

#### Machine Language, Assemblers and Linkers

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#### **Goals for this Lecture**



- Help you to learn about:
  - IA-32 machine language
  - The assembly and linking processes

# **Why Learn Machine Language**



- Last stop on the "language levels" tour
- A power programmer knows the relationship between assembly and machine languages
- A systems programmer knows how an assembler translates assembly to machine language

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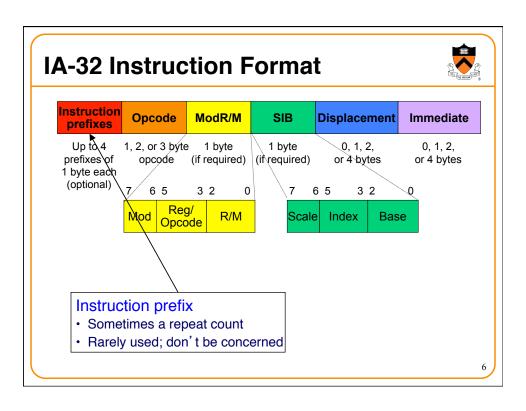


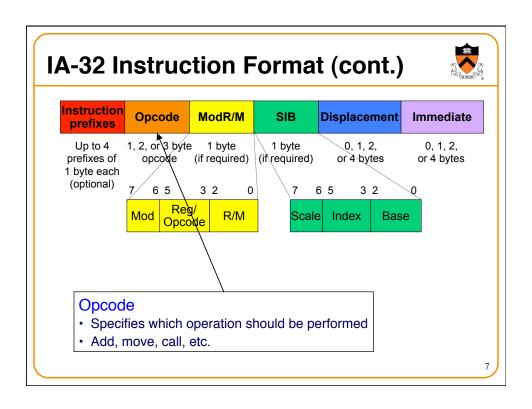
Part 1: Machine Language

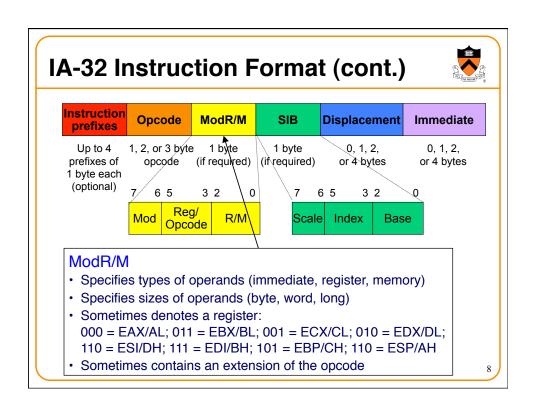
#### IA-32 Machine Language

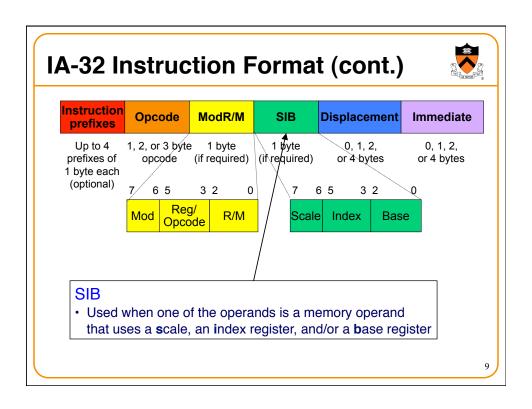


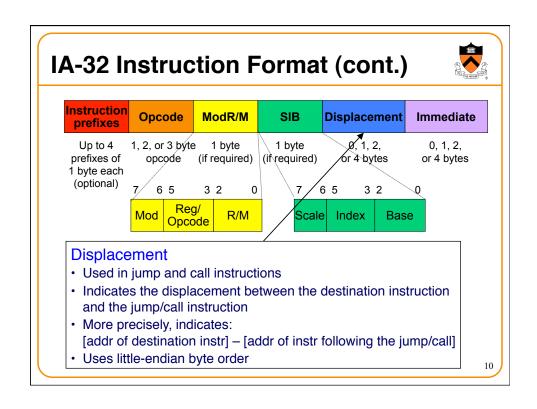
- IA-32 machine language
  - · Difficult to generalize about IA-32 instruction format
    - · Many (most) instructions are exceptions to the rules
  - · Generally, instructions use format in following slides
- · We'll go over
  - · The format of instructions
  - Two example instructions
- Just to convey a sense of how it works...

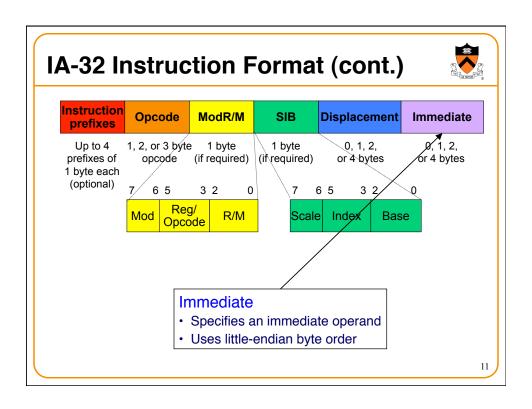












#### **Example: Push on to Stack**



· Assembly language:

#### pushl %edx

- · Machine code:
  - · IA32 has a separate opcode for push for each register operand
    - 50: pushl %eax
    - 51: pushl %ecx
    - 52: pushl %edx ——— 0101 0010

· Results in a one-byte instruction

 Sometimes one assembly language instruction can map to a group of different opcodes

#### **Example: Load Effective Address**



Assembly language:

leal (%eax,%eax,4), %eax

- Machine code:
  - Byte 1: 8D (opcode for "load effective address")
  - Byte 2: 04 (dest %eax, with scale-index-base)
  - Byte 3: 80 (scale=4, index=%eax, base=%eax)

1000 1101

0000 0100

1000 0000

Load the address %eax + 4 \* %eax into register %eax

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#### **CISC and RISC**



- IA-32 machine language instructions are **complex**
- IA-32 is a
  - Complex Instruction Set Computer (CISC)
- Alternative:
  - Reduced Instruction Set Computer (RISC)

#### **Characteristics of CISC and RISC**



- CISC
  - · Many instructions
  - Many addressing modes (direct, indirect, indexed, base-pointer)
  - · Hardware interpretation is complex
  - Few instructions required to accomplish a given job (expressive)
  - · Example: IA-32
- RISC
  - · Few instructions
  - Few addressing modes (typically only direct and indirect)
  - Hardware interpretation is simple
  - Many instructions required to accomplish a given job (not expressive)
  - · Examples: MIPS, SPARC

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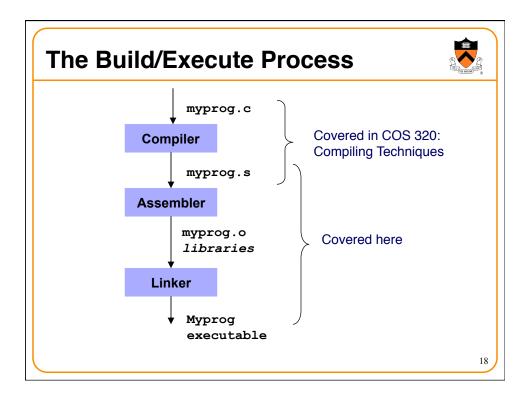
#### **Brief History of CISC and RISC**



- Stage 1: Programmers write assembly language
  - · Important that assembly/machine language be expressive
  - · CISC dominates (esp. Intel)
- Stage 2: Programmers write high-level language
  - Not important that assembly/machine language be expressive; the compiler generates it
  - Important that compilers work well => assembly/machine language should be simple
  - RISC takes a foothold (but CISC, esp. Intel, persists)
- Stage 3: Compilers get smarter
  - · Less important that assembly/machine language be simple
  - · Hardware is plentiful, enabling complex implementations
  - Much motivation for RISC disappears
  - · CISC (esp. Intel) dominates the computing world



# Part 2: The Assembly Process



#### Two Aspects of the Assembler/Linker



- Translating each instruction
  - Mapping an assembly-language instruction into the corresponding machinelanguage instruction
- Dealing with references across instructions
  - Jumps to other locations in same chunk (file) of code
  - Accesses a global variable by the name of its memory location
  - Calling to and returning from functions defined in other code

```
main:
        pushl
                 %ebp
        movl
                 %esp, %ebp
        call
                 getchar
                 $'A', %eax
        cmpl
        jne
                 skip
        pushl
                 $msg
        call
                 printf
        addl
                 $4, %esp
skip:
                 $0, %eax
        movl
                 %ebp, %esp
                 %ebp
        popl
        ret
```

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#### **References Across Instructions**



- Many instructions can be assembled independently
  - pushl %edx
  - leal (%eax, %eax, 4), %eax
  - movl \$0, %eax
  - · addl %ebx, %ecx
- But, some make references to other data or code
  - jne skip
  - · pushl \$msg
  - call printf
- Need to fill in those references
  - To generate a final executable binary

#### **The Forward Reference Problem**



Problem

...
jmp mylabel
...
mylabel:
...

Any assembler must deal with the forward reference problem

- Assembler must generate machine language code for "jmp mylabel"
- But assembler hasn't yet seen the definition of mylabel
  - I.e., the jmp instruction contains a forward reference to mylabel

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#### **The Forward Reference Solution**



- Solution
  - Assembler performs 2 passes over assembly language program
  - · One to record labels and addresses, e.g.
  - · Another to patch them in
- Different assemblers perform different tasks in each pass
- One straightforward design...

#### **Assembler Passes**



- Pass1
  - · Assembler traverses assembly program to create...
  - · Symbol table
    - Key: label
    - · Value: information about label
      - · Label name, which section, what offset within that section, ...
- Pass 2
  - Assembler traverses assembly program again to create...
  - · RODATA section
  - · DATA section
  - · BSS section
  - · TEXT section

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#### **References Outside the File**



- e.g. Call printf
  - printf is defined in another library; assembler doesn't know at what address that code will be placed in memory
- Defer to linker
  - Pass 1 defers symbol resolution to pass 2, pass 2 defers some to linker
- So, Pass 2 becomes...
  - · Assembler traverses assembly program again to create...
  - RODATA section
  - · DATA section
  - · BSS section
  - TEXT section
  - · Relocation record section
    - · Each relocation record indicates something that linker must patch

### **An Example Program**



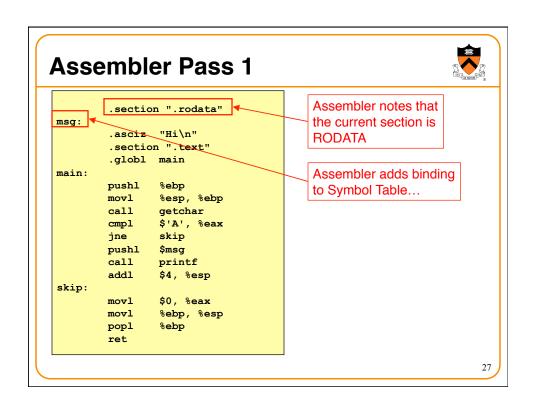
 A simple (nonsensical) program:

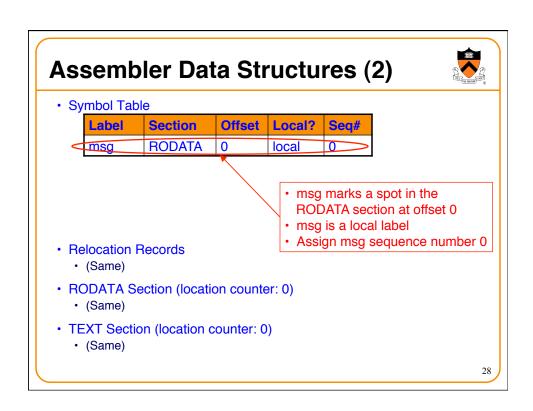
```
#include <stdio.h>
int main(void) {
   if (getchar() == 'A')
      printf("Hi\n");
   return 0;
}
```

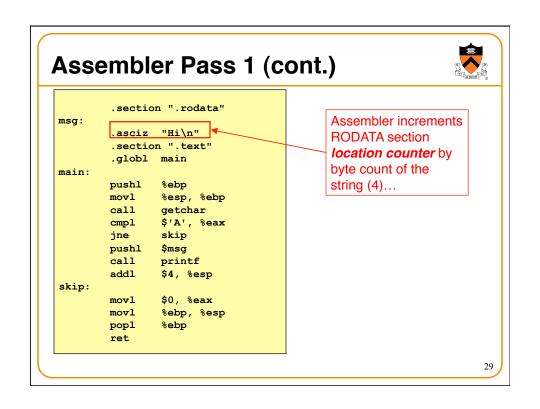
 Let's consider how the assembler handles that program...

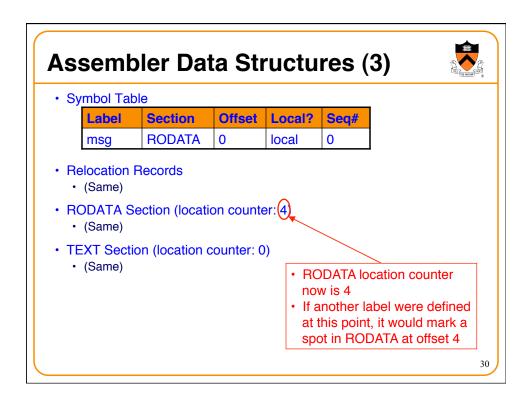
```
.section ".rodata"
msg:
        .asciz "Hi\n"
        .section ".text"
        .globl main
main:
        pushl
                %ebp
        movl
                %esp, %ebp
        call
                getchar
                $'A', %eax
        cmpl
                skip
        jne
        pushl
                $msg
                printf
        call
        addl
                $4, %esp
skip:
                $0, %eax
        movl
        movl
                %ebp, %esp
        popl
                %ebp
        ret
                                   25
```

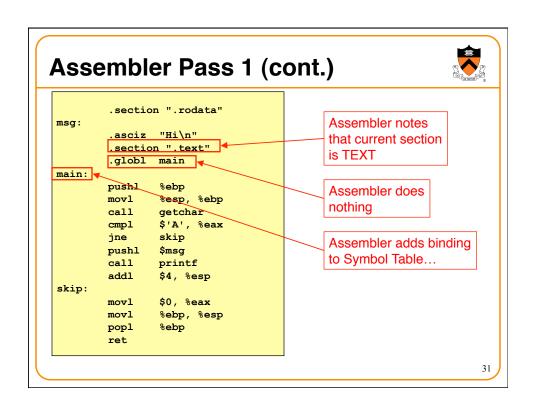
**Assembler Data Structures (1)**  Symbol Table Label **Section** Offset Local? Seq# · Relocation Records Seq# Section Offset **Rel Type** • RODATA Section (location counter: 0) No DATA or BSS Contents Explanation section in this program · Initially all sections are empty • TEXT Section (location counter: 0) Offset | Contents **Explanation** 26

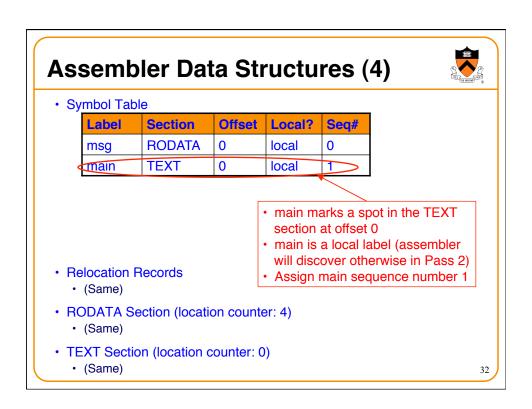


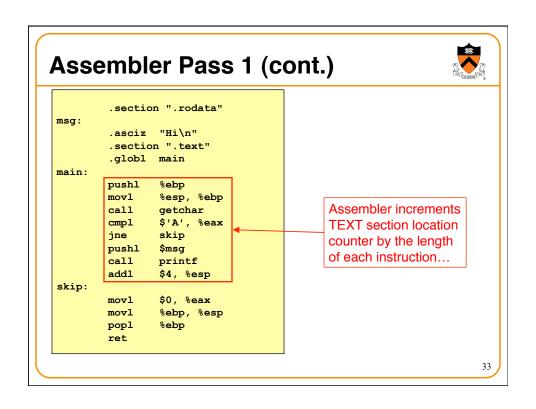


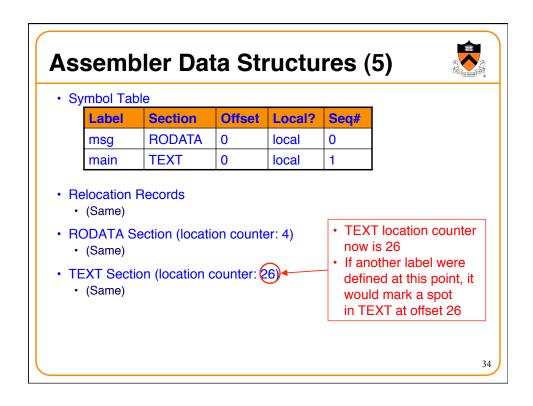


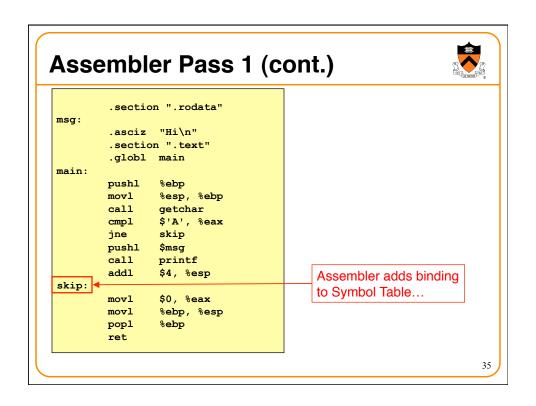


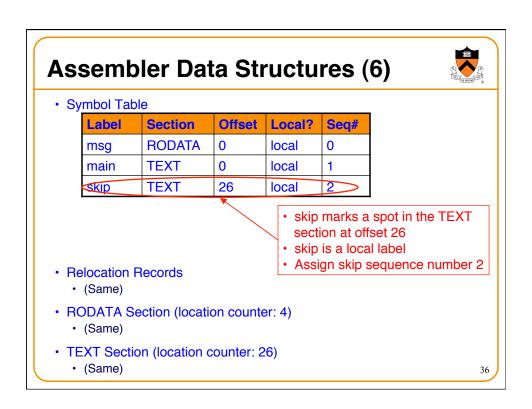


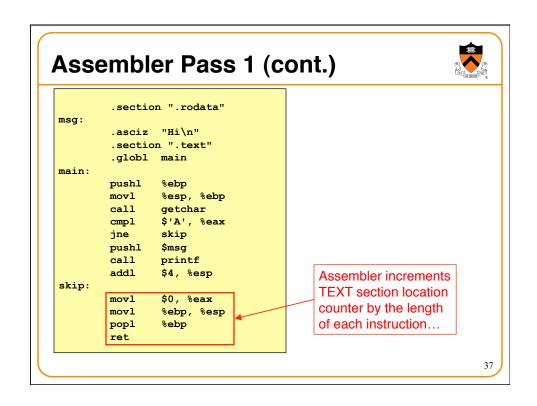












#### **Assembler Data Structures (7)**



· Symbol Table

Label	Section	Offset	Local?	Seq#
msg	RODATA	0	local	0
main	TEXT	0	local	1
skip	TEXT	26	local	2

- · Relocation Records
  - (Same)
- RODATA Section (location counter: 4)
  - (Same)
- TEXT Section (location counter: 35)
  - (Same)

- TEXT location counter now is 35
- If another label were defined at this point, it would mark a spot in TEXT at offset 35

# From Assembler Pass 1 to Pass 2



- End of Pass 1
  - · Assembler has (partially) created Symbol Table
  - · So assembler now knows which location each label marks
- Beginning of Pass 2
  - · Assembler resets all section location counters...

#### **Assembler Data Structures (8)**

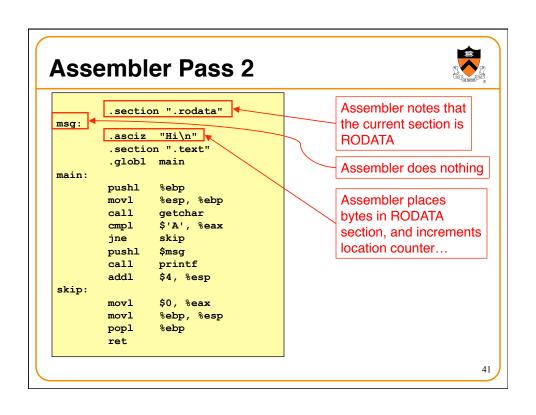


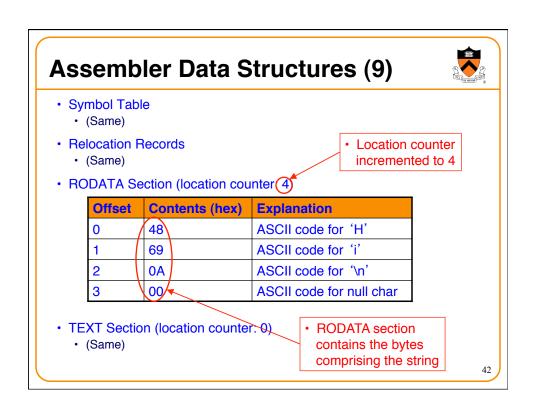
· Symbol Table

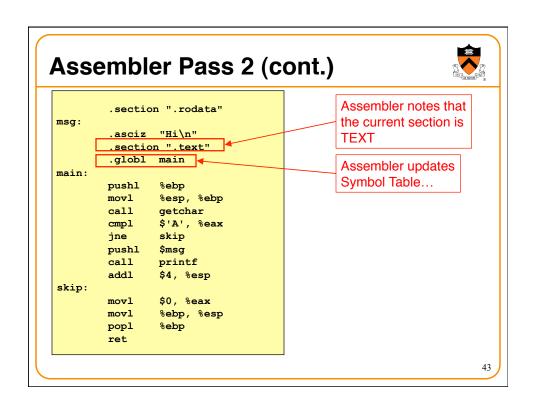
Label	Section	Offset	Local?	Seq#
msg	RODATA	0	local	0
main	TEXT	0	local	1
skip	TEXT	26	local	2

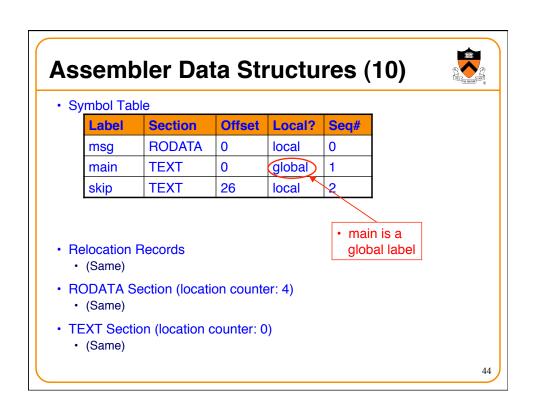
- Relocation Records
  - (Same)
- RODATA Section (location counter 0)
  - (Same)
- TEXT Section (location counter: 0)
  - (Same)

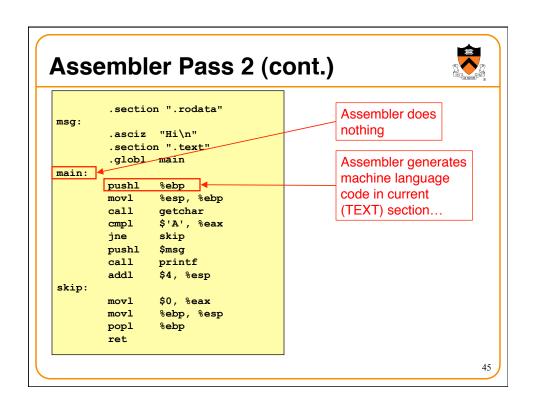
 Location counters reset to 0



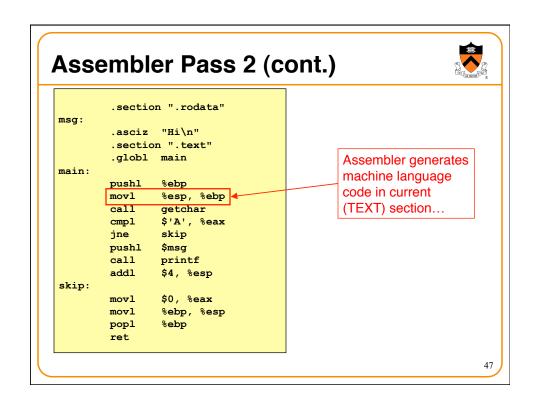








# Assembler Data Structures (11) • Symbol Table • (Same) • Relocation Records • (Same) • RODATA Section (location counter: 4) • (Same) • TEXT Section (location counter: 1) Offset Contents Explanation 0 55 pushl %ebp 01010101 This is a "pushl %ebp" instruction

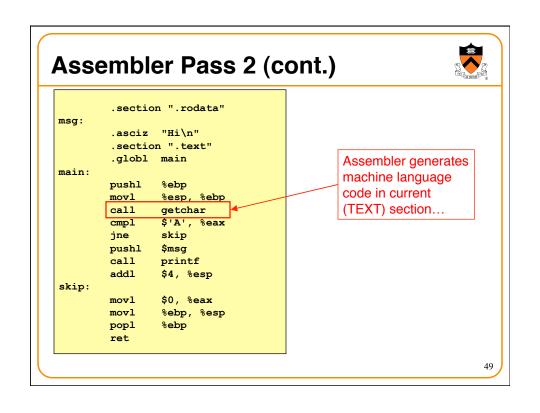


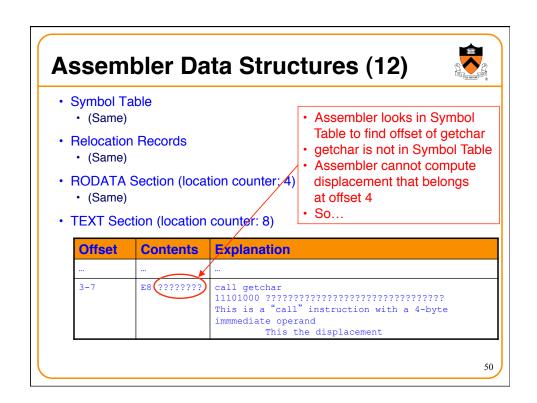
# **Assembler Data Structures (12)**

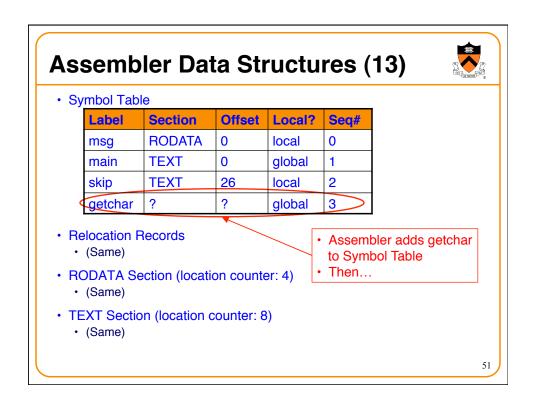


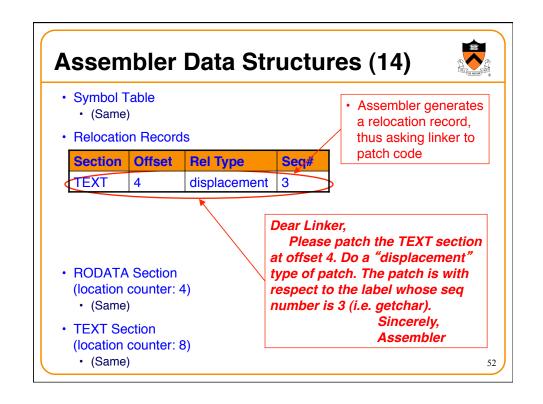
- Symbol Table
  - (Same)
- Relocation Records
  - · (Same)
- RODATA Section (location counter: 4)
  - · (Same)
- TEXT Section (location counter: 3)

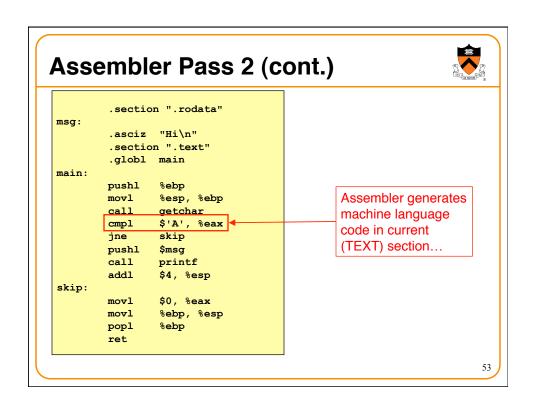
Offset	Contents	Explanation
1-2	89 E5	movl %esp, %ebp 10001001 11 100 101 This is a "movl" instruction whose source operand is a register The M field designates a register The source register is ESP The destination register is EBP









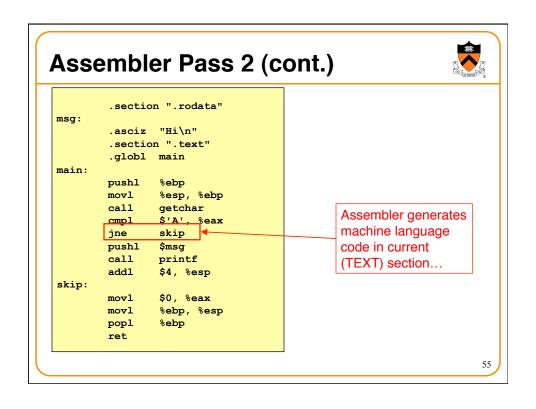


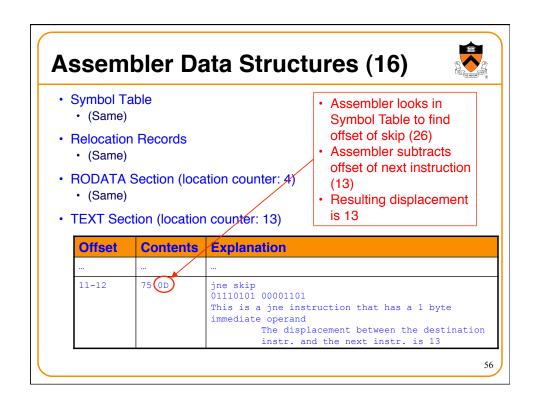
# **Assembler Data Structures (15)**

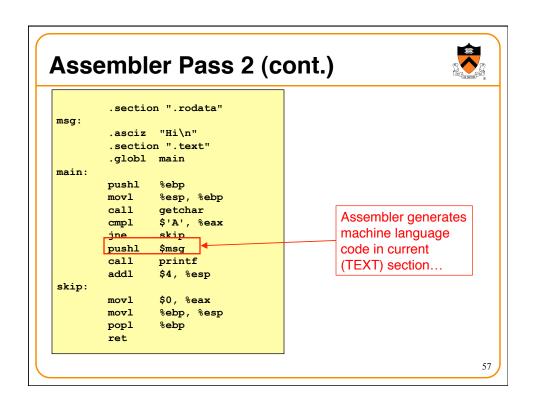


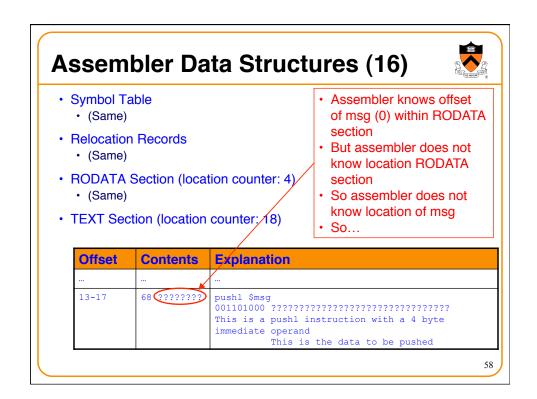
- Symbol Table
  - (Same)
- · Relocation Records
  - · (Same)
- RODATA Section (location counter: 4)
  - · (Same)
- TEXT Section (location counter: 11)

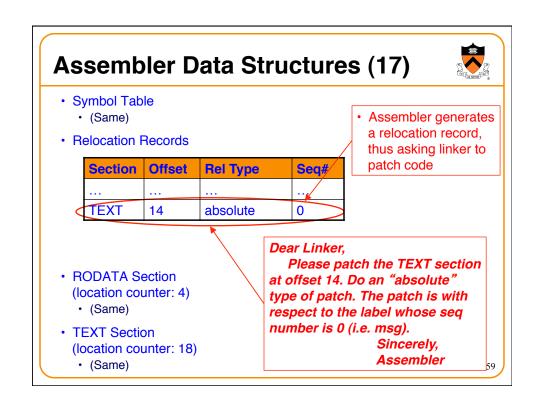
Offset	Contents	Explanation
8-10	83 F8 41	cmpl %'A',%eax 10000011 11 111 000 01000001 This is some "1" instruction that has a 1 byte immediate operand  The M field designates a register  This is a "cmp" instruction  The destination register is EAX  The immediate operand is 'A' 54

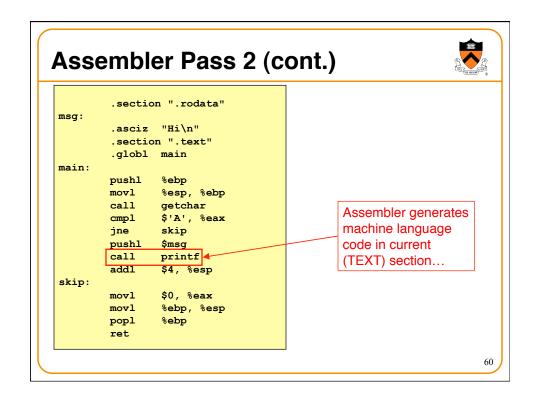


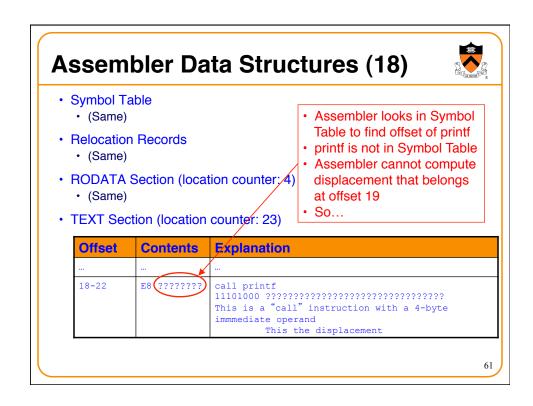




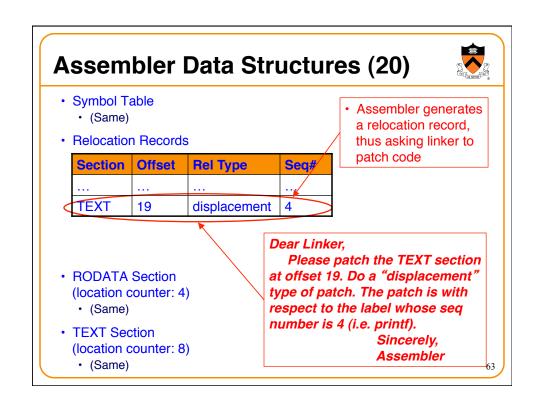


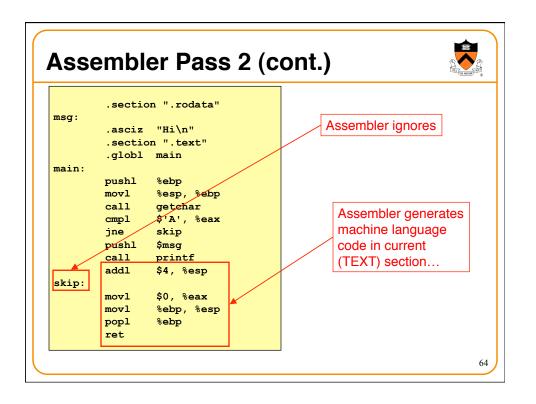






#### **Assembler Data Structures (19)** · Symbol Table Label **Section** Offset Local? Seq# **RODATA** 0 local 0 msg main **TEXT** 0 global 1 2 **TEXT** 26 skip local ? global 3 getchar ? ? printf global Relocation Records · Assembler adds printf • (Same) to Symbol Table RODATA Section (location counter: 4) Then... · (Same) • TEXT Section (location counter: 23) · (Same) 62





# **Assembler Data Structures (21)**



- Symbol Table, Relocation Records, RODATA Section
   (Same)
- TEXT Section (location counter: 31)

Offset	Contents	Explanation
23-25	83 C4 04	addl \$4,%esp 10000011 11 000 100 00000100 This is some "1" instruction that has a 1 byte immediate operand  The M field designates a register  This is an "add" instruction  The destination register is ESP  The immediate operand is 4
26-30	B8 00000000	movl \$0,%eax 10111000 00000000000000000000000000000

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# **Assembler Data Structures (22)**



- Symbol Table, Relocation Records, RODATA Section
   (Same)
- TEXT Section (location counter: 35)

Offset	Contents	Explanation
31-32	89 EC	movl %ebp,%esp 10001001 11 101 100 This is a "movl" instruction whose source operand is a register  The M field designates a register  The source register is EBP  The destination register is ESP
33	5D	popl %ebp 01011101 This is a "popl %ebp" instruction
34	С3	ret 11000011 This is a "ret" instruction

#### From Assembler to Linker



- · Assembler writes its data structures to .o file
- Linker:
  - · Reads .o file
  - Works in two phases: resolution and relocation

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#### **Linker Resolution**



- Resolution
  - · Linker resolves references
- For this program, linker:
  - · Notes that Symbol Table contains undefined labels
    - · getchar and printf
  - Fetches, from libc.a, machine language code defining getchar and printf
  - · Adds that code to TEXT section
    - (May add code to other sections too)
  - · Updates Symbol Table to note offsets of getchar and printf
  - · Adds column to Symbol Table to note addresses of all labels

#### **Linker Relocation**



- Relocation
  - Linker patches ("relocates") code
  - · Linker traverses relocation records, patching code as specified
- · For this program

Section	Offset	Rel Type	Seq#
TEXT	4	displacement	3
TEXT	14	absolute	0
TEXT	19	displacement	4

- Linker looks up offset of getchar
- Linker computes: [offset of getchar] – 8
- Linker places difference in TEXT section at offset 4

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# **Linker Relocation (cont.)**



• For this program

Section	Offset	Rel Type	Seq#
TEXT	4	displacement	3
TEXT	14	absolute	0
TEXT	19	displacement	4

- · Linker looks up addr of msg
- Linker places addr in TEXT section at offset 14

# Linker Relocation (cont.) • For this program

Section	Offset	Rel Type	Seq#
TEXT	4	displacement	3
TEXT	14	absolute	0
TEXT	19	displacement	4

- · Linker looks up offset of printf
- Linker computes:
   [offset of printf] 23
- Linker places difference in TEXT section at offset 19

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#### **Linker Finishes**

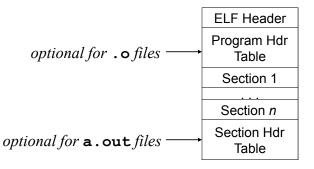


• Linker writes resulting TEXT, RODATA, DATA, BSS sections to executable binary file

#### **ELF: Executable and Linking Format**



- · Unix format of object and executable files
  - · Output by the assembler
  - · Input and output of linker



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#### **Conclusions**



- · Assembler: reads assembly language file
  - Pass 1: Generates Symbol Table
    - · Contains info about labels
  - Pass 2: Uses Symbol Table to generate code
    - TEXT, RODATA, DATA, BSS sections
    - · Relocation Records
  - Writes object file (ELF)
- Linker: reads object files
  - Resolution: Resolves references to make Symbol Table complete
  - Relocation: Uses Symbol Table and Relocation Records to patch code
  - Writes executable binary file (ELF)

#### **Appendix: Generating Machine Lang**



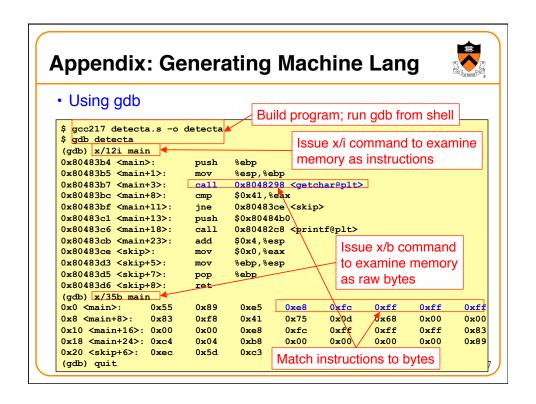
- Hint for Buffer Overrun assignment...
- Given an assembly language instruction, how can you find the machine language equivalent?
- Option 1: Consult IA-32 reference manuals
  - · See course Web pages for links to the manuals

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#### **Appendix: Generating Machine Lang**



- Option 2:
  - Compose an assembly language program that contains the given assembly language instruction
  - Then use gdb...



#### **Appendix: Generating Machine Lang**



- Option 3:
  - Compose an assembly language program that contains the given assembly language instruction
  - Then use objdump a special purpose tool...

