

## What is TOY?

## An imaginary machine similar to:

- Ancient computers.
- Today's microprocessors.





# 5. The TOY Machine



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## Why Study TOY?

## Machine language programming.

- How do Java programs relate to computer?
- Key to understanding Java references.
- Still situations today where it is really necessary.

multimedia, computer games, embedded devices, scientific computing, MMX, Altivec

## Computer architecture.

- How does it work?
- How is a computer put together?

TOY machine. Optimized for simplicity, not cost or performance.

#### Inside the Box

Switches. Input data and programs.

Lights. View data.

#### Memory.

- Stores data and programs.
- 256 16-bit "words."
- Special word for stdin / stdout.

#### Program counter (PC).

- An extra 8-bit register.
- Keeps track of next instruction to be executed.

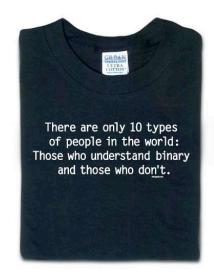
#### Registers.

- Fastest form of storage.
- Scratch space during computation.
- 16 16-bit registers.
- Register 0 is always 0.

Arithmetic-logic unit (ALU). Manipulate data stored in registers.

Standard input, standard output. Interact with outside world.

## Binary People



http://www.thinkgeek.com/tshirts/frustrations/5aa9/zoom/

## Data and Programs Are Encoded in Binary

#### Each bit consists of two states:

- 1 or 0: true or false.
- Switch is on or off; wire has high voltage or low voltage.

## Everything stored in a computer is a sequence of bits.

- Data and programs.
- Text, documents, pictures, sounds, movies, executables, ...



## Binary Encoding

## How to represent integers?

- Use binary encoding.
- Ex:  $6375_{10} = 0001100011100111_2$

Dec	Bin	Dec	Bin
	0000	8	1000
	0001	9	1001
2	0010	10	1010
3	0011	11	1011
	0100	12	1100
5	0101	13	1101
6	0110	14	1110
7	0111	15	1111

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	1	1	0	0	0	1	1	1	0	0	1	1	1

$$6375_{10} = +2^{12} +2^{11} +2^{7} +2^{6} +2^{5} +2^{2} +2^{1} +2^{0}$$

$$= 4096 +2048 +128 +64 +32 +4 +2 +1$$

## Hexadecimal Encoding

## Machine "Core" Dump

#### How to represent integers?

- Use hexadecimal encoding.
- Binary code, four bits at a time.
- Ex: 6375<sub>10</sub> = 0001100011100111<sub>2</sub> = 18E7<sub>16</sub>

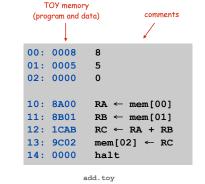
Dec	Bin	Hex	Dec	Bin	Hex
0	0000	0	8	1000	8
1	0001	1	9	1001	9
2	0010	2	10	1010	A
3	0011	3	11	1011	В
4	0100	4	12	1100	С
5	0101	5	13	1101	D
6	0110	6	14	1110	E
7	0111	7	15	1111	F

15 1	4 13	12	11	10	9	8	7	6	5	4	3	2	1	0
0 0	0	1	1	0	0	0	1	1	1	0	0	1	1	1
	1			8	3			]	Ε				7	
6375 <sub>10</sub>	= 1	× 16	3	+	8 ×	16 <sup>2</sup>		+	14 ×	16¹		+	7 ×	160
	= 40	96			т э	048			_	224				+ 7

## A Sample Program

A sample program. Adds 0008 + 0005 = 000D.

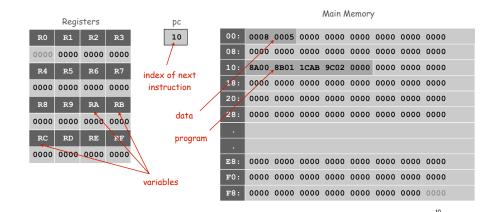




11

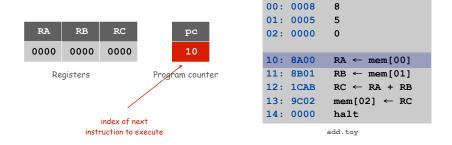
#### Machine contents at a particular place and time.

- Record of what program has done.
- Completely determines what machine will do.



## A Sample Program

Program counter. The pc is initially 10, so the machine interprets 8A00 as an instruction.

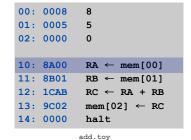


Load Load

#### Load. [opcode 8]

- Loads the contents of some memory location into a register.
- 8A00 means load the contents of memory cell 00 into register A.

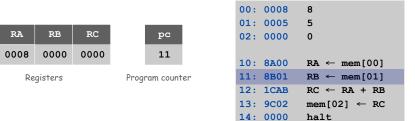




0 0 0 816  $A_{16}$ 0016 dest d addr opcode

### Load. [opcode 8]

- Loads the contents of some memory location into a register.
- $\bullet$  8B01 means load the contents of memory cell 01 into register B.



add.toy

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	0	0	0	1	0	1	1	0	0	0	0	0	0	0	1
	8	16			В	16					01	16			
	opc	ode			des	t d					ad	dr			

16

#### Add

### Add. [opcode 1]

- Add contents of two registers and store sum in a third.
- 1CAB means add the contents of registers  ${\tt A}$  and  ${\tt B}$  and put the result into register c.

RA	RB	RC
8000	0005	0000

Registers

12

Program counter

01: 0005 02: 0000 10: 8A00  $RA \leftarrow mem[00]$ 11: 8B01 mem[01] 12: 1CAB RC ← RA + RB 13: 9C02  $mem[02] \leftarrow RC$ 14: 0000 halt

00: 0008

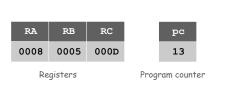
add.toy

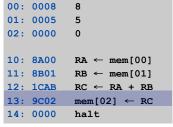
15	14	13	12	11	10	9		7	6	5	4	3	2	1			
0	0	0	1	1	1	0	0	1	0	1	0	1	0	1	1		
	1	16			С	16			А	16			В	16			
	opc	ode		dest d					sour	ce s		source t					

#### Store

### Store. [opcode 9]

- Stores the contents of some register into a memory cell.
- 9C02 means store the contents of register C into memory cell 02.





add.toy

15	14	13	12	11	10			7			4	3	2		0
1	0	0	1	1	1	0	0	0	0	0	0	0	0	1	0
	9	16			С	16					02	216			
	opc	ode			des	t d					ado	dr			

## Program and Data

Halt. [opcode 0] • Stop the machine.

RA	RB	RC
8000	0005	000D



Halt

Registers Program counter

```
00: 0008
01: 0005
             5
02: 000D
10: 8A00
             RA \leftarrow mem[00]
11: 8B01
             RB \leftarrow mem[01]
12: 1CAB
             RC ← RA + RB
13: 9C02
             mem[02] \leftarrow RC
14: 0000
             halt
```

add.toy

Program. Sequence of 16-bit integers, interpreted one way.

Data. Sequence of 16-bit integers, interpreted other way.

Program counter (pc). Holds memory address of the "next instruction" and determines which integers get interpreted as instructions.

16 instruction types. Changes contents of registers, memory, and pc in specified, well-defined ways.



jump and link

### TOY Instruction Set Architecture

### TOY instruction set architecture (ISA).

- Interface that specifies behavior of machine.
- 16 register, 256 words of main memory, 16-bit words.
- 16 instructions.

#### Each instruction consists of 16 bits.

- Bits 12-15 encode one of 16 instruction types or opcodes.
- Bits 8-11 encode destination register d.
- Bits 0-7 encode:

[Format 1] source registers s and t

[Format 2] 8-bit memory address or constant

	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	1	0	1	1	1	0	1	0	0	0	0	0	0	1	0	0
Format 1		opo	ode			des	t d			sour	ce s			sour	ce t	
Format 2		opo	ode			des	t d					ad	ldr			

### TOY Reference Card

	15			12	11			7	6			3	2		0
Format 1		opo	ode			des	t d		sour	ce s			sour	ce t	
Format 2		opc	ode			des	t d				ac	ldr			

#	Operation	Fmt	Pseudocode
0:	halt	1	exit(0)
1:	add	1	$R[d] \leftarrow R[s] + R[t]$
2:	subtract	1	$R[d] \leftarrow R[s] - R[t]$
3:	and	1	$R[d] \leftarrow R[s] \& R[t]$
4:	xor	1	$R[d] \leftarrow R[s] ^ R[t]$
5:	shift left	1	$R[d] \leftarrow R[s] \iff R[t]$
6:	shift right	1	$R[d] \leftarrow R[s] >> R[t]$
7:	load addr	2	R[d] ← addr
8:	load	2	R[d] ← mem[addr]
9:	store	2	mem[addr]
A:	load indirect	1	$R[d] \leftarrow mem[R[t]]$
B:	store indirect	1	$mem[R[t]] \leftarrow R[d]$
C:	branch zero	2	if $(R[d] == 0)$ pc $\leftarrow$ addr
D:	branch positive	2	if $(R[d] > 0)$ pc $\leftarrow$ addr
E:	jump register	2	pc ← R[d]
F:	jump and link	2	R[d] ← pc; pc ← addr

Register 0 always 0. Loads from mem[FF] from stdin. Stores to mem[FF] to stdout.

19

17

## TEQ on TOY 1

What is the interpretation of 1A75

A. as a TOY instruction?

B. as an integer value?

Using the TOY Machine: Run

## To run the program:

- Set 8 memory address switches to address of first instruction.
- Press Look to set pc to first instruction.
- Press Run button to repeat fetch-execute cycle until halt opcode.

## Fetch-execute cycle.

- Fetch: get instruction from memory.
- Execute: update pc move data to or from memory and registers, perform calculations.



21

## Interfacing with the TOY Machine

## To enter a program or data:

- Set 8 memory address switches.
- Set 16 data switches.
- Press Load: data written into addressed word of memory.

## To view the results of a program:

- Set 8 memory address switches.
- Press Look: contents of addressed word appears in lights.



Flow Control

#### Flow control.

- To harness the power of TOY, need loops and conditionals.
- Manipulate pc to control program flow.

## Branch if zero. [opcode C]

- Changes pc depending on whether value of some register is zero.
- Used to implement: for, while, if-else.

#### Branch if positive. [opcode D]

- ullet Changes pc depending on whether value of some register is positive.
- Used to implement: for, while, if-else.

## An Example: Multiplication

Multiply. Given integers a and b, compute  $c = a \times b$ .

TOY multiplication. No direct support in TOY hardware.

### Brute-force multiplication algorithm:

- Initialize c to 0.
- Add b to c, a times.

```
int a = 3;
int b = 9;
int c = 0;
while (a != 0) {
   c = c + b;
   a = a - 1;
}
```

brute force multiply in Java

Issues ignored. Slow, overflow, negative numbers.

## Step-By-Step Trace

```
_R1_
                                                                  RB
                                                         RA
                                                                            RC
10: 8A0A
                                                        0003
            RA \leftarrow mem[0A]
11: 8B0B
            RB ← mem[0B]
                                                                  0009
            RC ← mem[OD]
                                                                           0000
12: 8COD
13: 810E
            R1 \leftarrow mem[0E]
                                               0001
14: CA18
            if (RA == 0) pc \leftarrow 18
15: 1CCB
            RC ← RC + RB
                                                                           0009
                                                        0002
16: 2AA1
            RA \leftarrow RA - R1
            pc ← 14
17: C014
14: CA18
            if (RA == 0) pc \leftarrow 18
15: 1CCB
            RC ← RC + RB
                                                                            0012
                                                        0001
16: 2AA1
            RA \leftarrow RA - R1
            pc ← 14
17: C014
14: CA18
            if (RA == 0) pc ← 18
15: 1CCB
            RC ← RC + RB
                                                                           001B
16: 2AA1
            RA \leftarrow RA - R1
                                                        0000
17: C014
            pc ← 14
14: CA18
            if (RA == 0) pc \leftarrow 18
            mem[OC] ← RC
18: 9COC
19: 0000
            halt
```

## Multiply

```
0A: 0003
                  ← inputs
0B: 0009
OC: 0000
                 ← output
OD: 0000
                 ← constants
OE: 0001
10: 8A0A
             RA \leftarrow mem[0A]
11: 8B0B
             RB \leftarrow mem[0B]
12: 8C0D
             RC \leftarrow mem[0D]
                                             c = 0
13: 810E
             R1 \leftarrow mem[0E]
                                             always 1
14: CA18
             if (RA == 0) pc ← 18
                                             while (a != 0) {
15: 1CCB
             RC \leftarrow RC + RB
                                                c = c + b
16: 2AA1
             RA \leftarrow RA - R1
                                                 a = a - 1
17: C014
             pc ← 14
18: 9C0C
             mem[0C] \leftarrow RC
19: 0000
             halt
```

multiply.toy

26

## TEQ on TOY 2

What does the following TOY program leave in R2?

```
10: 7C0A
11: 7101
12: 7201
13: 5221
14: 2CC1
15: DC13
16: 0000
```

## A Little History

#### Electronic Numerical Integrator and Calculator (ENIAC).

- First widely known general purpose electronic computer.
- Conditional jumps, programmable.
- Programming: change switches and cable connections.
- Data: enter numbers using punch cards.

30 tons  $30 \times 50 \times 8.5 \text{ ft}$ 300 multiply/sec



John Mauchly (left) and J. Presper Eckert (right) http://cs.swau.edu/~durkin/articles/history\_computing.html



ENIAC, Ester Gerston (left), Gloria Gordon (right) US Army photo: http://ftp.arl.mil/ftp/historic-computers

# 17,468 vacuum tubes

## Harvard vs. Princeton

#### Harvard architecture.

- Separate program and data memories.
- Can't load game from disk (data) and execute (program).
- Used in some microcontrollers.

#### Von Neumann architecture.

- Program and data stored in same memory.
- Used in almost all computers.



- Q. What's the difference between Harvard and Princeton?
- A. At Princeton, data and programs are the same.

#### Basic Characteristics of TOY Machine

## TOY is a general-purpose computer.

- Sufficient power to perform ANY computation.
- Limited only by amount of memory and time.

## Stored-program computer. [von Neumann memo, 1944]

- Data and program encoded in binary.
- Data and program stored in SAME memory.
- Can change program without rewiring.

### Outgrowth of Alan Turing's work. (stay tuned)

All modern computers are general-purpose computers and have same (von Neumann) architecture.





Maurice Wilkes (left) EDSAC (right)