



# Assembly Language: IA-32 Instructions

1



## Goals of this Lecture

- Help you learn how to:
  - Manipulate data of various sizes
  - Leverage more sophisticated addressing modes
  - Use condition codes and jumps to change control flow... and thereby ...
  - Write more efficient assembly-language programs
  - Understand the relationship to data types and common programming constructs in high-level languages
- Focus is on the assembly-language code
  - Rather than layout of memory for storing data (precept)

2



## Handling Different Data Sizes

3

## Variable Sizes in High-Level Language



- C data types vary in size
  - Character: 1 byte
  - Short, int, and long: ??
  - Float and double: ??
  - Pointers: ??
- Programmer-created types
  - Struct: ??
- Arrays
  - Multiple consecutive elements of some fixed size
  - Where each element could be a struct

4

## Supporting Different Sizes in IA-32



- Three main data sizes
  - Byte (b): 1 byte
  - Word (w): 2 bytes
  - Long (l): 4 bytes
- Separate assembly-language instructions
  - E.g., addb, addw, and addl
- Separate ways to access (parts of) a register
  - E.g., for EAX register: %ah or %al, %ax, and %eax
- Larger sizes (e.g., struct)
  - Manipulated in smaller byte, word, or long units

5

## Byte Order in Multi-Byte Entities



- IA-32 is a **little endian** architecture
  - Least significant byte of multi-byte entity is stored at lowest memory address
  - “Little end goes first”
- Some other systems use **big endian**
  - Most significant byte of multi-byte entity is stored at lowest memory address
  - “Big end goes first”

The 4-byte int 5 (hex 00 00 00 05) at address 1000:

|      |          |
|------|----------|
| 1000 | 00000101 |
| 1001 | 00000000 |
| 1002 | 00000000 |
| 1003 | 00000000 |

The 4-byte int 5 (hex 00 00 00 05) at address 1000:

|      |          |
|------|----------|
| 1000 | 00000000 |
| 1001 | 00000000 |
| 1002 | 00000000 |
| 1003 | 00000101 |

6

## Little Endian Example



```
int main(void) {
    int i=0x003377ff, j;
    unsigned char *p = (unsigned char *) &i;

    for (j=0; j<4; j++)
        printf("Byte %d: %x\n", j, p[j]);
}
```

Output on a  
little-endian  
machine

Byte 0: ff  
Byte 1: 77  
Byte 2: 33  
Byte 3: 00

• Portable?

7

## IA-32 General Purpose Registers



| 31 | 15 | 8  | 7 | 0 | 16-bit | 32-bit | Common Use                        |
|----|----|----|---|---|--------|--------|-----------------------------------|
|    | AH | AL |   |   | AX     | EAX    | Accumulator                       |
|    | BH | BL |   |   | BX     | EBX    | Pointer to data                   |
|    | CH | CL |   |   | CX     | ECX    | Counter for loops                 |
|    | DH | DL |   |   | DX     | EDX    | I/O pointer                       |
|    | SI |    |   |   |        | ESI    | Pointers (string source and dest) |
|    | DI |    |   |   |        | EDI    |                                   |

General-purpose registers

- EBP: pointer to data on stack
- ESP: stack pointer

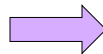
8

## C Example: One-Byte Data



Global *char* variable *i* is in *%al*,  
the *lower byte* of the “A” register.

```
char i;  
...  
if (i > 5) {  
    i++;  
else  
    i--;  
}
```



```
    cmpb $5, %al  
    jle else  
    incb %al  
    jmp endif  
else:  
    decb %al  
endif:
```

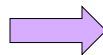
9

## C Example: Four-Byte Data



Global *int* variable *i* is in *%eax*,  
the *full 32 bits* of the “A” register.

```
int i;  
...  
if (i > 5) {  
    i++;  
else  
    i--;  
}
```



```
    cmpl $5, %eax  
    jle else  
    incl %eax  
    jmp endif  
else:  
    decl %eax  
endif:
```

10



## Memory Addressing Modes

11



## Loading and Storing Data

- Processors have many ways to access data
  - Known as “addressing modes”
  - Two simple ways seen in previous examples ...
- Immediate addressing
  - Example: `movl $0, %ecx`
  - Initialize register ECX with zero
  - Data (e.g., number “0”) embedded in the instruction
- Register addressing
  - Example: `movl %edx, %ecx`
  - Copy value in register EDX into register ECX
  - Choice of register(s) embedded in the instruction

12

## Accessing Memory



- Variables are stored in memory
  - Global and static local variables in Data or BSS section
  - Dynamically allocated variables in the heap
  - Function parameters and local variables on the stack
- Need to be able to load from and store to memory
  - To copy the data between main memory and registers
  - Or manipulate the data directly in memory
- IA-32 has many different addressing modes
  - Corresponding to common programming constructs
  - E.g., accessing a global variable, dereferencing a pointer, accessing a field in a struct, or indexing an array

13

## Direct Addressing



- Useful when the address is known in advance
  - Global variables in the Data or BSS sections
- Load or store from a particular memory location
  - Memory address is embedded in the instruction
  - Instruction reads from or writes to that address
- IA-32 example: `movl 2000, %ecx`
  - Four-byte variable located at address 2000
  - Read four bytes starting at address 2000
  - Load the value into the ECX register
- Can use a label for (human) readability
  - E.g., “i” to allow “`movl i, %eax`”

14

## Indirect Addressing



- Useful when address is not known in advance
  - Dereference a pointer, for dynamically allocated data
- Load or store from a previously-computed address
  - Register with the address is embedded in the instruction
  - Instruction reads from or writes to that address
- IA-32 example: `movl (%eax), %ecx`
  - EAX register stores a 32-bit address (e.g., 2000)
  - Read long-word variable stored at that address
  - Load the value into the ECX register
  - The “(%eax)” essentially dereferences the pointer stored in register %eax

15

## Base Pointer Addressing



- Useful when accessing part of a larger variable
  - Specific field within a “struct”
  - E.g., if “age” starts at the 8<sup>th</sup> byte of “student” record
- Load or store with an offset from a base address
  - `movl offset(r1), r2`
  - Register r1 stores the base address
  - Fixed offset also embedded in the instruction
  - Instruction computes the address and does access
- IA-32 example: `movl 8(%eax), %ecx`
  - EAX register stores a 32-bit base address (e.g., 2000)
  - Offset of 8 is added to compute address (e.g., 2008)
  - Load the value into the ECX register

16



## Indexed Addressing



- Load/store with offset made of register, multiplier
  - Fixed base address embedded in the instruction
  - Offset = register \* constant multiplier
- Useful to iterate through an array (e.g., `a[i]`)
  - Base is the start of the array (i.e., “a”)
  - Register is the index (i.e., “i”)
  - Multiplier is the size of the element (e.g., 4 for “int”)
- IA-32 example: `movl 2000(,%eax,4), %ecx`
  - Index register EAX (say, with value of 10)
  - Multiplied by a multiplier of 1, 2, 4, or 8 (here, 4)
  - Added to a fixed base of 2000 (to get 2040)

17

## Indexed Addressing Example



```
int a[20]; ← global variable
int i, sum=0;
for (i=0; i<20; i++)
    sum += a[i];
```

EAX: temporary  
EBX: sum  
ECX: i

```
movl $0, %ecx
movl $0, %ebx
sumloop:
    movl a(,%ecx,4), %eax
    addl %eax, %ebx
    incl %ecx
    cmpl $19, %ecx
    jle sumloop
```

## Effective Address: More Generally



$$\text{Offset} = \begin{pmatrix} \text{eax} \\ \text{ebx} \\ \text{ecx} \\ \text{edx} \\ \text{esp} \\ \text{ebp} \\ \text{esi} \\ \text{edi} \end{pmatrix} + \begin{pmatrix} \text{eax} \\ \text{ebx} \\ \text{ecx} \\ \text{edx} \\ \text{esp} \\ \text{ebp} \\ \text{esi} \\ \text{edi} \end{pmatrix} * \begin{pmatrix} 1 \\ 2 \\ 4 \\ 8 \end{pmatrix} + \begin{pmatrix} \text{None} \\ 8\text{-bit} \\ 16\text{-bit} \\ 32\text{-bit} \end{pmatrix}$$

Base                  Index          scale          displacement

- Displacement `movl foo, %ebx`
- Base `movl (%eax), %ebx`
- Base + displacement `movl foo(%eax), %ebx`  
`movl 1(%eax), %ebx`
- (Index \* scale) + displacement `movl (%edx,%eax,4), %ebx`
- Base + (index \* scale) + displacement `movl foo(%edx,%eax,4), %ebx`

19

## Data Access Methods: Summary



- **Immediate addressing:** data stored in the instruction itself
  - `movl $10, %ecx`
- **Register addressing:** data stored in a register
  - `movl %eax, %ecx`
- **Direct addressing:** address stored in instruction
  - `movl foo, %ecx`
- **Indirect addressing:** address stored in a register
  - `movl (%eax), %ecx`
- **Base pointer addressing:** indirect plus offset
  - `movl 4(%eax), %ecx`
- **Indexed addressing:** instruction contains base address, and specifies an index register and a multiplier (1, 2, 4, or 8)
  - `movl 2000(%eax,1), %ecx`
  - Can also have an additional displacement register

20



## Condition Codes and Control Flow

21



### Control Flow

- Common case
  - Execute code sequentially
  - One instruction after another
- Sometimes need to change control flow
  - If-then-else
  - Loops
  - Switch
- Two key ingredients
  - Testing a condition
  - Selecting what to run next based on result

```
    cmpl $5, %eax
    jle else
    incl %eax
    jmp endif
else:
    decl %eax
endif:
```

## Condition Codes



- 1-bit registers set by arithmetic & logic instructions
  - ZF: Zero Flag
  - SF: Sign Flag
  - CF: Carry Flag
  - OF: Overflow Flag
- Example: “addl Src, Dest” (“t = a + b”)
  - ZF: set if t == 0
  - SF: set if t < 0
  - CF: set if carry out from most significant bit (unsigned)
    - *Unsigned* overflow
  - OF: set if two's complement overflow
    - (a>0 && b>0 && t<0) || (a<0 && b<0 && t>=0)

23

## Condition Codes (continued)



- Example: “cmpl Src2,Src1” (compare b,a)
  - Like computing a-b without setting destination
  - ZF: set if a == b
  - SF: set if (a-b) < 0
  - CF: set if carry out from most significant bit (unsigned)
  - OF: set if two's complement overflow
    - (a>0 && b<0 && (a-b)<0) || (a<0 && b>0 && (a-b)>0)
- Flags are *not* set by lea, inc, or dec instructions
  - Hint: this is useful for the extra-credit part of the assembly-language programming assignment

24

## Jumps after Comparison (cmpl)



- Equality
  - Equal: je (ZF is set)
  - Not equal: jne ( $\sim$ ZF)
- Below/above (e.g., unsigned arithmetic)
  - Below: jb (CF is set)
  - Above or equal: jae ( $\sim$ CF)
  - Below or equal: jbe (CF  $\vee$  ZF)
  - Above: ja ( $\sim$ (CF  $\vee$  ZF))
- Less/greater (e.g., signed arithmetic)
  - Less: jl (SF  $\wedge$  OF)
  - Greater or equal: jge ( $\sim$ (SF  $\wedge$  OF))
  - Less or equal: jle ((SF  $\wedge$  OF)  $\vee$  ZF)
  - Greater: jg ( $\sim$ ((SF  $\wedge$  OF)  $\vee$  ZF))

26

## Branch Instructions



- Conditional jump
  - $\{j\{l,g,e,ne,\dots\} \text{ target} \quad \text{if (condition) \{eip = target\}}$

| Comparison     | Signed | Unsigned |                  |
|----------------|--------|----------|------------------|
| $=$            | e      | e        | "equal"          |
| $\neq$         | ne     | ne       | "not equal"      |
| $>$            | g      | a        | "greater, above" |
| $\geq$         | ge     | ae       | "...-or-equal"   |
| $<$            | l      | b        | "less, below"    |
| $\leq$         | le     | be       | "...-or-equal"   |
| overflow/carry | o      | c        |                  |
| no ovf/carry   | no     | nc       |                  |

- Unconditional jump
  - jmp target
  - jmp \*register

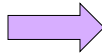
27

## Jumping



- Simple model of a “goto” statement
  - Go to a particular place in the code
  - Based on whether a condition is true or false
  - Can represent if-the-else, switch, loops, etc.
- Pseudocode example: If-Then-Else

```
if (Test) {  
    then-body;  
} else {  
    else-body;
```



```
if (!Test) jump to Else;  
then-body;  
jump to Done;  
Else:  
    else-body;  
Done:
```

28

## Jumping (continued)



- Pseudocode example: Do-While loop

```
do {  
    Body;  
} while (Test);
```



```
loop:  
    Body;  
    if (Test) then jump to loop;
```

- Pseudocode example: While loop

```
while (Test)  
    Body;
```



```
jump to middle;  
loop:  
    Body;  
middle:  
    if (Test) then jump to loop;
```

29

## Jumping (continued)



- Pseudocode example: For loop

```
for (Init; Test; Update)  
  Body
```



```
Init;  
  if (!Test) jump to done;  
loop:  
  Body;  
  Update;  
  if (Test) jump to loop;  
done:
```

30

## Example Instruction Types



31

## Arithmetic Instructions



- Simple instructions

- |                               |                      |
|-------------------------------|----------------------|
| • add{b,w,l} source, dest     | dest = source + dest |
| • sub{b,w,l} source, dest     | dest = dest - source |
| • Inc{b,w,l} dest             | dest = dest + 1      |
| • dec{b,w,l} dest             | dest = dest - 1      |
| • neg{b,w,l} dest             | dest = ~dest + 1     |
| • cmp{b,w,l} source1, source2 | source2 - source1    |

- Multiply

- mul (unsigned) or imul (signed)

- Divide

- div (unsigned) or idiv (signed)

- Many more in Intel manual (volume 2)

- adc, sbb, decimal arithmetic instructions

32

## Bitwise Logic Instructions



- Simple instructions

- |                                      |                       |
|--------------------------------------|-----------------------|
| and{b,w,l} source, dest              | dest = source & dest  |
| or{b,w,l} source, dest               | dest = source   dest  |
| xor{b,w,l} source, dest              | dest = source ^ dest  |
| not{b,w,l} dest                      | dest = ~dest          |
| sar{b,w,l} source, dest (arithmetic) | dest = dest << source |
| sar{b,w,l} source, dest (arithmetic) | dest = dest >> source |

- Many more in Intel Manual (volume 2)

- Logic shift
- Rotation shift
- Bit scan
- Bit test
- Byte set on conditions

33



## Data Transfer Instructions

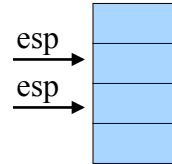


- **mov{b,w,l} source, dest**

- General move instruction

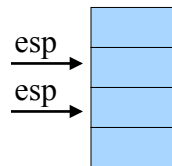
- **push{w,l} source**

`pushl %ebx`    # equivalent instructions  
                  `subl $4, %esp`  
                  `movl %ebx, (%esp)`



- **pop{w,l} dest**

`popl %ebx`    # equivalent instructions  
                  `movl (%esp), %ebx`  
                  `addl $4, %esp`



- **Many more in Intel manual (volume 2)**

- Type conversion, conditional move, exchange, compare and exchange, I/O port, string move, etc.

34

## Conclusions



- **Accessing data**

- Byte, word, and long-word data types
- Wide variety of addressing modes

- **Control flow**

- Common C control-flow constructs
- Condition codes and jump instructions

- **Manipulating data**

- Arithmetic and logic operations

- **Next time**

- Calling functions, using the stack

35