

**Due:** Wednesday October 12, in class.

1. The goal of this exercise is to work out an axiomatic definition of entropy. In other words, we start with some desirable properties of the entropy function, and show that  $H(\cdot)$  is the only function satisfying these properties — and hence the formula for  $H(\cdot)$  follows from the specified axioms. Let  $f$  be a function that takes a random variable  $X$  on finite support and outputs a real number with the following properties:

- $f(X)$  only depends on the frequencies of the different values of  $X$ .
- If  $X$  is a uniformly random point from a set of size  $M$ , and  $Y$  is a uniformly random point from a set of size  $M' > M$ , then  $f(M') > f(M)$ .
- If  $X, Y$  are independent, then  $f(X, Y) = f(X) + f(Y)$ .
- If  $B_q$  is such that  $\Pr[B_q = 1] = q$  and  $\Pr[B_q = 0] = 1 - q$ , then  $f(B_q)$  is a continuous function of  $q$ .
- If  $B$  is a random variable taking 0/1 values, and  $X$  is another random variable, then  $f(BX) = f(B) + \Pr[B = 1] \cdot f(X|B = 1) + \Pr[B = 0] \cdot f(X|B = 0)$ .
- $f(B_{1/2}) = 1$ .

Show that  $f(X) = H(X)$  for all finitely supported  $X$ . *Hint:* start with uniform  $X$ 's, then proceed to  $f(B_q)$ .

2. (Problem 2.25 in the book). There isn't really a notion of mutual information common to three random variables. Here is one attempt at a definition: Using Venn diagrams, we can see that the mutual information common to three random variables  $X, Y$ , and  $Z$  can be defined by

$$I(X; Y; Z) = I(X; Y) - I(X; Y|Z).$$

The quantity is symmetric in  $X, Y$ , and  $Z$  despite the preceding asymmetric definition (this fact will follow from the identities below). Unfortunately,  $I(X; Y; Z)$  is not necessarily non-negative. Find  $X, Y$ , and  $Z$  such that  $I(X; Y; Z) < 0$ , and prove the following two identities.

- (a)  $I(X; Y; Z) = H(XYZ) - H(X) - H(Y) - H(Z) + I(X; Y) + I(Y; Z) + I(Z; X)$ .  
 (b)  $I(X; Y; Z) = H(XYZ) - H(XY) - H(YZ) - H(ZX) + H(X) + H(Y) + H(Z)$ .

The first identity can be understood using the Venn diagram analogy for entropy and mutual information. The second identity follows easily from the first.

3. (Problem 2.39 in the book, extended). Let  $X, Y, Z$  be three Bernoulli(1/2) random variables that are pairwise independent:  $I(X; Y) = I(Y; Z) = I(Z; X) = 0$ .
- (a) Under this constraint, what is the minimum value for  $H(XYZ)$ ?  
 (b) Give an example achieving this minimum.  
 (c) What would the minimum value for  $H(XYZ)$  be if the constraint above was replaced with  $I(X; Y) = I(Y; Z) = I(Z; X) = \alpha$ , for some  $0 \leq \alpha \leq 1$ ?  
 (d) Show (by giving an example, or otherwise) that your bound from the previous part of the question is tight.