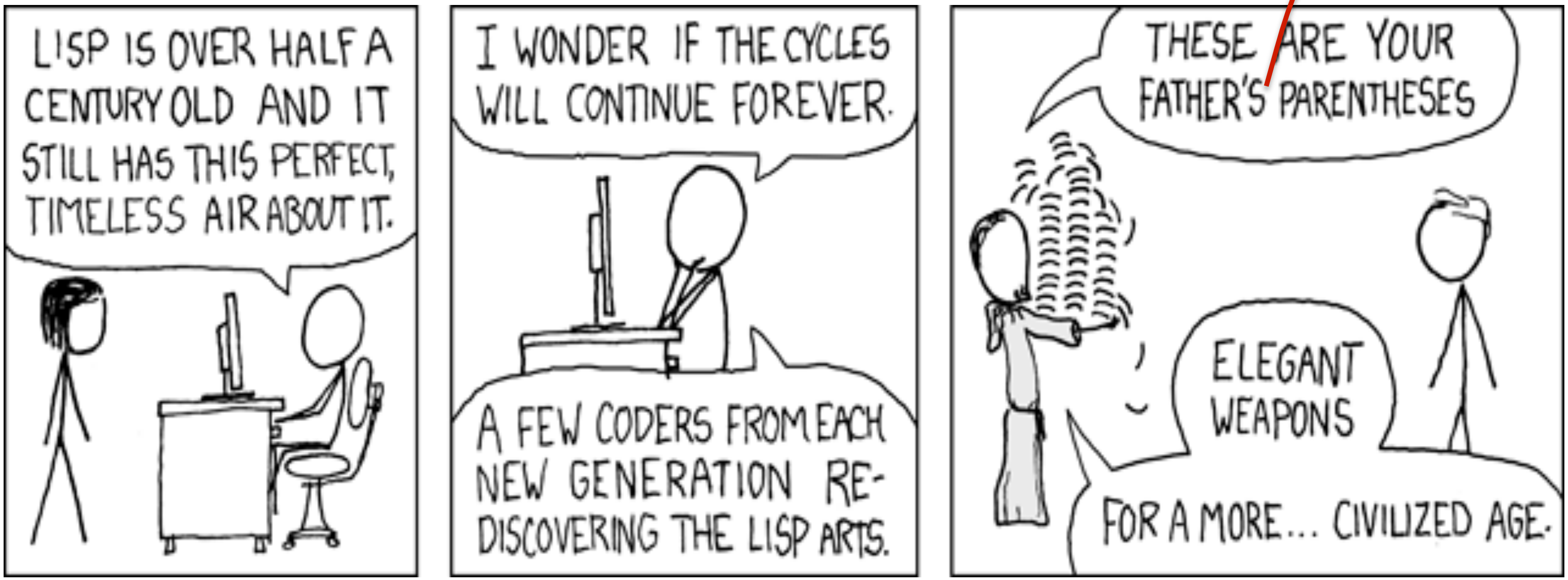

or mother's

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COS 326 Functional Programming: An elegant weapon for the modern age

David Walker

Princeton University

Alonzo Church, 1903-1995
Princeton Professor, 1929-1967

In 1936, Alonzo Church invented the lambda calculus. He called it a logic, but it was a language of pure functions -- the world's first programming language.

He said:

"*There may, indeed, be other applications of the system than its use as a logic.*"

Greatest technological understatement of the 20th century?

He said:

"*There may, indeed, be other applications of the system than its use as a logic.*"

Alonzo Church, 1903-1995
Princeton Professor, 1929-1967

Alonzo Church
1934 -- developed lambda calculus

Programming Languages

Alan Turing (PhD Princeton 1938)
1936 -- developed Turing machines

Computers

Optional reading: ***The Birth of Computer Science at Princeton in the 1930s***
by Andrew W. Appel, 2012. http://press.princeton.edu/chapters/s9780.pdf

A few designers of functional programming languages

Alonzo Church:
λ-calculus, 1934

John McCarthy
(PhD Princeton 1951)
LISP, 1958

Guy Steele & Gerry Sussman:
Scheme, 1975

A few designers of functional programming languages

Alonzo Church:
λ-calculus, 1934

Robin Milner
ML, 1978

Appel & MacQueen: SML/NJ, 1988

Xavier Leroy: Ocaml, 1990's

They were younger than they appear…

Alonzo Church:
λ-calculus, 1934
Photo ~1960

John McCarthy
(PhD Princeton 1951)
LISP, 1958
Photo ~1975

Guy Steele & Gerry Sussman:
Scheme, 1975
Photo ~1995 Photo ~1995

Robin Milner
ML, 1978
Photo ~2005

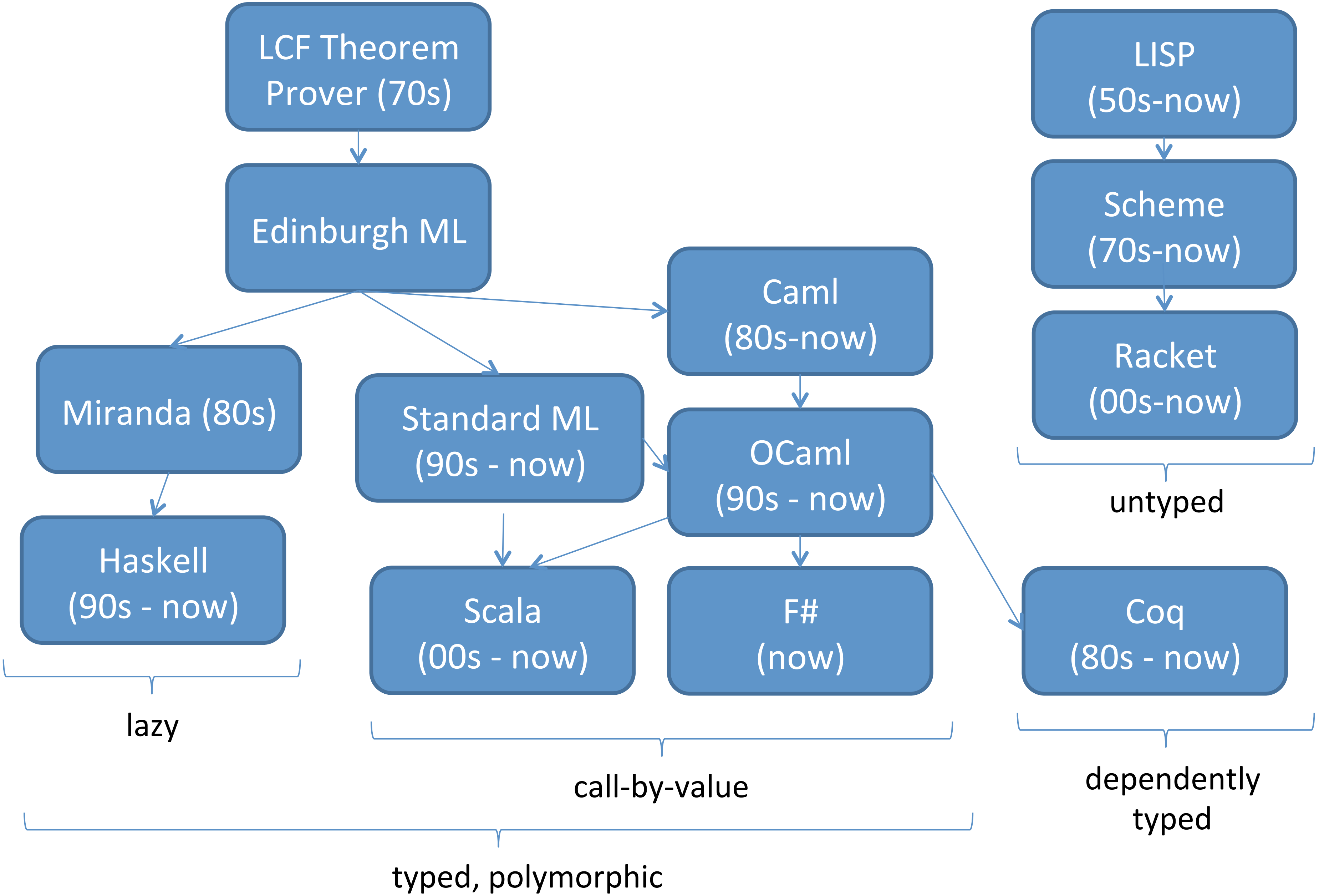
Luca Cardelli
Edinburgh ML, 1981
Photo ~1995

Implementations:

Appel & MacQueen: SML/NJ, 1988
Photo 2005 Photo ~2000

Xavier Leroy:
Ocaml, 1990’s
Photo ~2005

Vastly Abbreviated FP Geneology



LCF Theorem Prover (70s)

Edinburgh ML

Caml (80s-now)

Miranda (80s)

Standard ML (90s - now)

OCaml (90s - now)

Haskell (90s - now)

Scala (00s - now)

F# (now)

Coq (80s - now)

LISP (50s-now)

Scheme (70s-now)

Racket (00s-now)

untyped

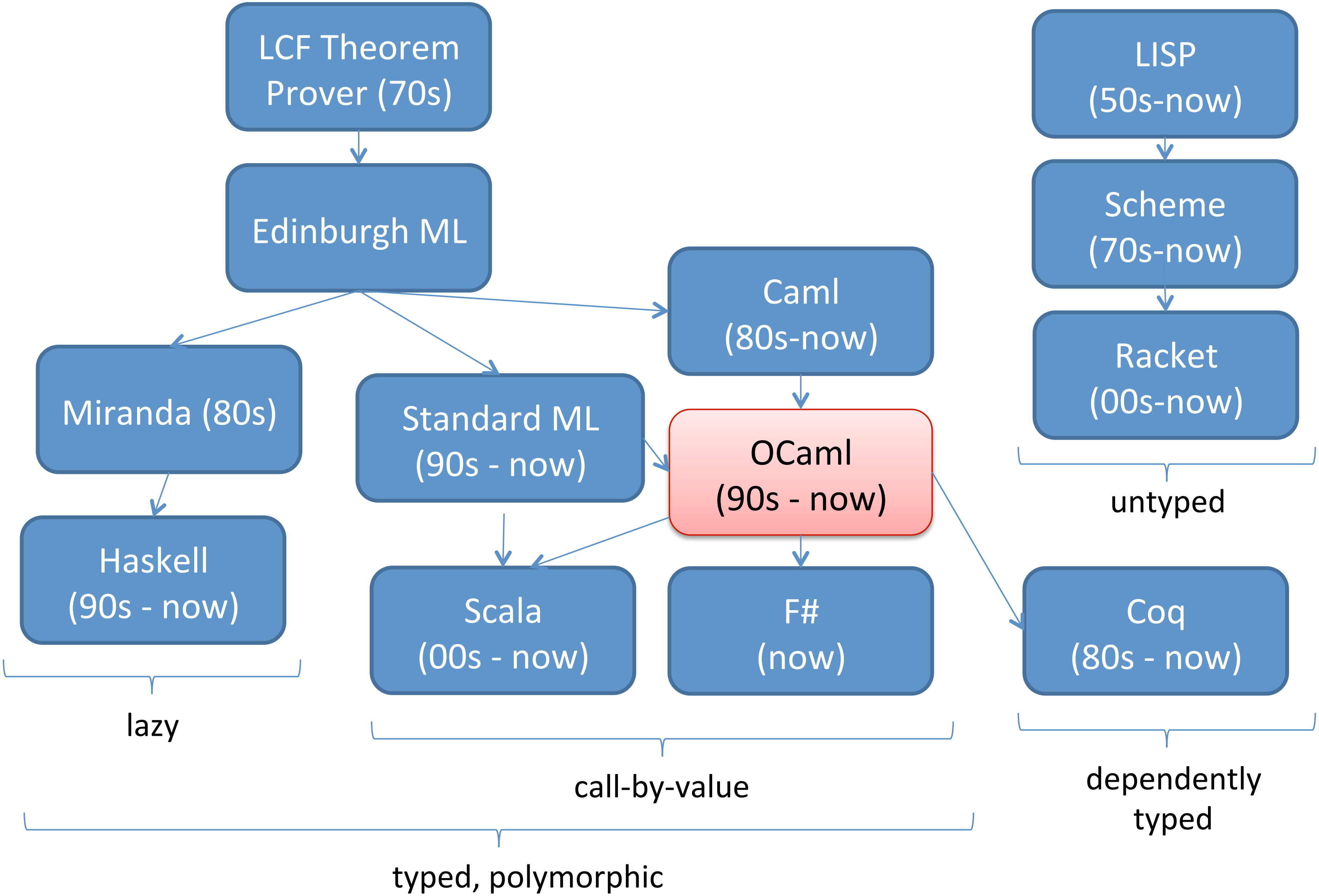
lazy

call-by-value

dependently typed

typed, polymorphic

Vastly Abbreviated FP Geneology



LCF Theorem Prover (70s)

Edinburgh ML

Miranda (80s)

Haskell (90s - now)

Standard ML (90s - now)

Caml (80s-now)

OCaml (90s - now)

Scala (00s - now)

F# (now)

Coq (80s - now)

LISP (50s-now)

Scheme (70s-now)

Racket (00s-now)

untyped

lazy

call-by-value

dependently typed

typed, polymorphic

But Why Functional Programming *Now*?

- Functional programming will introduce you to new ways to *think about* and *structure* your programs:
  - new reasoning principles
  - new abstractions
  - new design patterns
  - new algorithms
  - elegant code

- Technology trends point to increasing parallelism:
  - multicore, gpu, data center
  - functional programming techniques such as map-reduce provide a plausible way forward for many applications

Functional Languages: Who's using them?

twitter

map-reduce in their data centers

Google

Scala for
correctness, maintainability, flexibility

facebook

Erlang for
concurrency,
Haskell for
managing PHP

Microsoft Be what's next.

F# in Visual Studio

mathematicians

Coq (re)proof of
4-color theorem

bluespec

Haskell to
synthesize hardware

BARCLAYS

Haskell
for specifying
equity derivatives

www.artima.com/scalazine/articles/twitter_on_scala.html
gregosuri.com/how-facebook-uses-erlang-for-real-time-chat
www.janestcapital.com/technology/ocaml.php
msdn.microsoft.com/en-us/fsharp/cc742182
labs.google.com/papers/mapreduce.html
www.haskell.org/haskellwiki/Haskell_in_industry

Functional Languages: Join the crowd

- Elements of functional programming are showing up all over
  - F# in Microsoft Visual Studio
  - Scala combines ML (a functional language) with Objects
    - runs on the JVM
  - C# includes “delegates”
    - delegates == functions
  - Python includes “lambdas”
    - lambdas == more functions
  - Javascript
    - find tutorials online about using functional programming techniques to write more elegant code
  - C++ libraries for map-reduce
    - enabled functional parallelism at Google
  - Java has generics and GC
  - ...

COURSE LOGISTICS

Course Staff

David Walker
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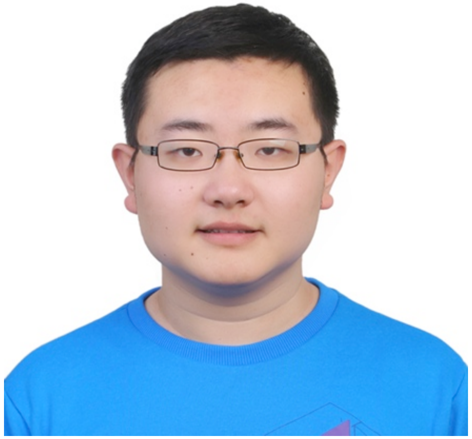


Robin Qiu
Grad Student
office: Fishbowl
email: yqiu@cs

Resources

- Web:
  - http://www.cs.princeton.edu/~cos326
- Lecture schedule and readings:
  - $(coursehome)/lectures.php
- Assignments:
  - $(coursehome)/assignments.php
- Precepts
  - useful if you want to do well on exams and homeworks
- Install OCaml: $(coursehome)/resources.php

Collaboration Policy

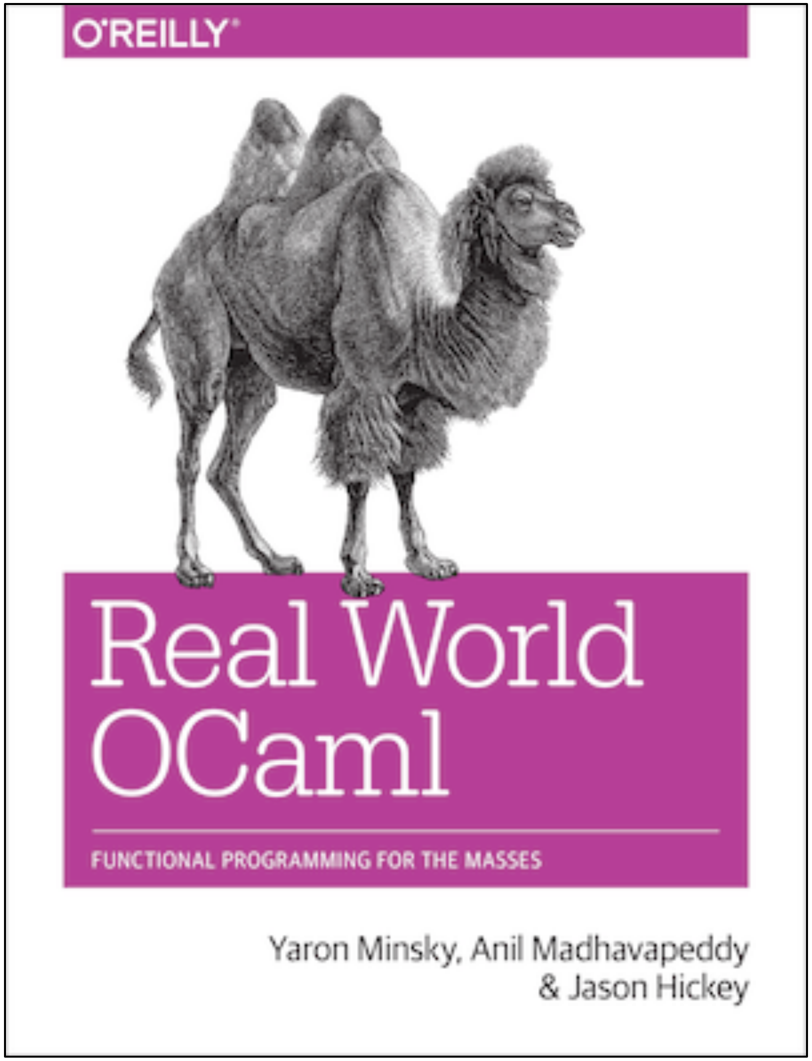
The COS 326 collaboration policy can be found here:

http://www.cs.princeton.edu/~cos326/info.php#collab

Read it in full prior to beginning the first assignment.

Please ask questions whenever anything is unclear, at any time during the course.

Course Textbook



O'REILLY

Real World OCaml

FUNCTIONAL PROGRAMMING FOR THE MASSES

Yaron Minsky, Anil Madhavapeddy
& Jason Hickey

http://realworldocaml.org/

Exams

Minterm: Wednesday of midterm week:

Wednesday, October 25

There will be a final exam, in exam period

(January — Make your travel plans accordingly)

Assignment 0

Figure out how to download and install the latest version of OCaml
on your machine by the time precept begins tomorrow.
(or, how to use OCaml by ssh to Princeton University servers)

Resources Page:

http://www.cs.princeton.edu/~cos326/resources.php

Hint:

ocaml.org

Public Service Announcement

The Pen is Mighter than the Keyboard: Advantages of Longhand Over Laptop Note Taking

Pam Mueller (Princeton University)
Daniel Oppenheimer (UCLA)
Journal of Psychological Science, June 2014, vol 25, no 6

http://pss.sagepub.com/content/25/6/1159.fullkeytype=ref&siteid=sppss&ijkey=CjRAwmrlURGNw

- You learn conceptual topics better by taking notes by hand.
- Instagram and World of Warcraft distract your classmates.

A Functional Introduction

Thinking Functionally

In Java or C, you get (most) work done by *changing* something

```
temp = pair.x;
pair.x = pair.y;
pair.y = temp;
```

commands *modify* or *change* an existing data structure (like pair)

In ML, you get (most) work done by *producing something new*

```
let
  (x,y) = pair
in
  (y,x)
```

you *analyze* existing data (like pair) and you *produce* new data (y,x)

This simple switch in perspective can change the way you

think

about programming and problem solving.

Thinking Functionally

pure, functional code:

```
let (x,y) = pair in
(y,x)
```

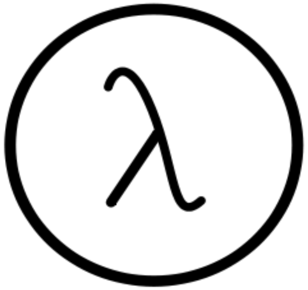
- *outputs are everything!*
- *output is function of input*
- *data properties are stable*
- *repeatable*
- *parallelism apparent*
- *easier to test*
- *easier to compose*

imperative code:

```
temp = pair.x;
pair.x = pair.y;
pair.y = temp;
```

- *outputs are irrelevant!*
- *output is not function of input*
- *data properties change*
- *unrepeatable*
- *parallelism hidden*
- *harder to test*
- *harder to compose*

Why OCaml?



Small, *orthogonal* core based on the *lambda calculus*.

- Control is based on (recursive) functions.
- Instead of for-loops, while-loops, do-loops, iterators, etc.
  - can be defined as library functions.
- Makes it easy to define semantics

Supports *first-class, lexically-scoped, higher-order* procedures

- a.k.a. first-class functions or closures or lambdas.
- first-class: functions are data values like any other data value
  - like numbers, they can be stored, defined anonymously, ...
- lexically-scoped: meaning of variables determined statically.
- higher-order: functions as arguments and results
  - programs passed to programs; generated from programs

These features also found in Racket, Haskell, SML, F#, Clojure, ...

Why OCaml?

Statically typed: debugging and testing aid

- compiler catches many silly errors before you can run the code.
  - A type is worth a thousand tests (start at 6:20):
    - https://www.youtube.com/watch?v=q1Yi-WM7XqQ
- Java is also strongly, statically typed.
- Scheme, Python, Javascript, etc. are all strongly, *dynamically typed* – type errors are discovered while the code is running.

Strongly typed: compiler enforces type abstraction.

- cannot cast an integer to a record, function, string, etc.
  - so we can utilize *types as capabilities*; crucial for local reasoning
- C/C++ are *weakly-typed* (statically typed) languages. The compiler will happily let you do something smart (*more often stupid*).

Type inference: compiler fills in types for you

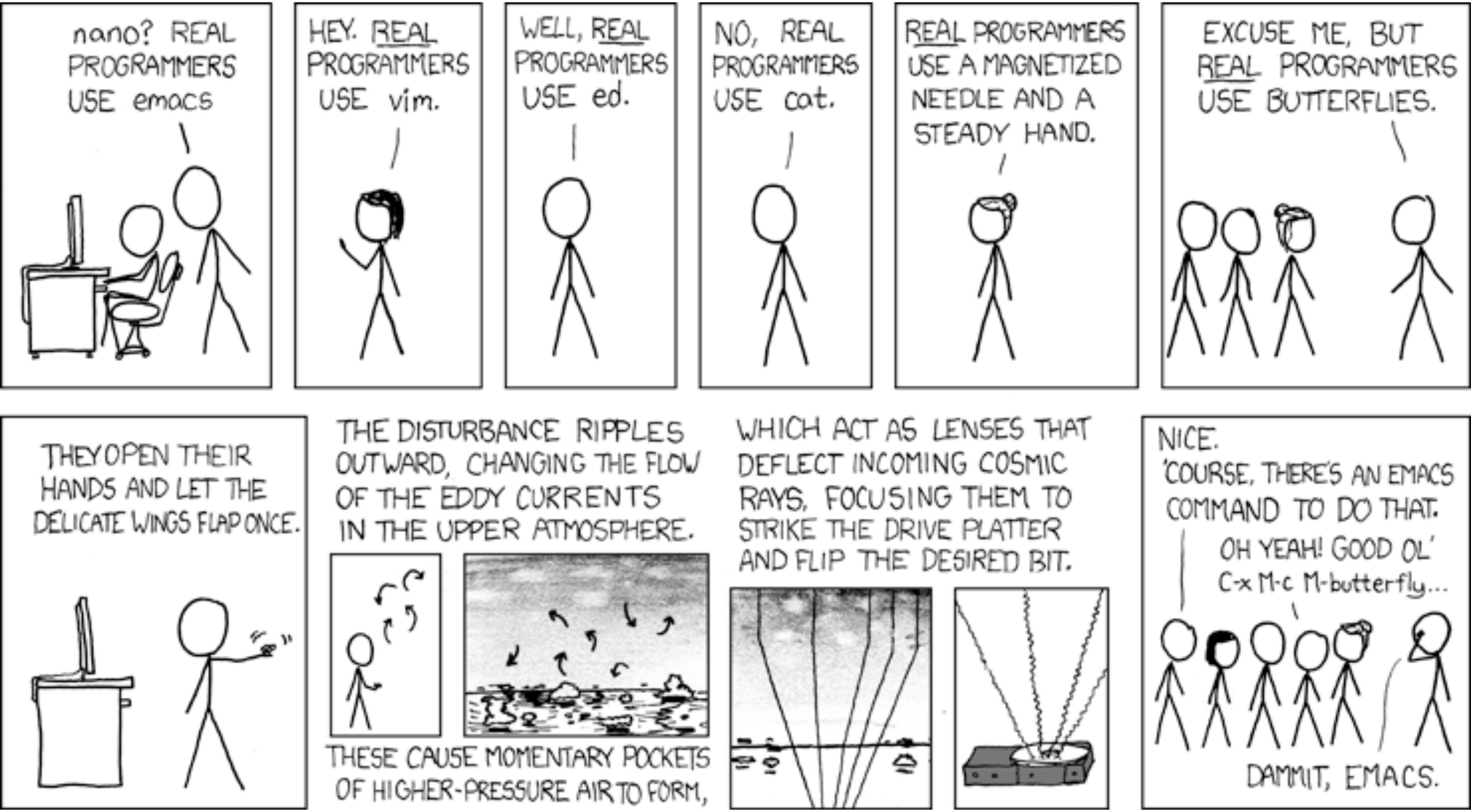
Installing, running Ocaml

- OCaml comes with compilers:
  - “ocamlc” – fast bytecode compiler
  - “ocamlopt” – optimizing, native code compiler
  - “ocamlbuild – a nice wrapper that computes dependencies
- And an interactive, top-level shell:
  - occasionally useful for trying something out.
  - “ocaml” at the prompt.
  - *but use the compiler most of the time*
- And many other tools
  - e.g., debugger, dependency generator, profiler, etc.
- See the course web pages for installation pointers
  - also OCaml.org

Editing Ocaml Programs

- Many options: pick your own poison
  - Emacs
    - what I'll be using in class.
    - good but not great support for OCaml.
    - I like it because it's what I'm used to
    - (extensions written in elisp – a functional language!)
  - OCaml IDE
    - integrated development environment written in Ocaml.
    - haven't used it much, so can't comment.
  - Eclipse
    - I've put up a link to an Ocaml plugin
    - I haven't tried it but others recommend it
  - Sublime, atom
    - A lot of students seem to gravitate to this

XKCD on Editors



AN INTRODUCTORY EXAMPLE (OR TWO)

OCaml Compiler and Interpreter

- Demo:
  - emacs
  - ml files
  - writing simple programs: hello.ml, sum.ml
  - simple debugging and unit tests
  - ocamlc compiler

A First OCaml Program

`hello.ml:`

```
print_string “Hello COS 326!!\n";;
```

A First OCaml Program

`hello.ml:`

```
print_string “Hello COS 326!!\n"
```

a function

its string argument enclosed in `"..."`

a program can be nothing more than just a single expression (but that is uncommon)

no parens. normally call a function f like this:

```
f arg
```

(parens are used for grouping, precedence only when necessary)

A First OCaml Program

`hello.ml:`

```
print_string “Hello COS 326!!\n"
```

`compiling and running hello.ml:`

```
$ ocamlbuild hello.d.byte
$ ./hello.d.byte
Hello COS 326!!
$
```

.d for debugging
(other choices .p for profiled; or none)

.byte for interpreted bytecode
(other choices .native for machine code)

A First OCaml Program

`hello.ml:`

```
print_string “Hello COS 326!!\n"
```

`interpreting and playing with hello.ml:`

```
$ ocaml
        Objective Caml Version 3.12.0
#
```

A First OCaml Program

hello.ml:

```
print_string “Hello COS 326!!\n"
```

interpreting and playing with hello.ml:

```
$ ocaml
        Objective Caml Version 3.12.0
# 3 + 1;;
- : int = 4
#
```

A First OCaml Program

`hello.ml:`

```
print_string “Hello COS 326!!\n"
```

`interpreting and playing with hello.ml:`

```
$ ocaml
      Objective Caml Version 3.12.0
# 3 + 1;;
- : int = 4
# #use "hello.ml";;
hello cos326!!
- : unit = ()
#
```

A First OCaml Program

hello.ml:

```
print_string “Hello COS 326!!\n"
```

interpreting and playing with hello.ml:

```
$ ocaml
        Objective Caml Version 3.12.0
# 3 + 1;;
- : int = 4
# #use "hello.ml";;
hello cos326!!
- : unit = ()
# #quit;;
$
```

A Second OCaml Program

a comment
(* … *)

sumTo8.ml:

```
(* sum the numbers from 0 to n
   precondition: n must be a natural number
*)
let rec sumTo (n:int) : int =
  match n with
    0 -> 0
  | n -> n + sumTo (n-1)

let _ =
  print_int (sumTo 8);
  print_newline()
```

A Second OCaml Program

the name of the function being defined

sumTo8.ml:

```
(* sum the numbers from 0 to n
   precondition: n must be a natural number
*)
let rec sumTo (n:int) : int =
  match n with
    0 -> 0
  | n -> n + sumTo (n-1)

let _ =
  print_int (sumTo 8);
  print_newline()
```

the keyword “let” begins a definition; keyword “rec” indicates recursion

A Second OCaml Program

`sumTo8.ml:`

```
(* sum the numbers from 0 to n
   precondition: n must be a natural number
*)
let rec sumTo (n:int) : int =
  match n with
    0 -> 0
  | n -> n + sumTo (n-1)

let _ =
  print_int (sumTo 8);
  print_newline()
```

result type int

argument
named n
with type int

A Second OCaml Program

deconstruct the value n using pattern matching

`sumTo8.ml:`

```
(* sum the numbers from 0 to n
   precondition: n must be a natural number
*)
let rec sumTo (n:int) : int =
  match n with
    0 -> 0
  | n -> n + sumTo (n-1)

let _ =
  print_int (sumTo 8);
  print_newline()
```

data to be deconstructed appears between key words “match” and “with”

A Second OCaml Program

vertical bar "|" separates the alternative patterns

`sumTo8.ml:`

```
(* sum the numbers from 0 to n
   precondition: n must be a natural number
*)
let rec sumTo (n:int) : int =
  match n with
    0 -> 0
  | n -> n + sumTo (n-1)

let _ =
  print_int (sumTo 8);
  print_newline()


_
```

deconstructed data matches one of 2 cases:
(i) the data matches the pattern 0, or (ii) the data matches the variable pattern n

A Second OCaml Program

Each branch of the match statement constructs a result

`sumTo8.ml:`

```
(* sum the numbers from 0 to n
   precondition: n must be a natural number
*)
let rec sumTo (n:int) : int =
  match n with
    0 -> 0
  | n -> n + sumTo (n-1)

let _ =
  print_int (sumTo 8);
  print_newline()
```

construct the result 0

construct a result using a recursive call to sumTo

A Second OCaml Program

`sumTo8.ml:`

```
(* sum the numbers from 0 to n
   precondition: n must be a natural number
*)
let rec sumTo (n:int) : int =
  match n with
    0 -> 0
  | n -> n + sumTo (n-1)

let _ =
  print_int (sumTo 8);
  print_newline()
```

print the result of calling sumTo on 8

print a new line

OCAML BASICS: EXPRESSIONS, VALUES, SIMPLE TYPES

Terminology: Expressions, Values, Types

- Expressions are computations
  - 2 + 3 is a computation

- Values are the results of computations
  - 5 is a value

- Types describe collections of values and the computations that generate those values
  - int is a type

  - values of type int include
    - 0, 1, 2, 3, …, max_int
    - -1, -2, …, min_int

Some simple types, values, expressions

Type:	Values:	Expressions:
`int`	`-2, 0, 42`	`42 * (13 + 1)`
`float`	`3.14, -1., 2e12`	`(3.14 +. 12.0) *. 10e6`
`char`	`’a’, ’b’, ’&’`	`int_of_char ’a’`
`string`	`"moo", "cow"`	`"moo" ^ "cow"`
`bool`	`true, false`	`if true then 3 else 4`
`unit`	`()`	`print_int 3`

For more primitive types and functions over them, see the OCaml Reference Manual here:

http://caml.inria.fr/pub/docs/manual-ocaml/libref/Pervasives.html

Not every expression has a value

Expression:

`42 * (13 + 1)`	**evaluates to**	`588`
`(3.14 +. 12.0) *. 10e6`	↦	`151400000.`
`int_of_char ’a’`	↦	`97`
`"moo" ^ "cow"`	↦	`“moocow"`
`if true then 3 else 4`	↦	`3`
`print_int 3`	↦	`()`

`1 + "hello"` **does not evaluate!**

Language Definition

- There are a number of ways to define a programming language
- In this class, we will briefly investigate:
  - Syntax
  - Evaluation
  - Type checking
- Standard ML, a very close relative of OCaml, has a full definition of each of these parts and a number of proofs of correctness
  - For more on this theme, see COS 441/510
- The OCaml Manual fleshes out the syntax, evaluation and type checking rules informally

OCAML BASICS:
CORE EXPRESSION SYNTAX

Core Expression Syntax

The simplest OCaml expressions e are:

• values	*numbers, strings, bools, ...*
• id	*variables (x, foo, ...)*
• e_1 op e_2	*operators (x+3, ...)*
• id e_1 e_2 … e_n	*function call (foo 3 42)*
• **let** id = e_1 **in** e_2	*local variable decl.*
• **if** e_1 **then** e_2 **else** e_3	*a conditional*
• (e)	*a parenthesized expression*
• (e : t)	*an expression with its type*

A note on parentheses

In most languages, arguments are parenthesized & separated by commas:

```
f(x,y,z)      sum(3,4,5)
```

In OCaml, we don’t write the parentheses or the commas:

```
f x y z       sum 3 4 5
```

But we do have to worry about *grouping*. For example,

```
f x y z
f x (y z)
```

The first one passes three arguments to f (x, y, and z)

The second passes two arguments to f (x, and the result of applying the function y to z.)

OCAML BASICS: TYPE CHECKING

Type Checking

- Every value has a type and so does every expression
- This is a concept that is familiar from Java but it becomes more important when programming in a functional language
- The type of an expression is determined by the type of its subexpressions
- We write (e : t) to say that expression e has type t. eg:

2 : int	"hello" : string
2 + 2 : int	"I say " ^ "hello" : string

Type Checking Rules

- There are a set of simple rules that govern type checking
  - programs that do not follow the rules will not type check and O'Caml will refuse to compile them for you (the nerve!)
  - at first you may find this to be a pain …

- But types are a great thing:
  - they *help us think* about *how to construct* our programs
  - they help us *find stupid programming errors*
  - they help us track down compatibility errors quickly when we edit and *maintain our code*
  - they allow us to *enforce powerful invariants* about our data structures

Type Checking Rules

- Example rules:

(1) 0 : int (and similarly for any other integer constant n)

(2) "abc" : string (and similarly for any other string constant "…")

Type Checking Rules

- Example rules:

(1) 0 : int (and similarly for any other integer constant n)

(2) "abc" : string (and similarly for any other string constant "…")

(3) if e1 : int and e2 : int
then e1 + e2 : int

(4) if e1 : int and e2 : int
then e1 * e2 : int

Type Checking Rules

- Example rules:

(1) 0 : int (and similarly for any other integer constant n)

(2) "abc" : string (and similarly for any other string constant "…")

(3) if e1 : int and e2 : int
then e1 + e2 : int

(4) if e1 : int and e2 : int
then e1 * e2 : int

(5) if e1 : string and e2 : string
then e1 ^ e2 : string

(6) if e : int
then string_of_int e : string

Type Checking Rules

- Example rules:

(1) 0 : int (and similarly for any other integer constant n)

(2) "abc" : string (and similarly for any other string constant "…")

(3) if e1 : int and e2 : int
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then e1 ^ e2 : string

(6) if e : int
then string_of_int e : string

- Using the rules:

2 : int and 3 : int. (By rule 1)

Type Checking Rules

- Example rules:

(1) 0 : int (and similarly for any other integer constant n)

(2) "abc" : string (and similarly for any other string constant "…")

(3) if e1 : int and e2 : int
then e1 + e2 : int

(4) if e1 : int and e2 : int
then e1 * e2 : int

(5) if e1 : string and e2 : string
then e1 ^ e2 : string

(6) if e : int
then string_of_int e : string

- Using the rules:

2 : int and 3 : int. (By rule 1)
Therefore, (2 + 3) : int (By rule 3)

Type Checking Rules

- Example rules:

(1) 0 : int (and similarly for any other integer constant n)

(2) "abc" : string (and similarly for any other string constant "…")

(3) if e1 : int and e2 : int
then e1 + e2 : int

(4) if e1 : int and e2 : int
then e1 * e2 : int

(5) if e1 : string and e2 : string
then e1 ^ e2 : string

(6) if e : int
then string_of_int e : string

- Using the rules:

2 : int and 3 : int. (By rule 1)
Therefore, (2 + 3) : int (By rule 3)
5 : int (By rule 1)

Type Checking Rules

- Example rules:

(1) 0 : int (and similarly for any other integer constant n)

(2) "abc" : string (and similarly for ")

(3) if e1 : int and e2 : int
then e1 + e2 : int

(5) if e1 : string and e2 : string
then e1 ^ e2 : string

(

tring_of_int e : string

FYI: This is a ***formal proof*** that the expression is well-typed!

- Using the rules:

2 : int and 3 : int. (By rule 1)
Therefore, (2 + 3) : int (By rule 3)
5 : int (By rule 1)
Therefore, (2 + 3) * 5 : int (By rule 4 and our previous work)

Type Checking Rules

- Example rules:

(1) 0 : int (and similarly for any other integer constant n)

(2) "abc" : string (and similarly for any other string constant "…")

(3) if e1 : int and e2 : int
then e1 + e2 : int

(4) if e1 : int and e2 : int
then e1 * e2 : int

(5) if e1 : string and e2 : string
then e1 ^ e2 : string

(6) if e : int
then string_of_int e : string

- Another perspective:

???? * ???? : int

rule (4) for typing expressions
says I can put any expression
with type int in place of the ????

Type Checking Rules

- Example rules:

(1) 0 : int (and similarly for any other integer constant n)

(2) "abc" : string (and similarly for any other string constant "…")

(3) if e1 : int and e2 : int
then e1 + e2 : int

(4) if e1 : int and e2 : int
then e1 * e2 : int

(5) if e1 : string and e2 : string
then e1 ^ e2 : string

(6) if e : int
then string_of_int e : string

- Another perspective:

7 * ???? : int

rule (4) for typing expressions
says I can put any expression
with type int in place of the ????

Type Checking Rules

- Example rules:

(1) 0 : int (and similarly for any other integer constant n)

(2) "abc" : string (and similarly for any other string constant "…")

(3) if e1 : int and e2 : int
then e1 + e2 : int

(4) if e1 : int and e2 : int
then e1 * e2 : int

(5) if e1 : string and e2 : string
then e1 ^ e2 : string

(6) if e : int
then string_of_int e : string

- Another perspective:

7 * (add_one 17) : int

rule (4) for typing expressions
says I can put any expression
with type int in place of the ????

Type Checking Rules

- You can always start up the OCaml interpreter to find out a type of a simple expression:

```
$ ocaml
         Objective Caml Version 3.12.0
#
```

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- : int = 4
#
```

press
return
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```
$ ocaml
       Objective Caml Version 3.12.0
# 3 + 1;;
- : int = 4
# “hello ” ^ “world”;;
- : string = “hello world”
#
```

press return and you find out the type and the value

Type Checking Rules

- You can always start up the OCaml interpreter to find out a type of a simple expression:

```
$ ocaml
        Objective Caml Version 3.12.0
# 3 + 1;;
- : int = 4
# “hello ” ^ “world”;;
- : string = “hello world”
# #quit;;
$
```

Type Checking Rules

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- Violating the rules:

"hello" : string (By rule 2)
1 : int (By rule 1)
1 + "hello" : ?? (NO TYPE! Rule 3 does not apply!)

Type Checking Rules

- Violating the rules:

```
# "hello" + 1;;
Error: This expression has type string but an
expression was expected of type int
```

- The type error message tells you the type that was expected and the type that it inferred for your subexpression
- By the way, this was one of the nonsensical expressions that did not evaluate to a value
- It is a ***good thing*** that this expression does not type check!

“Well typed programs do not go wrong”

Robin Milner, 1978

Type Checking Rules

- Violating the rules:

```
# "hello" + 1;;
Error: This expression has type string but an
expression was expected of type int
```

- A possible fix:

```
# "hello" ^ (string_of_int 1);;
- : string = "hello1"
```

- *One of the keys to becoming a good ML programmer is to understand type error messages.*

Type Checking Rules

- What about this expression:

```
# 3 / 0 ;;
Exception: Division_by_zero.
```

- Why doesn't the ML type checker do us the favor of telling us the expression will raise an exception?

Type Checking Rules

- What about this expression:

```
# 3 / 0 ;;
Exception: Division_by_zero.
```

- Why doesn't the ML type checker do us the favor of telling us the expression will raise an exception?
  - In general, detecting a divide-by-zero error requires we know that the divisor evaluates to 0.
  - In general, deciding whether the divisor evaluates to 0 requires solving the halting problem:

```
# 3 / (if turing_machine_halts m then 0 else 1);;
```

- There are type systems that will rule out divide-by-zero errors, but they require programmers supply proofs to the type checker

Isn't that cheating?

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(3 / 0) is well typed. Does it "go wrong?" Answer: No.

"Go wrong" is a technical term meaning, "have no defined semantics." Raising an exception is perfectly well defined semantics, which we can reason about, which we can handle in ML with an exception handler.

So, it's not cheating.

(Discussion: why do we make this distinction, anyway?)

Type Soundness

“Well typed programs do not go wrong”

Programming languages with this property have *sound* type systems. They are called *safe* languages.

Safe languages are generally *immune* to buffer overrun vulnerabilities, uninitialized pointer vulnerabilities, etc., etc. (but not immune to all bugs!)

Safe languages: ML, Java, Python, …

Unsafe languages: C, C++, Pascal

Well typed programs do not go wrong

Robin Milner

Turing Award, 1991

“For three distinct and complete achievements:

1. LCF, the mechanization of Scott's Logic of Computable Functions, probably the first theoretically based yet practical tool for machine assisted proof construction;

2. ML, the first language to include polymorphic type inference together with a type-safe exception-handling mechanism;

3. CCS, a general theory of concurrency.

In addition, he formulated and strongly advanced full abstraction, the study of the relationship between operational and denotational semantics.”

“Well typed programs do not go wrong”

Robin Milner, 1978

OVERALL SUMMARY:
A SHORT INTRODUCTION TO FUNCTIONAL PROGRAMMING

OCaml

OCaml is a *functional* programming language

- Java gets most work done by *modifying* data
- OCaml gets most work done by producing *new*, *immutable* data

OCaml is a *typed* programming language

- the *type* of an expression *correctly predicts* the kind of *value* the expression will generate when it is executed
- types help us *understand* and *write* our programs
- the type system is *sound*; the language is *safe*