# Written Exam 1

#### 0. Miscellaneous.

Don't forget to write your name, NetID, precept, and exam room.

# 1. Java basics.

| Java expression                   | value             |
|-----------------------------------|-------------------|
| x                                 | 1                 |
| 1.0 * x + y / z                   | 1.0               |
| ((x > y)    (y > z)) && (z > x)   | false             |
| Math.min(Math.max(x, y), y*y % z) | 1                 |
| Math.sqrt(x - y - z)              | NaN or Double.NaN |
| Double.parseDouble(x + "2" + y)   | 122.0             |

# 2. Properties of arrays and functions.

- (a) T T F F F
- (b) T T F F T
- 3. Debugging and arrays.
  - $\mathbf{D} \neq \mathbf{G} \neq \mathbf{C}$

#### 4. Functions.

(a) CKDKEK

The letters C, D, and E can be permuted in any order. EKGKHK is an alternative solution, but poorer style.

```
public static boolean oddParity(boolean x, boolean y, boolean z) {
    int count = 0;
    if (x) count++;
    if (y) count++;
    if (z) count++;
    return (count % 2) != 0;
}
```

(b) GEHFE

GEHFJ, GIHFE, GIHFJ are alternative solutions, but poorer style.

```
public static boolean oddParity(boolean x, boolean y, boolean z) {
    if (x && y) return z;
    else if (x || y) return !z;
    else return z;
}
```

This part is tricker than it might appear. It is tempting to start with IAGB, but then you quickly get stuck.

#### 5. Recursion.

(a)

(b) **T T T** 

- Let  $L_n$  be the length of the return value of f(n). Then,  $L_n = L_{n-1} + 1 + L_{n-2} + L_{n-1}$ , with  $L_0 = L_1 = 1$ . From above,  $L_2 = 4$  and  $L_3 = 10$ . Thus,  $L_4 = 10 + 1 + 4 + 10 = 25$ and  $L_5 = 25 + 1 + 10 + 25 = 61$ .
- Exchanging statements 2 and 3 can have an effect only when n is 1. When n is 1, the modified function still returns "1"—it just performs the unnecessary work of calling f(0) before doing so.
- The function f() has no side effects (such as printing to standard output). Thus first = f(n-1) and third = f(n-1) will always be equal, and the return value in the original function (first + n + second + third) will always equal the return value in the modified function (first + n + second + first).

# 6. Powers of 2.

# $\mathbf{D} \; \mathbf{B} \; \mathbf{F} \; \mathbf{D} \; \mathbf{C} \; \mathbf{N} \; \mathbf{E} \; \mathbf{E}$

### 7. TOY.

The TOY program computes the smallest power of 2 that is strictly greater than a given integer.

| 10: 8A0 | 00 R[A] <- M[OO]      | load a from M[OO]      |
|---------|-----------------------|------------------------|
| 11: 710 | )1 R[1] <- 1          | power = 1              |
| 12: 221 | .A R[2] <- R[1] - R   | [A] while (power <= a) |
| 13: D21 | .6 if $(R[2] > 0) PC$ | ; <- 16    {           |
| 14: 111 | .1 R[1] <- R[1] + R   | [1] power = 2*power    |
| 15: CO1 | .2 goto 12            | }                      |
| 16: 910 | )1 M[01] <- R[1]      | store power to M[01]   |
| 17: 000 | 00 halt               |                        |

- (a) **0002**
- (b) **0008**
- (c) **0010**

Remember that everything is in hex:  $10_{16} = 16_{10}$ .

(d) **2000** 

You will get this part only by reasoning about what the TOY program does (or wasting an extraordinary amount of time tracing code).

(e) **0001** 

FACE is a negative integer (two's complement representation). Thus, the loop is skipped and R[1] remains 0001.