COS 318: Operating Systems

Message Passing

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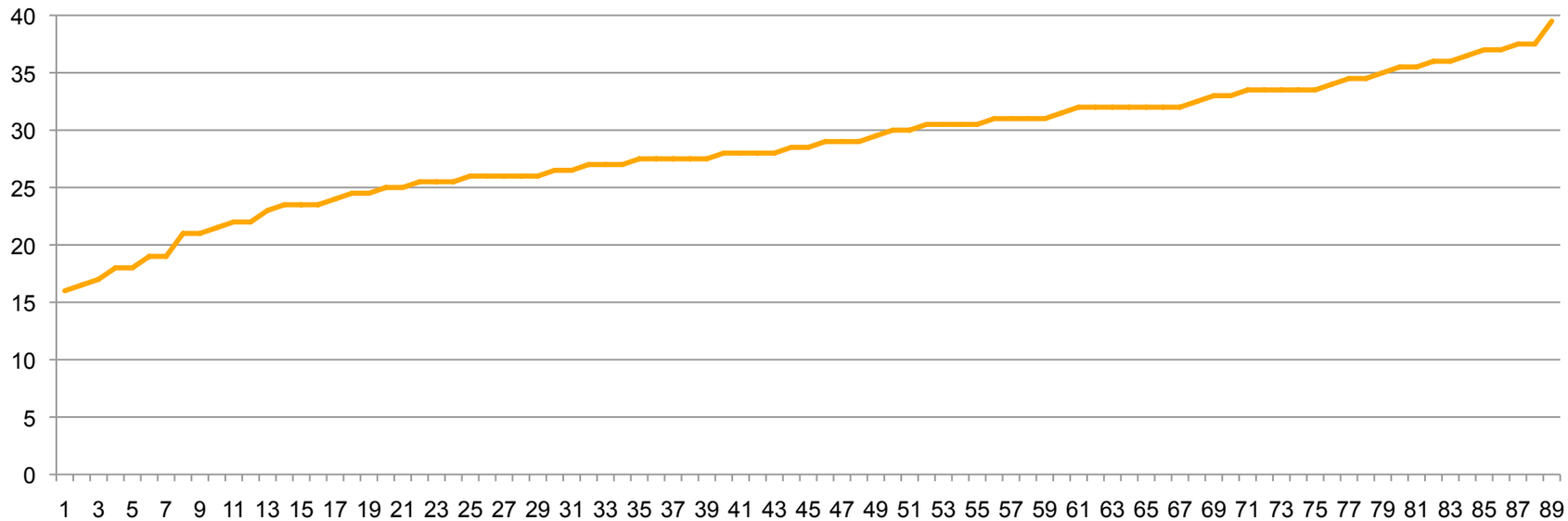
(http://www.cs.princeton.edu/courses/cos318/)

Midterm Results



- Average: 28.58
- Median: 28.5
- See suggested solutions online

Midterm Grading

- Problem 1 (Aaron)
  - Main issue: did not realize that monitors can cause deadlocks
- Problem 2 (Yida)
  - Main issue: definition of turnaround time
- Problem 3 (Scott)
  - Main issue: not thinking about device driver
- Problem 4 (Yida & Aaron)
  - Main issue: making multi multilock atomic can avoid deadlock but inefficient
- Problem 5 (Me)
  - Main issue: not knowing how to program with Mesa-style monitor
- Look for graders outside the classroom

Revisit Idiom of Mesa-style Monitor

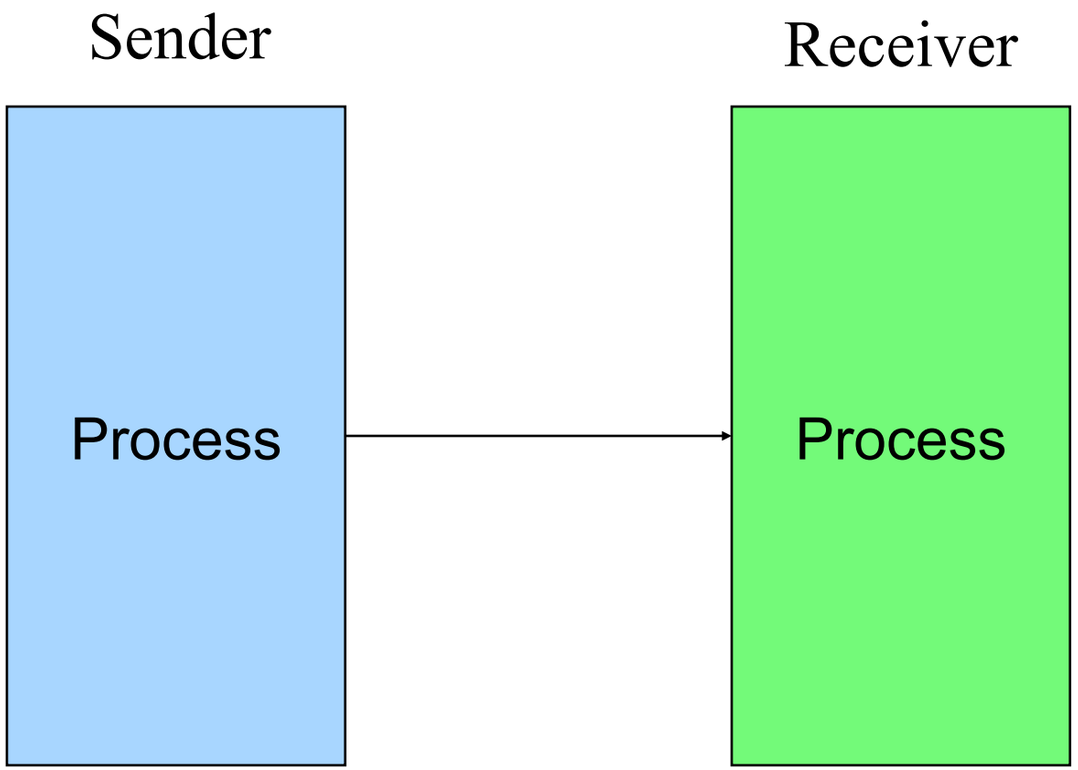
```
Acquire(mutex);
while (“condition not true”)
  wait(mutex, cond);
...
Release(mutex);
```

```
Acquire(mutex);
...
Signal(cond);
/* or Broadcast(cond);
Release(mutex);
```

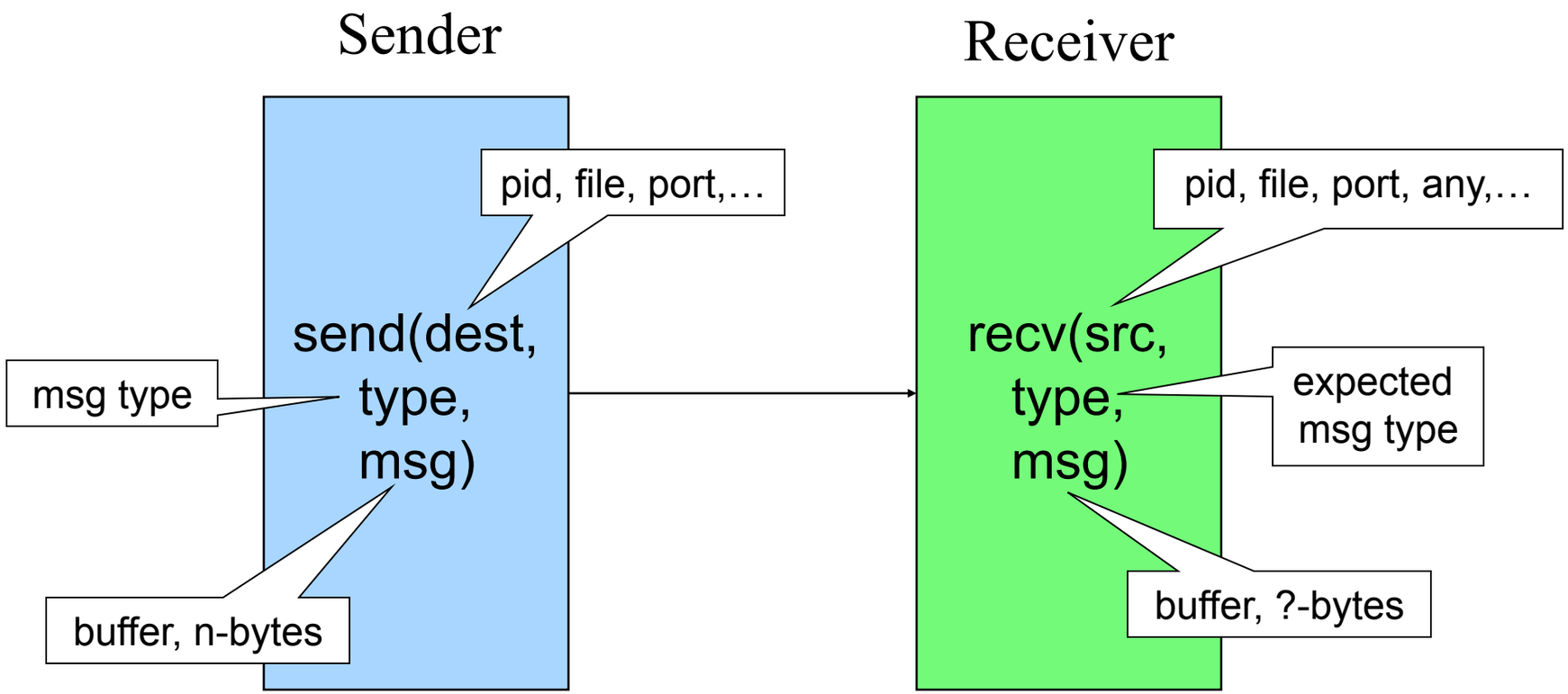
Today's Topics

- Message passing
- Implementation issues

Big Picture

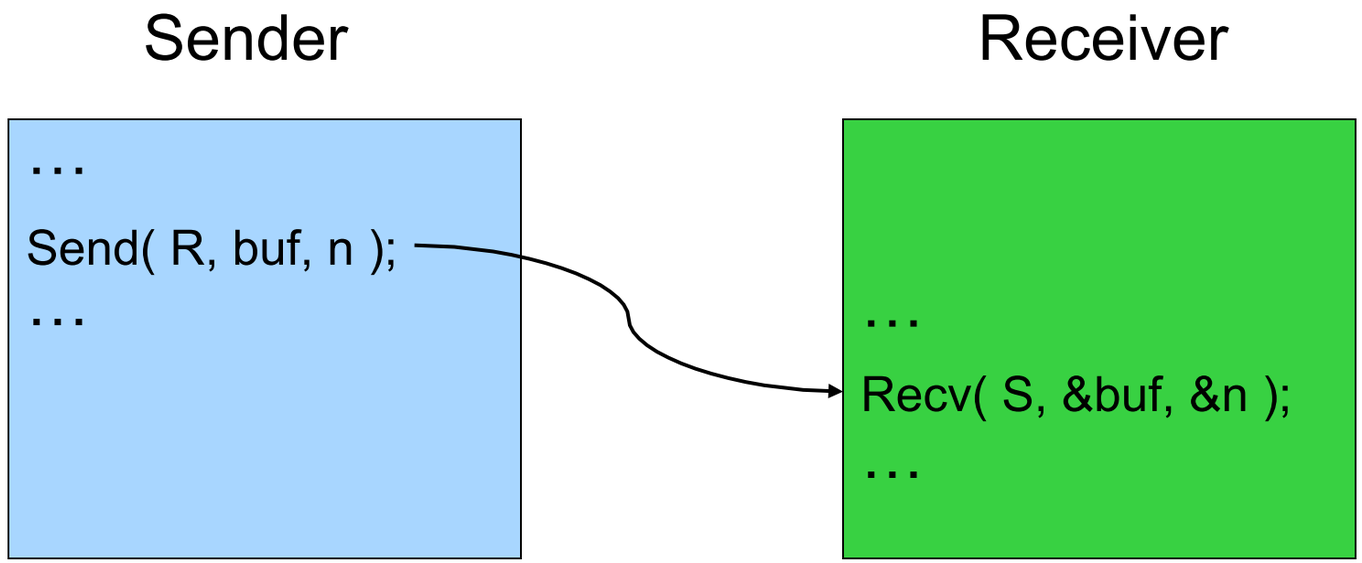


Send and Receive Primitives



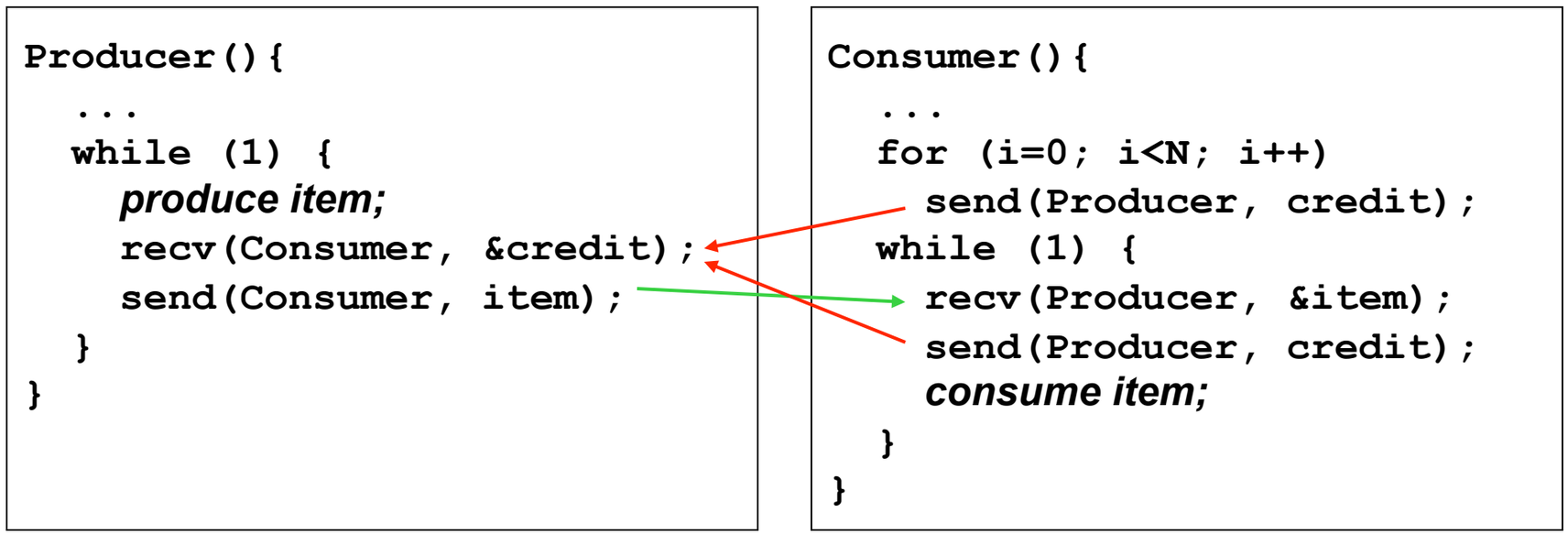
Many ways to design the message passing API

Synchronous Message Passing



- Move data between processes
  - Sender: when data is ready, send it to the receiver process
  - Receiver: when the data has arrived and when the receive process is ready to take the data, move the data
- Synchronization
  - Sender: signal the receiver process that a particular event happens
  - Receiver: block until the event has happened

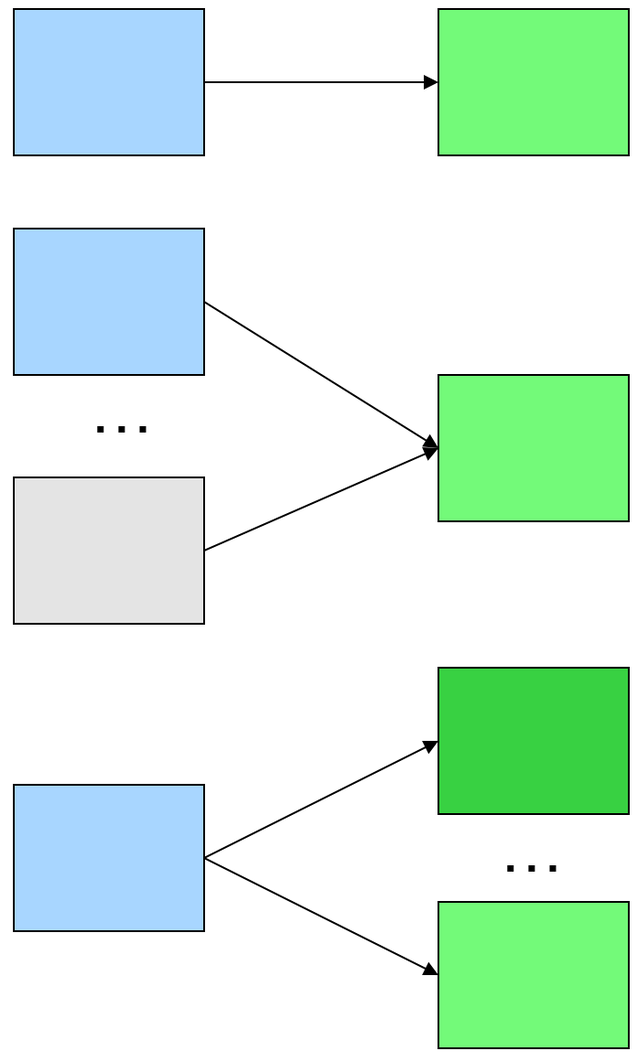
Example: Producer-Consumer



- Does this work?
- Would it work with multiple producers and 1 consumer?
- Would it work with 1 producer and multiple consumers?
- What about multiple producers and multiple consumers?

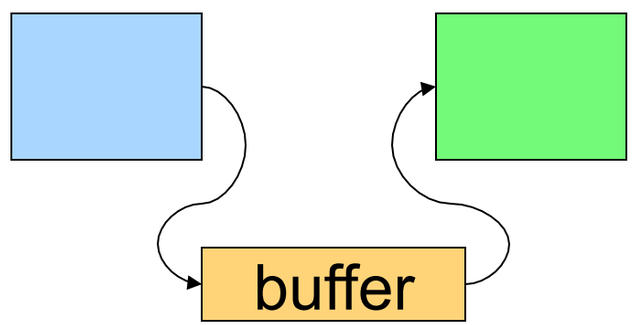
Implementation Issues

- Buffering messages
- Direct vs. indirect
- Unidirectional vs. bidirectional
- Asynchronous vs. synchronous
- Event handler vs. receive
- Handle exceptions



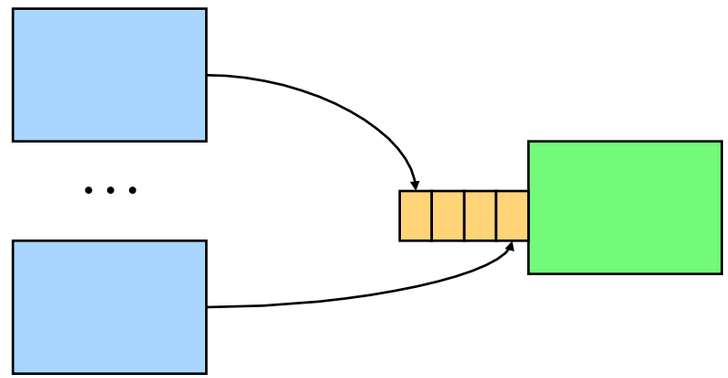
Buffering Messages

- No buffering
  - Sender must wait until the receiver receives the message
  - Rendezvous on each message
- Bounded buffer
  - Finite size
  - Sender blocks on buffer full
  - Use mesa-monitor to solve the problem
- Unbounded buffer
  - “Infinite” size
  - Sender never blocks

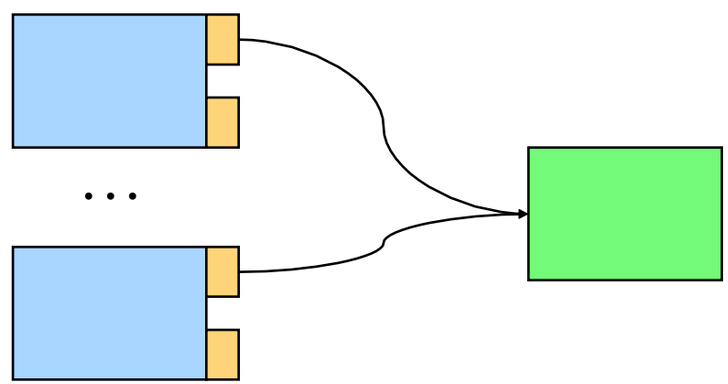


Direct Communication

- A single buffer at the receiver
  - More than one process may send messages to the receiver
  - To receive from a specific sender, it requires searching through the whole buffer

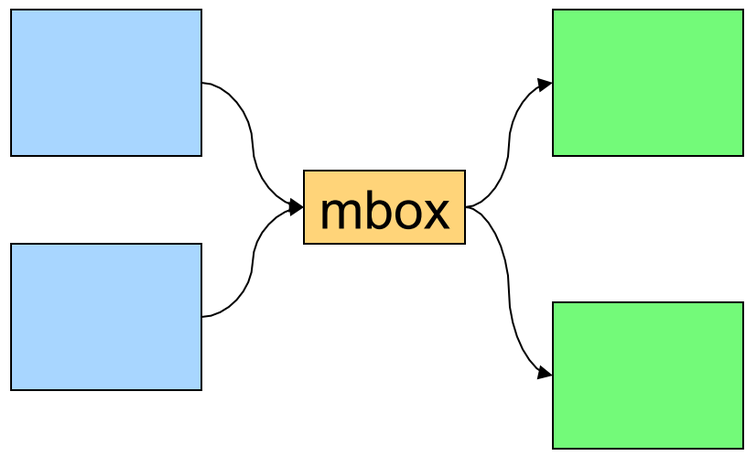


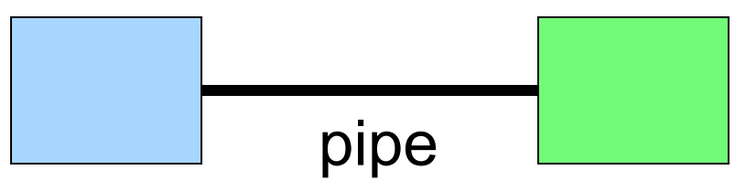
- A buffer at each sender
  - A sender may send messages to multiple receivers
  - To get a message, it also requires searching through the whole buffer



Indirect Communication

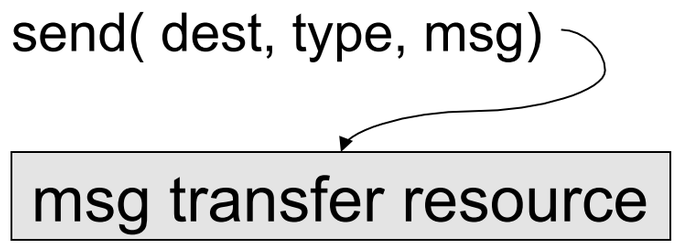
- Use mailbox as the abstraction
  - Allow many-to-many communication
  - Require open/close a mailbox
- Buffering
  - A buffer, its mutex and condition variables should be at the mailbox
- Message size
  - Not necessarily. One can break a large message into packets
- Mailbox vs. pipe
  - A mailbox allows many to many communication
  - A pipe implies one sender and one receiver





Synchronous vs. Asynchronous: Send

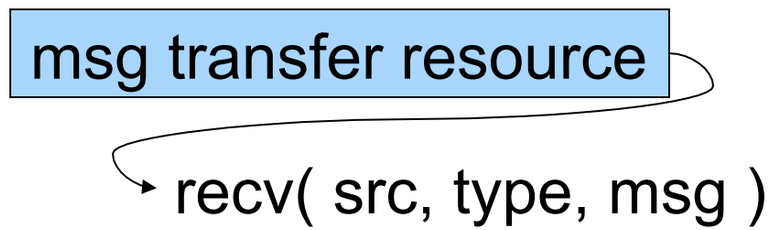
- Synchronous
  - Block on if resource is busy
  - Initiate data transfer
  - Block until data is out of its source memory
- Asynchronous
  - Block if resource is busy
  - Initiate data transfer and return
  - Completion
    - Require applications to check status
    - Notify or signal the application




```
status = async_send( dest, type, msg )
…
if !send_complete( status )
    wait for completion;
…
use msg data structure;
…
```

Synchronous vs. Asynchronous: Receive

- Synchronous
  - Return data if there is a message



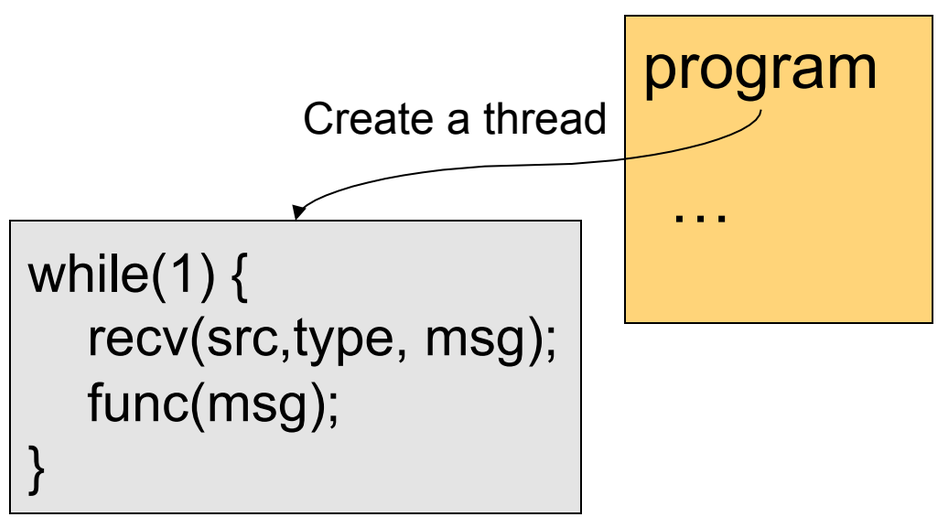
- Asynchronous
  - Return data if there is a message
  - Return status if there is no message (probe)

```
status = async_recv( src, type, msg );
if ( status == SUCCESS )
    consume msg;

while ( probe(src) != HaveMSG )
    wait for msg arrival
recv( src, type, msg );
consume msg;
```

Event Handler vs. Receive

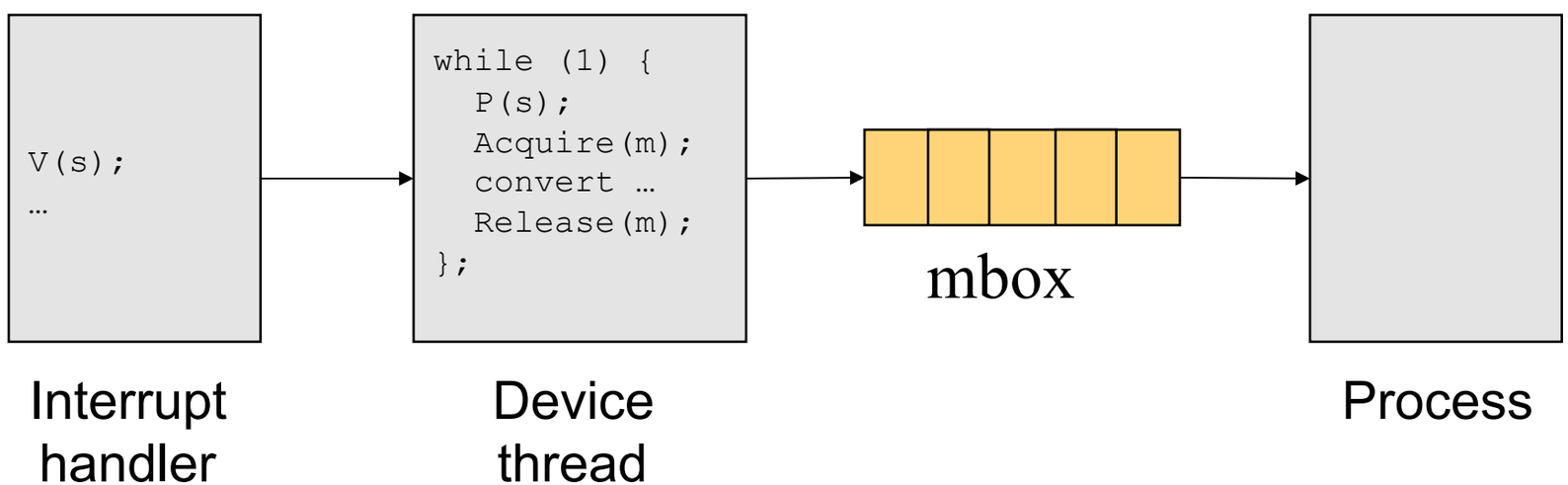
- hrecv(src, type, msg, func)
  - msg is an arg of func
  - Execute “func” on a message arrival
- Which one is more powerful?
  - Recv with a thread can emulate a Handler
  - Handler can be used to emulate recv by using Monitor
- Pros and Cons

```
void func( char * msg ) {
   …
}

…
hrecv( src, type, msg, func)
…
```



Example: Keyboard Input

- How do you implement keyboard input?
  - Need an interrupt handler
  - Generate a mbox message from the interrupt handler
- Suppose a keyboard device thread converts input characters into an mbox message
  - How would you synchronize between the keyboard interrupt handler and device thread?
  - How can a device thread convert input into mbox messages?



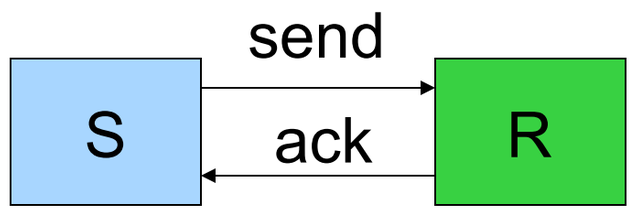
Exception: Process Termination

- R waits for a message from S, but S has terminated
  - R may be blocked forever

- S sends a message to R, but R has terminated
  - S has no buffer and will be blocked forever

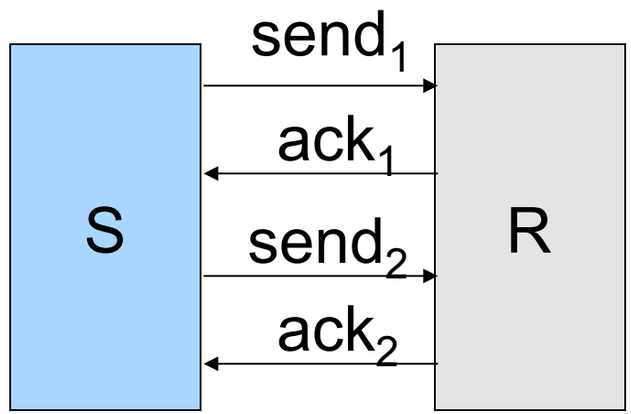
Exception: Message Loss

- Use ack and timeout to detect and retransmit a lost message
  - Receiver sends an ack for each msg
  - Sender blocks until an ack message is back or timeout
    status = send(dest, msg, timeout);
  - If timeout happens and no ack, then retransmit the message
- Issues
  - Duplicates
  - Losing ack messages

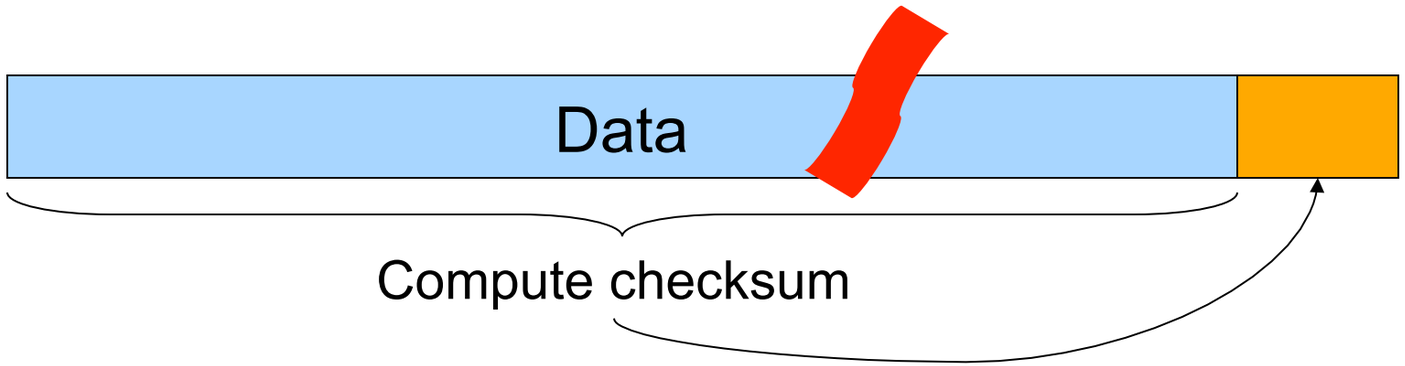


Exception: Message Loss, cont'd

- Retransmission must handle
  - Duplicate messages on receiver side
  - Out-of-sequence ack messages on sender side
- Retransmission
  - Use sequence number for each message to identify duplicates
  - Remove duplicates on receiver side
  - Sender retransmits on an out-of-sequence ack
- Reduce ack messages
  - Bundle ack messages
  - Receiver sends noack messages: can be complex
  - Piggy-back acks in send messages



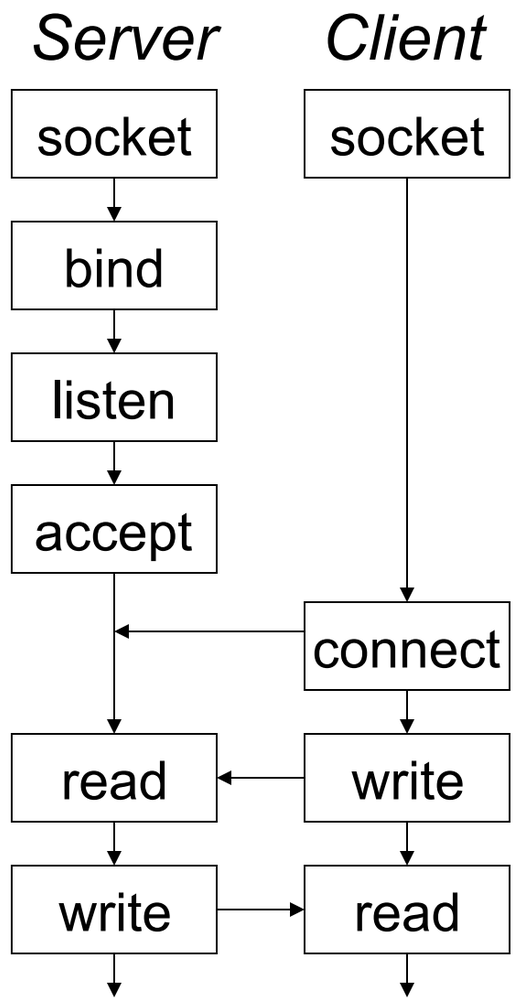
Exception: Message Corruption



- Detection
  - Compute a checksum over the entire message and send the checksum (e.g. CRC code) as part of the message
  - Recompute a checksum on receive and compare with the checksum in the message
- Correction
  - Trigger retransmission
  - Use correction codes to recover

Example: Sockets API

- Abstraction for TCP and UDP
- Addressing
  - IP address and port number (2^{16} ports available for users)
- Create and close a socket
  - `sockid = socket(af, type, protocol);`
  - `Sockerr = close(sockid);`
- Bind a socket to a local address
  - `sockerr = bind(sockid, localaddr, addrlength);`
- Negotiate the connection
  - `listen(sockid, length);`
  - `accept(sockid, addr, length);`
- Connect a socket to destimation
  - `connect(sockid, destaddr, addrlength);`



Summary

- Message passing
  - Move data between processes
  - Implicit synchronization
  - API design is important

- Implementation issues
  - Synchronous method is most common
  - Asynchronous method provides overlapping but requires careful design considerations
  - Indirection makes implementation flexible
  - Exception needs to be carefully handled