the rep invariant:

(1) valid (fst v) for type t:
    ie: inv (fst v)

(2) valid (snd v) for type t -> t:
    ie: for all v1:t,
        if Inv(v1) then
        Inv ((snd v) v1)

Representation Invariants: More Types

What is a valid option? v is valid for type s1 option if

- (1) v is None, or
- (2) v is Some u, and u is valid for type s1

eg: for abstract type t, consider: val op : t * t -> t option

```
must prove to satisfy rep invariant:
 for all x : t * t,
    if Inv(fst x) and Inv(snd x)
    then
       either:
          (1) op x is None or
          (2) op x is Some u and Inv u
```

Representation Invariants: More Types

Suppose we are defining an abstract type t.

Consider happens when the type int shows up in a signature.

The type int does not involve the abstract type t at all, in any way.

eg: in our set module, consider: val size : t -> int

When is a value v of type int valid?

all values v of type int are valid

val size : t -> int ← must prove nothing

val const : int ← must prove nothing

val create : int -> t ← for all v:int, assume nothing about v, must prove Inv (create v)

Representation Invariants: More Types

What is a valid function? Value f is valid for type t1 -> t2 if

- for all inputs arg that are valid for type t1,
- it is the case that f arg is valid for type t2

eg: for abstract type t, consider: val op : t * t -> t option

```
must prove to satisfy rep invariant:
 for all x : t * t,
    if Inv(fst x) and Inv(fst x)
    then
       either:
          (1) op x is None or
          (2) op x is Some u and Inv u
```

valid for type t * t (the argument) — points to "if Inv(fst x) and Inv(fst x)"

valid for type t option (the result) — points to "(1) op x is None or (2) op x is Some u and Inv u"

Representation Invariants: More Types

What is a valid function? Value f is valid for type t1 -> t2 if

- for all inputs arg that are valid for type t1,
- it is the case that f arg is valid for type t2

eg: for abstract type t, consider: val op : (t -> t) -> t

must prove to satisfy rep invariant:

```
 for all x : t -> t,
    if
      {for all arguments arg:t,
         if Inv(arg) then Inv(x arg) }
    then
       Inv (op x)
```

valid for type t -> t (the argument)

valid for type t (the result)

Representation Invariants: More Types

```
sig
  type t
  val create : int -> t
  val incr : t -> t
  val apply : t * (t -> t) -> t
  val check_t : t -> t
end
```

```
struct
  type t = int
  let create n = abs n
  let incr n = if n<maxint then n + 1
               else raise Overflow
  let apply (x, f) = f x
  let check_t x = assert (x >= 0); x
end
```

representation invariant:
let inv x = x >= 0

function apply, must prove:
  for all x:t,
  for all f:t -> t
    if x valid for t
    and f valid for t -> t
    then f x valid for t

function apply, must prove:
  for all x:t,
  for all f:t -> t
    if (1) inv(x)
    and (2) for all y:t, if inv(y) then inv(f y)
    then inv(f x)

Proof: By (1) and (2), inv(f x)

ANOTHER EXAMPLE

Natural Numbers

```
module type NAT =
  sig

    type t

    val from_int : int -> t

    val to_int : t -> int

    val map : (t -> t) -> t -> t list

 end
```

Natural Numbers

```
module type NAT =
  sig

    type t

    val from_int : int -> t

    val to_int : t -> int

    val map : (t -> t) -> t -> t list

 end
```

```
module Nat : NAT =
  struct

    type t = int

    let from_int (n:int) : t =
      if n <= 0 then 0 else n

    let to_int (n:t) : int = n

    let rec map f n =
      if n = 0 then []
      else f n :: map f (n-1)

end
```

Natural Numbers

```
module type NAT =
  sig

    type t

    val from_int : int -> t

    val to_int : t -> int

    val map : (t -> t) -> t -> t list

 end
```

```
let inv n : bool =
  n >= 0
```

```
module Nat : NAT =
  struct

    type t = int

    let from_int (n:int) : t =
      if n <= 0 then 0 else n

    let to_int (n:t) : int = n

    let rec map f n =
      if n = 0 then []
      else f n :: map f (n-1)

end
```

Look to the signature to figure out what to verify

```
module type NAT =
 sig

   type t

   val from_int : int -> t

   val to_int : t -> int

   val map : (t -> t) -> t -> t list

 end
```

```
type t = int
```

since function result has type t, must prove the output satisfies inv()

can assume inv(x) for all inputs; don't need to prove anything of the outputs with type int

for map f x, assume:
(1) inv(x), and
(2) f's results satisfy inv() when it's inputs satisfy inv().

then prove that all elements of the output list satisfy inv()

```
let inv n : bool =
  n >= 0
```

Verifying The Invariant

In general, we use a type-directed proof methodology:

- Let t be the abstract type and inv() the representation invariant
- For each value v with type s in the signature, we must check that v is valid for type s as follows:
  - v is valid for t if
    - inv(v)
  - (v1, v2) is valid for s1 * s2 if
    - v1 is valid for s1, and
    - v2 is valid for s2
  - v is valid for type s option if
    - v is None or,
    - v is Some u and u is valid for type s
  - v is valid for type s1 -> s2 if
    - for all arguments a, if a is valid for s1, then v a is valid for s2
  - v is valid for int if
    - always
  - [v1; ...; vn] is valid for type s list if
    - v1 ... vn are all valid for type s

Natural Numbers

```
module type NAT =
 sig

   type t

   val from_int : int -> t

   ...

 end
```

```
module Nat : NAT =
 struct

   type t = int

   let from_int (n:int) : t =
     if n <= 0 then 0 else n

   ...

end
```

```
let inv n : bool =
  n >= 0
```

Must prove:

```
for all n,
  inv (from_int n) == true
```

Proof strategy: Split into 2 cases.
(1) n > 0, and (2) n <= 0

Natural Numbers

```
module type NAT =
 sig

   type t

   val from_int : int -> t

   ...

 end
```

Must prove:

```
for all n,
  inv (from_int n) == true
```

```
module Nat : NAT =
 struct

   type t = int

   let from_int (n:int) : t =
     if n <= 0 then 0 else n

   ...

end
```

```
let inv n : bool =
  n >= 0
```

Case: n > 0

```
   inv (from_int n)
== inv (if n <= 0 then 0 else n)
== inv n
== true
```

Natural Numbers

```
module type NAT =
  sig

    type t

    val from_int : int -> t

    ...

  end
```

Must prove:

```
for all n,
  inv (from_int n) == true
```

```
module Nat : NAT =
  struct

    type t = int

    let from_int (n:int) : t =
      if n <= 0 then 0 else n

    ...

  end
```

```
let inv n : bool =
  n >= 0
```

Case: n <= 0

```
     inv (from_int n)
== inv (if n <= 0 then 0 else n)
== inv 0
== true
```

Natural Numbers

```
module type NAT =
  sig

    type t

    val to_int : t -> int

    ...

  end
```

```
module Nat : NAT =
  struct

    type t = int

    let to_int (n:t) : int = n

    ...

end
```

```
let inv n : bool =
  n >= 0
```

Must prove:

for all n,
  if inv n then
  we must show ... nothing ...
  since the output type is int

Natural Numbers

```
module type NAT =
 sig

   type t

    val map : (t -> t) -> t -> t list

   ...


 end
```

```
module Nat : NAT =
 struct

   type t = int

    let rep map f n =
     if n = 0 then []
     else f n :: map f (n-1)

   ...
end
```

```
let inv n : bool =
  n >= 0
```

Must prove:

```
for all f valid for type t -> t
for all n valid for type t
  map f n is valid for type t list
```

Proof: By induction on n.

Natural Numbers

```
module type NAT =
 sig

  type t

   val map : (t -> t) -> t -> t list

  ...


 end
```

Must prove:

for all f valid for type t -> t
for all n valid for type t
  map f n is valid for type t list

Proof: By induction on nat n.

```
module Nat : NAT =
 struct

  type t = int

   let rep map f n =
    if n = 0 then []
    else f n :: map f (n-1)

  ...
end
```

```
let inv n : bool =
  n >= 0
```

Case: n = 0

map f n == []

(Note: each value v in [ ] satisfies inv(v))

Natural Numbers

```
module type NAT =
 sig

   type t

    val map : (t -> t) -> t -> t list

   ...

 end
```

Must prove:

```
for all f valid for type t -> t
for all n valid for type t
  map f n is valid for type t list
```

Proof: By induction on nat n.

```
module Nat : NAT =
 struct

   type t = int

    let rep map f n =
     if n = 0 then []
     else f n :: map f (n-1)

   ...
end
```

```
let inv n : bool =
  n >= 0
```

Case: n > 0

```
map f n  == f n :: map f (n-1)
```

Natural Numbers

```
module type NAT =
 sig

  type t

   val map : (t -> t) -> t -> t list

  ...

 end
```

Must prove:

```
for all f valid for type t -> t
for all n valid for type t
  map f n is valid for type t list
```

Proof: By induction on nat n.

```
module Nat : NAT =
 struct

  type t = int

   let rep map f n =
    if n = 0 then []
    else f n :: map f (n-1)

  ...
end
```

```
let inv n : bool =
  n >= 0
```

Case: n > 0

map f n == f n :: map f (n-1)

By IH, map f (n-1) is valid for t list.

Natural Numbers

```
module type NAT =
 sig

   type t

    val map : (t -> t) -> t -> t list

   ...

 end
```

Must prove:

> for all f valid for type t -> t
> for all n valid for type t
>   map f n is valid for type t list

Proof: By induction on nat n.

```
module Nat : NAT =
 struct

   type t = int

    let rep map f n =
     if n = 0 then []
     else f n :: map f (n-1)

   ...
end
```

```
let inv n : bool =
  n >= 0
```

Case: n > 0

map f n == f n :: map f (n-1)

By IH, map f (n-1) is valid for t list.
Since f valid for t -> t and n valid for t
  f n::map f (n-1) is valid for t list

Natural Numbers

```
module type NAT =
  sig

    type t

     val map : (t -> t) -> t -> t list

    ...

  end
```

```
module Nat : NAT =
  struct

    type t = int

     let rep map f n =
      if n = 0 then []
      else f n :: map f (n-1)

    ...
end
```

Hooray! n is never negative so we don't infinite loop

End result: We have proved a strong property (n >= 0) of every value with abstract type Nat.t

Summary for Representation Invariants

- The signature of the module tells you what to prove

- Roughly speaking:
  - assume invariant holds on values with abstract type *on the way in*
  - prove invariant holds on values with abstract type *on the way out*

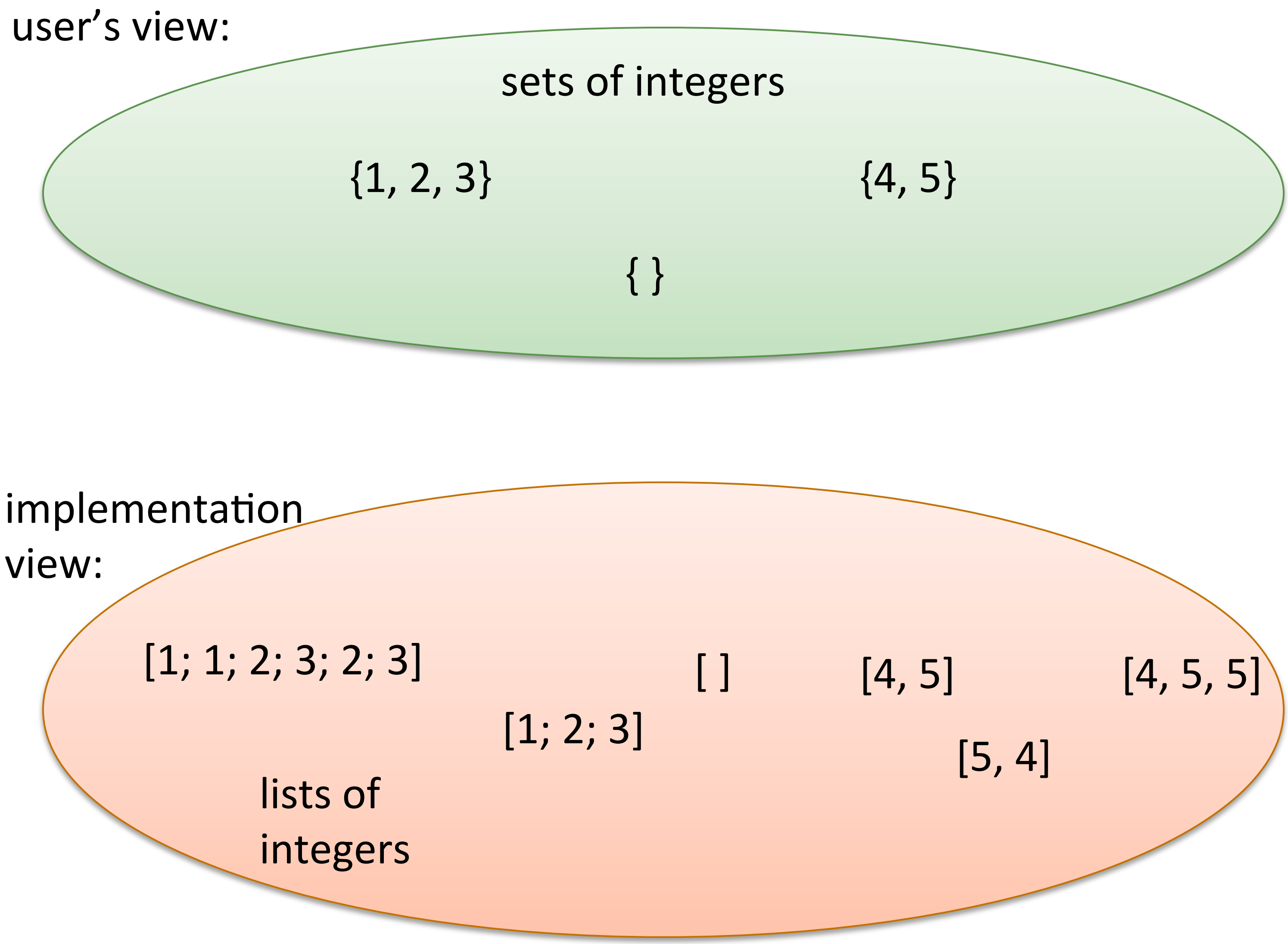
ABSTRACTION FUNCTIONS

Abstraction

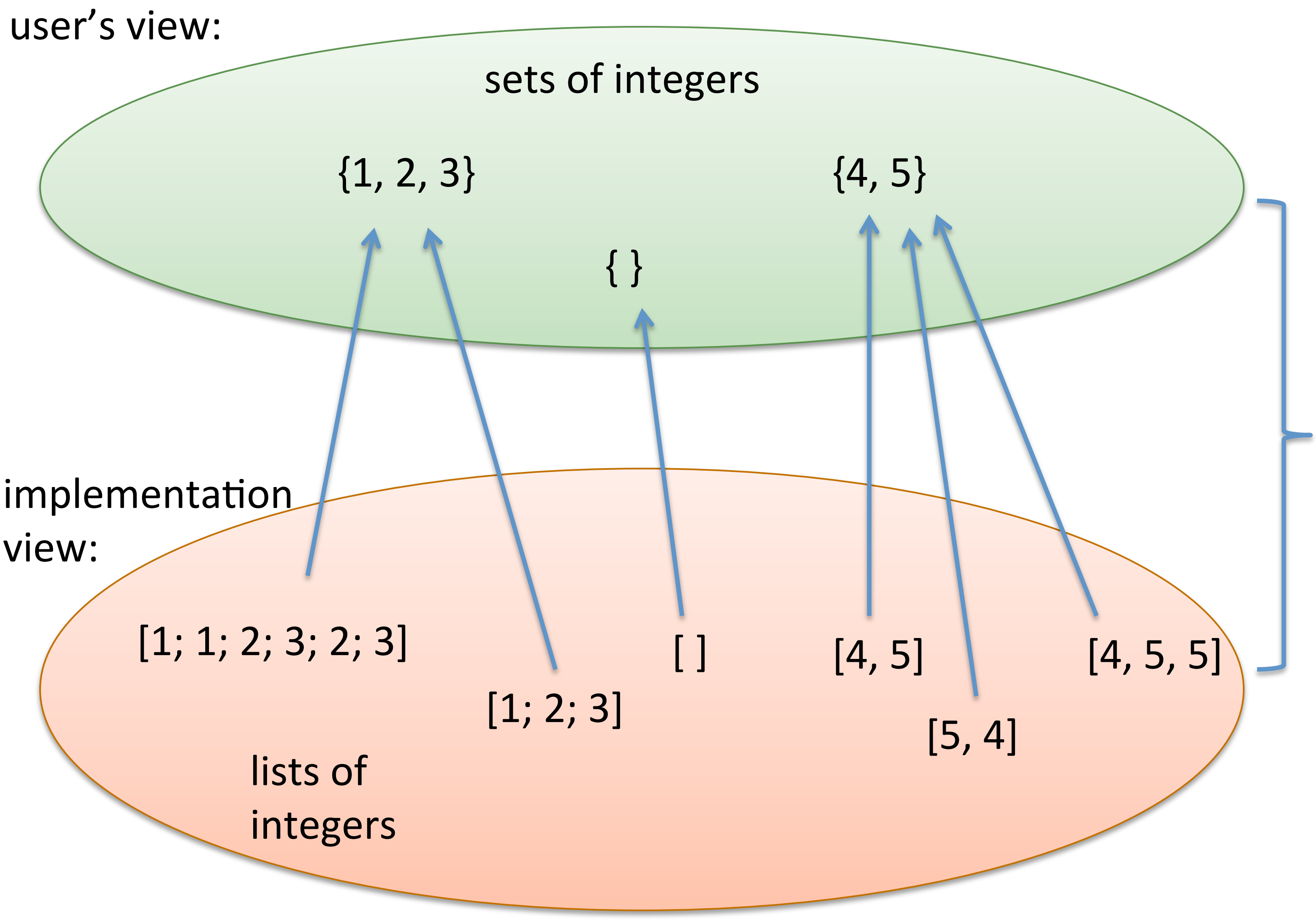
```
module type SET =
  sig
    type 'a set
    val empty : 'a set
    val mem : 'a -> 'a set -> bool
    ...
end
```

- When explaining our modules to clients, we would like to explain them in terms of *abstract values*
  - sets, not the lists (or maybe trees) that implement them
- From a client's perspective, operations act on abstract values
- Signature comments, specifications, preconditions and post-conditions in terms of those abstract values
- *How are these abstract values connected to the implementation?*

Abstraction



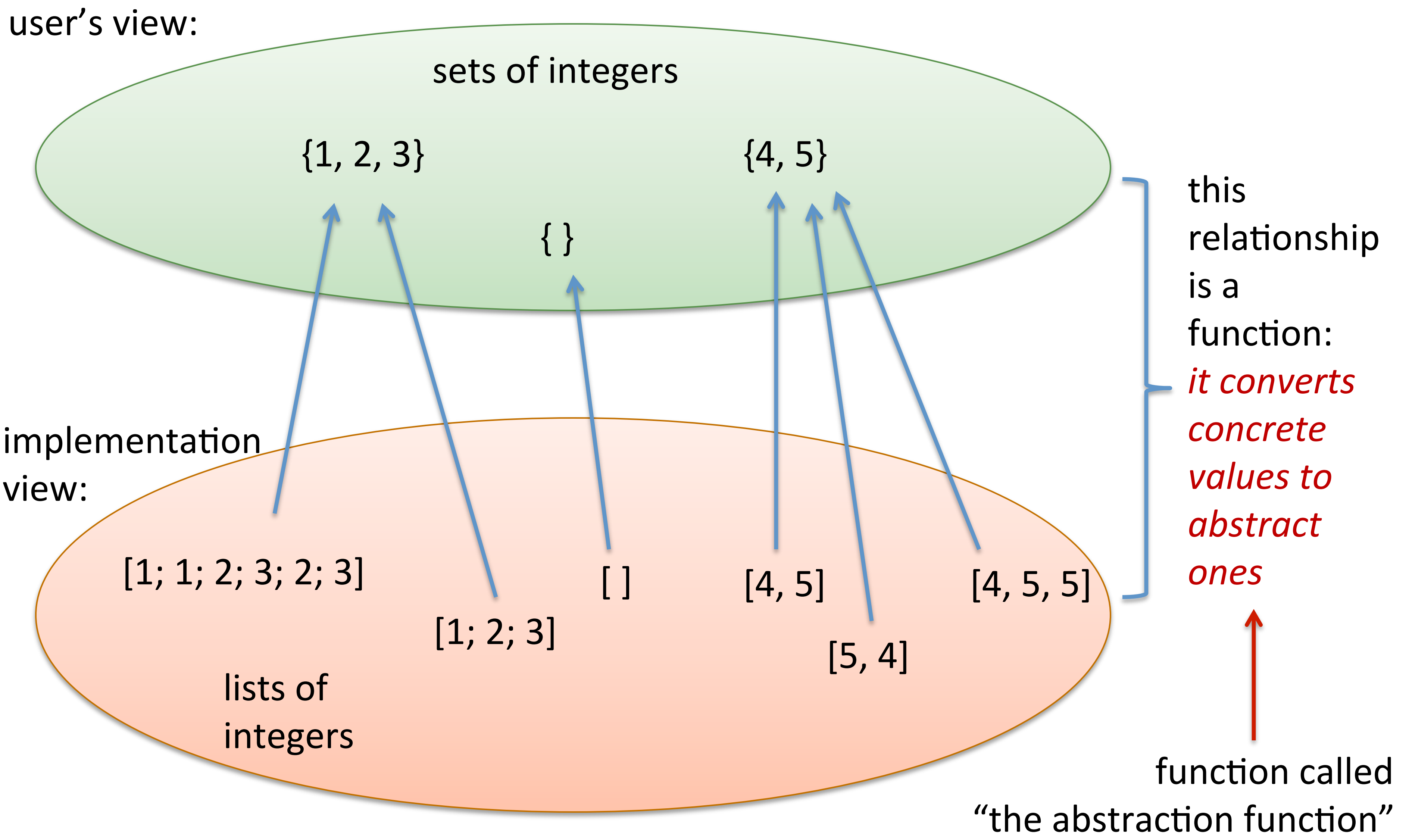
Abstraction



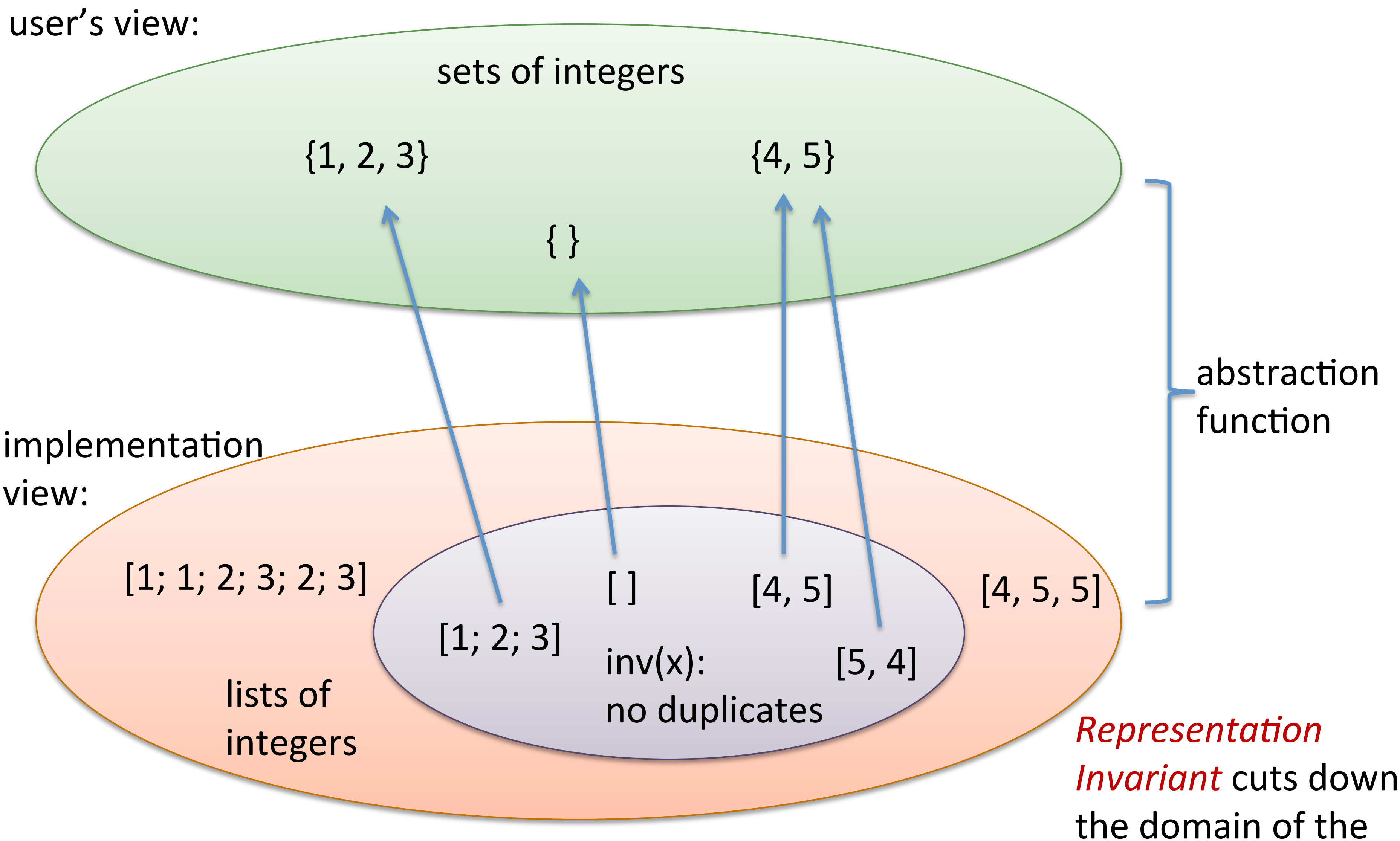
there's a relationship here, of course!

we are trying to *implement* the *abstraction*

Abstraction



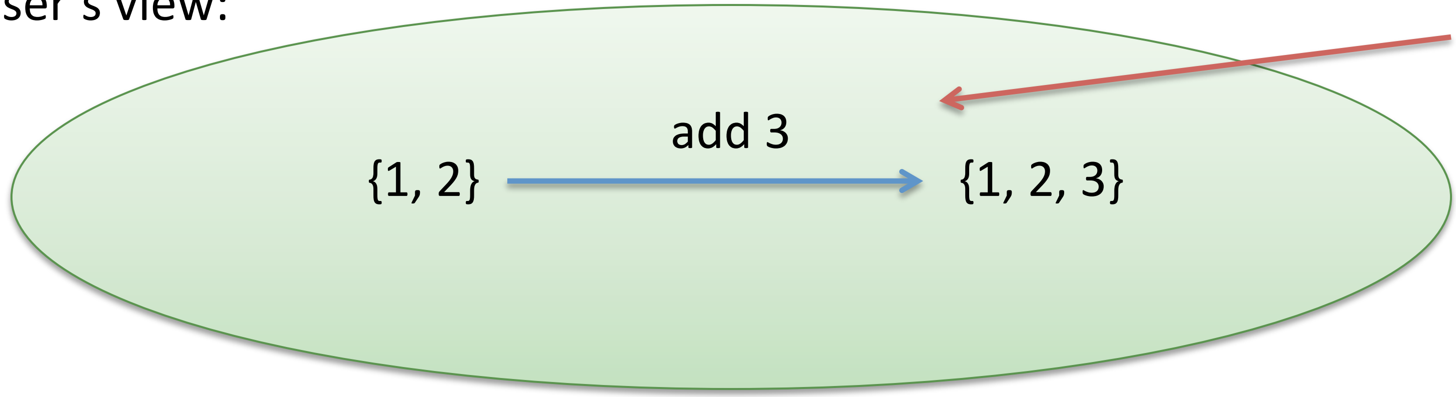
Abstraction



Representation Invariant cuts down the domain of the abstraction function

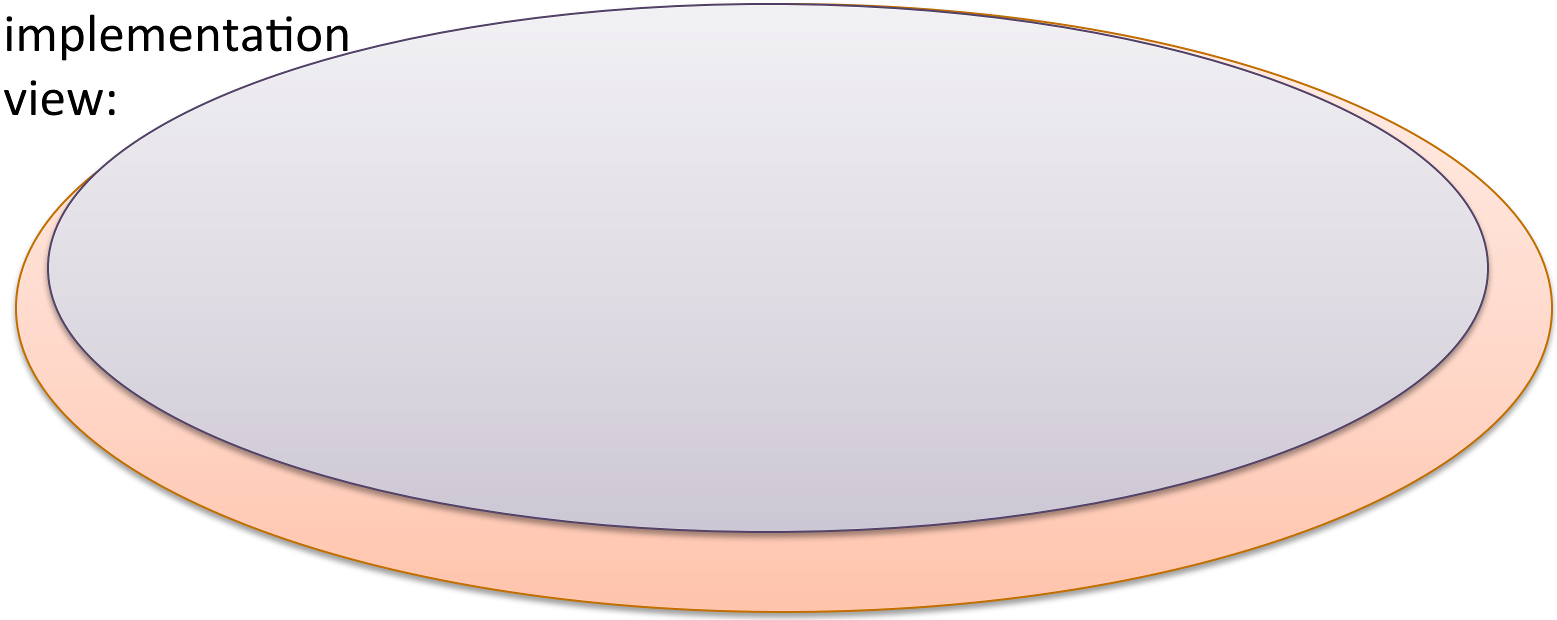
Specifications

user's view:

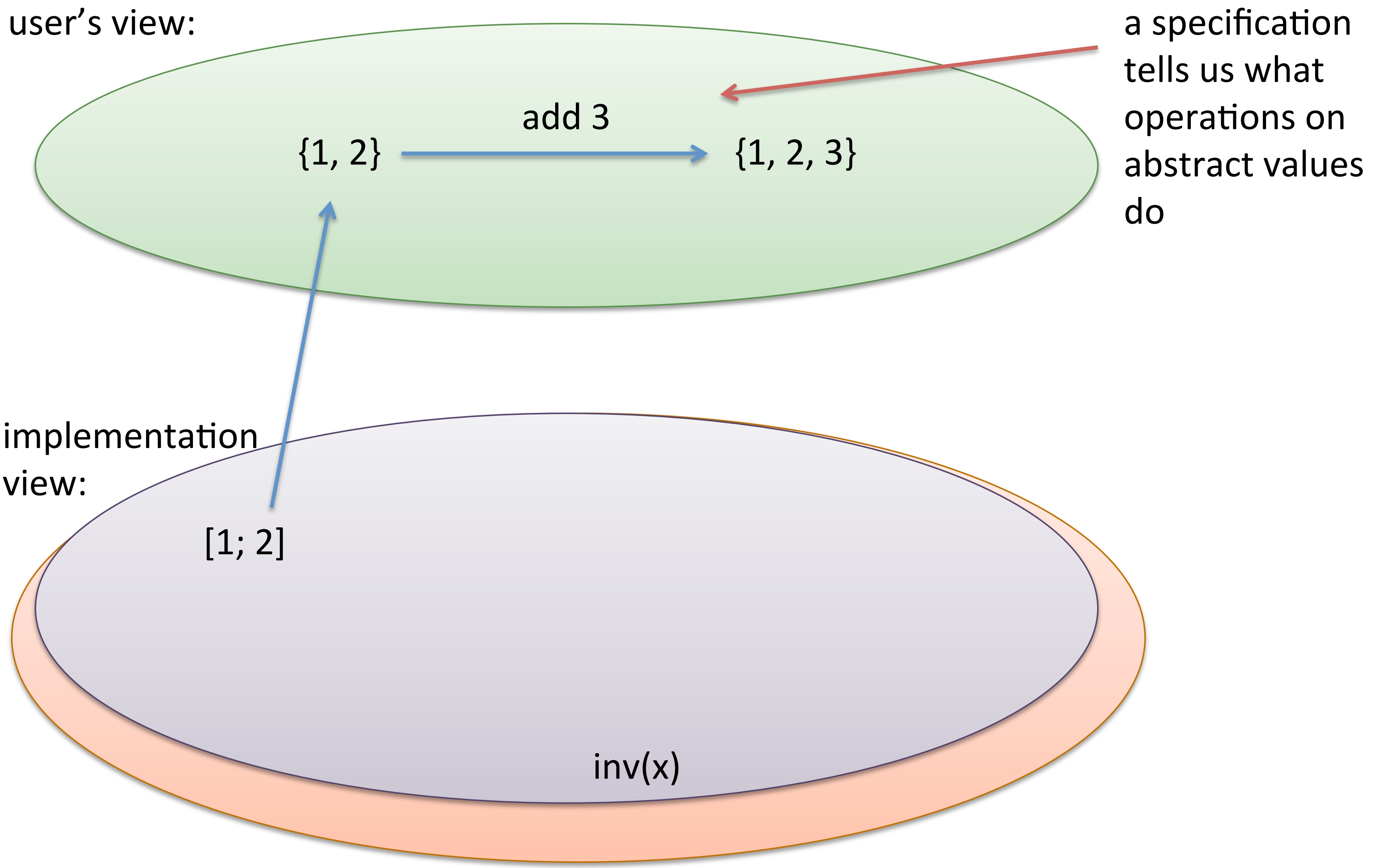


a specification tells us what operations on abstract values do

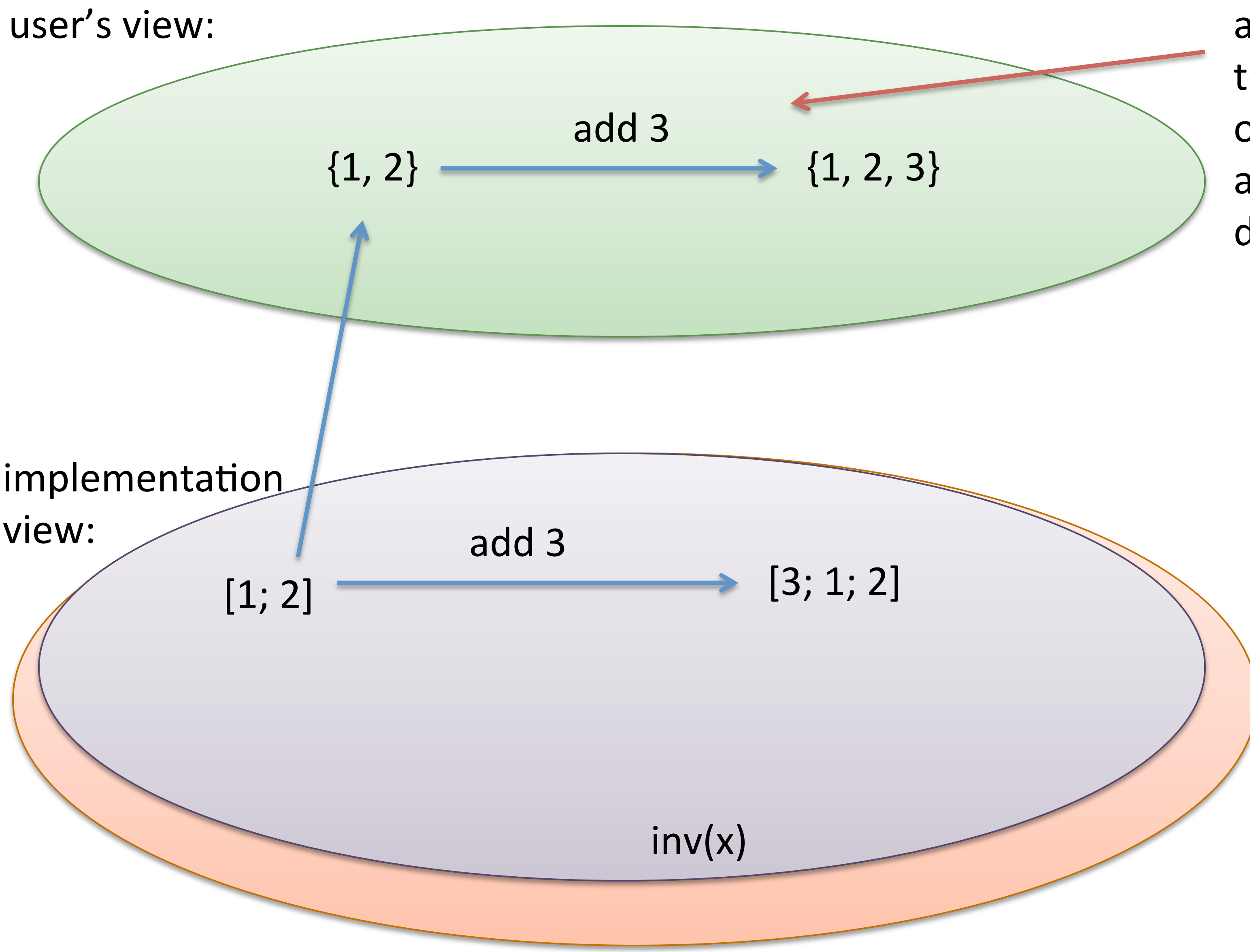
implementation view:



Specifications

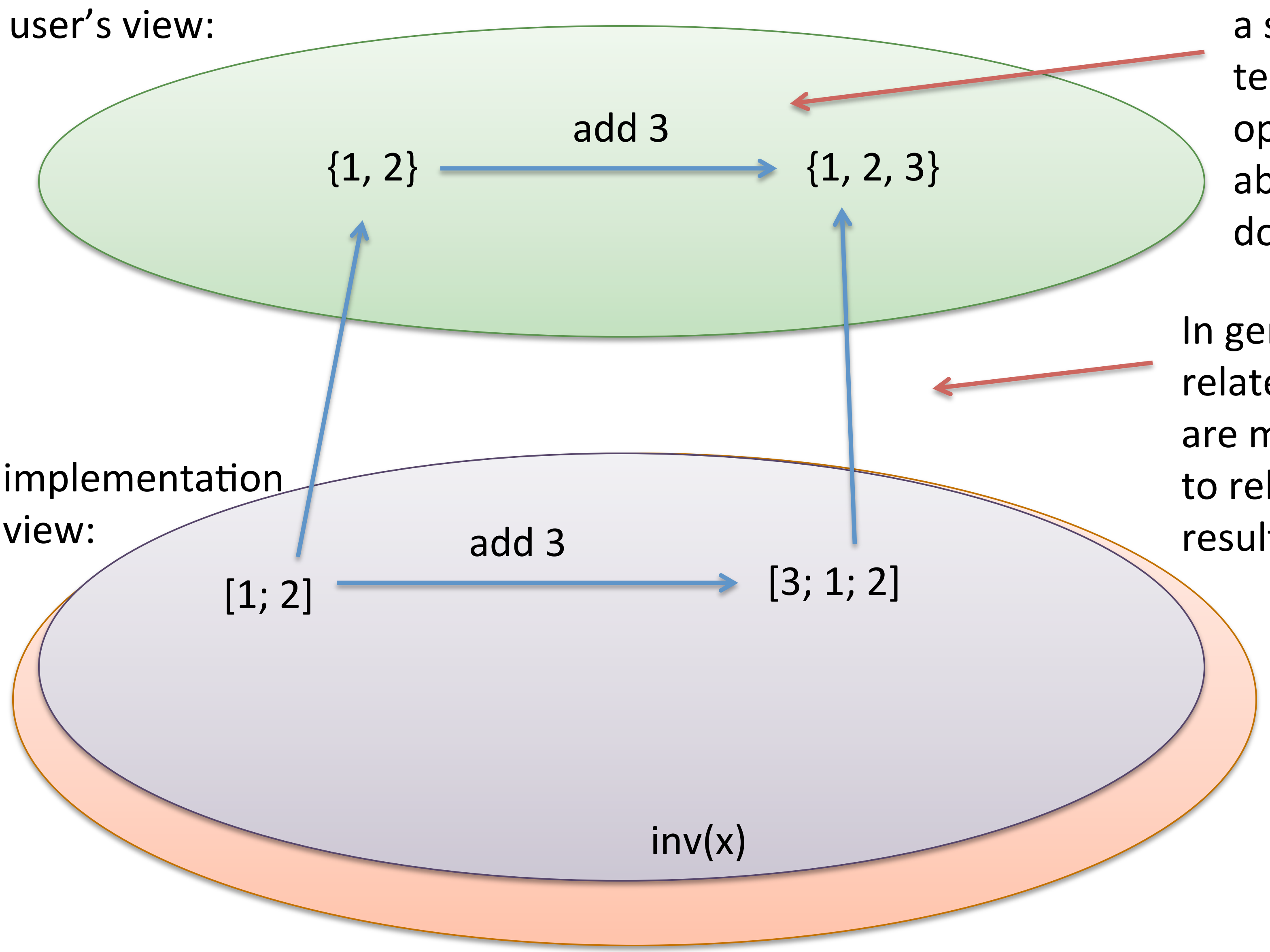


Specifications



a specification tells us what operations on abstract values do

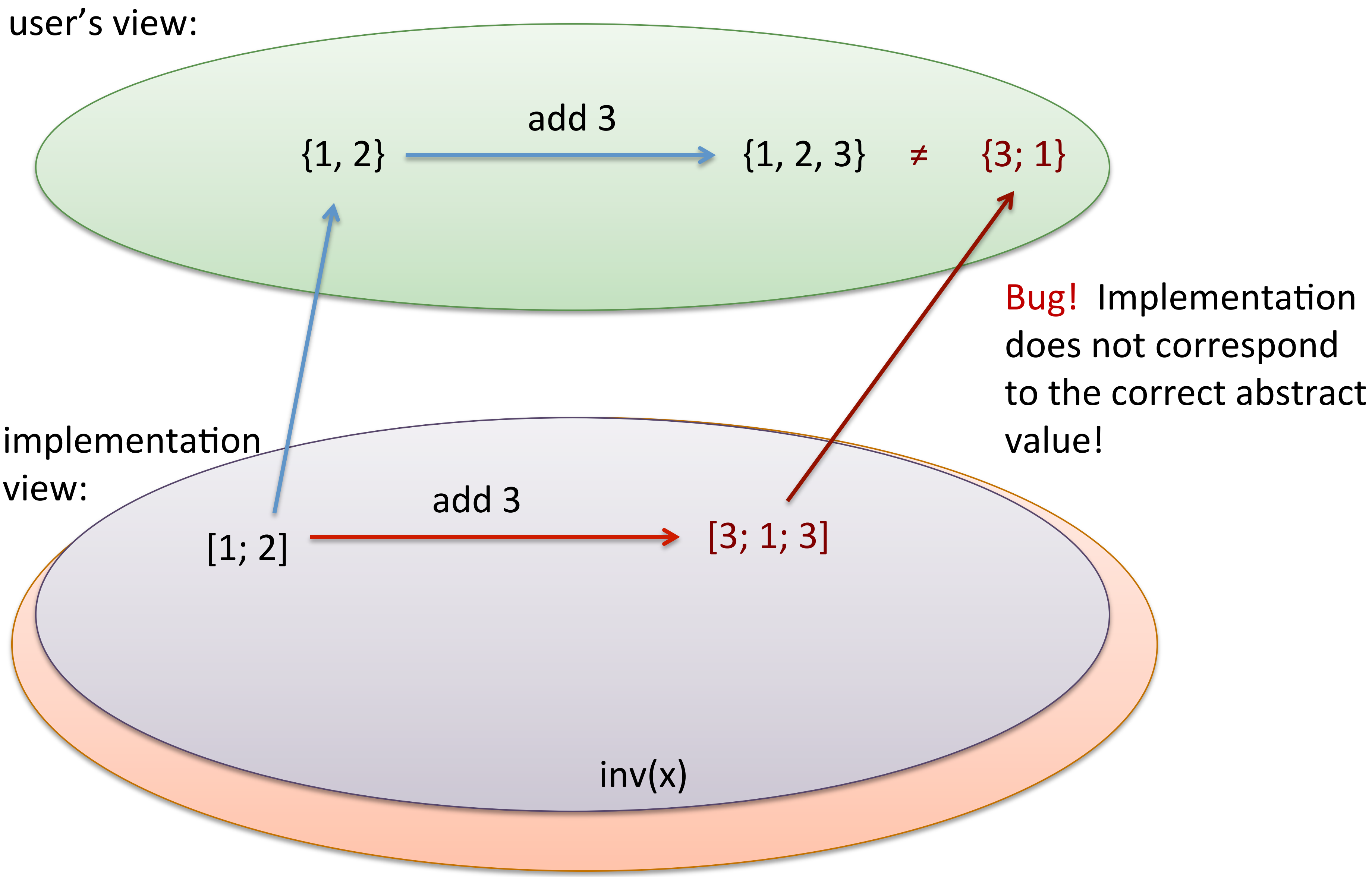
Specifications



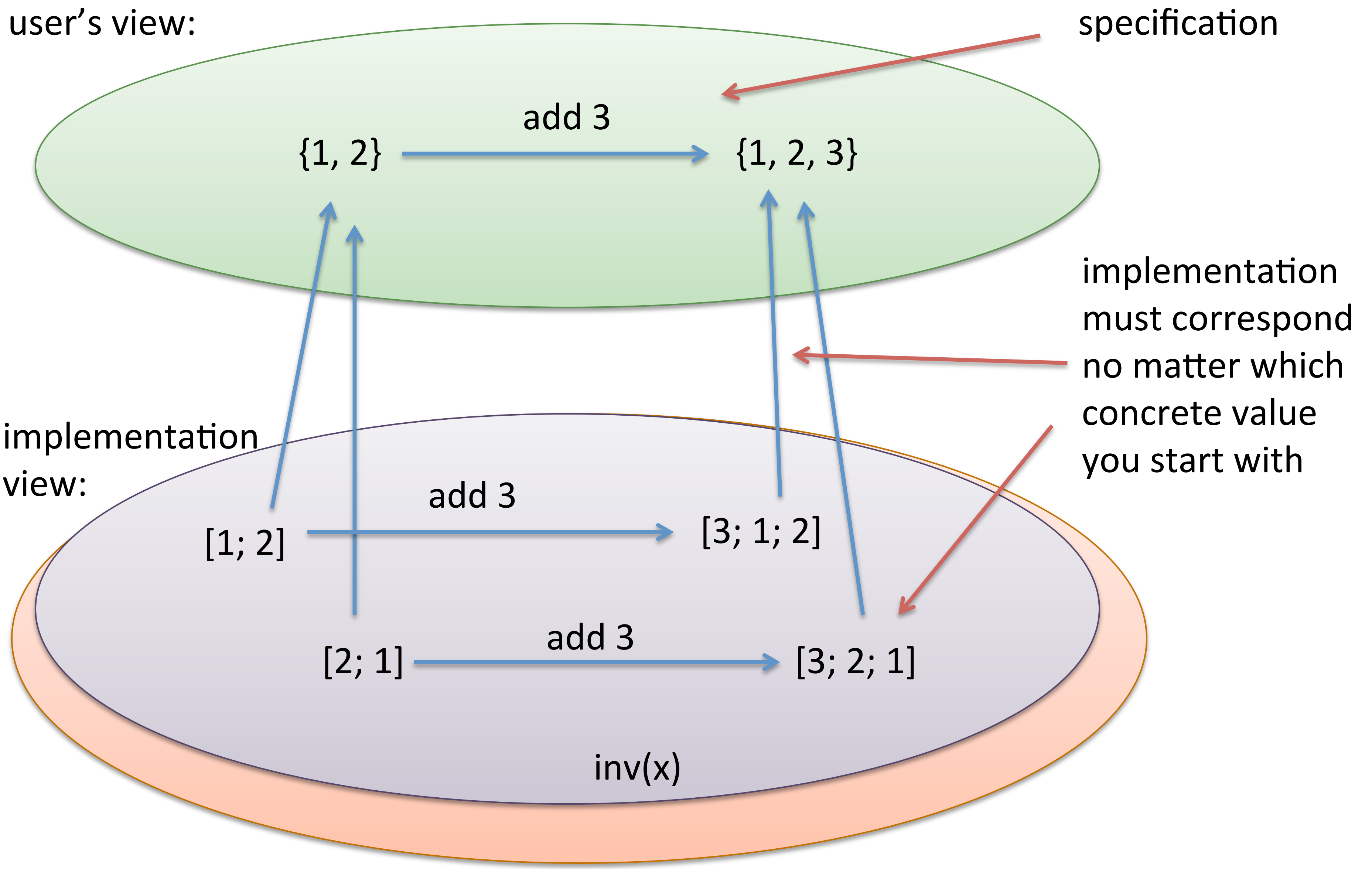
a specification tells us what operations on abstract values do

In general: related arguments are mapped to related results

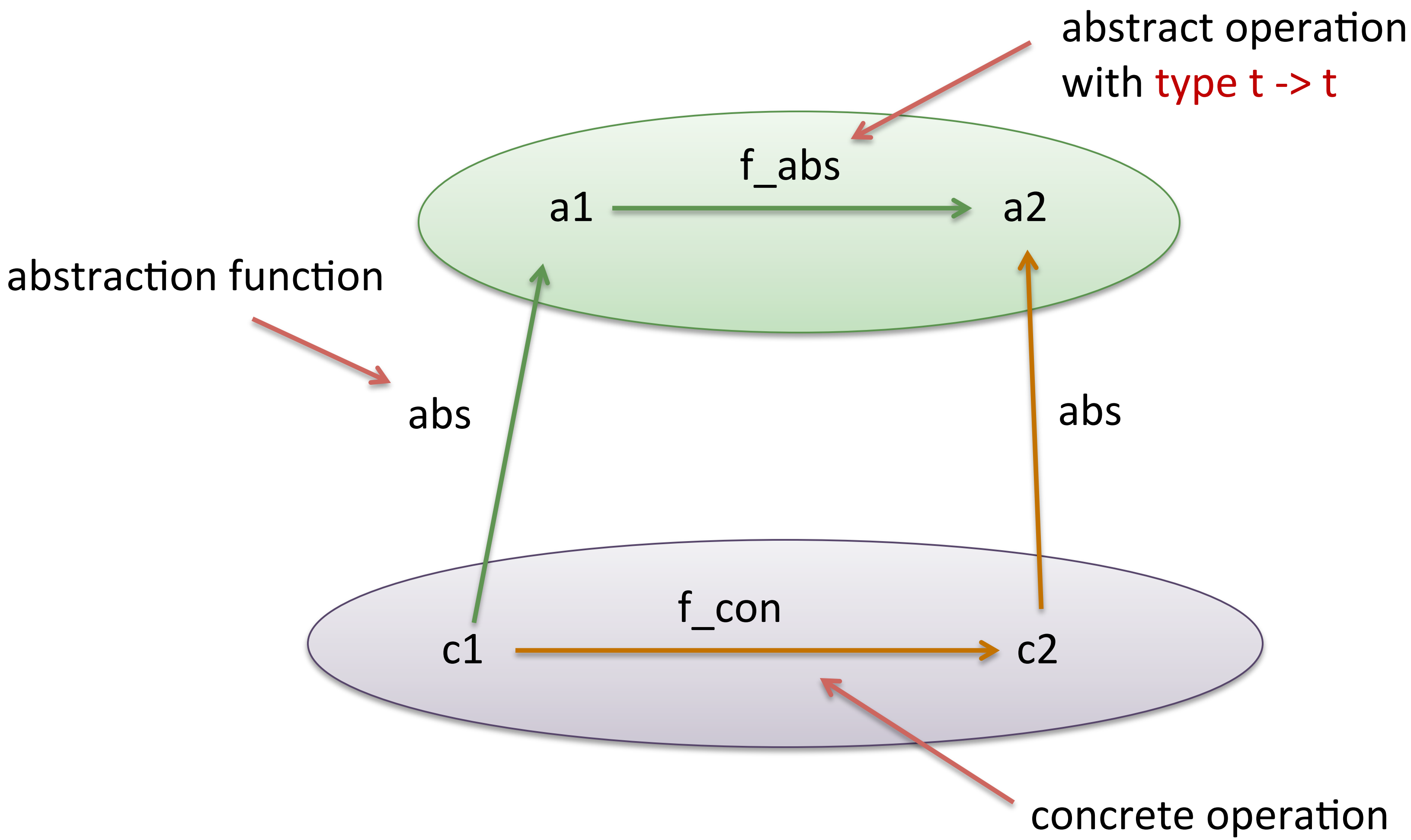
Specifications



Specifications



A more general view



to prove:
for all c1:t, if inv(c1) then f_abs (abs c1) == abs (f_con c1)

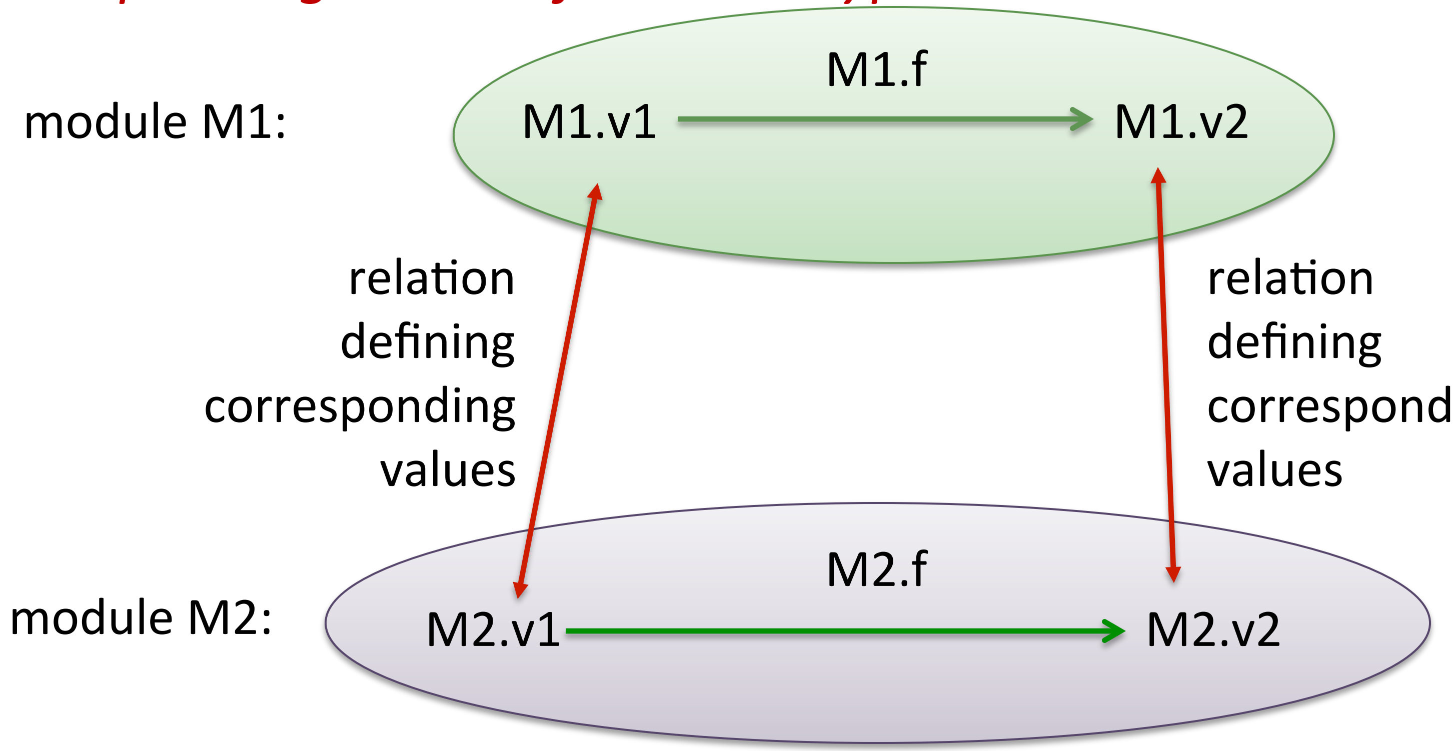
abstract then apply the abstract op == apply concrete op then abstract

Another Viewpoint

A specification is really just another implementation (in this viewpoint)

- but it's often simpler ("more abstract")

We can use similar ideas to compare *any two implementations of the same signature. Just come up with a relation between corresponding values of abstract type.*



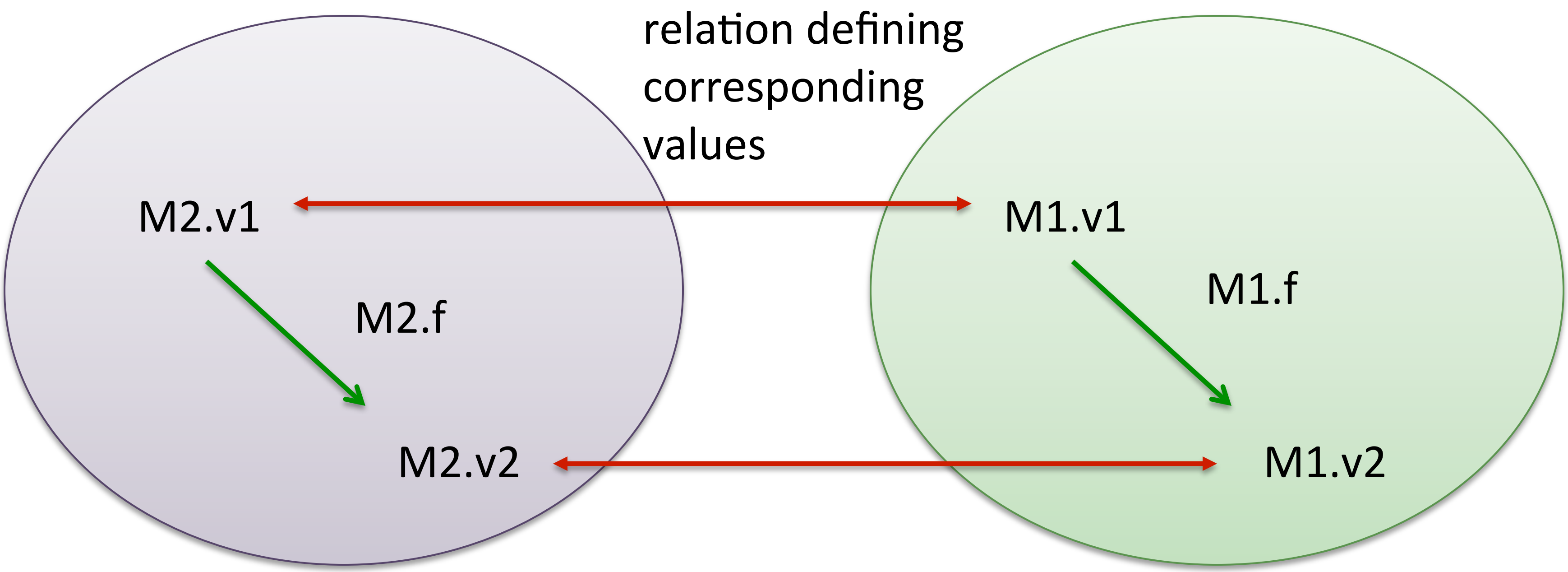
We ask: Do operations like f take related arguments to related results?

What is a specification?

It is really just another implementation

- but it's often simpler ("more abstract")

We can use similar ideas to compare *any two implementations of the same signature. Just come up with a relation between corresponding values of abstract type.*



One Signature, Two Implementations

```
module type S =
  sig
    type t
    val zero : t
    val bump : t -> t
    val reveal : t -> int
  end
```

```
module M1 : S =
  struct
    type t = int
    let zero = 0
    let bump n = n + 1
    let reveal n = n
  end
```

```
module M2 : S =
  struct
    type t = int
    let zero = 2
    let bump n = n + 2
    let reveal n = n/2 - 1
  end
```

Consider a client that might use the module:

```
let x1 = M1.bump (M1.bump (M1.zero)
```

```
let x2 = M2.bump (M2.bump (M2.zero)
```

What is the relationship?

```
is_related (x1, x2) =
  x1   ==   x2/2 - 1
```

And it persists: Any sequence of operations produces related results from M1 and M2!

How do we prove it?

One Signature, Two Implementations

```
module type S =
  sig
    type t
    val zero : t
    val bump : t -> t
    val reveal : t -> int
  end
```

```
module M1 : S =
  struct
    type t = int
    let zero = 0
    let bump n = n + 1
    let reveal n = n
  end
```

```
module M2 : S =
  struct
    type t = int
    let zero = 2
    let bump n = n + 2
    let reveal n = n/2 - 1
  end
```

Recall: A representation invariant is a property that holds for all values of abs. type:

- if M.v has abstract type t,
  - we want inv(M.v) to be true

Inter-module relations are a lot like representation invariants!

- if M1.v and M2.v have abstract type t,
  - we want is_related(M1.v, M2.v) to be true

It's just a relation between two modules instead of one

One Signature, Two Implementations

```
module type S =
  sig
    type t
    val zero : t
    val bump : t -> t
    val reveal : t -> int
  end
```

```
module M1 : S =
  struct
    type t = int
    let zero = 0
    let bump n = n + 1
    let reveal n = n
  end
```

```
module M2 : S =
  struct
    type t = int
    let zero = 2
    let bump n = n + 2
    let reveal n = n/2 - 1
  end
```

Recall: To prove a rep. inv., assume it holds on inputs & prove it holds on outputs:

- if M.f has type t -> t, we prove that:
  - if inv(v) then inv(M.f v)

Likewise for inter-module relations:

- if M1.f and M2.f have type t -> t, we prove that:
  - if is_related(v1, v2) then
  - is_related(M1.f v1, M2.f v2)

related functions produce related results from related arguments

One Signature, Two Implementations

```
module type S =
  sig
    type t
    val zero : t
    val bump : t -> t
    val reveal : t -> int
  end
```

```
module M1 : S =
  struct
    type t = int
    let zero = 0
    let bump n = n + 1
    let reveal n = n
  end
```

```
module M2 : S =
  struct
    type t = int
    let zero = 2
    let bump n = n + 2
    let reveal n = n/2 - 1
  end
```

Consider zero, which has abstract type t.

```
is_related (x1, x2) =
  x1   ==   x2/2 - 1
```

Must prove: is_related (M1.zero, M2.zero)

Equivalent to proving: M1.zero == M2.zero/2 – 1

Proof:

```
   M1.zero
== 0                  (substitution)
== 2/2 – 1            (math)
== M2.zero/2 – 1      (subsitution)
```

One Signature, Two Implementations

```
module type S =
  sig
    type t
    val zero : t
    val bump : t -> t
    val reveal : t -> int
  end
```

```
module M1 : S =
  struct
    type t = int
    let zero = 0
    let bump n = n + 1
    let reveal n = n
  end
```

```
module M2 : S =
  struct
    type t = int
    let zero = 2
    let bump n = n + 2
    let reveal n = n/2 - 1
  end
```

```
is_related (x1, x2) =
  x1   ==   x2/2 - 1
```

Consider bump, which has abstract type t -> t.

Must prove for all v1:int, v2:int
if is_related(v1,v2) then is_related (M1.bump v1, M2.bump v2)

Proof:
(1) Assume is_related(v1, v2).
(2) v1 == v2/2 – 1 (by def)

Next, prove:
(M2.bump v2)/2 – 1 == M1.bump v1

(M2.bump v2)/2 - 1	
== (v2 + 2)/2 – 1	(eval)
== (v2/2 – 1) + 1	(math)
== v1 + 1	(by 2)
== M1.bump v1	(eval, reverse)

One Signature, Two Implementations

```
module type S =
  sig
    type t
    val zero : t
    val bump : t -> t
    val reveal : t -> int
 end
```

```
module M1 : S =
  struct
    type t = int
    let zero = 0
    let bump n = n + 1
    let reveal n = n
 end
```

```
module M2 : S =
  struct
    type t = int
    let zero = 2
    let bump n = n + 2
    let reveal n = n/2 - 1
 end
```

```
is_related (x1, x2) =
  x1   ==   x2/2 - 1
```

Consider reveal, which has abstract type t -> int.

Must prove for all v1:int, v2:int
if is_related(v1,v2) then M1.reveal v1 == M2.reveal v2

Proof:
(1) Assume is_related(v1, v2).
(2) v1 == v2/2 – 1 (by def)

Next, prove:
(M2.reveal v2 == M1.reveal v1

(M2.reveal v2)	
== v2/2 – 1	(eval)
== v1	(by 2)
== M1.reveal v1	(eval, reverse)

Summary of Proof Technique

To prove M1 == M2 relative to signature S,

- Start by defining a relation "is_related":
  - is_related (v1, v2) should hold for values with abstract type t when v1 comes from module M1 and v2 comes from module M2

- Extend "is_related" to types other than just abstract t. For example:
  - if v1, v2 have type int, then they must be exactly the same
    - ie, we must prove: v1 == v2
  - if v1, v2 have type s1 -> s2 then we consider arg1, arg2 such that:
    - if is_related(arg1, arg2) then we prove
    - is_related(v1 arg1, v2 arg2)
  - if v1, v2 have type s option then we must prove:
    - v1 == None and v2 == None, or
    - v1 == Some u1 and v2 == Some u2 and is_related(u1, u2) at type s

- For each val v:s in S, prove is_related(M1.v, M2.v) at type s